Algorithms \& Models of Computation
CS/ECE 374, Fall 2020
23.3

NP-Completeness of Hamiltonian Cycle

Algorithms \& Models of Computation

## CS/ECE 374, Fall 2020

23.3.1

Reduction from 3SAT to Hamiltonian Cycle:
Basic idea

## Directed Hamiltonian Cycle

Input Given a directed graph $\boldsymbol{G}=(\boldsymbol{V}, \boldsymbol{E})$ with $\boldsymbol{n}$ vertices
Goal Does $\boldsymbol{G}$ have a Hamiltonian cycle?

- A Hamiltonian cycle is a cycle in the graph that visits every vertex in G exactly once



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Is the following graph Hamiltonian?
(A) Yes.

(B) No.

## Directed Hamiltonian Cycle is NP-Complete

- Directed Hamiltonian Cycle is in NP: exercise
- Hardness: We will show 3SAT $\leq_{p}$ Directed Hamiltonian Cycle .


## Reduction construction

## From 3SAT to Hamiltonian cycle in directed graph

1. To show reduction, we next describe an algorithm:

- Input: 3SAT formula $\varphi$
- Output: A graph $\boldsymbol{G}_{\varphi}$.
- Running time is polynomial.
- Requirement: $\boldsymbol{\varphi}$ is satisfiable $\Longleftrightarrow \boldsymbol{G}_{\varphi}$ is Hamiltonian.

2. Given 3SAT formula $\varphi$ create a graph $G_{\varphi}$ such that

- $G_{\varphi}$ has a Hamiltonian cycle if and only if $\varphi$ is satisfiable
- $G_{\varphi}$ should be constructible from $\varphi$ by a polynomial time algorithm $\mathcal{A}$

3. Notation: $\varphi$ has $n$ variables $x_{1}, x_{2}, \ldots, x_{n}$ and $m$ clauses $C_{1}, C_{2}, \ldots, C_{m}$.

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## Encoding assignments

## Converting $\varphi$ to a graph

Given a formula with $n$ variables, we need a graph with $2^{n}$ different Hamiltonian paths, that can encode their assignments.


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## THE END

(for now)

