## Algorithms \& Models of Computation

CS/ECE 374, Fall 2020

## 18.5

Summary of shortest path algorithms

## Summary of results on shortest paths

| Single source |  |  |
| :--- | :--- | :--- |
| No negative edges | Dijkstra | $\boldsymbol{O}(\boldsymbol{n} \log \boldsymbol{n}+\boldsymbol{m})$ |
| Edge lengths can be negative | Bellman Ford | $\boldsymbol{O}(\boldsymbol{n m})$ |

## All Pairs Shortest Paths

| No negative edges | $\boldsymbol{n}^{*}$ Dijkstra | $\boldsymbol{O}\left(\boldsymbol{n}^{2} \log \boldsymbol{n}+\boldsymbol{n m}\right)$ |
| :--- | :--- | :--- |
| No negative cycles | $\boldsymbol{n}^{*}$ Bellman Ford | $\boldsymbol{O}\left(\boldsymbol{n}^{2} \boldsymbol{m}\right)=\boldsymbol{O}\left(\boldsymbol{n}^{4}\right)$ |
| No negative cycles $\left(^{*}\right)$ | BF $+\boldsymbol{n}^{*}$ Dijkstra | $\boldsymbol{O}\left(\boldsymbol{n m}+\boldsymbol{n}^{2} \log \boldsymbol{n}\right)$ |
| No negative cycles | Floyd-Warshall | $\boldsymbol{O}\left(\boldsymbol{n}^{3}\right)$ |
| Unweighted | Matrix multiplication | $\boldsymbol{O}\left(\boldsymbol{n}^{2.38}\right), \boldsymbol{O}\left(\boldsymbol{n}^{2.58}\right)$ |

## Summary of results on shortest paths

## More details

$\left(^{*}\right)$ : The algorithm for the case that there are no negative cycles, and doing all shortest paths, works by computing a potential function using Bellman-Ford and then doing Dijkstra. It is mentioned for the sake of completeness, but it outside the scope of the class.

## THE END

(for now)

