## 9.5 <br> Turing complete

## Equivalent to a program

## Definition

A system is Turing complete if one can simulate a Turing machine using it.
(1) Programming languages (yey!)
(2) $\mathrm{C}++$ templates system (boo)
(3) John Conway's game of life
© Many games (Minesweeper)
© Post's correspondence problem

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- Many games (Minesweeper).
- Post's correspondence problem.


## Post's correspondence problem

$S$ : set of domino tiles.
abb
bc
domino piece a string at the top and a string at the bottom.
Example:

$$
S=\left\{\begin{array}{c}
\hline b \\
\hline c a \\
\hline a b \\
\hline a b \\
\hline a \\
\hline a \\
\hline a b c \\
\hline a
\end{array}\right\} .
$$

## Matching dominos

$$
S=\left\{\begin{array}{c}
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\hline c a \\
\left.\left.\hline \frac{a}{a b}, \begin{array}{|c|}
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\hline a \\
\hline a b c \\
\hline c
\end{array}\right\} . . . \begin{array}{c}
a b c
\end{array} . . \begin{array}{c} 
\\
\hline
\end{array}\right] \\
\hline
\end{array}\right.
$$

match for $S$ : ordered list of dominos from $S$, such that top strings make same string as bottom strings. Example:

(1) Can use same domino more than once
(2) Do not have to use all pieces of $S$

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| $a$ | $b$ | $c a$ | $a$ | $a b c$ |
| :---: | :---: | :---: | :---: | :---: |
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## Post's Correspondence Problem

Post's Correspondence Problem (PCP) is deciding whether a set of dominos has a match or not. modified Post's Correspondence Problem (MPCP): PCP + a special tile. Matches for MPCP have to start with the special tile.

## Theorem

The MPCP problem is undecidable.

## THE END

## (for now)

