
Submission instructions as in previous [homeworks](#).

17 (100 PTS.) Shortest legible walk.

You are given a directed graph $G = (V, E)$ with n vertices and $m \geq n$ edges, and a pair of vertices s and t . In addition, every edge $e \in E$, has a *weight* $w(e) > 0$, and an associated character $c(e) \in \Sigma$. Thus, a walk $\pi = e_1 e_2 e_3 \dots e_k$ encodes the string $c(\pi) = c(e_1) c(e_2) \dots c(e_k)$.

You are also given a DFA $M = (Q, \Sigma, \delta, q_0, A)$, where $N = |Q|$, and $|\Sigma| = O(1)$.

Describe **formally** an algorithm, as fast as possible (the running time depends on n, m and N), that computes the shortest walk π from s to t , such that $c(\pi) \in L(M)$. Formally, let

$$\Pi = \{c(\pi) \in L(M) \mid \pi \text{ is a walk in } G \text{ from } s \text{ to } t \}.$$

The task is to compute $\arg \min_{\pi \in \Pi} w(\pi)$.

18 (100 PTS.) Paths and cycles.

Let G be a directed graph with n vertices and m edges, with positive numbers as weights on the edges.

- 18.A.** (50 PTS.) You are given in addition to G , two vertices $s, t \in V(G)$, and a parameter k . Describe an algorithm, as fast as possible, that computes the shortest path in G between s and t , where you are allowed to travel on k edges for free. What is the running time of your algorithm?
- 18.B.** (50 PTS.) The *heft* of a cycle C in G , is the minimum weight edge in C . Describe an algorithm, as fast as possible, that computes the maximum heft cycle in G .