

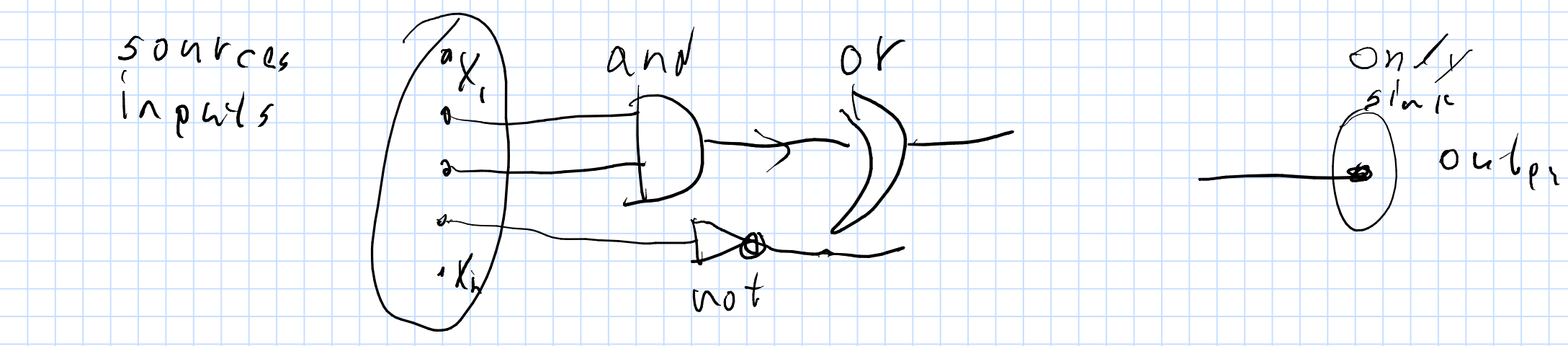
27: more NP problems

COOK Levin theorem

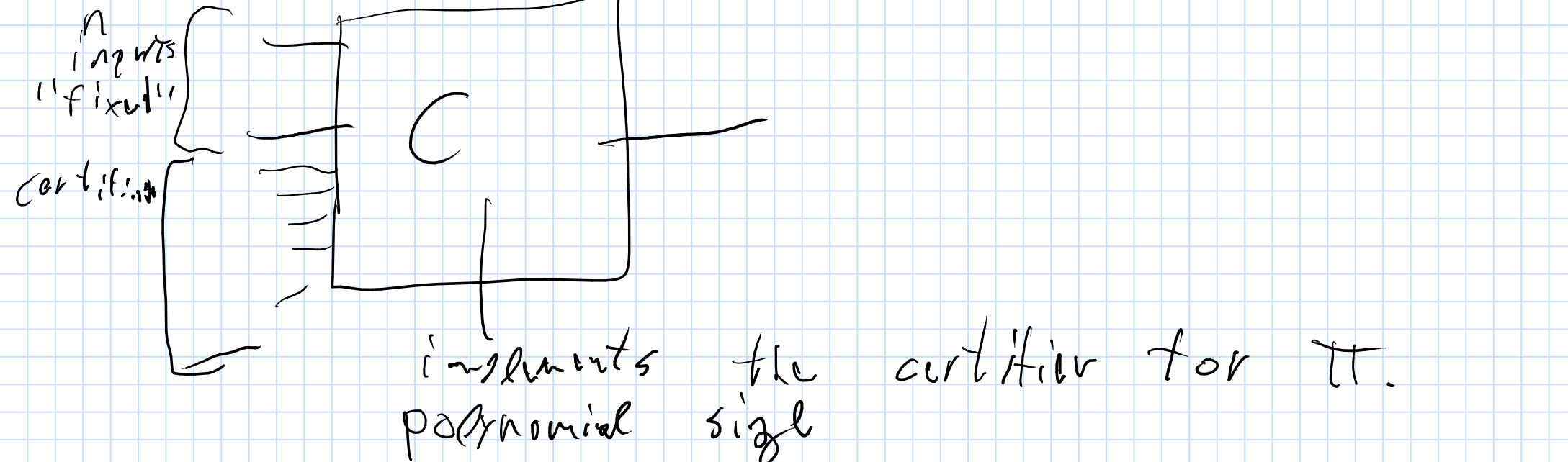
CSAT

Input: Boolean circuit with n inputs and one output.
 Q: Is there an input to the circuit that the output is 1/TRUE.

Circuit $\equiv$ DAG



A problem Π is in NP if $\exists$ a poly time verifier for it.
 $\Rightarrow$ $\text{verify}_{\Pi}(\text{instance}, \text{certificate})$ runs in poly time
 $\Rightarrow$ TM that implements the verifiers and runs in poly time.
 $\Rightarrow$ "compile TM to a circuit" (Cook-Levin Th)



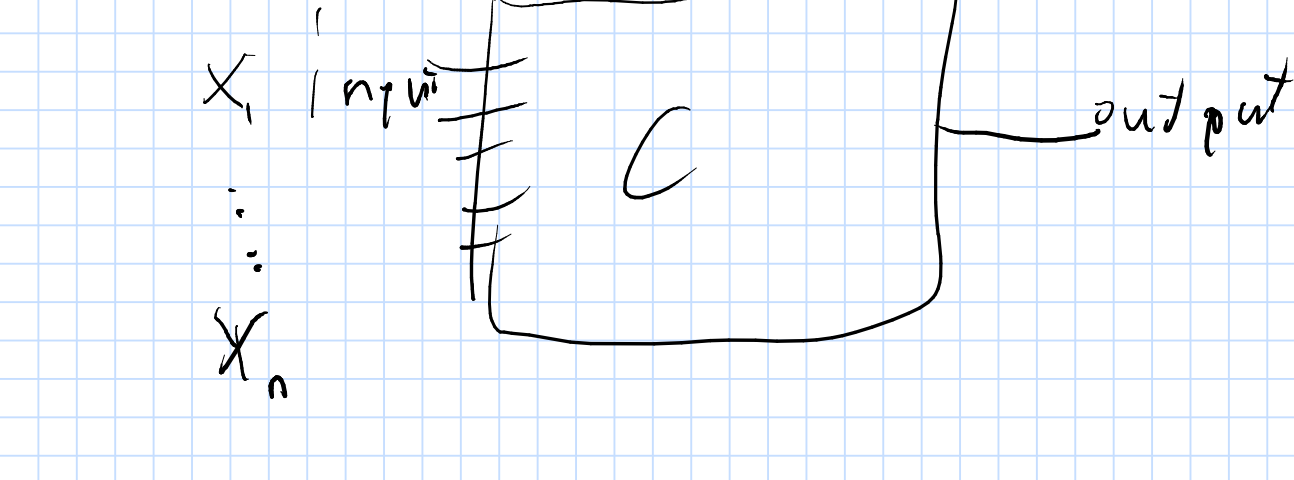
CSAT is NPC. (Cook-Levin)

claim SAT is NPC.

proof

- SAT is in NP. $\checkmark$
 - need to show reduction from CSAT to SAT.
 $\text{CSAT} \leq_p \text{SAT}$.

input is a circuit C

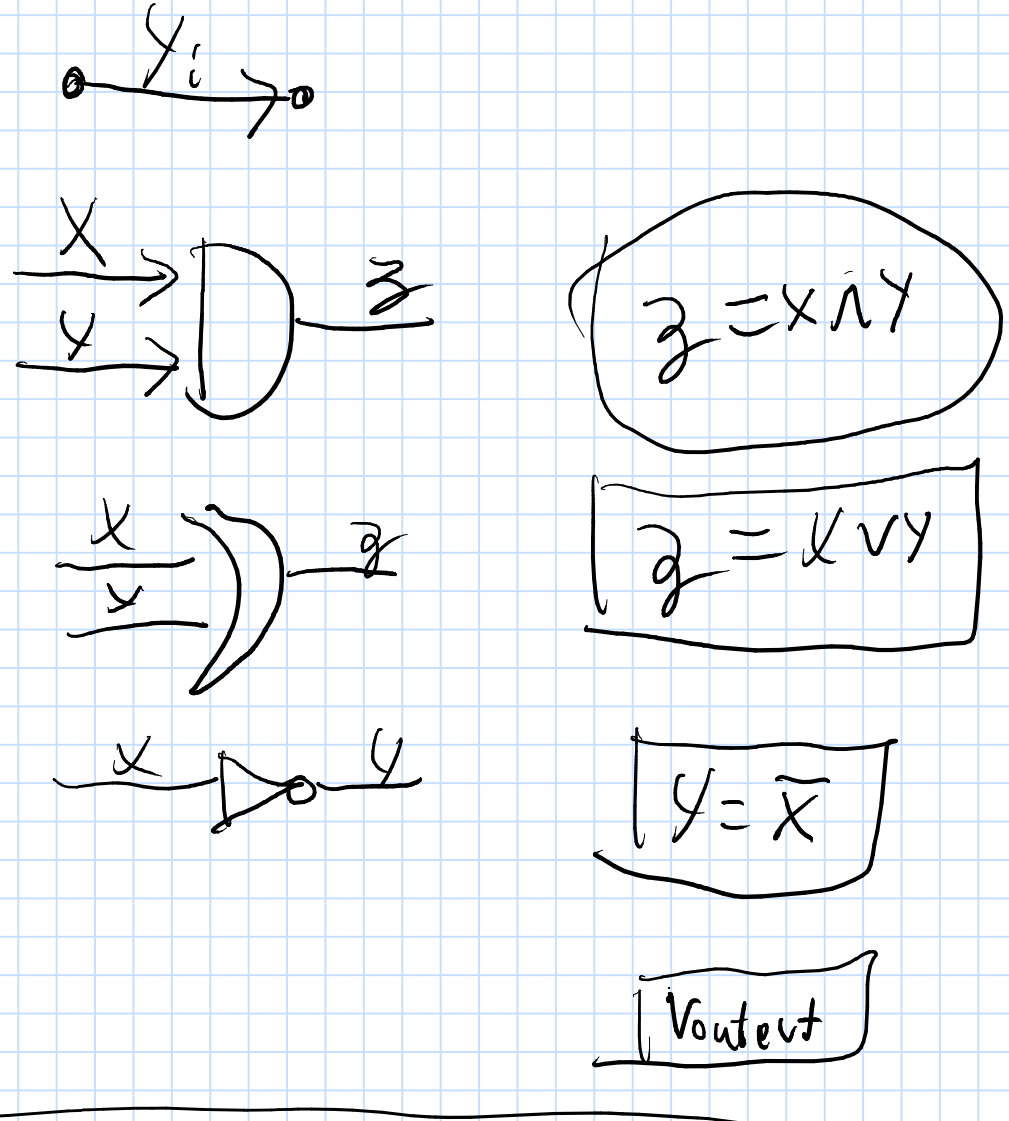


SAT
 Input: CNF formula
 Q: Is there a satisfying assignment

CNF formula

$$\bigwedge_{j=1}^m C_j$$

$$C_j = \bigvee_{v=1}^k x_v$$



x	y	z	$z = x \wedge y$	$x \vee y \vee \bar{z}$	$x \vee \bar{y} \vee \bar{z}$
0	0	0	0	0	1
0	0	1	0	0	1
0	1	0	0	1	0
0	1	1	0	1	1
1	0	0	0	1	0
1	0	1	0	1	1
1	1	0	0	1	1
1	1	1	1	1	1

$$\text{CSAT} \leq_p \text{SAT} \leq_p \exists \text{SAT} \leq_p \text{IS} \leq_p \text{CLIQUE} \leq_p \text{VC} \leq_p \text{HAMILTONIAN CYCLE/PATH}$$

Graph coloring $\langle G, k \rangle$

Q: Can the graph be colored with k colors.

3COLOR: $\langle G \rangle$

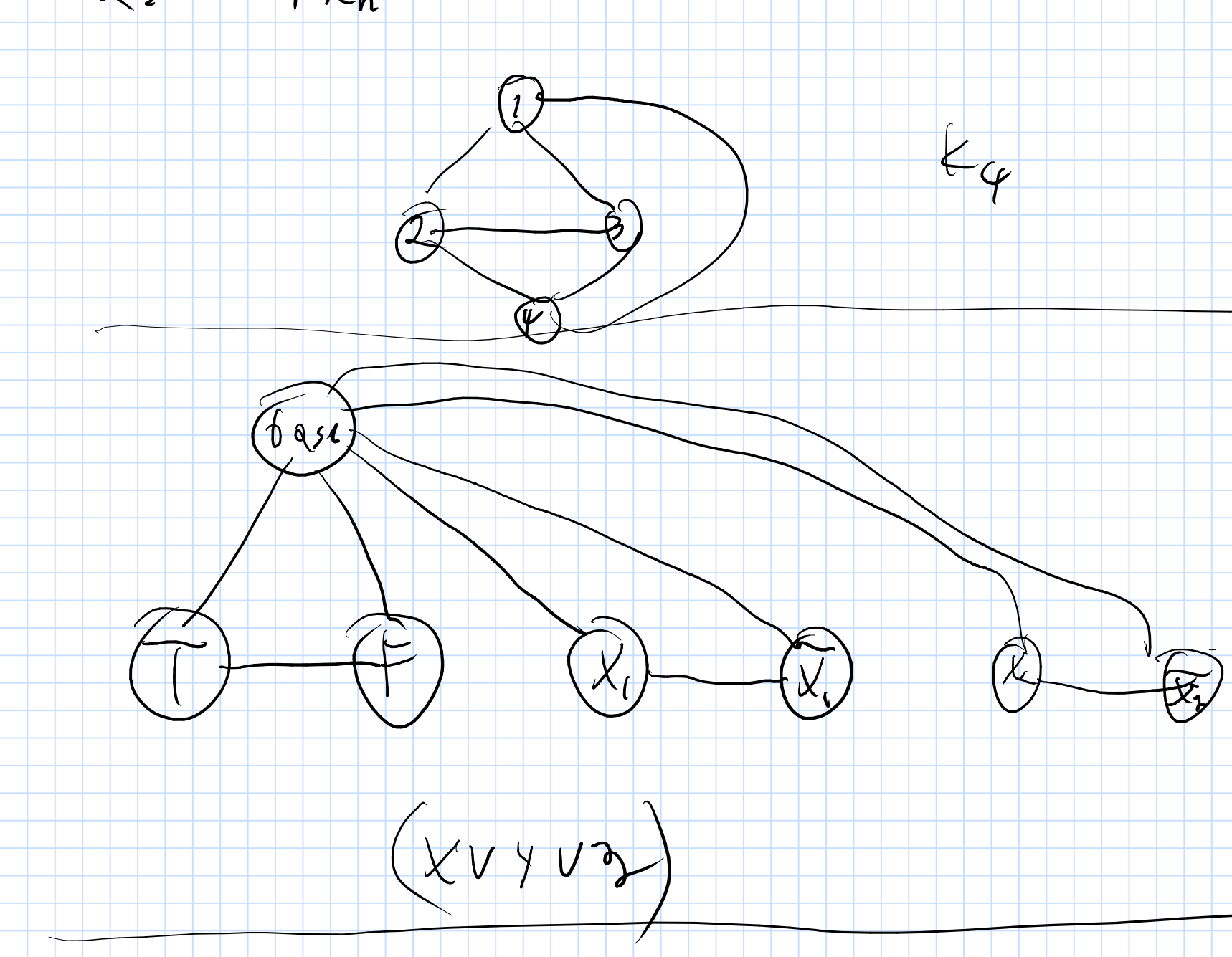
Q: Is G three colorable?

- 3COLOR is in NP. $\Rightarrow$ 3COLOR is NP-Complete.
 - $\exists \text{SAT} \leq_p \text{3COLOR}$.

$O(n^{1/4})$ colors

F is a 3SAT formula

$x_1, x_2, \dots, x_n$



subset sum

$i_1, i_2, \dots, i_n$ integer positive numbers
 t : target number

Q: Is there a subset of the numbers that add up to t ?

subset sum is NPC

Exercise

If all the n numbers in subset sum are at most $O(n^3)$ then the problem can be solved in poly time.

PARTITION

$\{i_1, i_2, \dots, i_n\} = I$

$$\exists S \subseteq I \quad \sum_{s \in S} s = \sum_{s \in I \setminus S} s$$