

5 (100 PTS.) String flip.

Let $\Sigma = \{0, 1\}$, and let $L \subseteq \Sigma^*$ be a regular language. For a string $s = s_1 \dots s_n \in \Sigma^*$, let $s^R = s_n s_{n-1} \dots s_1$ be the *reverse* of s . Consider the following language

$$L_8 = \{xy^Rz \mid x, y, z \in \Sigma^*, |y| \leq 8, \text{ and } xyz \in L\}.$$

Thus, if $0101000011110101 \in L$, then $0101111100000101 \in L_8$ as is $0101001001110101 \in L_8$. Prove that L_8 is a regular language.

To this end, you are given a DFA M for L – provide an NFA N for L_8 . Describe formally how you construct N from M , and argue why your construction is correct. A formal proof that your construction works is not required.

Hints: (A) Your NFA should use its ability to guess things, and remember constant amount of information (how?). (B) To build up to the solution consider special cases, and solve them first, such as: (i) $x = \varepsilon$, (ii) $z = \varepsilon$, and (iii) $|y| = 2$.

6 (100 PTS.) Highly irregular.

For each of the following languages prove that they are not regular using fooling sets. Here $\Sigma = \{0, 1\}$.

- 6.A.** (30 PTS.) For a string $w = w_1 w_2 \dots w_k$, let $\text{odd}(w) = w_1 w_3 w_5 \dots$ be the string formed by the odd characters of w . Consider the language $L_A = \{w \in (0 + 1)^* \mid \text{odd}(w) \text{ is a palindrome}\}$.
- 6.B.** (30 PTS.) $L_B = \{w \in \Sigma^* \mid 10^n 10^n 1 \text{ is a substring of } w, \text{ where } n \text{ is an integer}\}$.
- 6.C.** (40 PTS.) $L_C = \{0^i 1^j \mid i + j = k^2, \text{ where } k \text{ is an integer}\}$.