



S. SARMATA RUM.

AMAXOBILIS ARMA T

Tricornuo. Monte aureo. Margum. Viminatio. Lederata. Punicum. Municipio. Louis pago. Gabuleo.

Bistue. Houa. Stanecli. Argentaria. Sallunto. Halata. Berfumno. Sinna. Scio BR.

adpicaria. Creuent. Uyratio. Diftum. Herabum. Baletum. Vhintum. Veretum.

Castra. Minerue. ydrunte. Balentum. Lupia. Manduris. Herabum. Baletum. Vhintum. Veretum.

Scammum. Tarento. Melochoro. Grater. Heraclesia. Sennum. Turs. Pelelia. Frontona. Lacennium.

Herulos. Hierampio. Capraia. Tanno. Vibona. Balentia. Tauriana. drude.

Clampeta. Tanno. Vibona. Balentia. Tauriana. drude. Leucopetra.

Leucopetra. Leucopetra. Leucopetra. Leucopetra. Leucopetra.

Leucopetra. Leucopetra. Leucopetra. Leucopetra. Leucopetra.

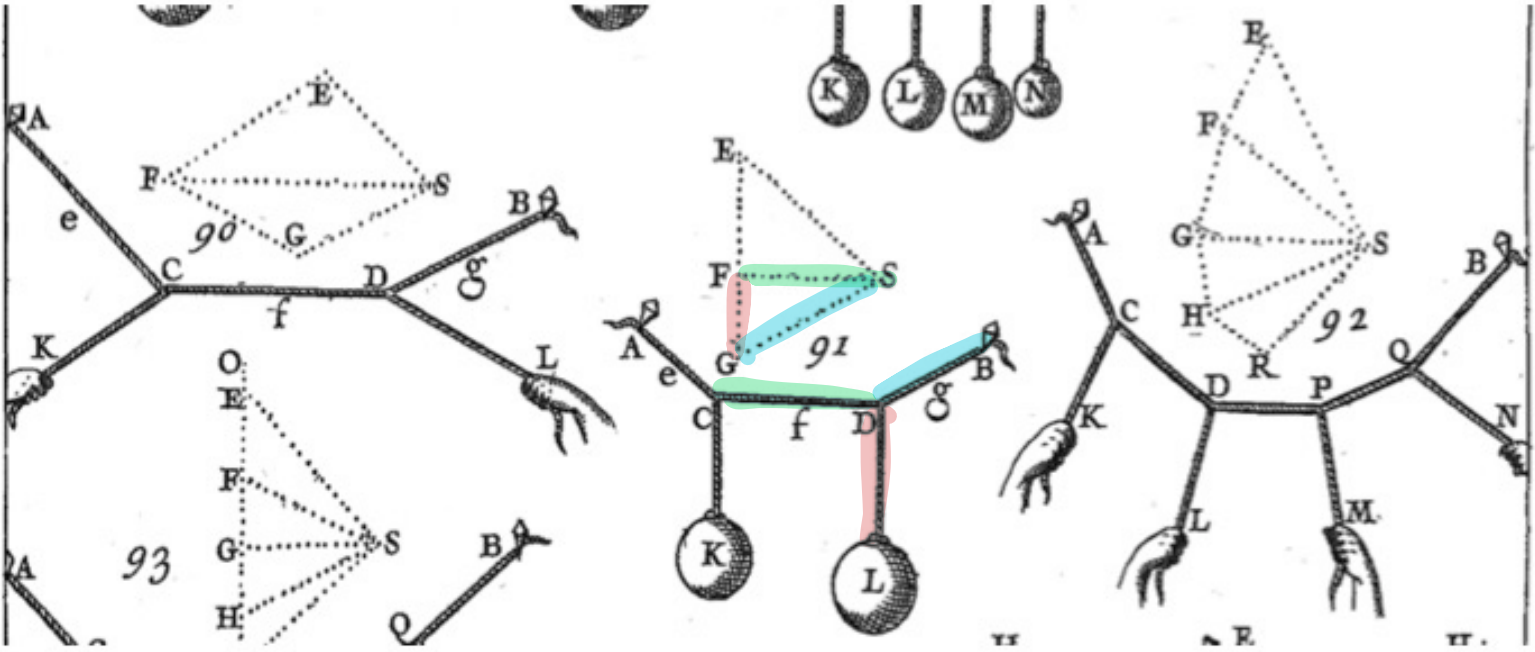
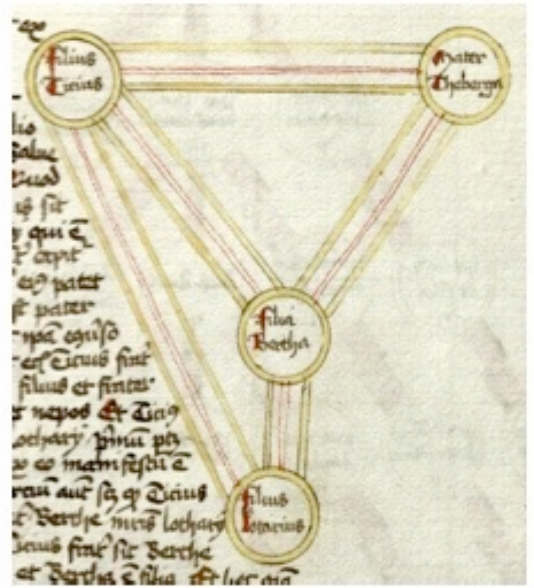
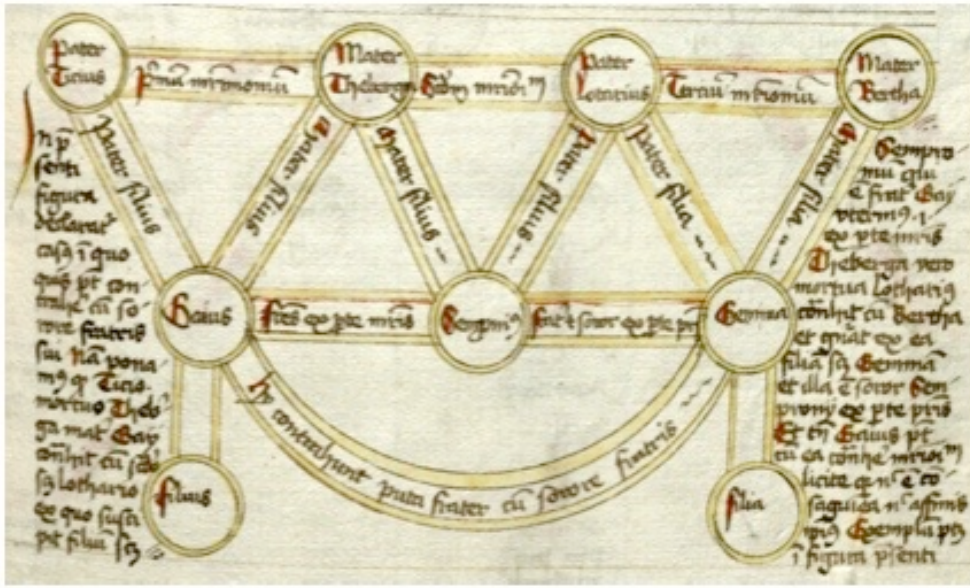
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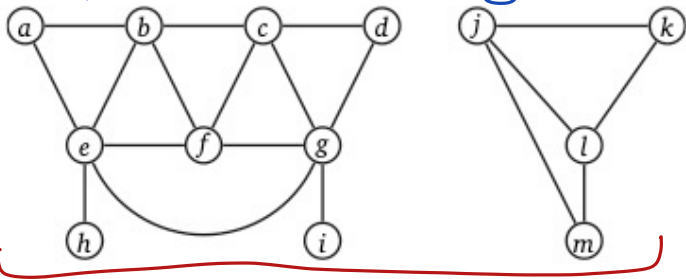
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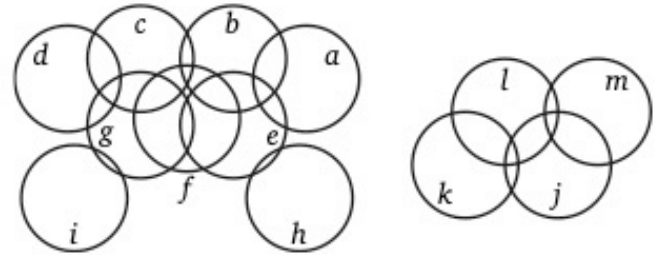
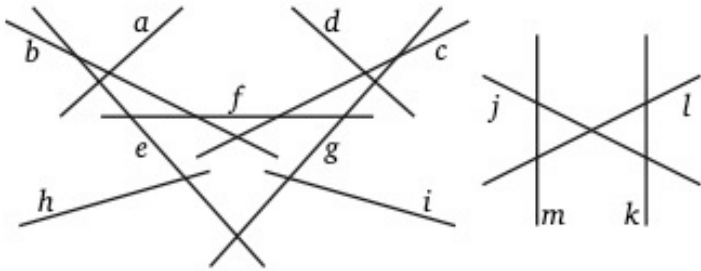
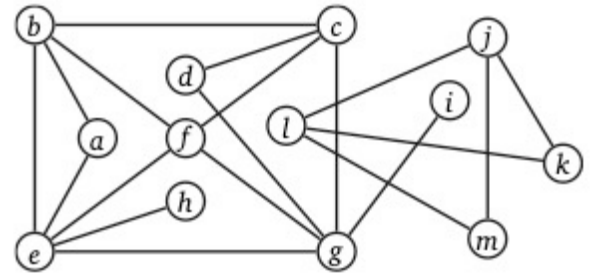
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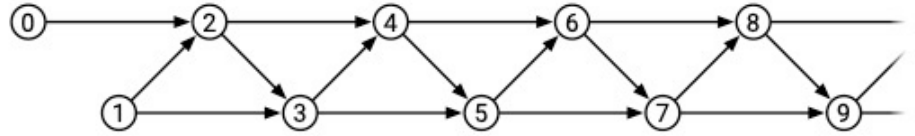
planar embeddings



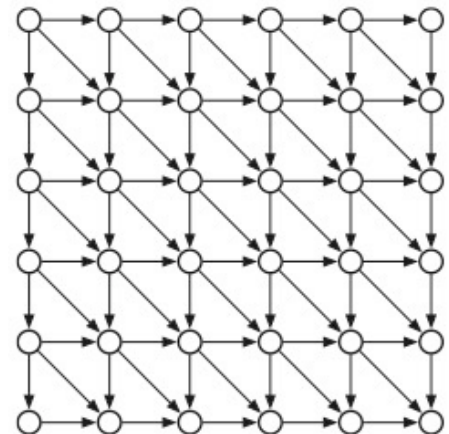
13 vertices
edges



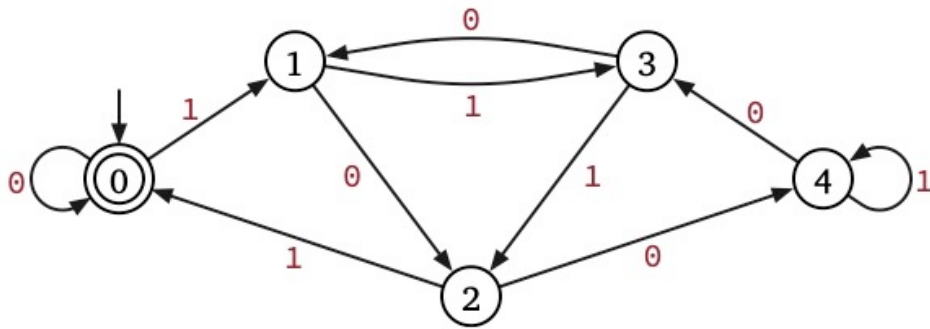
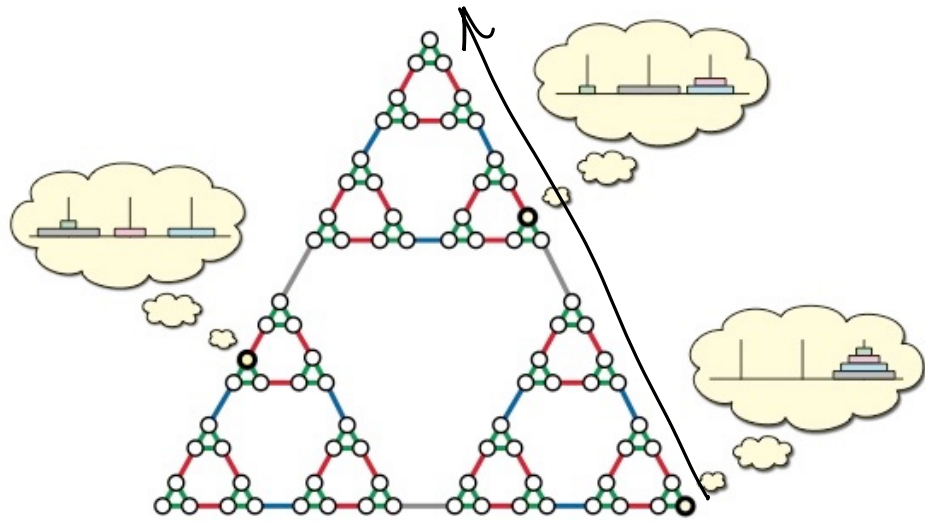
$$F_n = \begin{cases} 0 & \text{if } n = 0, \\ 1 & \text{if } n = 1, \\ F_{n-1} + F_{n-2} & \text{otherwise,} \end{cases}$$



$$Edit(i, j) = \begin{cases} i & \text{if } j = 0 \\ j & \text{if } i = 0 \\ \min \begin{cases} Edit(i-1, j) + 1 \\ Edit(i, j-1) + 1 \\ Edit(i-1, j-1) + [A[i] \neq B[j]] \end{cases} & \text{otherwise} \end{cases}$$



directed acyclic graphs



NFA \rightarrow DFA conversion
incremental subset

Vertices?

Edges?

Problem?

Algo?

Time?

Graph = (V, E)

V = vertices
ANY nonempty
finite set

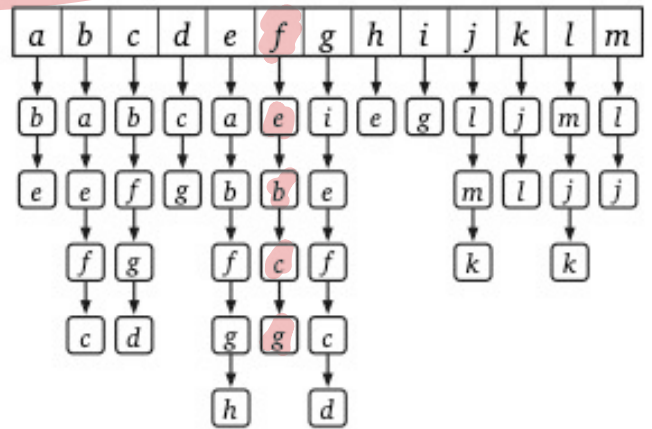
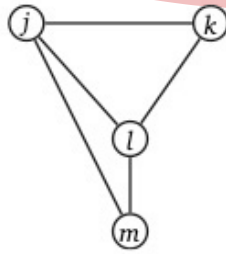
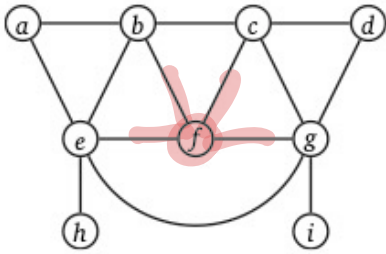
E = edges
= pairs of vertices

$\{u, v\} = uv$ undirected

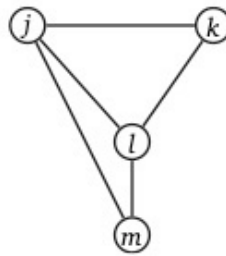
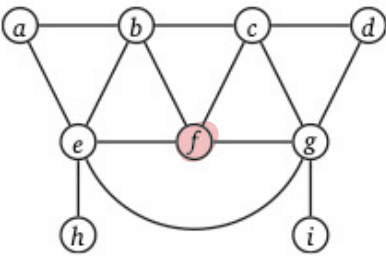
$(u, v) = u \rightarrow v$ directed

DATA STRUCTURES

"adjacency list"



adjacency matrix



	a	b	c	d	e	f	g	h	i	j	k	l	m
a	0	1	0	0	1	0	0	0	0	0	0	0	0
b	1	0	1	0	1	1	0	0	0	0	0	0	0
c	0	1	0	1	0	1	1	0	0	0	0	0	0
d	0	0	1	0	0	0	1	0	0	0	0	0	0
e	1	1	0	0	0	1	1	1	0	0	0	0	0
f	0	1	1	0	1	0	1	0	0	0	0	0	0
g	0	0	1	1	1	1	0	0	1	0	0	0	0
h	0	0	0	0	1	0	0	0	0	0	0	0	0
i	0	0	0	0	0	0	1	0	0	0	0	0	0
j	0	0	0	0	0	0	0	0	0	0	1	1	1
k	0	0	0	0	0	0	0	0	0	1	0	1	0
l	0	0	0	0	0	0	0	0	0	1	1	0	1
m	0	0	0	0	0	0	0	0	0	1	0	1	0

Reachability

v can reach $w \iff$ there is a path from v to w .

RECURSIVEDFS(v):

```
if  $v$  is unmarked
  mark  $v$ 
  for each edge  $vw$ 
    RECURSIVEDFS( $w$ )
```

ITERATIVEDFS(s):

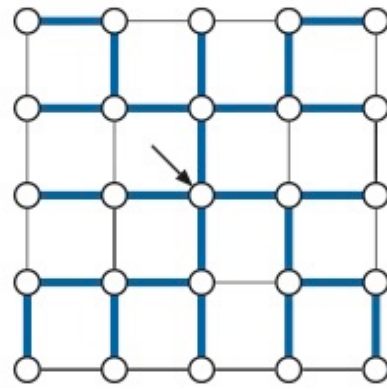
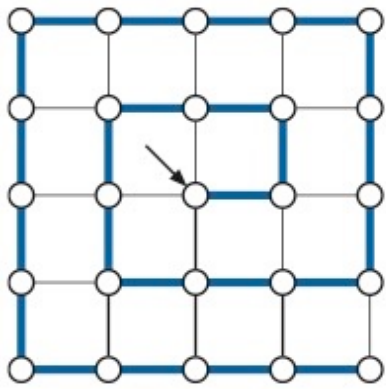
```
PUSH( $s$ )
while the stack is not empty
   $v \leftarrow$  POP
  if  $v$  is unmarked
    mark  $v$ 
    for each edge  $vw$ 
      PUSH( $w$ )
```

WHATEVERFIRSTSEARCH(s):

```
put  $s$  into the bag
while the bag is not empty
  take  $v$  from the bag
  if  $v$  is unmarked
    mark  $v$ 
    for each edge  $vw$ 
      put  $w$  into the bag
```

WHATEVERFIRSTSEARCH(s):

```
put  $(\emptyset, s)$  in bag
while the bag is not empty
  take  $(p, v)$  from the bag (*)
  if  $v$  is unmarked
    mark  $v$ 
     $parent(v) \leftarrow p$ 
    for each edge  $vw$  (†)
      put  $(v, w)$  into the bag (**)
```



WFSALL(G):

for all vertices v
unmark v

for all vertices v
if v is unmarked

WHATEVERFIRSTSEARCH(v)

COUNTCOMPONENTS(G):

count $\leftarrow 0$

for all vertices v
unmark v

for all vertices v
if v is unmarked

count \leftarrow *count* + 1

WHATEVERFIRSTSEARCH(v)

return count

COUNTANDLABEL(G):

count $\leftarrow 0$

for all vertices v
unmark v

for all vertices v
if v is unmarked

count \leftarrow *count* + 1

LABELONE($v, count$)

return count

⟨⟨Label one component⟩⟩

LABELONE($v, count$):

while the bag is not empty

take v from the bag

if v is unmarked

mark v

comp(v) \leftarrow *count*

for each edge vw

put w into the bag