

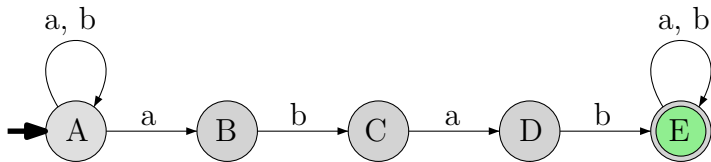
7.3

Converting regular languages into CFL

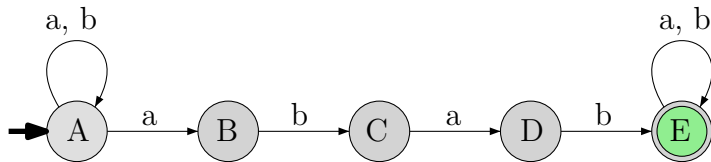
Converting regular languages into CFL

$M = (Q, \Sigma, \delta, s, A)$: DFA for regular language L .

$$G = \left(\underbrace{Q}_{\text{Variables}}, \underbrace{\Sigma}_{\text{Terminals}}, \underbrace{\left\{ q \rightarrow a\delta(q, a) \mid q \in Q, a \in \Sigma \right\} \cup \left\{ q \rightarrow \varepsilon \mid q \in A \right\}}_{\text{Productions}}, \underbrace{s}_{\text{Start var}} \right)$$



Conversion continued...



$$G = \left(\{A, B, C, D, E\}, \{a, b\}, \left\{ \begin{array}{l} A \rightarrow aA, A \rightarrow bA, A \rightarrow aB, \\ B \rightarrow bC, \\ C \rightarrow aD, \\ D \rightarrow bE, \\ E \rightarrow aE, E \rightarrow bE, E \rightarrow \epsilon \end{array} \right\}, A \right)$$

The result...

Lemma

For an regular language L , there is a context-free grammar (CFG) that generates it.

THE END

...

(for now)