## $\frac{\mathrm{Quiz}\ 5}{\mathrm{CS}\ 373}$ : Theory of Computation

Date: November 16, 2010. Lecture Section AL1. Time limit: 15 minutes.

| Name       |           |           |           |          |          |
|------------|-----------|-----------|-----------|----------|----------|
| netid      |           |           |           |          |          |
| Discussion | Tu 2-2:50 | Tu 3-3:50 | Tu 4-4:50 | W 4-4:50 | W 5-5:50 |

Pick the correct alternative from among the choices (A), (B), and (C) provided for each question below. Each question is worth 1 point.

$$\begin{array}{ll} S & \rightarrow AX \mid YC \\ A & \rightarrow aA \mid \epsilon \\ B & \rightarrow bB \mid \epsilon \\ C & \rightarrow cC \mid \epsilon \\ X & \rightarrow bXc \mid bB \mid cC \\ Y & \rightarrow aYb \mid aA \mid bB \end{array}$$

Figure 1: Context Free Grammar  $G_* = (V = \{S, A, B, C, X, Y\}, \Sigma = \{a, b, c\}, R, S)$  for problems 1, 2, and 3. The rules for  $G_*$  are given above.

- 1. In grammar  $G_*$ , which of the following strings can be derived from S in zero or more steps?
  - (A) aaba
  - (B) aabbcc
  - (C) aaAbXc
- 2. The set of strings (over  $\Sigma$ ) derivable from A in  $G_*$  is
  - (A)  $L(a^*)$
  - (B) Strings with an even number of as
  - (C)  $\emptyset$  because A is not the start symbol
- 3. Which of the following is true about grammar  $G_*$ ?
  - (a)  $G_*$  is ambiguous because abbccc has two derivations from S
  - (b)  $G_*$  is ambiguous because abbccc has two parse trees
  - (c)  $G_*$  is not ambiguous because we can transform the grammar using the techniques discussed in class

$$\begin{array}{ll} \delta(q_0,[,\epsilon) &= \{(q_0,A)\} \\ \delta(q_0,\epsilon,\epsilon) &= \{(q_1,\epsilon)\} \\ \delta(q_1,],A) &= \{(q_1,\epsilon)\} \\ \delta(q,a,x) &= \emptyset \quad \text{otherwise} \end{array}$$

Figure 2: Pushdown automata  $P_* = (Q = \{q_0, q_1\}, \Sigma = \{[,]\}, \Gamma = \{A\}, q_0, F = \{q_1\}, \delta)$  for problems 4 and 5. The transition function  $\delta$  is given above.

- 4. Suppose the current instantaneous description of  $P_*$  is  $\langle q_1, AAAAA \rangle$  and the unread portion of the input is |||. The instantaneous description after one step is
  - (A) The machine crashes
  - (B)  $\langle q_1, AAAAAA \rangle$
  - (C)  $\langle q_1, AAAA \rangle$
- 5. The language recognized by PDA  $P_*$  is
  - (A)  $\{[n]^n \mid n \ge 0\}$
  - (B)  $\{[i]^j \mid i \ge j \ge 0\}$
  - (C)  $\{[i]^j \mid j \ge i \ge 0\}$
- 6. Suppose L is recognized by a linear bounded automata A. Then,
  - (A) A always halts on all inputs because L is decidable
  - (B) L maybe undecidable because A need not halt on all inputs
  - (C) L need not be a context-free language