## $\frac{\text{Quiz 1}}{\text{CS 373: Theory of Computation}}$

Date: September 23, 2010. Lecture Section AL2. Time limit: 15 minutes.

Name					
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Discussion	Tu 2-2:50	Tu 3-3:50	Tu 4-4:50	W 4-4:50	W 5-5:50

Pick the correct alternative from among the choices (A), (B), and (C) provided for each question below. Each question is worth 1 point.

- 1. Let  $D = (Q, \{0, 1\}, \delta, q_0, F)$  be an DFA such that  $L(D) = \{0, 1\}^*$ . Then,
  - (A) Every state must be a final state, i.e., F = Q.
  - (B) No state is a final state, i.e.  $F = \emptyset$ .
  - (C) Neither of the above
- 2. Consider  $r = a(ab^*a \cup b^*)^*$ . Which of the following is true about L(r)?
  - (A)  $a \in L(r)$
  - (B)  $aa \in L(r)$
  - (C) There is at least one b in every string belonging to L(r)
- 3. For  $n \ge 0$ , let  $L_n = \{a^i b^k \mid i \ge n, \ 0 < k < n\}$ .
  - (A)  $L_n$  is regular, independent of the value of n
  - (B)  $L_n$  is not regular, independent of the value of n
  - (C)  $L_n$  is regular only for small values of n
- 4. Let  $L_1$  be an infinite regular language. Let  $L_2$  be an infinite set such that  $L_1 \subseteq L_2$ .
  - (A)  $L_2$  is definitely regular because  $L_1$  is regular
  - (B)  $L_2$  is never regular because  $L_2$  is infinite
  - (C)  $L_2$  may or may not be regular

- 5. Consider  $L_1, L_2 \subseteq \Sigma^*$  such that  $L_1$  is a finite language and  $L_1 \cup L_2$  is regular.
  - (A)  $L_2$  is definitely regular
  - (B)  $L_2$  may not be regular
  - (C)  $L_2 = (L_1 \cup L_2) \setminus L_1$
- 6. Recall from homework 3, for a string w,  $w^R$  denotes the reverse of w. For  $L\subseteq \Sigma^*$ , recall that  $L^R=\{w^R\mid w\in L\}$ . Suppose  $L^R$  is not regular. Then,
  - (A) L is definitely regular
  - (B) L may or may not be regular
  - (C) L is definitely not regular