

Algorithms and Data Structures for Data Science

Binary Search Tree

CS 277
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Department of Computer Science

Learning Objectives

Review understanding of Binary Trees

Extend ADT to Binary Search Trees

Implement BST operations

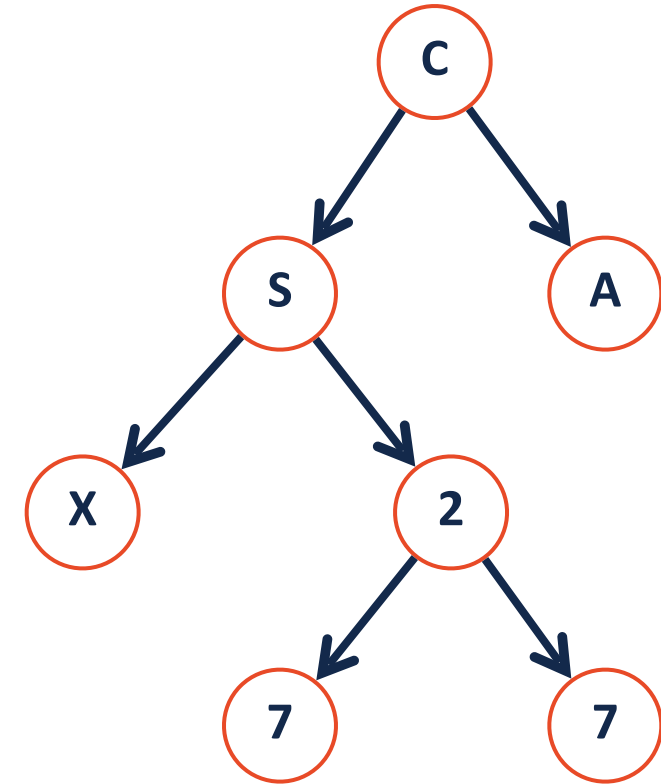
(Binary) Tree Recursion

A **binary tree** is a tree T such that:

$T = \text{None}$

or

$T = \text{treeNode}(\text{val}, T_L, T_R)$



```
1 class treeNode:
2     def __init__(self, val, left=None, right=None):
3         self.val = val
4         self.left = left
5         self.right = right
```

```
1 class binaryTree:
2     def __init__(self):
3         self.root = None
4
5
```

Tree ADT

Constructor: Build a new (empty) tree

Insert: Add an object into tree

Remove: Remove a specific object from tree

Traverse: Visit every node in tree (all objects)

Search: Find a specific object in the tree

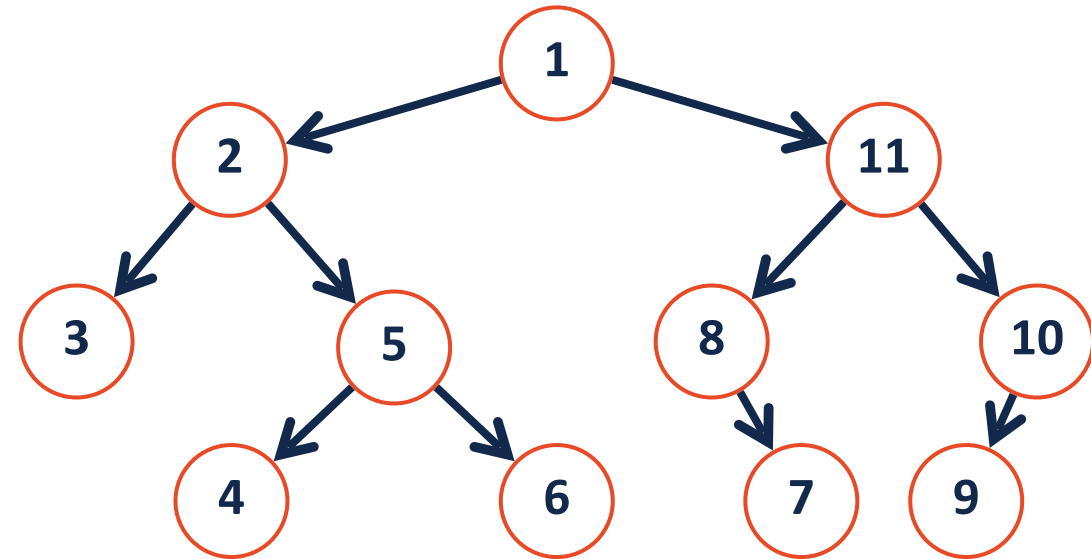
Binary Tree Traversal



Last class we implemented traversals using recursion, stacks, and queues.

What implementations led to a **depth first search traversal**?

Which lead to **breadth first search**?



Binary Tree Utility

This week we will deep dive into useful implementations of binary trees

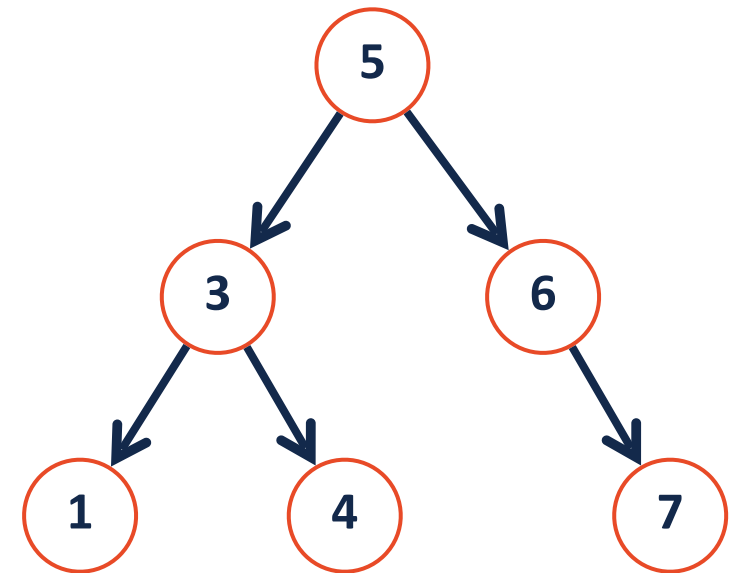
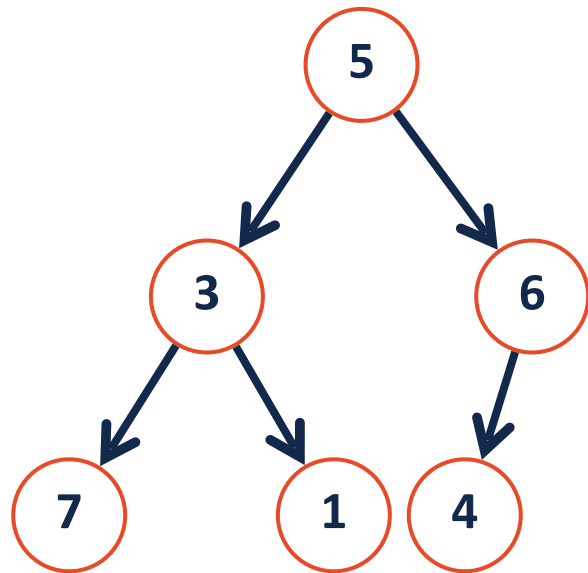
Binary Search Tree: An efficient implementation of a dictionary

Huffman Tree: A binary tree used to define an optimal text encoding

Improved search on a binary tree

5	3	6	7	1	4
---	---	---	---	---	---

1	3	4	5	6	7
---	---	---	---	---	---

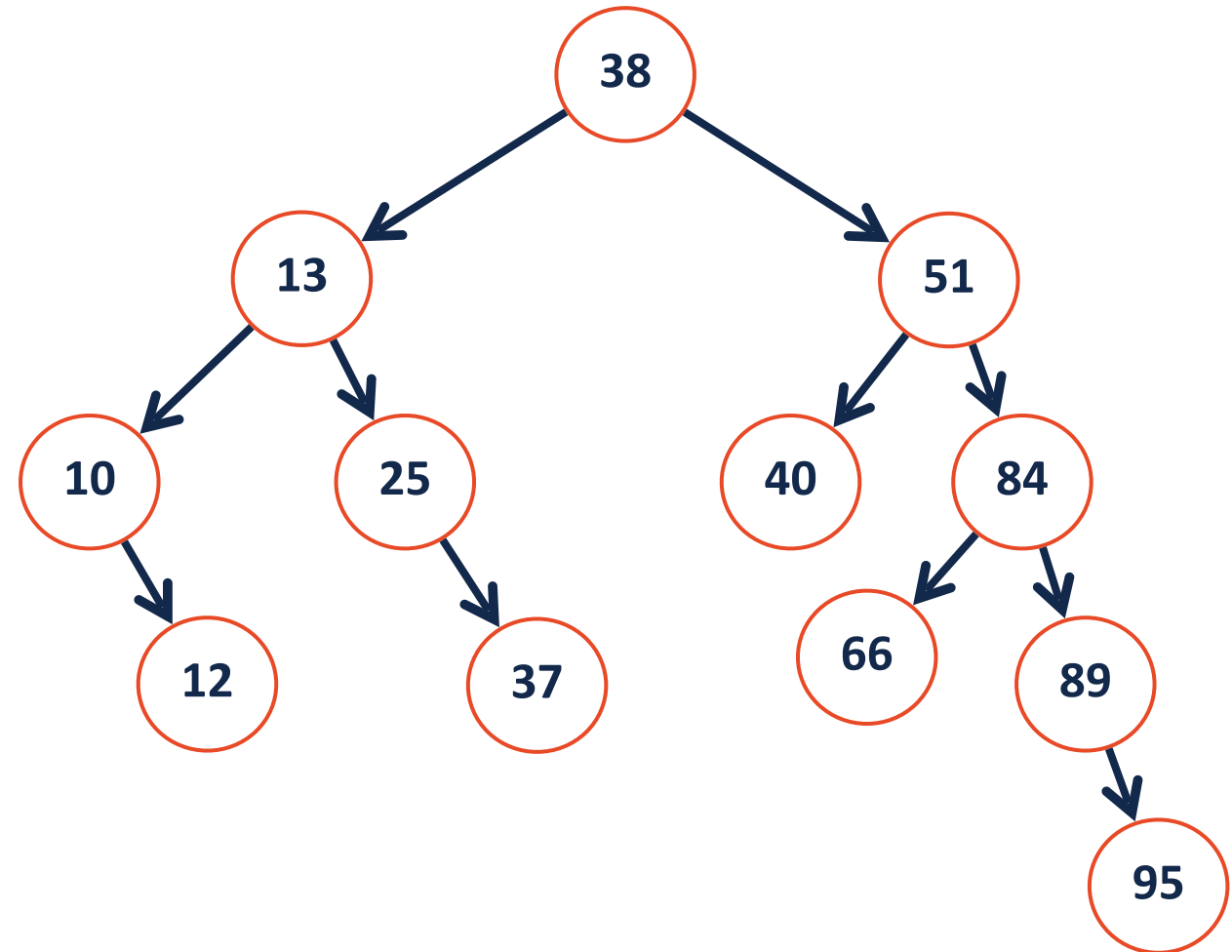


Binary Search Tree (BST)

A **BST** is a binary tree $T = \text{treeNode}(\text{val}, T_L, T_r)$ such that:

$\forall n \in T_L, n.\text{val} < T.\text{val}$

$\forall n \in T_R, n.\text{val} > T.\text{val}$



Dictionary ADT

Data is often organized into key/value pairs:

Word → Definition

Course Number → Lecture/Lab Schedule

Node → Edges

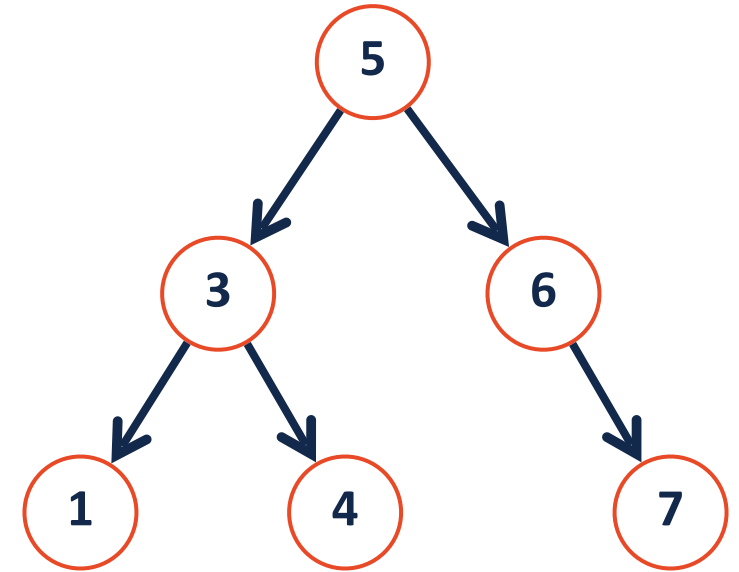
Flight Number → Arrival Information

URL → HTML Page

Average Image Color → File Location of Image

Dictionary as a Binary Search Tree

```
1 class bstNode:  
2     def __init__(self, key, val, left=None, right=None):  
3         self.key = key  
4         self.val = val  
5         self.left = left  
6         self.right = right
```



Key	5	3	6	7	1	4
Value	A	B	C	D	E	F

Binary **Search** Tree ADT — **what changed?**



Constructor: Build a new (empty) tree

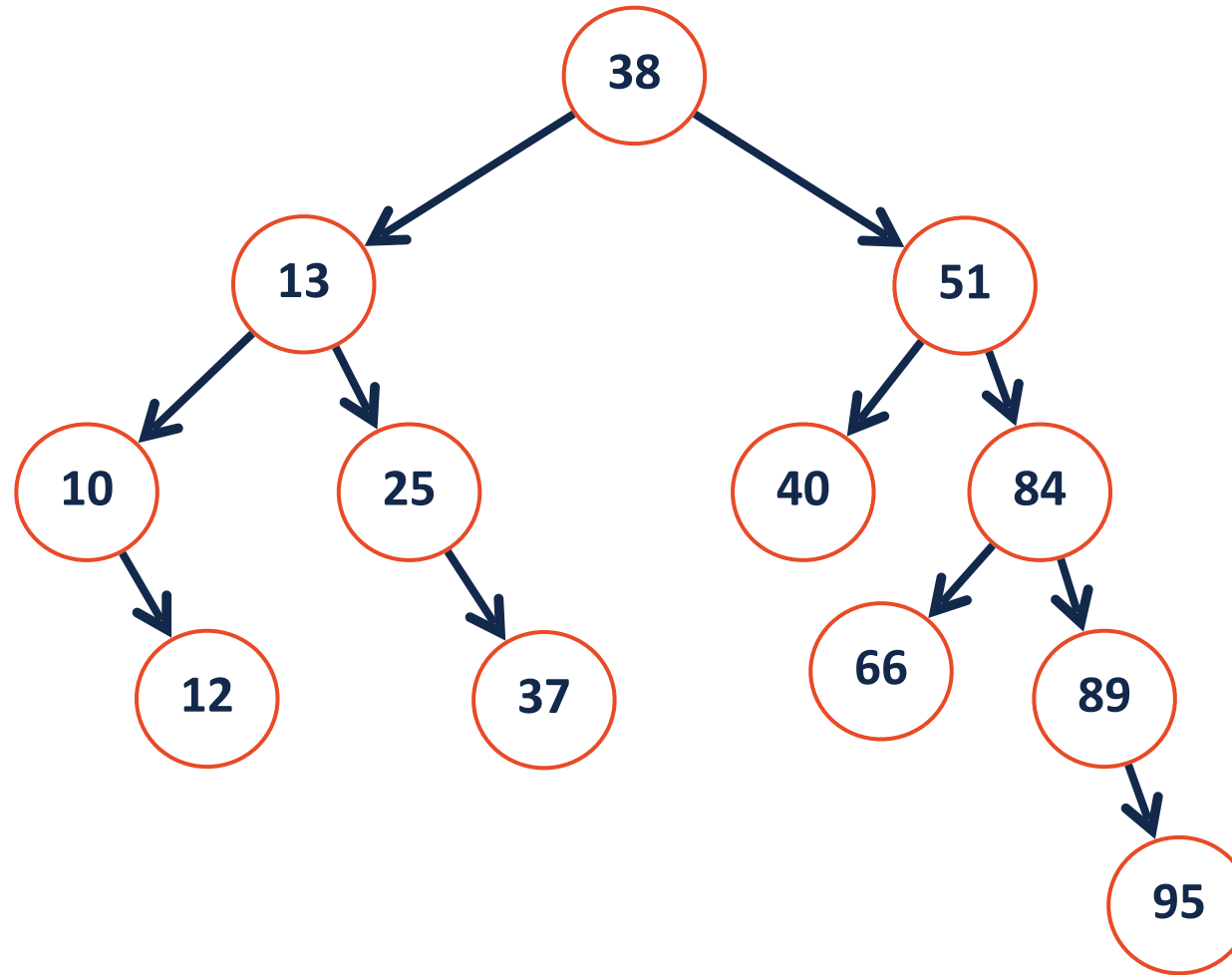
Insert: Add an object into tree

Remove: Remove a specific object from tree

Traverse: Visit every node in tree (all objects)

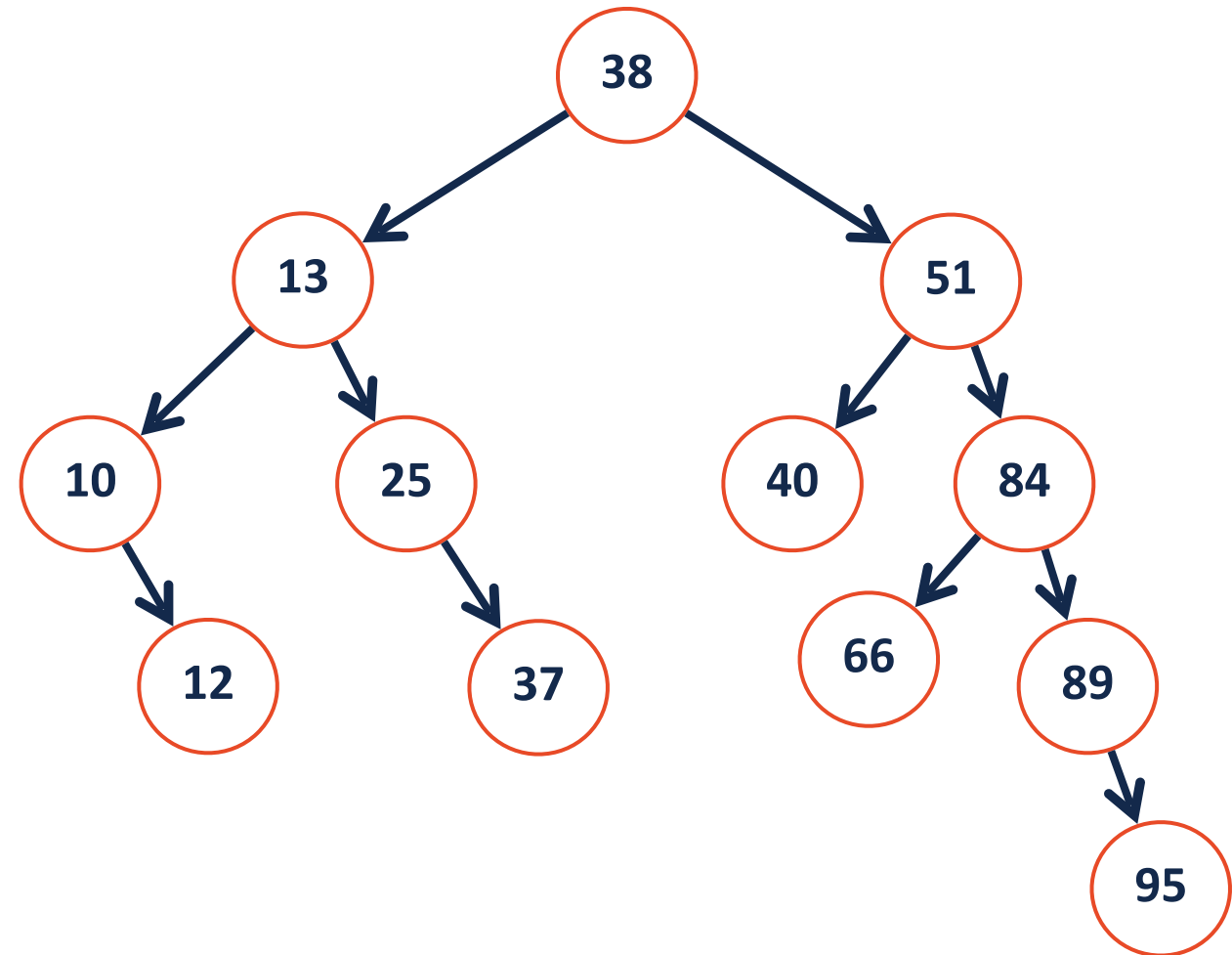
Find
Search: Find a specific **key** in the tree, **return value**

BST In-Order Traversal



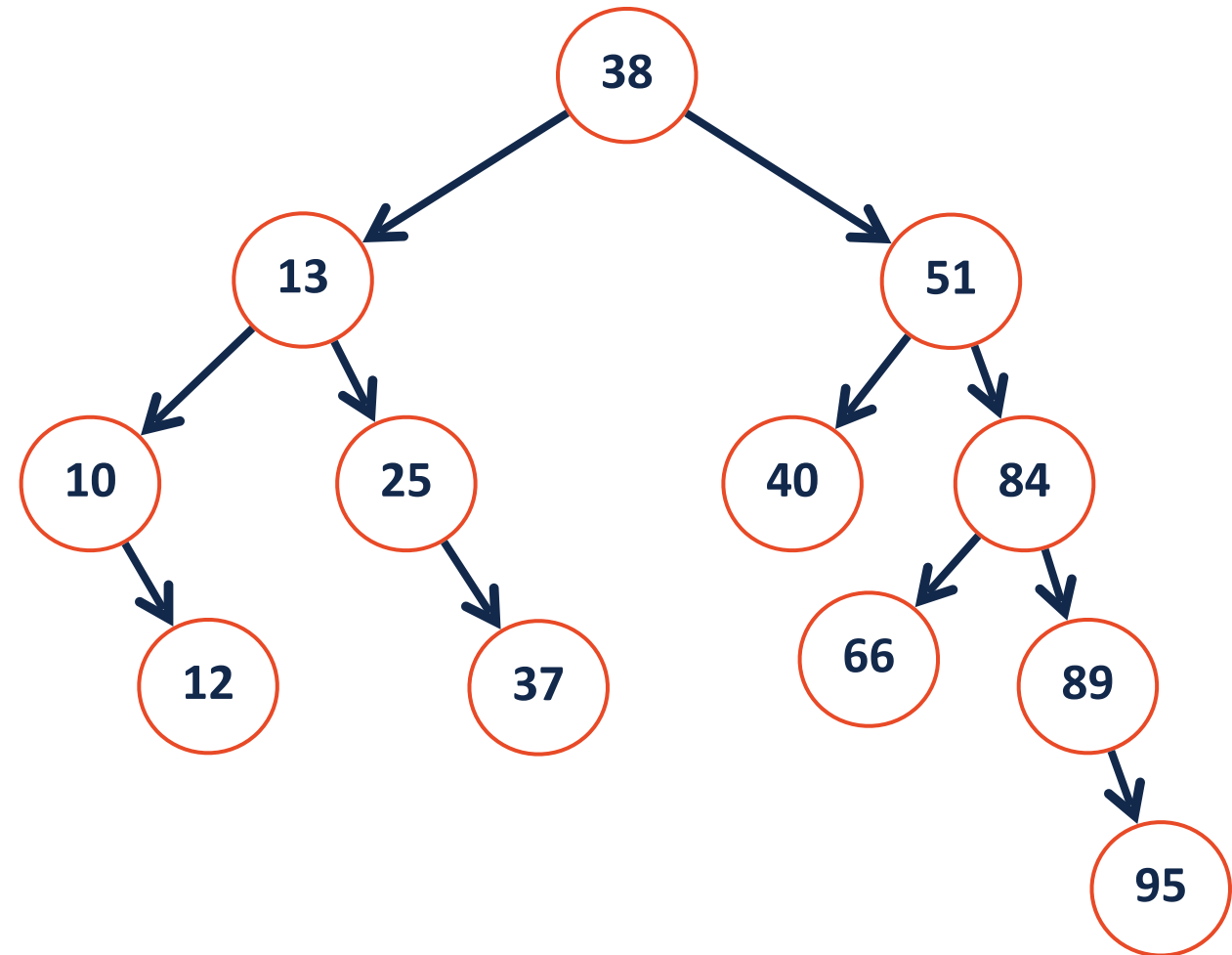
BST Find

find(66)



BST Find

find(9)

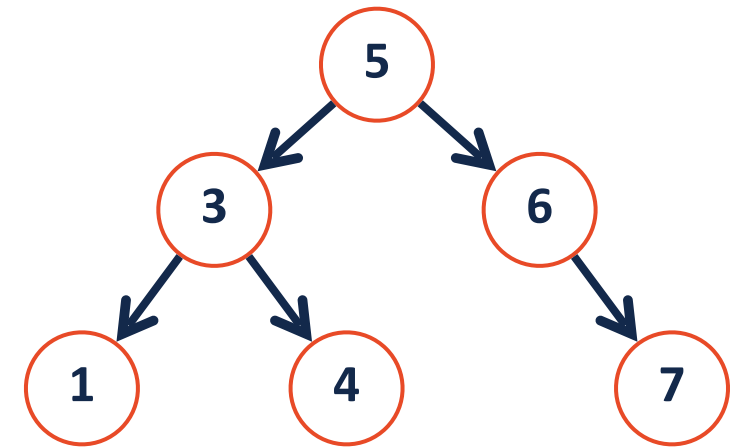


BST Find

Base Case:

Recursive Step:

Combining:



BST Find



A recursive function based around value of root:

Base Case: If root is None, return root

If root.key is query, return root

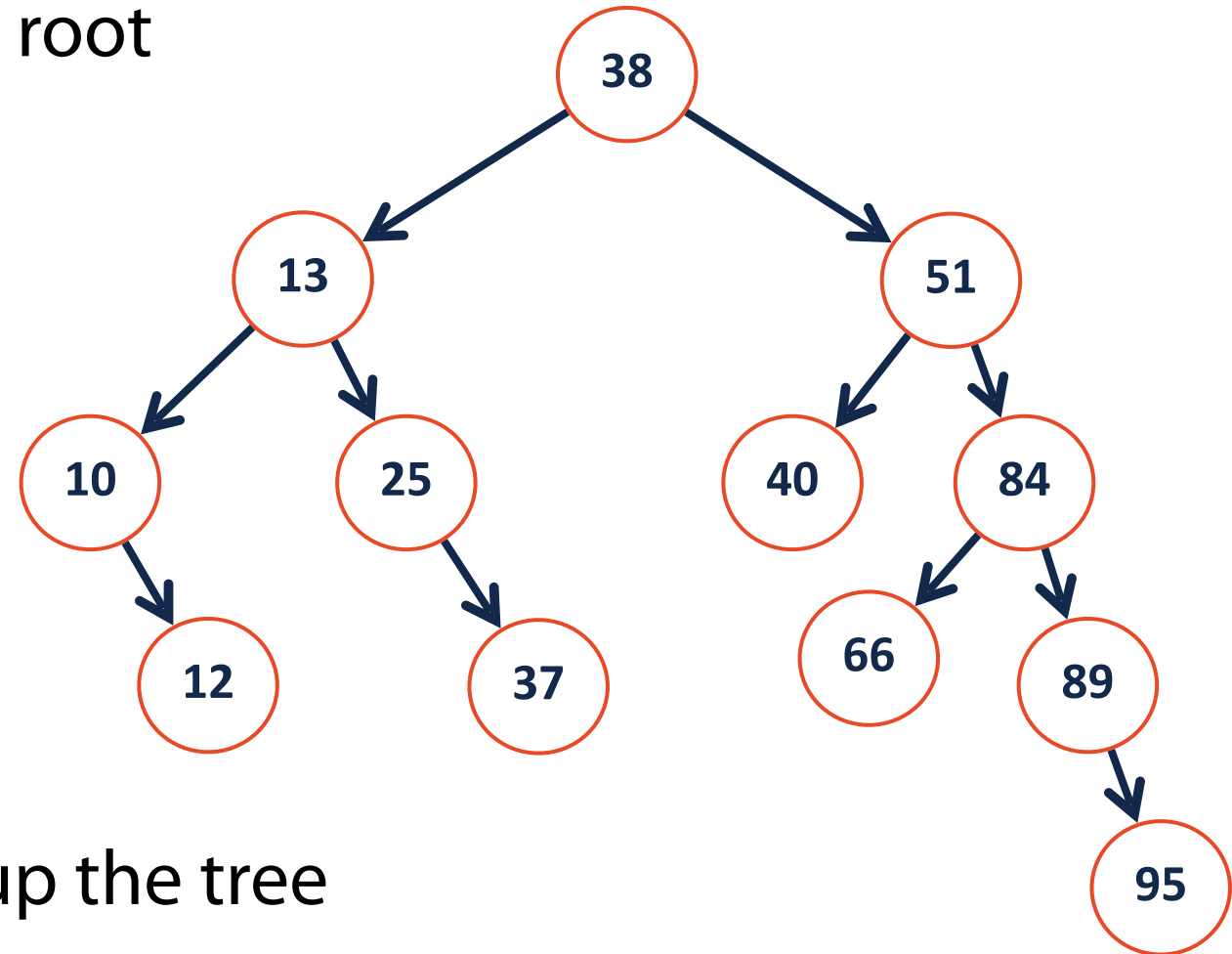
Recursion:

root.key ____ query, recurse right

root.key ____ query, recurse left

Combining:

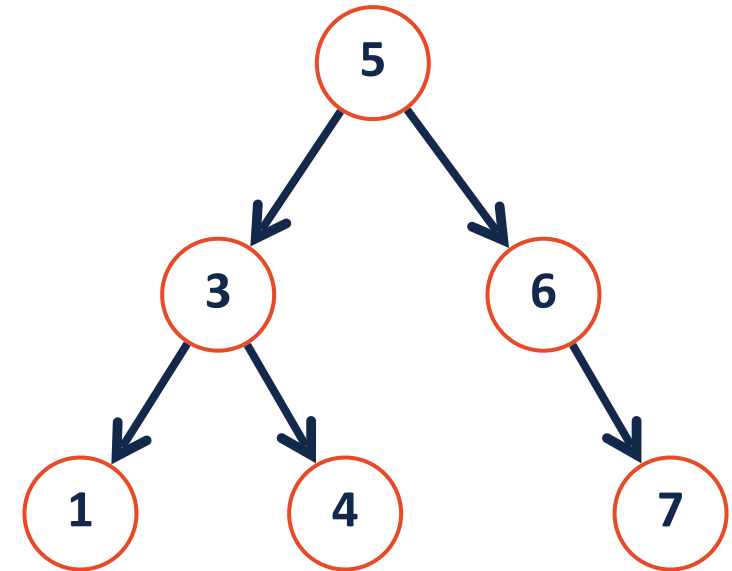
Return the recursive value back up the tree



BST Find



```
1 #inside class bst
2 def find(self, key):
3
4
5
6
7
8
9 def find_helper(self, node, key):
10
11
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19
20
21
22
23
```

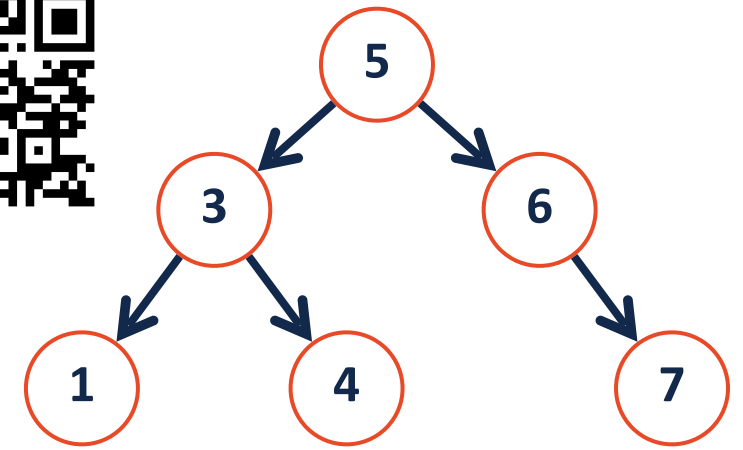
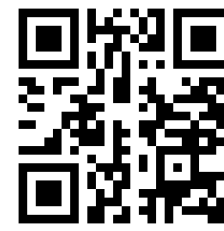


BST Insert

Base Case:

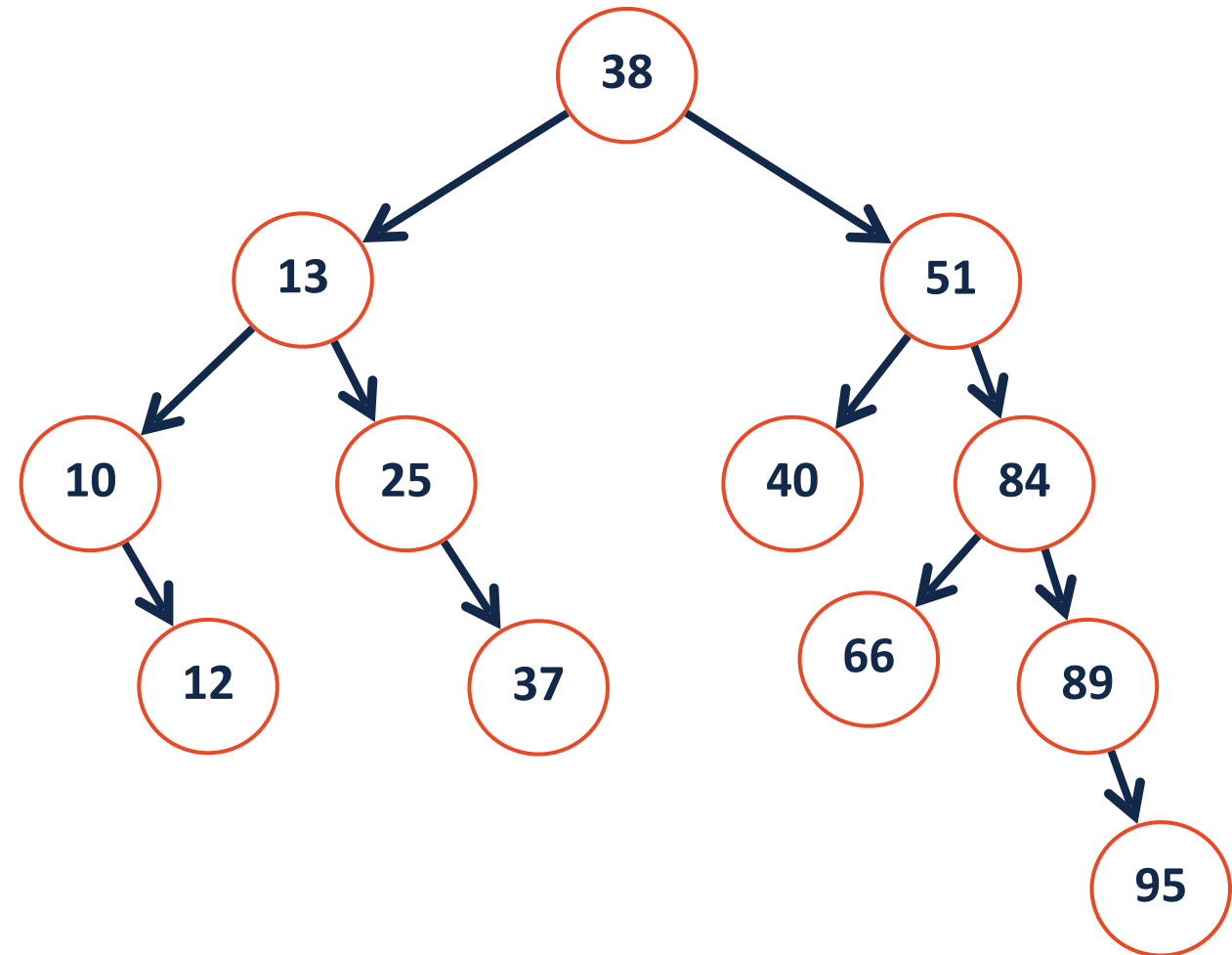
Recursive Step:

Combining:

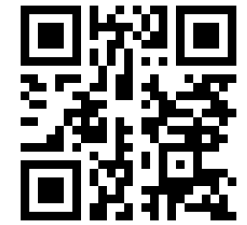


BST Insert

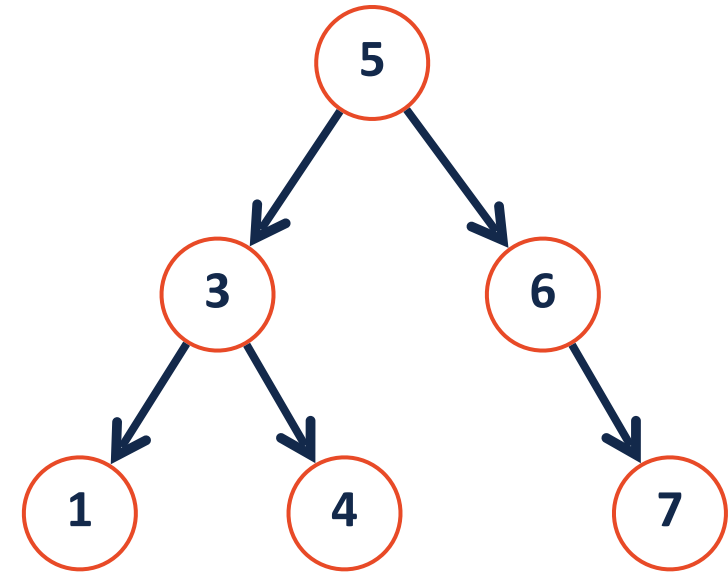
insert(33)



BST Insert



```
1 # inside class bst
2 def insert(self, key, val):
3     self.root = self.insert_helper(self.root, key, val)
4
5 def insert_helper(self, node, key, val):
6
7
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23
```

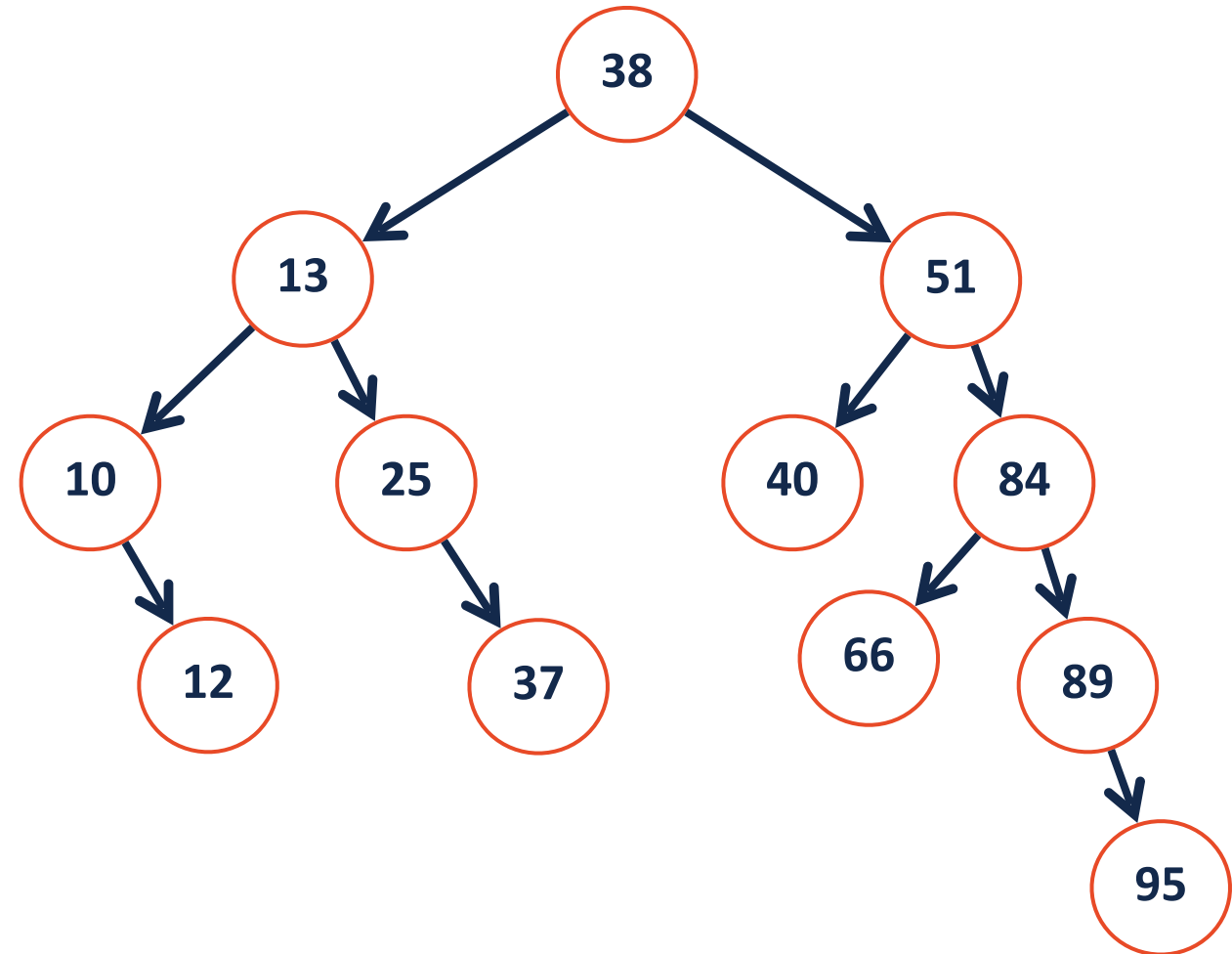


BST Insert

What binary tree would be formed by inserting the following sequence of integers: [3, 7, 2, 1, 4, 8, 0]

BST Remove

remove (40)



BST Remove

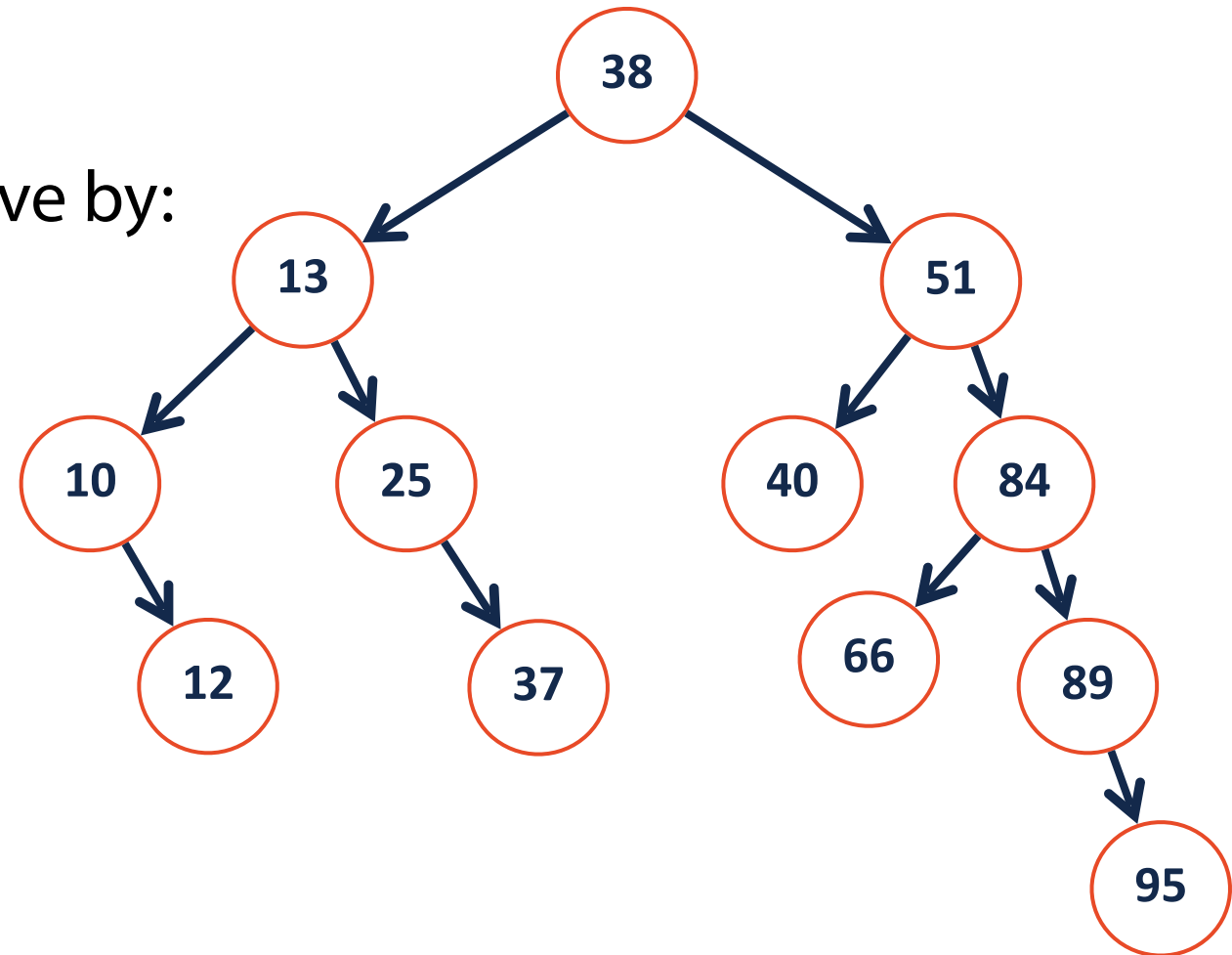
remove (12)

- 1) Find item being removed
- 2) Identify number of children

When we have zero children, remove by:

Set parent.**next** = None

next is either **left** or **right**

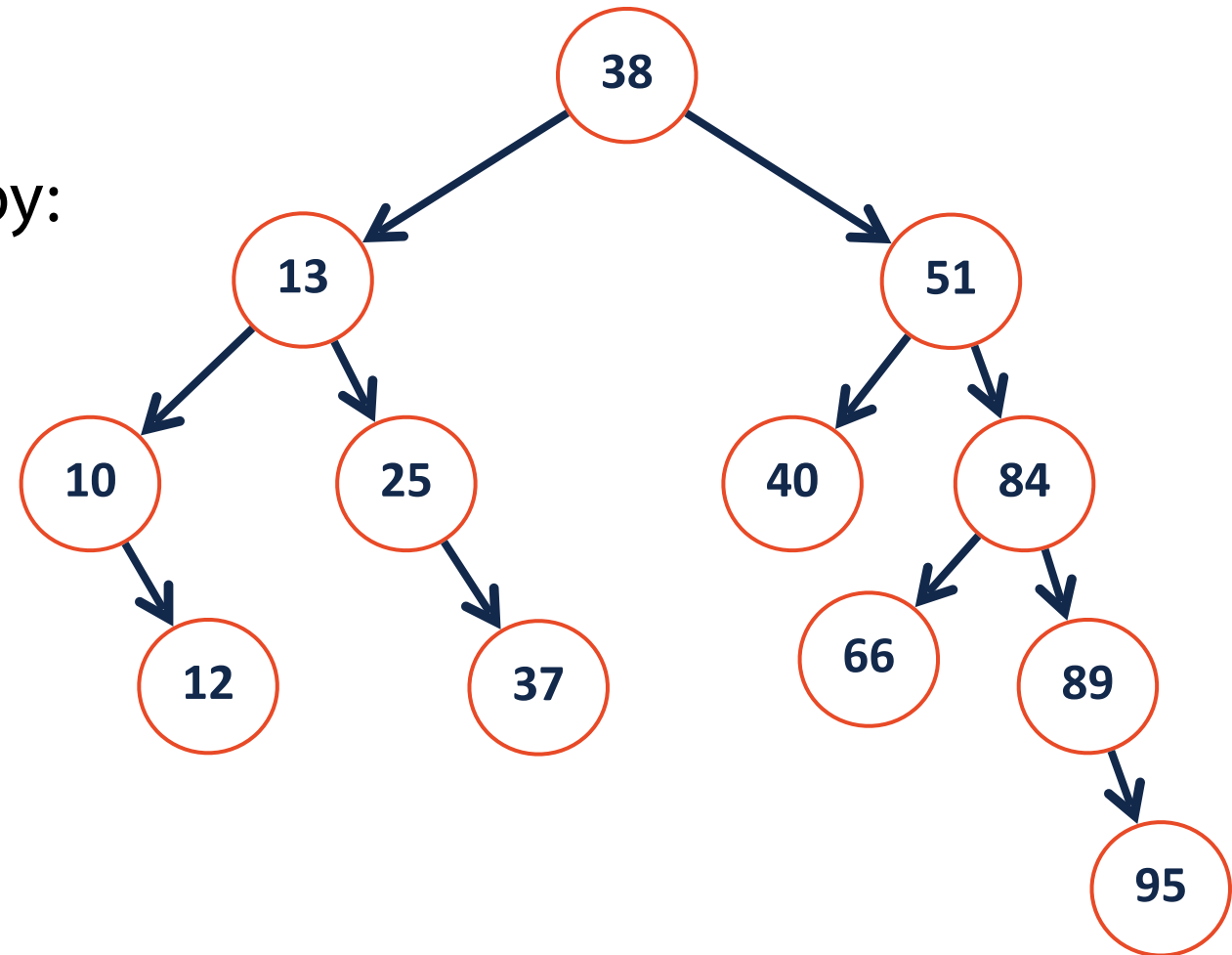


BST Remove

remove (25)

- 1) Find item being removed
- 2) Identify number of children

When we have one child, remove by:



BST Remove

remove(10)

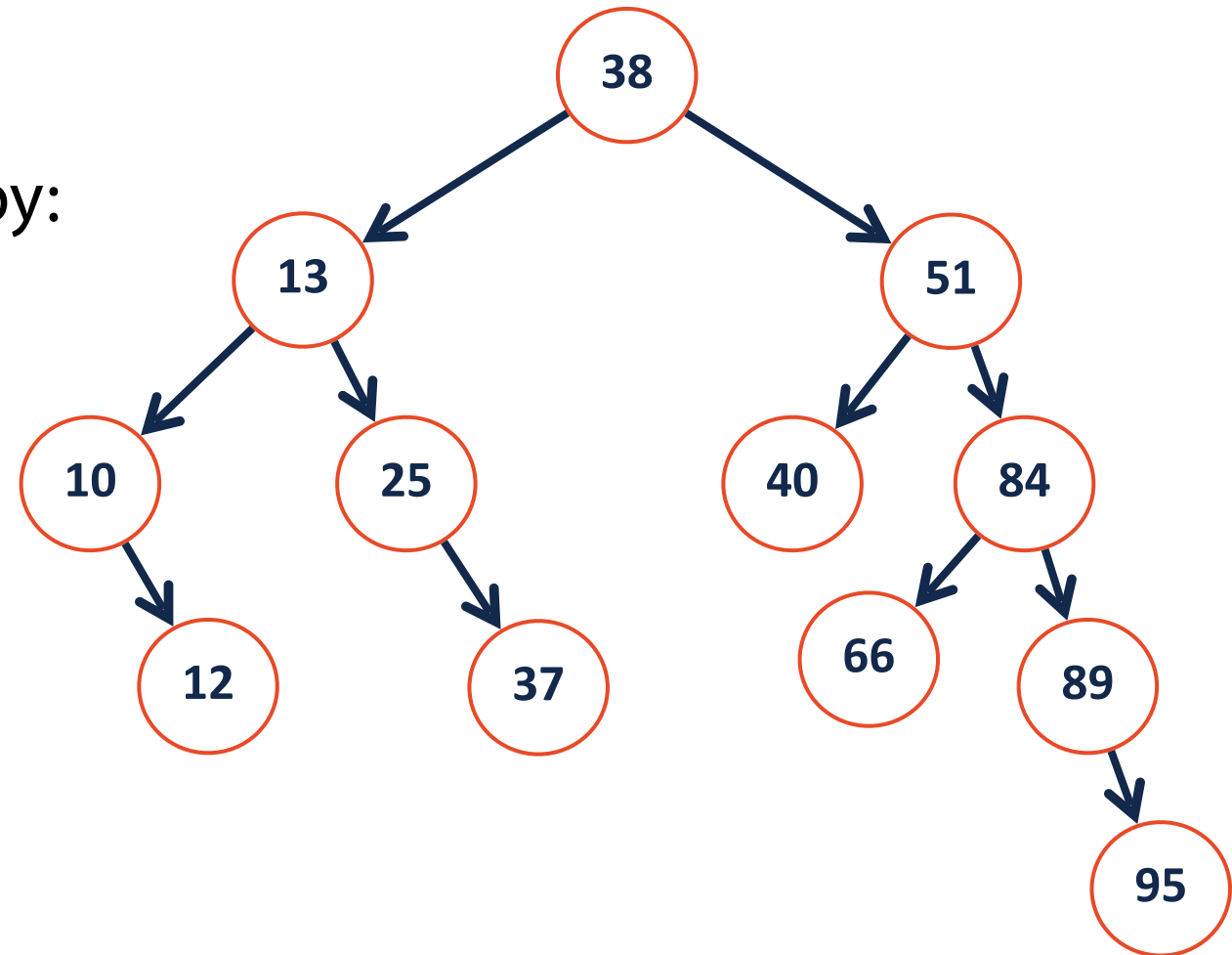
- 1) Find item being removed
- 2) Identify number of children

When we have one child, remove by:

Set parent.**next1** = target.**next2**

next1 is either **left** or **right**

next2 is either **left** or **right**

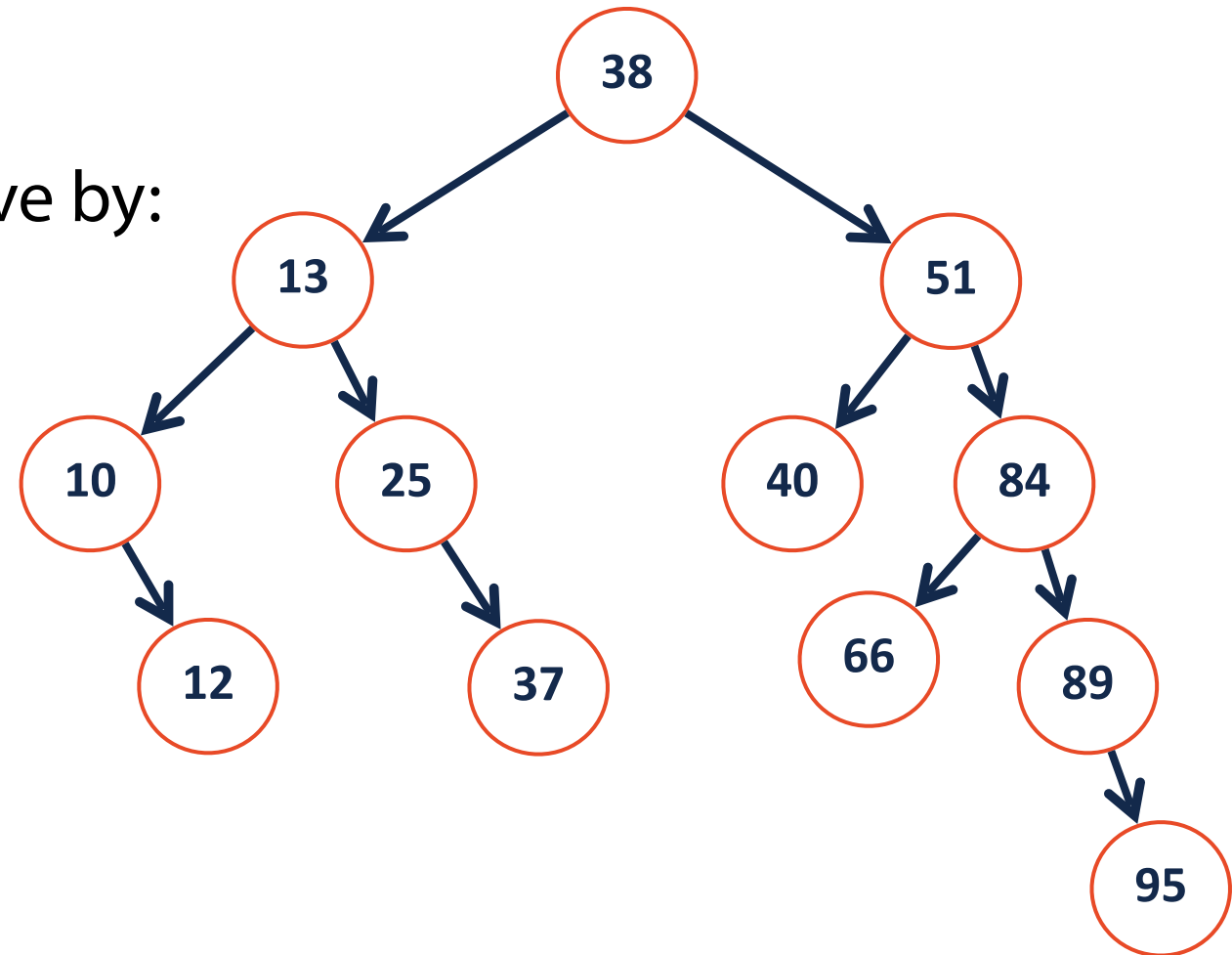


BST Remove

remove (13)

- 1) Find item being removed
- 2) Identify number of children

When we have two children, remove by:



BST In-Order _____



In-Order Predecessor

Rightmost left child

IOP(38) =

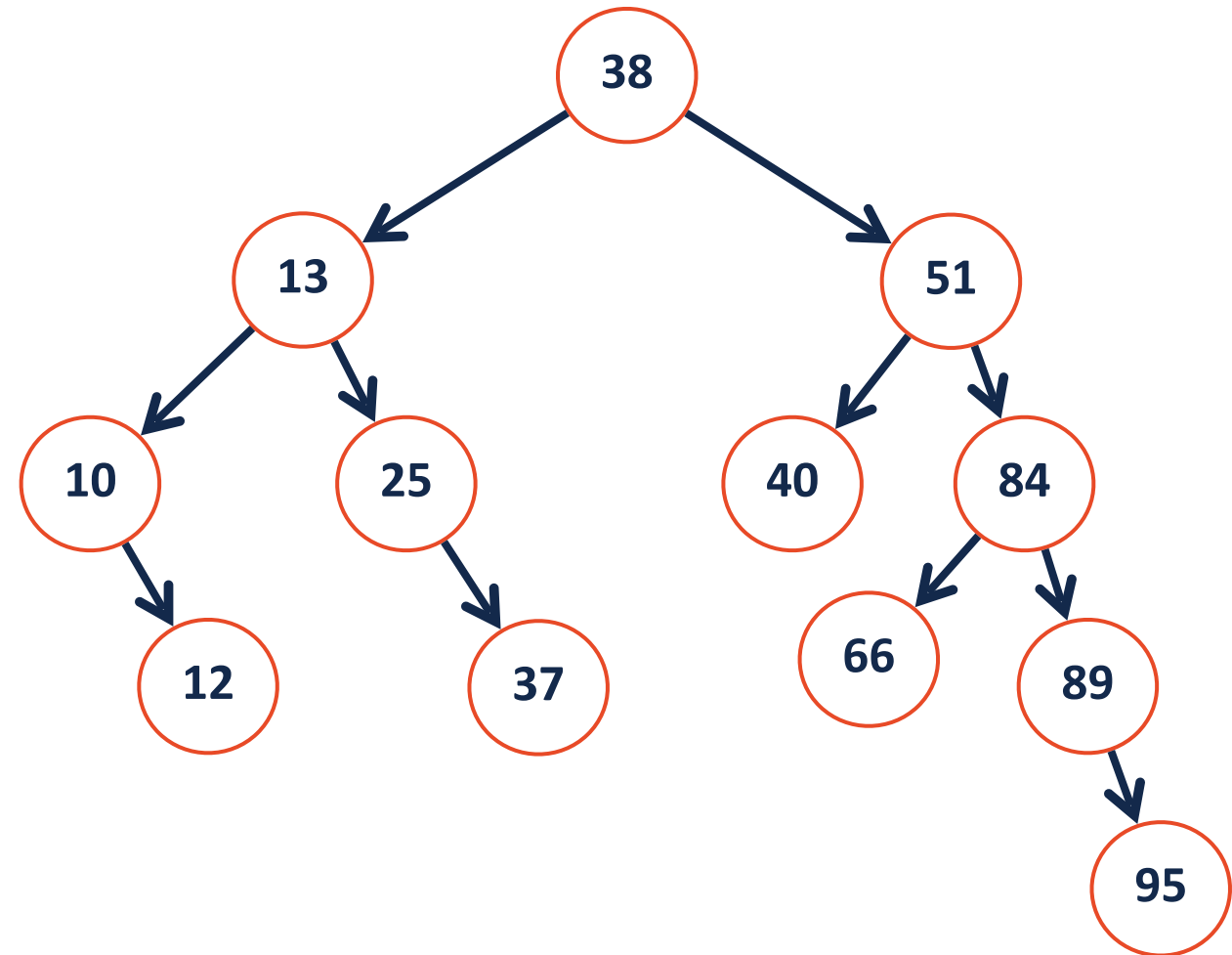
IOP(84) =

In-Order Successor

Leftmost right child

IOS(38) =

IOS(84) =



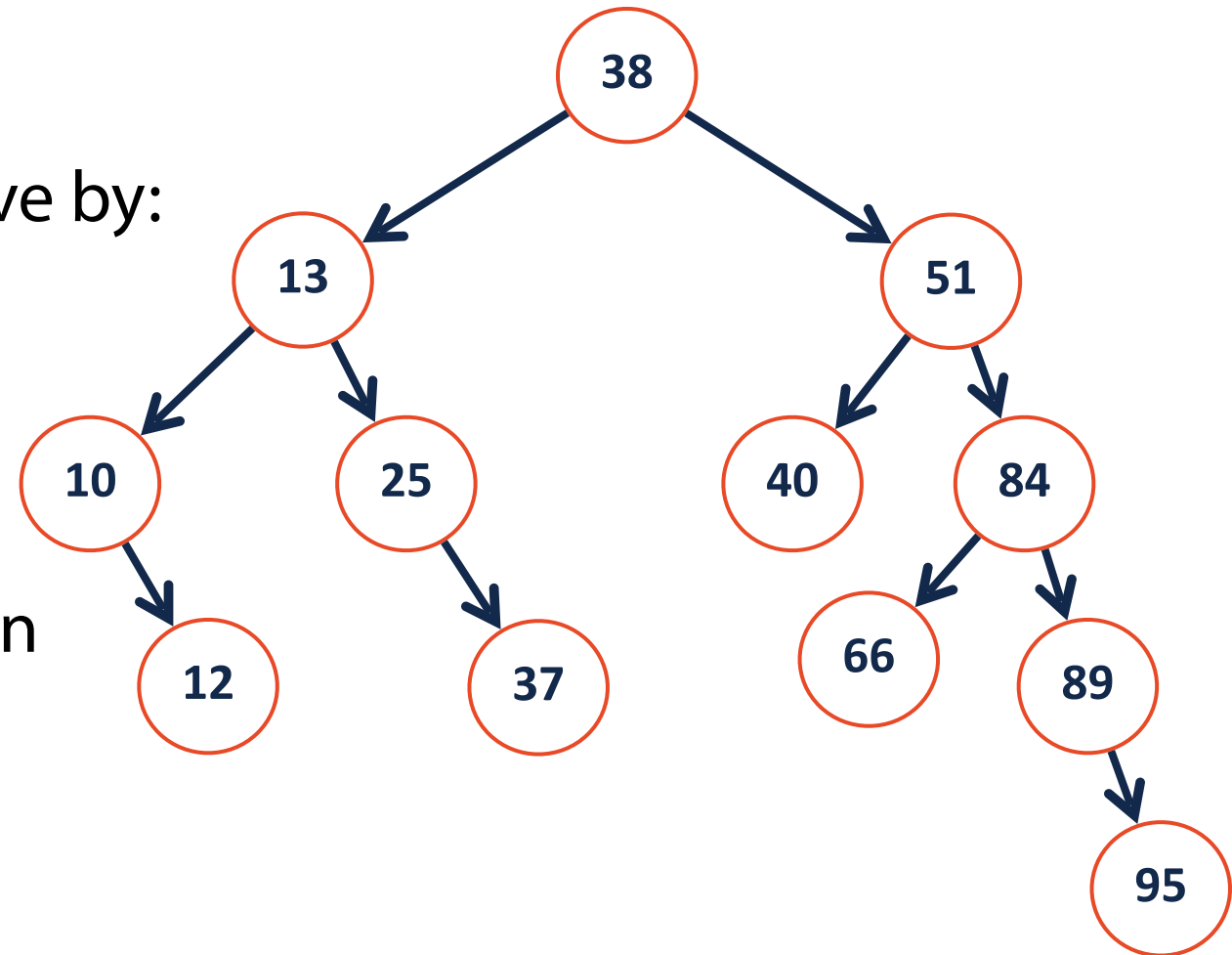
BST Remove

remove (13)

- 1) Find item being removed
- 2) Identify number of children

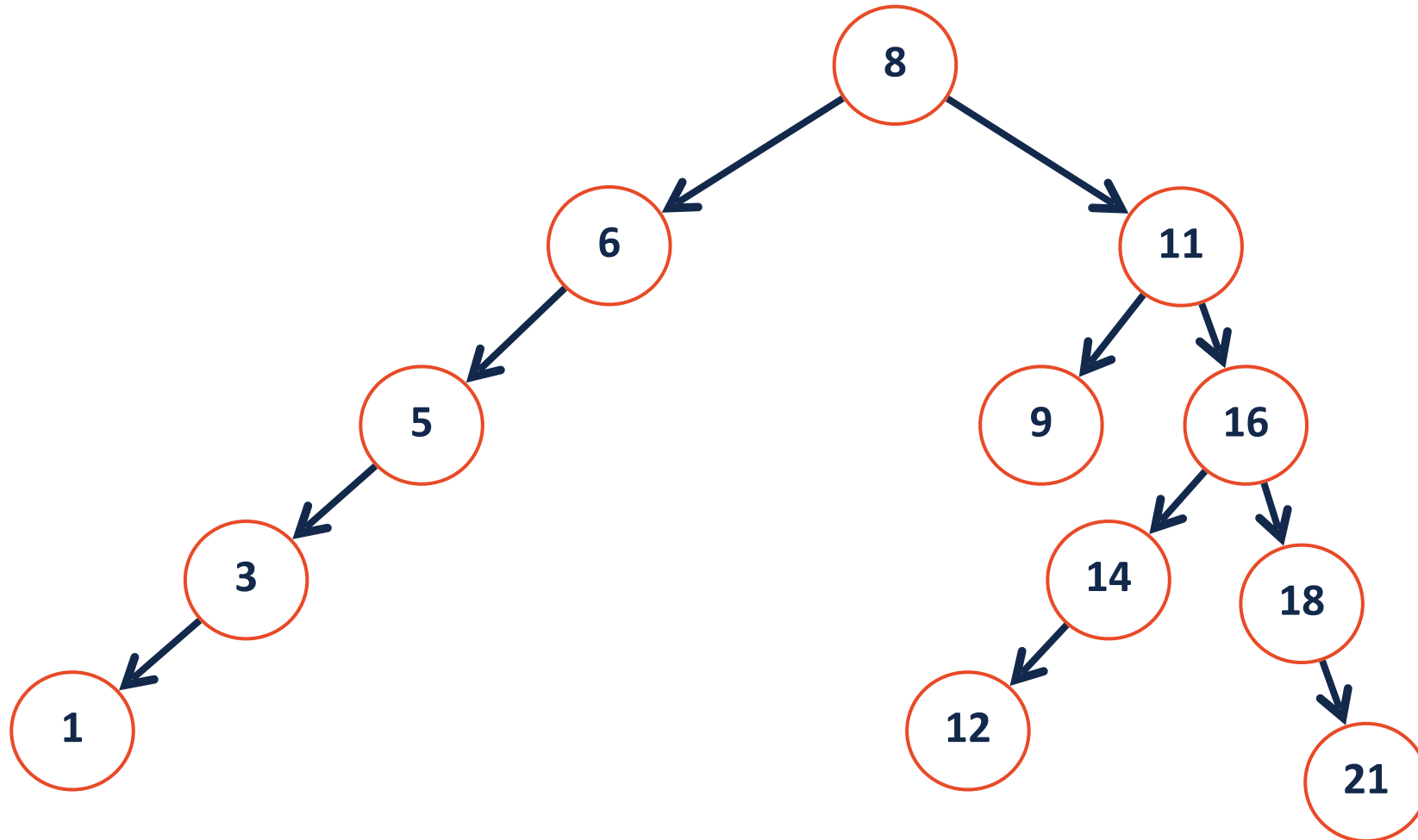
When we have two children, remove by:

- 3) Find the IOP / IOS
- 4) Swap target with IOP / IOS
- 5) Remove target at its new location



BST Remove

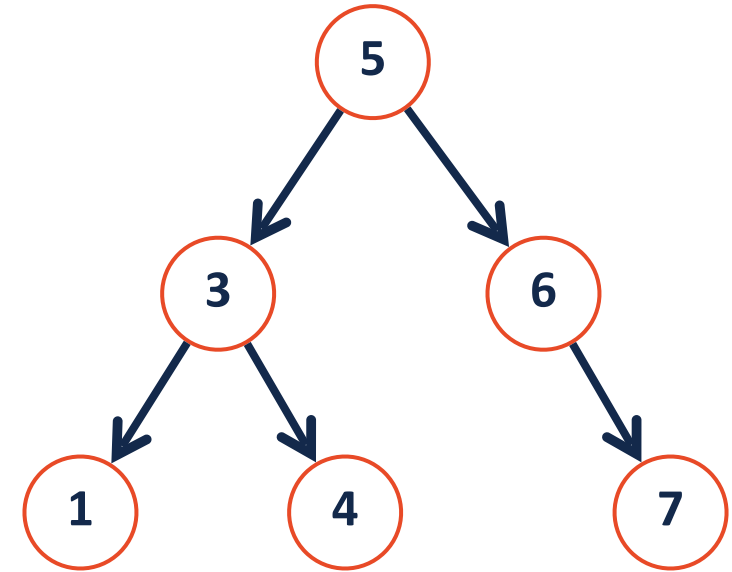
What will the tree structure look like if we remove node 16 using IOS?



BST Remove



```
1 def remove(self, key):
2     self.root = self.remove_helper(self.root, key)
3
4 def remove_helper(self, node, key):
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25
```



BST Analysis – Running Time



Operation	BST Worst Case
find	
insert	
delete	
traverse	

BST Analysis

Every operation on a BST depends on the **height** of the tree.

... how do we relate $O(h)$ to n , the size of our dataset?

BST Analysis



What is the max number of nodes in a tree of height h ?

BST Analysis

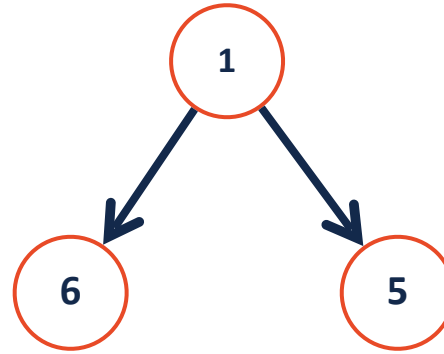


What is the min number of nodes in a tree of height h ?

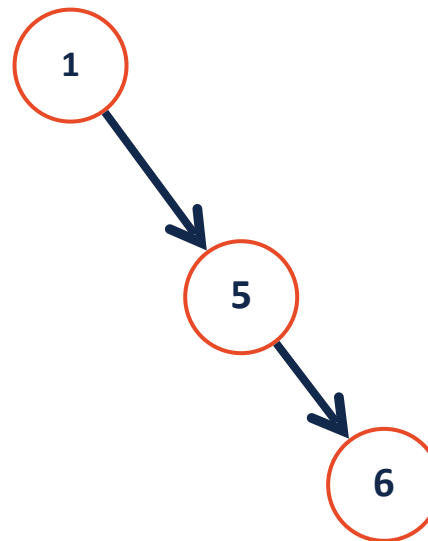
BST Analysis

A BST of n nodes has a height between:

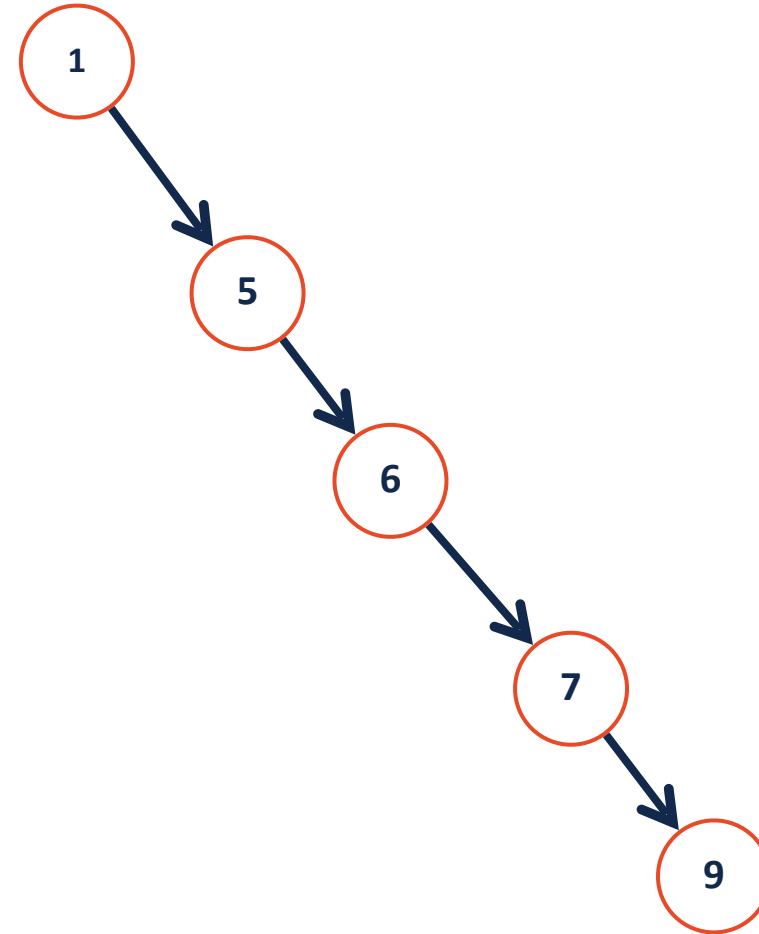
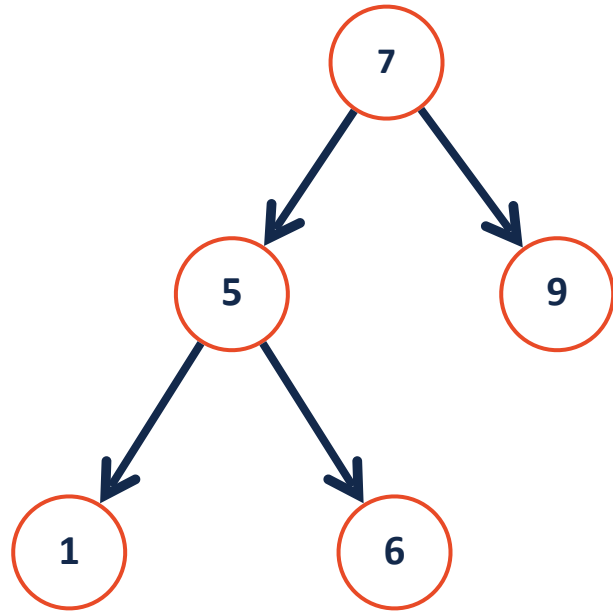
Lower-bound: $O(\log n)$



Upper-bound: $O(n)$



Limiting the height of a tree



Option A: Correcting bad insert order

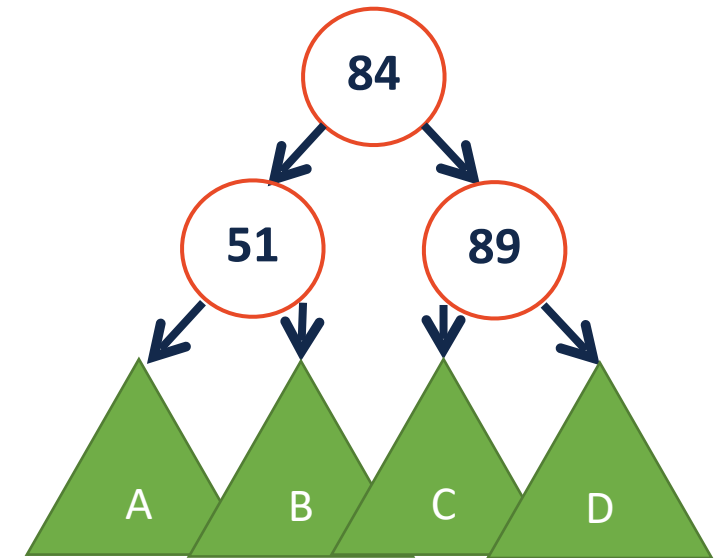
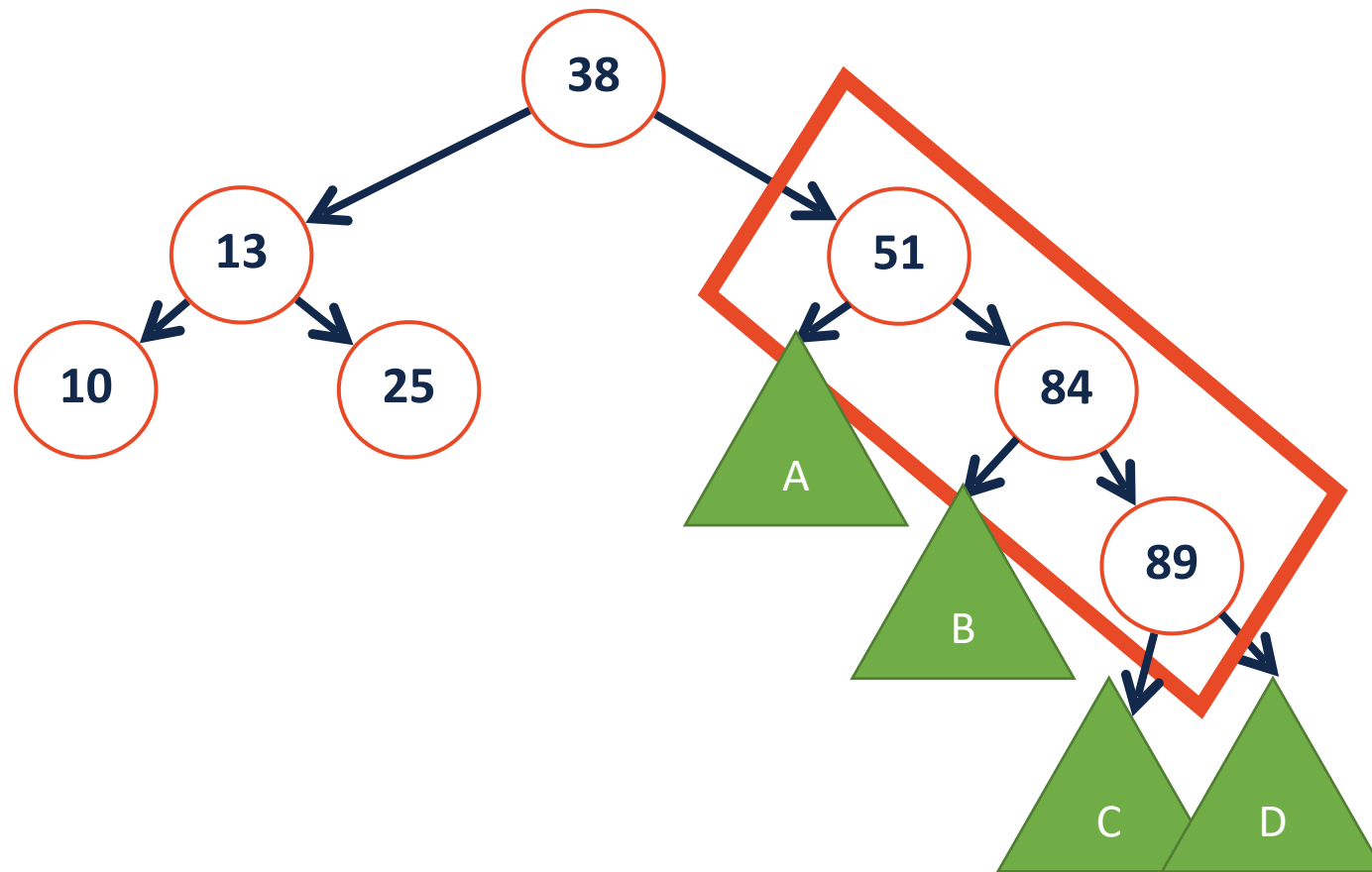
The height of a BST depends on the order in which the data was inserted

Insert Order: [1, 3, 2, 4, 5, 6, 7]

Insert Order: [4, 2, 3, 6, 7, 1, 5]

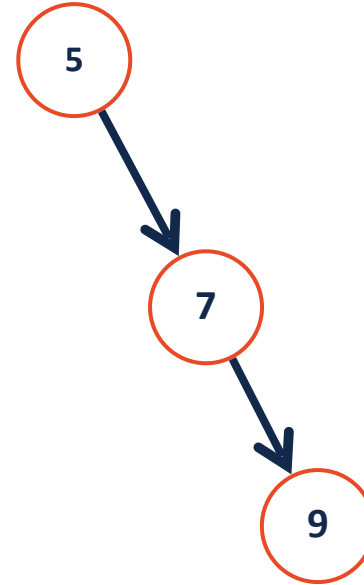
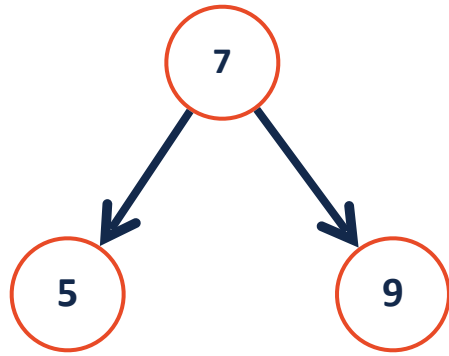
AVL-Tree: A self-balancing binary search tree

Rather than fixing an insertion order, just correct the tree as needed!



Height-Balanced Tree

What tree is better?

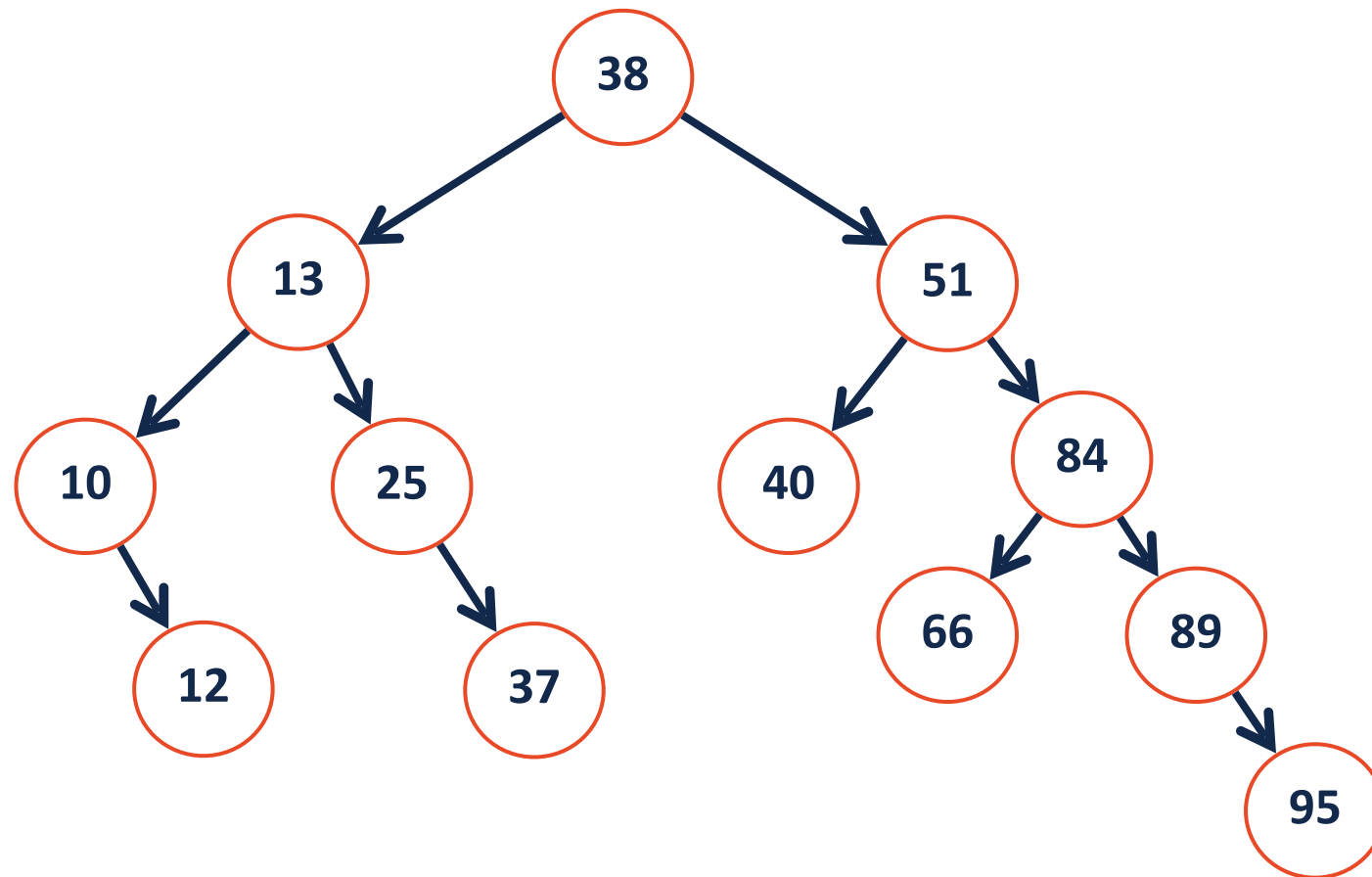


Height balance: $b = \text{height}(T_R) - \text{height}(T_L)$

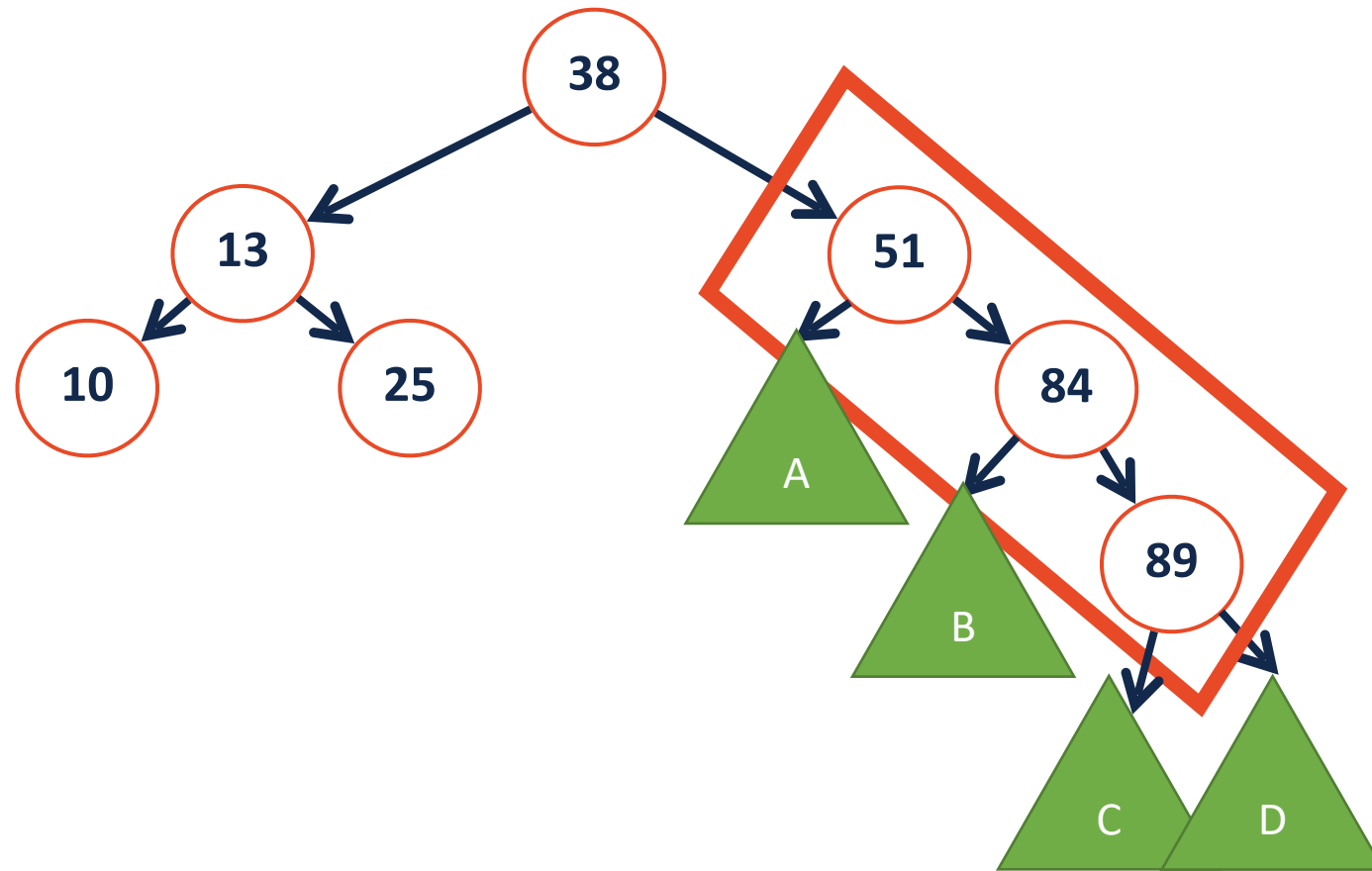
A tree is “balanced” if:

BST Rotations (The AVL Tree)

We can adjust the BST structure by performing **rotations**.



Left Rotation



Left Rotation

