

Algorithms and Data Structures for Data Science

Stacks and Queues and (maybe) 2D Lists

CS 277

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My white
whale
😊



The purpose of assessments in CS 277



Reminder: MP 0 due today!

Informal Early Feedback

An anonymous survey about the class

If 70% of class completes, everyone gets bonus points

Please provide constructive criticism and positive feedback

Learning Objectives

Observe data structure tradeoffs between data access and speed

Understand the fundamentals of the stack and queue

Introduce NumPy and practice 2D lists



Array Implementation

Given Node in CL
↓
⊕ → 0



| | Singly Linked List | Array |
|-------------------------------------|--------------------|----------|
| Look up arbitrary location | $O(n)$ | $O(1)$ ← |
| Insert after given element | $O(1)$ | $O(n)$ |
| Remove after given element | $O(1)$ | $O(n)$ |
| Insert at arbitrary location | $O(n)$ | $O(n)$ |
| Remove at arbitrary location | $O(n)$ | $O(n)$ |
| Search for an input value | $O(n)$ | $O(n)$ |

} find
:)



Thinking critically about lists: tradeoffs

Can we make our lists better at some things? What is the cost?

Imagine I want to do a very complex calculation.

$$\hookrightarrow O(n^3)$$

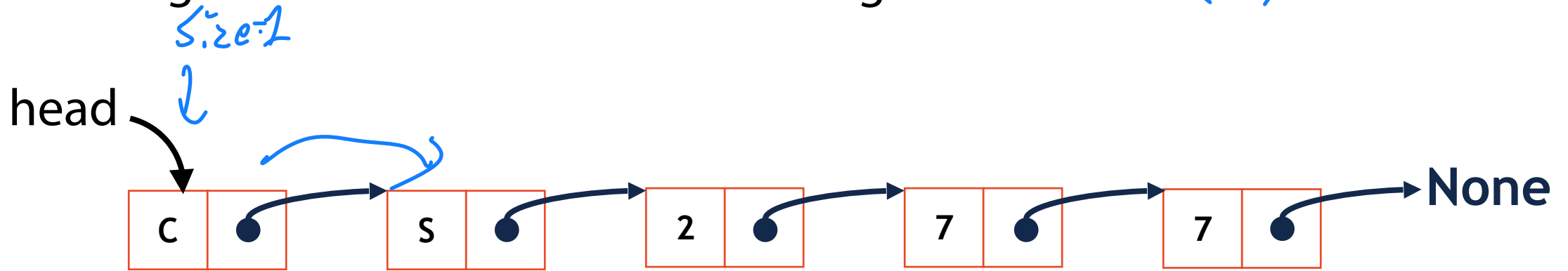
trivial way to become $O(1)$ \rightarrow preprocess every input

\hookrightarrow w/ infinite time & memory

store all possible answers

Thinking critically about lists: tradeoffs

Getting the size of a linked list has a Big O of: $O(n)$



Size = 1

Added `LinkedList.size` (a variable)

$\hookrightarrow O(1)$ ← cost is string size

Thinking critically about lists: tradeoffs

Does knowing our list is sorted change our Array Big O?

Specifically
in this
case
Insert here
is
 $O(1)$ *

unsorted

| | | | | | | | | | |
|---|---|---|---|---|----|---|---|---|---|
| 2 | 7 | 5 | 9 | 7 | 14 | 1 | 0 | 8 | X |
|---|---|---|---|---|----|---|---|---|---|

X → $O(n)$

sorted

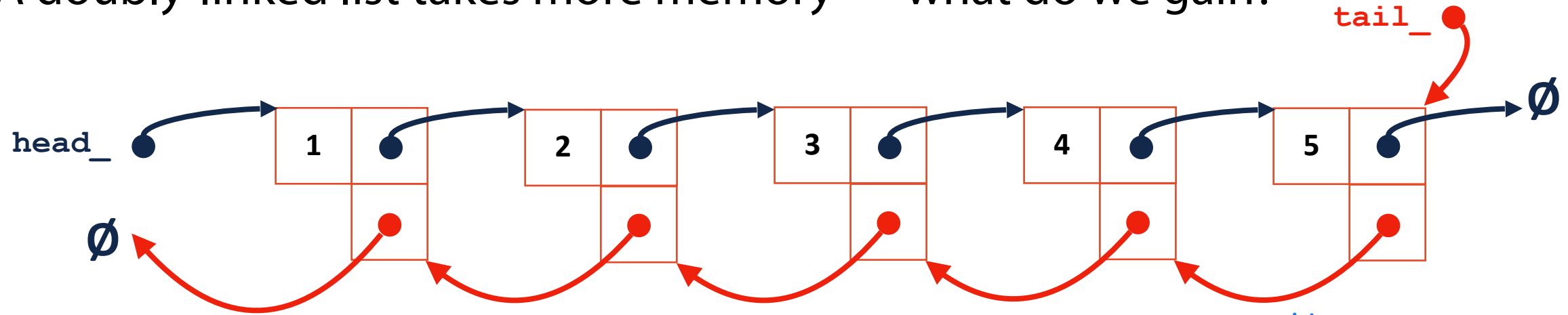
find Min $O(1)$

| | | | | | | | | | |
|---|---|---|---|---|---|---|---|----|--|
| 0 | 1 | 2 | 5 | 7 | 7 | 8 | 9 | 14 | |
|---|---|---|---|---|---|---|---|----|--|

find Max → $O(1)$

Thinking critically about lists: tradeoffs

A doubly-linked list takes more memory — what do we gain?



Node

↳ Data

↳ Next

↳ Prev

Increase in size

walk
backwards!

Thinking critically about lists: tradeoffs

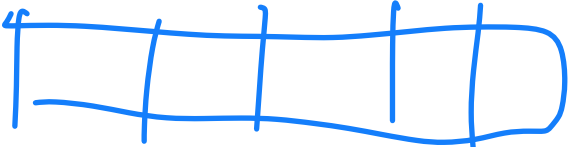
Consider carefully how data structures can be modified for a problem!

Lets see two examples of this:

I want a data structure that can add and remove in $O(1)$.

I'm willing to 'trade away' a lot of utility to do this.

↳ I don't care about (random access)

$X =$ 

$X[i]$

The Stack ADT

A **stack** stores an ordered collection of objects (like a list)

However you can only do three operations:

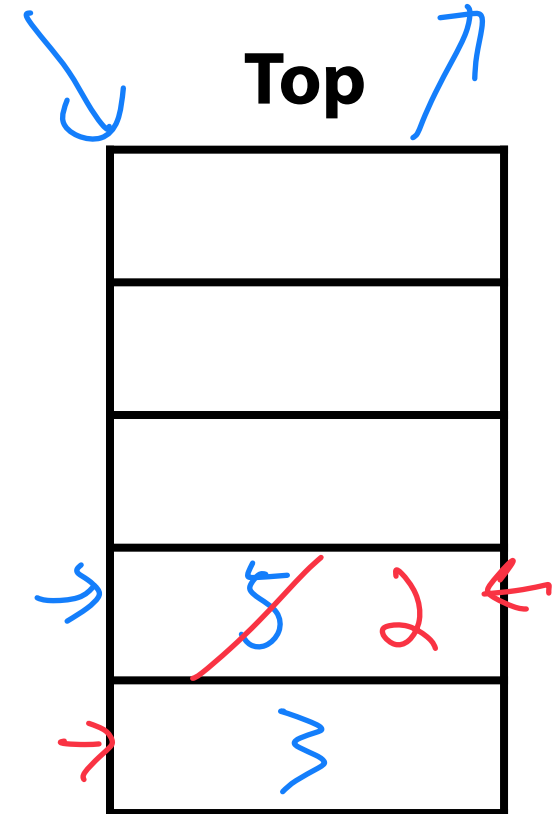
Push: Put an item on top of the stack

Pop: Remove the top item of the stack (and return it)

Top: Look at the value of the top item

`push (3) ; push (5) ; top () ; pop () ; push (2)`

↳ 5





Programming Toolbox: Stack

The stack is a **last in — first out** data structure (LIFO)

```
1 def reverse(inList):  
2     s = stack() ←  
3     for v in inList:  
4         s.push(v) ]  
5  
6     out = []  
7     while not s.empty():  
8         out.append(s.pop())  
9     return out
```

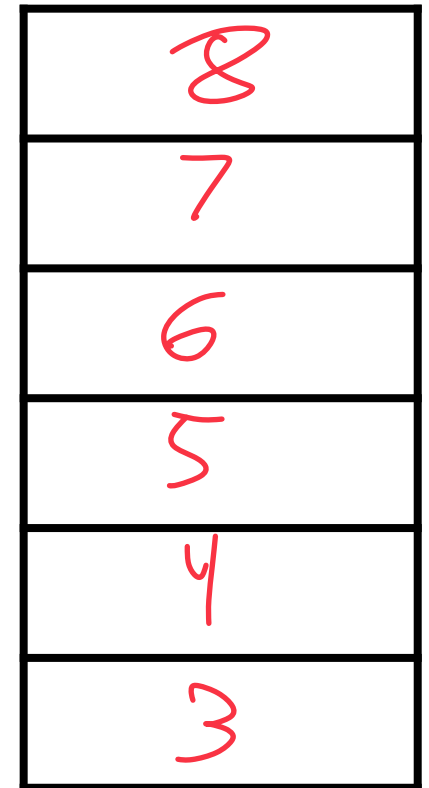
Not a Python built-in!

`reverse([3, 4, 5, 6, 7, 8])`

out



stack s



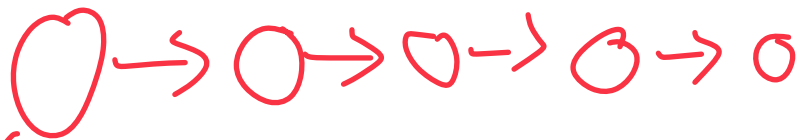
Stack Implementation (using Lists)

good at random access

| | Singly Linked List | Array |
|-------------------------------------|--------------------------|--------------------------|
| Look up arbitrary location | $O(n)$ | $O(1)$ |
| Insert after given element | $O(1)$ | $O(n)$ |
| Remove after given element | $O(1)$ | $O(n)$ |
| Insert at arbitrary location | $O(n)$ | $O(n)$ |
| Remove at arbitrary location | $O(n)$ | $O(n)$ |
| Search for an input value | $O(n)$ | $O(n)$ |

good at insert/remove in specific cases

Insert



remove

Stack as a Linked List

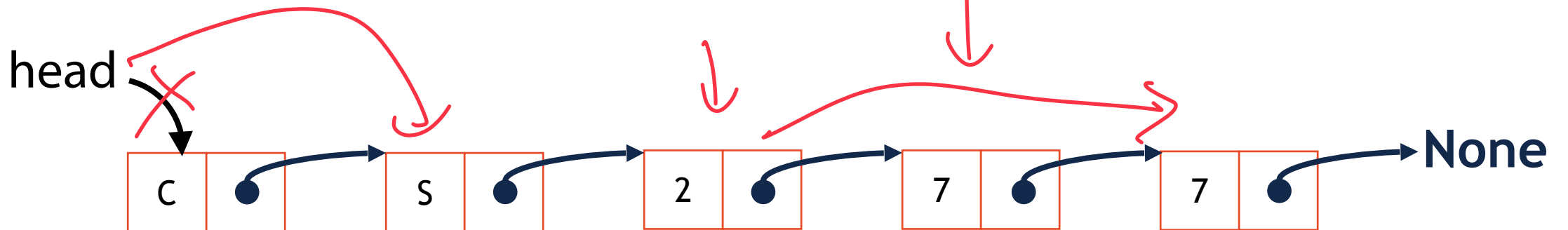
pop ()

Pop() removes the top item from my list

1) $head = head.next$ $O(1)$

$O(1)$ as well if we have curr

$curr.next = curr.next.next$



Stack as a Linked List

top ()

Top() looks at the top item of the list

head, data

$O(1)$

-val



Stack Implementation



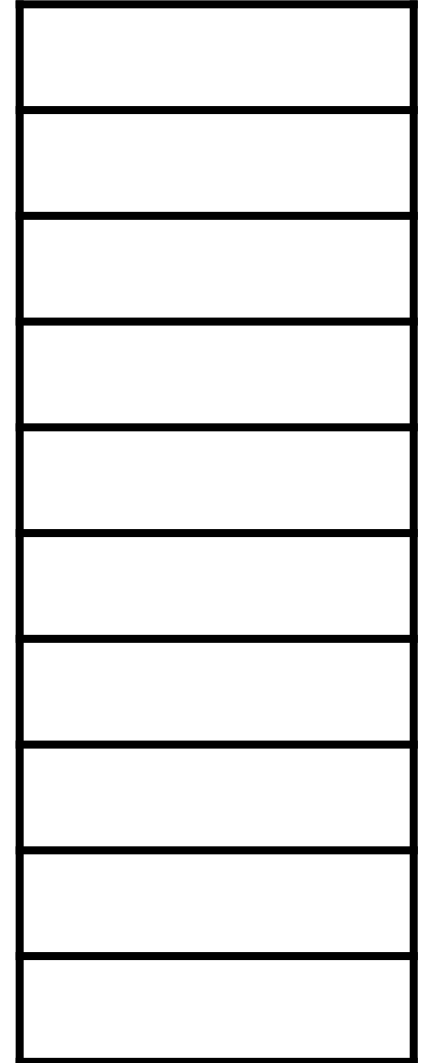
```
1 class Node:
2     def __init__(self, value, next = None):
3         self.val = value
4         self.next = next
5
6 class stack:
7     def __init__(self):
8         self.head = None
9         self.length = 0
10
11    def push(self, value):
12        self.length += 1
13        newNode = Node(value)
14        newNode.next = self.head
15        self.head = newNode
16
17    def __len__(self):
18        return self.length
19
20
21
22
23
24 # class stack:
25     def top(self):
26         if self.length > 0:
27             return self.head.val
28         return None
29
30     def pop(self):
31         if (self.length > 0):
32             self.length -= 1
33             popped = self.head
34             self.head = self.head.next
35             return popped.val
36         return None
37
38 # Some other support functions in code base
39
40
41
42
43
44
45
46
```

On your own: Stack Practice

What will the stack look like as you run the following code?

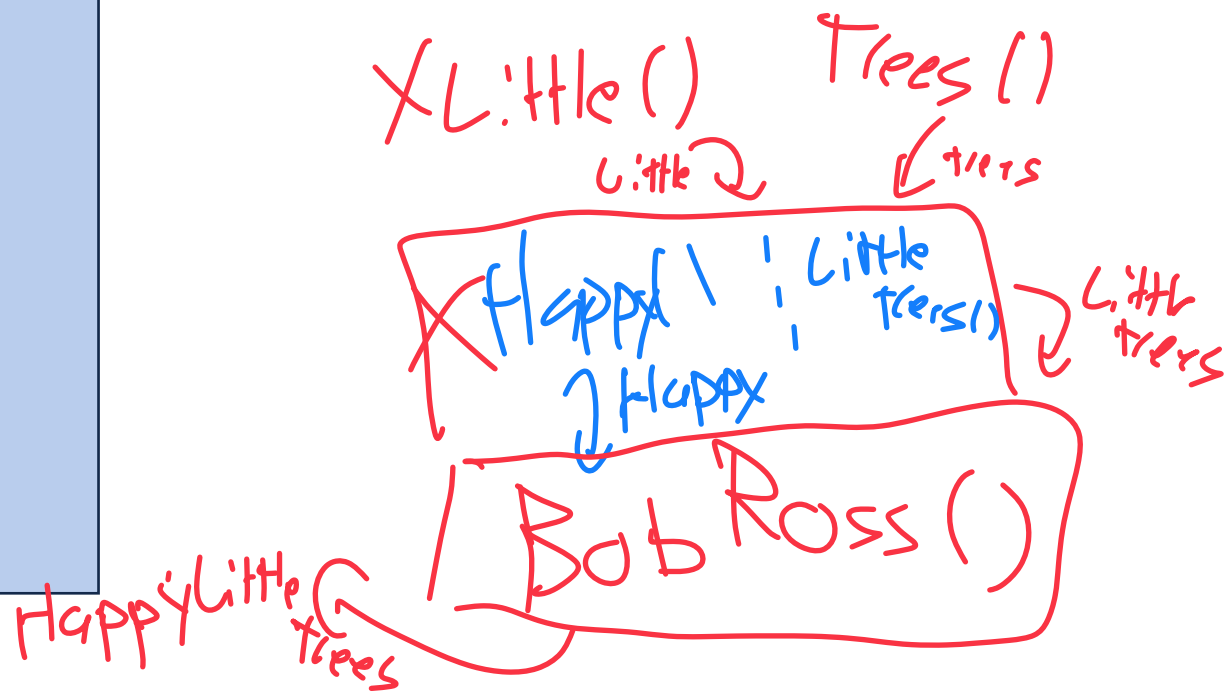
Try it by hand and run the code to check!

```
1 s = stack()
2
3 print(s.empty())
4
5 for i in range(0, 10, 2):
6     s.push(i)
7
8 print(s)
9
10 x = s.pop()
11 print(x, s)
12
13 print(len(s))
14
15 print(s.top())
16 s.pop()
17 print(s.top())
18
19 print(s.empty())
```



The Call Stack

```
1 def Happy():
2     print("Calling Happy!")
3     return "Happy"
4
5 def Little():
6     print("Calling Little!")
7     return "Little"
8
9 def Trees():
10    print("Calling Trees!")
11    return "Trees"
12
13 def LittleTrees():
14    print("Calling LittleTrees!")
15    return Little() + Trees()
16
17 def BobRoss():
18    print("Calling BobRoss!")
19    return Happy() + LittleTrees()
20
21 print(BobRoss())
22
23
```



Queue Data Structure

A **queue** stores an ordered collection of objects (like a list)

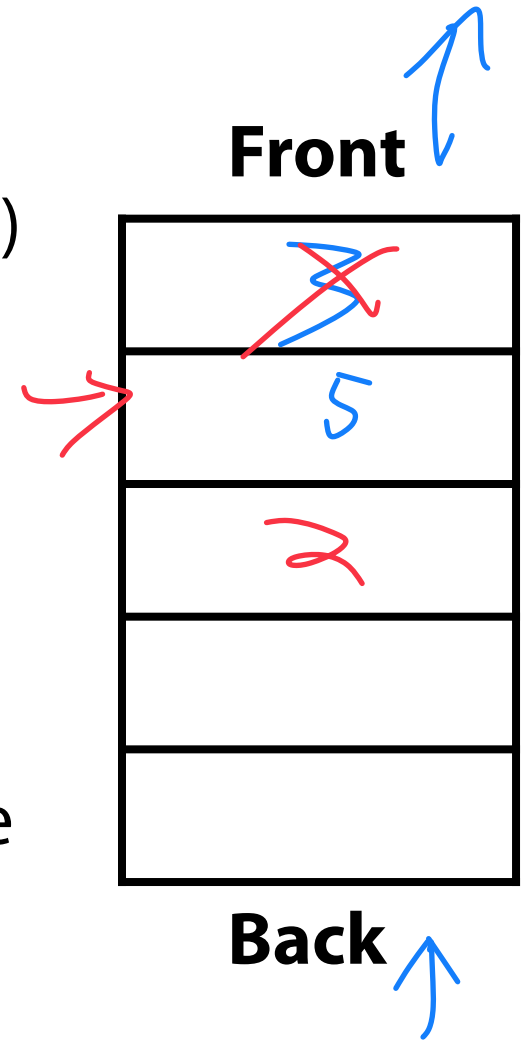
However you can only do ~~three~~ two operations:

Enqueue: Put an item at the back of the queue

Dequeue: Remove and return the front item of the queue

Front: Get value at front of queue

enqueue (3) ; enqueue (5) ; dequeue () ; enqueue (2)



Queue Data Structure

The queue is a **first in — first out** data structure (FIFO)

What data structure excels at removing from the front?

Linked list



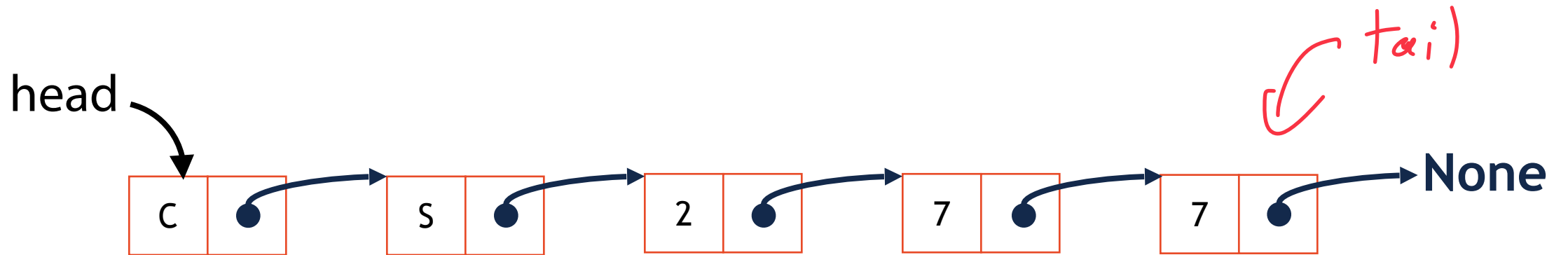
Can we make that same data structure good at inserting at the end? ??

↳ can we make head variable for end of list?

Queue Data Structure

Once again, a linked list is a great implementation of a queue!

But we need one modification...



Linked List

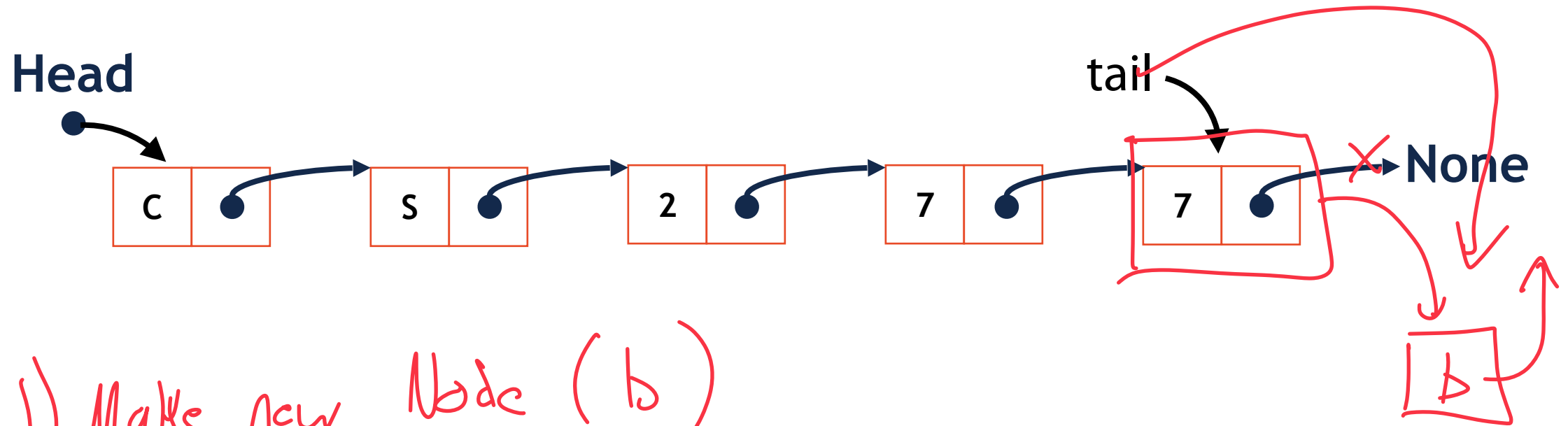
↳ head

↳ length

↳ tail - points to last item

Queue as Linked List

enqueue (b)



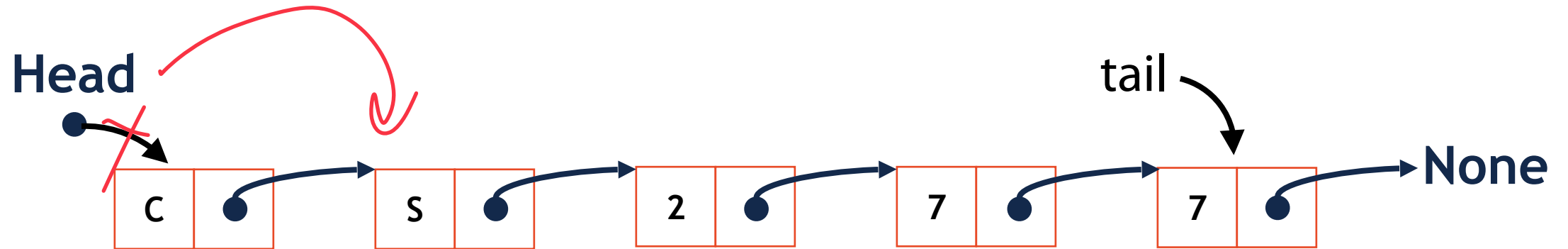
1) Make new Node (b)

2) $\text{tail.next} = \text{new Node}$

3) $\text{tail} = \text{new Node}$ # $\text{tail} = \text{tail.next}$

Queue as Linked List

dequeue ()



↳ This is start pop()

Queue Implementation



```
1 class Node:
2     def __init__(self, value, next = None):
3         self.val = value
4         self.next = next
5
6 class queue:
7     def __init__(self):
8         self.length = 0
9         self.head = None
10        self.tail = None
11
12    def enqueue(self, value):
13        self.length += 1
14        item = Node(value)
15        if self.head == None:
16            self.head = item
17            self.tail = item
18        else:
19            self.tail.next = item
20            self.tail = self.tail.next
21
22
23
```

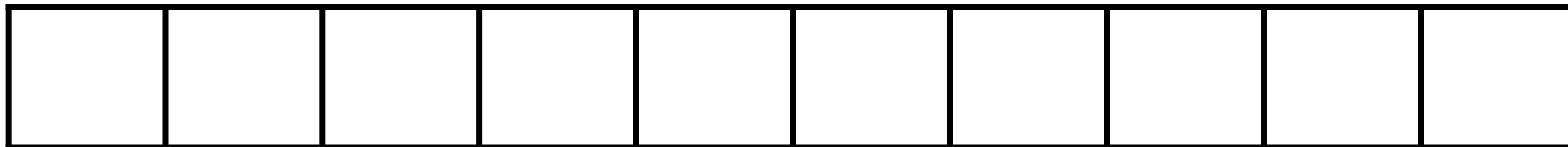
```
24 # class queue:
25     def dequeue(self):
26         if self.length > 0:
27             self.length -= 1
28             item = self.head
29             self.head = item.next
30
31             if self.head == None:
32                 self.tail = None
33             return item.val
34
35         return None
36
37     def front(self):
38         if self.length > 0:
39             return self.head.val
40         return None
41
42 # Other support functions in code base
43
44
45
46
```

On your own: Queue Practice

```
1 q = queue()
2
3 print(q.empty())
4
5 for i in range(0,20, 2):
6     q.enqueue(i)
7 print(q)
8
9 x = q.dequeue()
10 print(x, q)
11
12 print(len(q))
13
14 print(q.front())
15 q.dequeue()
16 print(q.front())
17
18 print(q.empty())
```

What will this code output?

Try it by hand and run the code to check!



The CS 277 Queue

Now you know how the CS 277 queue works!

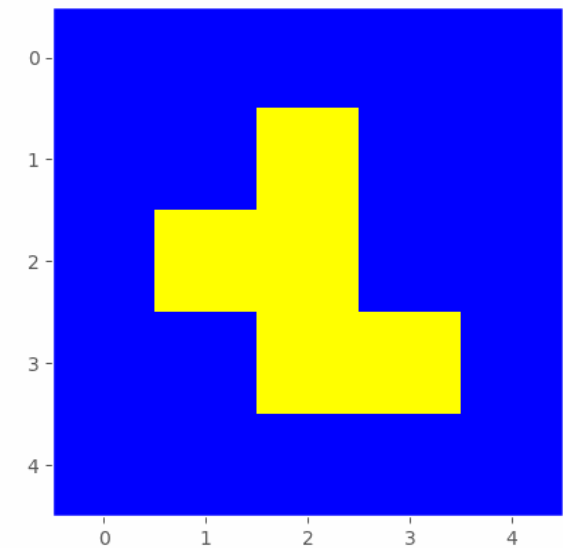
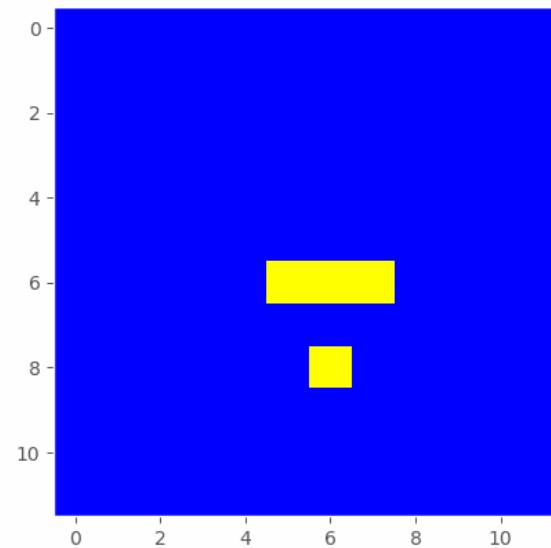
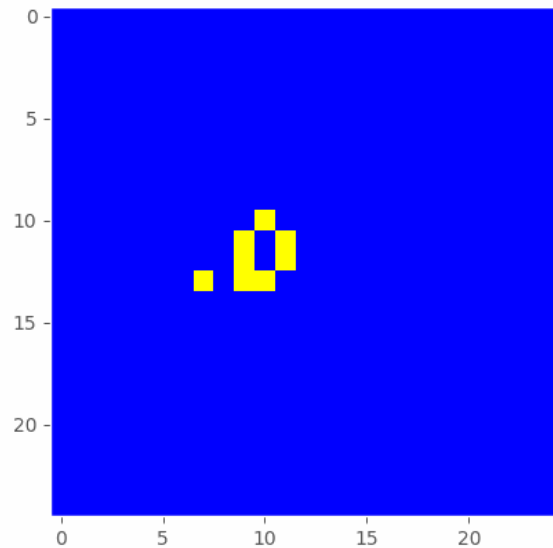
<https://queue.illinois.edu/q/course/185>

Cellular Automata — probably pushed back?

A computational model consisting of a **matrix** and a **set of rules**

Each iteration, the matrix changes based on its current state

There are a number of emergent behaviors that can be discovered!



Programming Practice: 2D Lists

Given a function that passes in three lists, make a 2D array storing all three.

How would I search a 2D list for a specific value?

See

code

Programming Practice: 2D Lists

What shape will this code produce?

5 x 5 list why?

of rows

```
1 outerList = []
2
3 for i in range(5):
4     innerList = []
5
6     for j in range(5):
7         innerList.append(i+j)
8
9     outerList.append(innerList)
10
11
12
13
14
15
16
17
18
```

size of inner list # of rows



Programming Practice: 2D Lists

What values will this list produce?

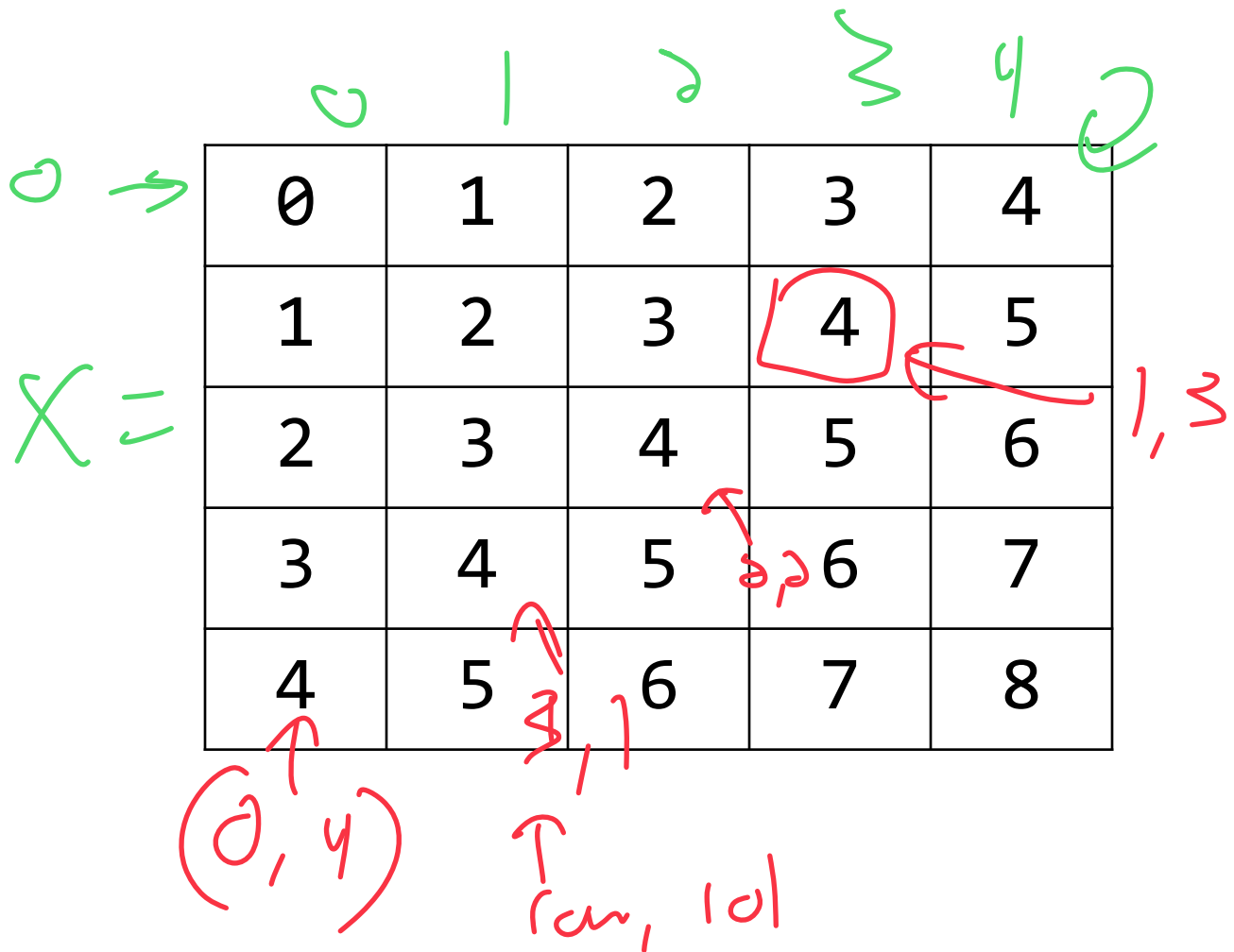
Handwritten diagram illustrating a 2D list structure. The grid is 5 rows by 5 columns. The first row is circled in green, and a green arrow points to the first cell (0,0). The first row is labeled with $i=0$ on the left and $j=0$ above the first cell. The columns are labeled with $j=0, 1, 2, 3, 4$ above the grid. The rows are labeled with $i=0, 1, 2, 3, 4$ on the left. The values in the first row are 0, 1, 2, 3, 4. The values in the second row are 1, 2, 3, 4, 5. The values in the third row are 2, 3, 4, 5, 6. The fourth and fifth rows are empty, with a dash and three dots in the fourth row indicating continuation.

| | | | | | |
|-------|---|---|---|---|---|
| $i=0$ | 0 | 1 | 2 | 3 | 4 |
| $i=1$ | 1 | 2 | 3 | 4 | 5 |
| $i=2$ | 2 | 3 | 4 | 5 | 6 |
| $i=3$ | | - | . | . | . |
| $i=4$ | | | | | |

```
1 outerList = []
2
3 for i in range(5):
4     innerList = []
5
6     for j in range(5):
7         innerList.append(i+j)
8
9     outerList.append(innerList)
10
11
12
13
14
15
16
17
18
```


Programming Practice: 2D Lists

What are the indices of every value of 4 in this list? $X[0][4]$



```
1 outerList = []
2
3 for i in range(5):
4     innerList = []
5
6     for j in range(5):
7         innerList.append(i+j)
8
9     outerList.append(innerList)
10
11
12
13
14
15
16
17
18
```

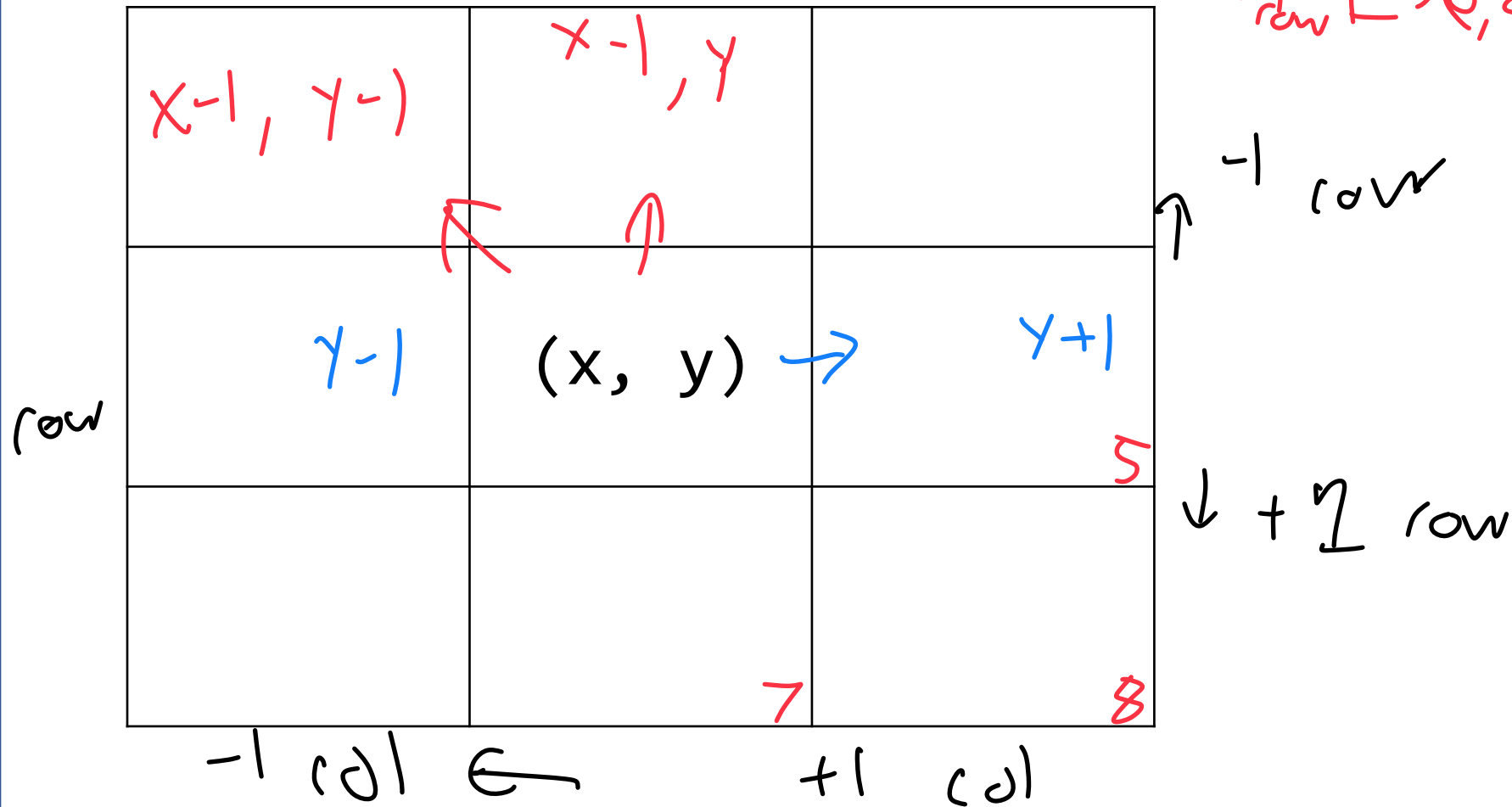


Programming Practice: 2D Lists

What are the indices of the square around the point (x, y) ?

$0, 0$ $x, y \equiv$ row, col col

Opt value
row $\rightarrow (0, 0)$



Programming Tip: Make a quick example!

I want to make a fast matrix with the values:

| | | |
|---|---|---|
| 1 | 2 | 3 |
| 4 | 5 | 6 |
| 7 | 8 | 9 |

```
1 count = 0
2 out = []
3 for i in range(3):
4     tmp = []
5     for j in range(3):
6         tmp.append(count)
7         count+=1
8     out.append(tmp)
9
10 print(out[2][2])
11
12 print(out[2][1])
13
14 print(out[1][2])
15
16
17
18
```

```
1 import numpy as np
2
3 x = np.arange(9).reshape(3, 3)
4
5 print(x)
6
7
8
9
```

Programming Toolbox: NumPy

NumPy is optimized for multidimensional arrays of numbers

```
1 import numpy as np
2
3 # Convert list to np list
4 n1 = np.array([1, 2, 3, 4, 5, 6])
5 print(n1)
6
7 # See list shape
8 print(n1.shape)
9
10 # Modify list shape
11 n12 = n1.reshape(3, 2)
12
13 print(n1)
14 print(n12)
15
16 # Create a new list
17 n13 = np.arange(15).reshape(5, 3)
18 n14 = np.zeros((2, 5))
19
20 print(n13)
21 print(n14)
22
23
```

`np.array(<list>)`

Convert a built-in list to a NumPy List

`<narray>.shape`

Get the shape of the NumPy array

`np.reshape(<row>, <col>)`

If the list contains exactly row x col items
reshape the list to those dimensions

`np.arange()`

returns a NumPy array with range() values

`np.zeros((<row>, col>))`

Create a list of 0s of the provided shape

Programming Toolbox: NumPy

Basic operations are applied **elementwise** (to each item of a list)

```
1 n1 = np.arange(4).reshape(2, 2)
2
3 print(n1)
4
5 n12 = n1 * 2
6
7 print(n12)
8
```

NumPy lists of equal size can also be added / subtracted / multiplied / etc...

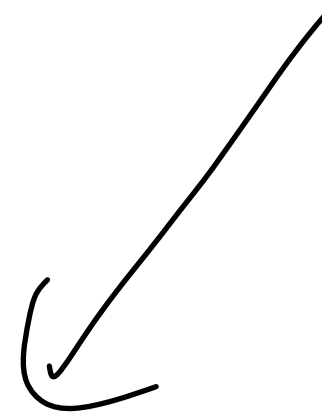
```
1 m1 = np.arange(9).reshape(3,3)
2 print(m1)
3
4 m2 = np.arange(9, 0, -1).reshape(3,3)
5 print(m2)
6
7 print(m1+m2)
8
```

Programming Toolbox: NumPy



NumPy is huge and can do many things

```
1 n1 = np.arange(4).reshape(2, 2)
2 n12 = n1 * 2
3
4 # Matrix multiplication
5 # 0*0+1*4      0*0+1*6
6 # 2*0+3*4      2*2+3*6
7 print(n1.dot(n12))
8
9 # Unique items
10 x = np.array([1, 1, 1, 2, 3, 4, 4, 5, 5, 6])
11 print(np.unique(x))
12
```



Explore on your own: <https://numpy.org/devdocs/>

Next Week: Copying 2D Lists

What happens when we run the following code? Why?

```
1 orig = [ [1,2,3], [4, 5, 6] ]
2
3 copy = orig[:]
4
5 orig[1][1]=9
6 copy[0][2]=7
7
8 print(orig)
9 print(copy)
10
11
12
13
14
15
16
17
18
```

