

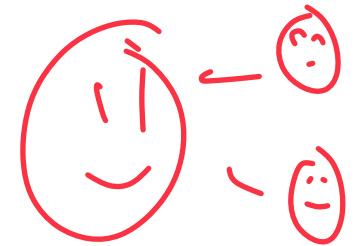
Algorithms and Data Structures for Data Science

Graph Implementations 2

CS 277

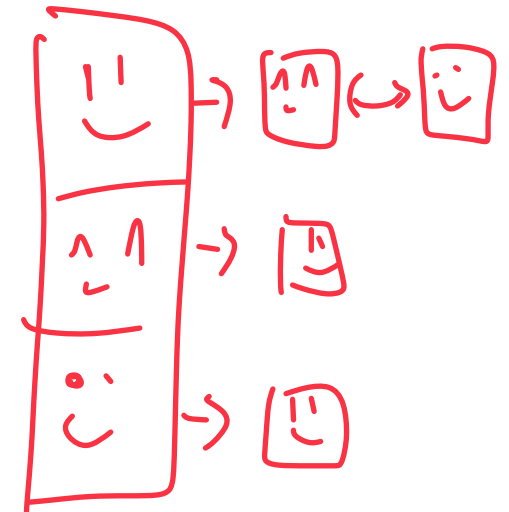
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Department of Computer Science



Learning Objectives

Practice implementing complex data structures (graphs)

Compare and contrast different implementations

Review Big O concepts in the context of graphs

Graph ADT

Find

`getVertices()` — return the list of vertices in a graph

`getEdges(v)` — return the list of edges that touch the vertex v

`areAdjacent(u, v)` — returns a bool based on if an edge from u to v exists

Insert

`insertVertex(v)` — adds a vertex to the graph

`insertEdge(u, v)` — adds an edge to the graph

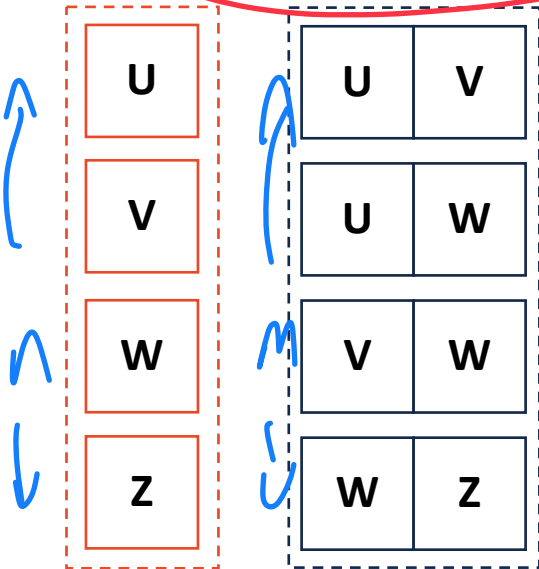
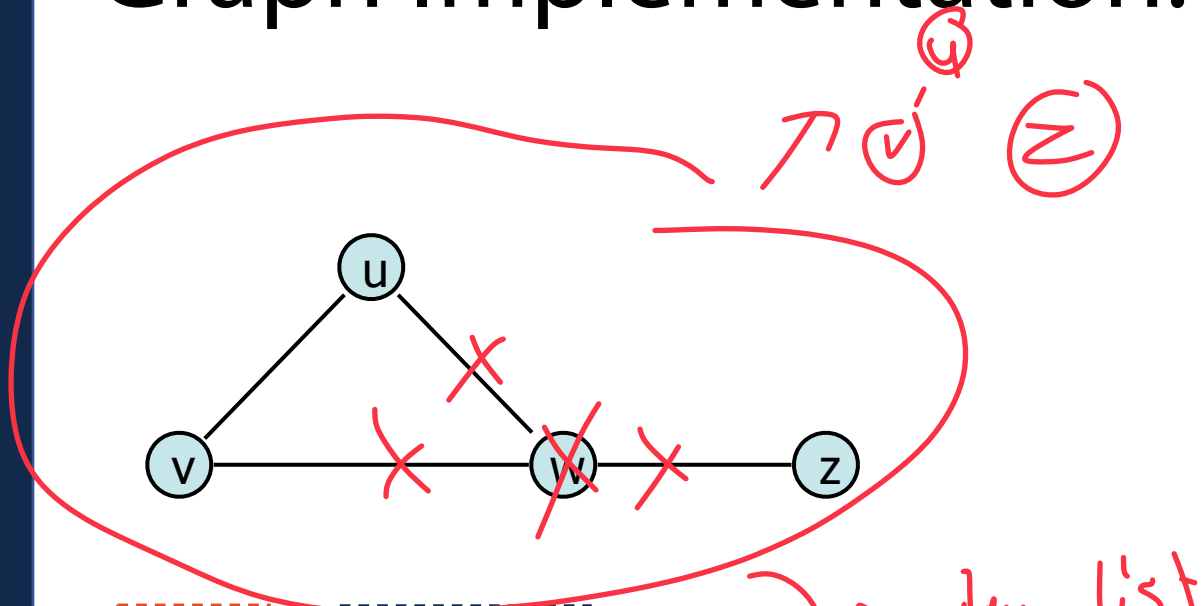
Remove

`removeVertex(v)` — removes a vertex from the graph

`removeEdge(u, v)` — removes an edge from the graph

Graph Implementation: Edge List

$n = |V|, m = |E|$



vertex list
edge list

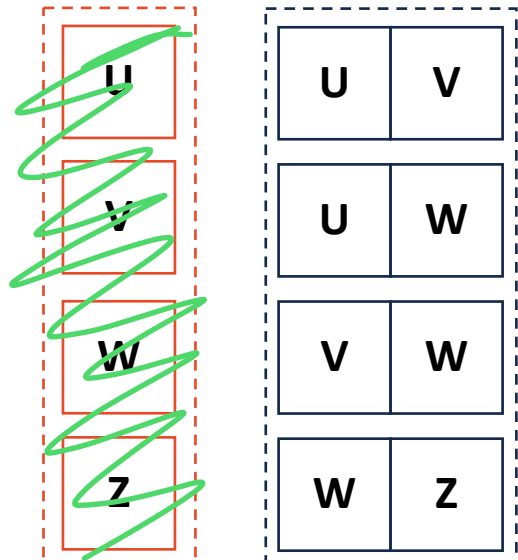
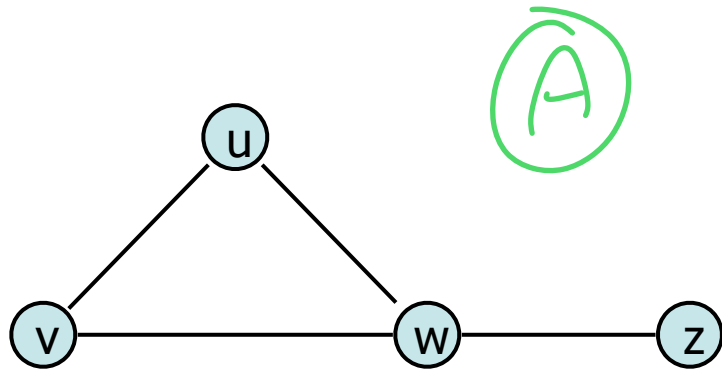
list insert

1) Remove from list twice

list insert

Expressed as O(f)	Edge List
Space	$n+m$
insertVertex(v)	$O(1)^*$
removeVertex(v)	$O(n+m)$
insertEdge(u, v)	$O(1)^*$
removeEdge(u, v)	$O(m)$
getEdges(v)	$O(m)$
areAdjacent(u, v)	$O(m)$

Graph Implementation: Edge List

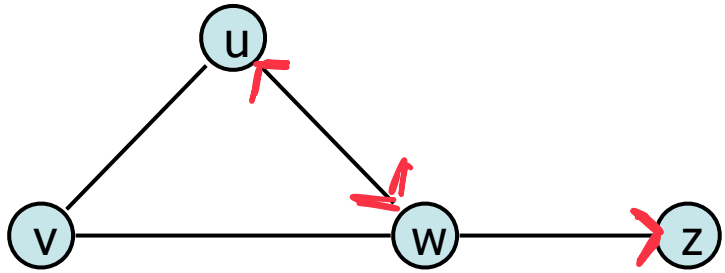


All vertices are in edge list

Expressed as $O(f)$	Edge List
Space	$n+m$
insertVertex(v)	1^*
removeVertex(v)	$n+m$
insertEdge(u, v)	1
removeEdge(u, v)	m
getEdges(v)	m
areAdjacent(u, v)	m

Graph Implementation: Adjacency Matrix

$$[u, v] = [v, u]$$



Vertex Storage:

A dictionary (key = label, value = index)

Edge Storage:

A 2D matrix storing edge existence

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

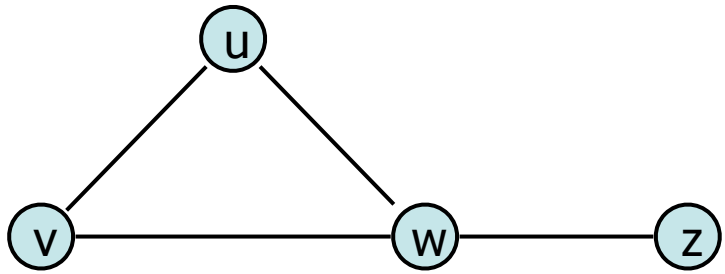
$i = \text{am. vertices ["v"]} \neq 2$

$j = \text{am. vertices ["w"]} \neq 2$

am. edges $[i][j]$ $[j][i]$

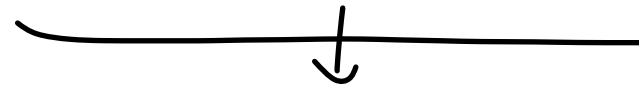
← This is directed edge

Graph Implementation: Adjacency Matrix



getVertices():

```
return self.vertices.keys()
```



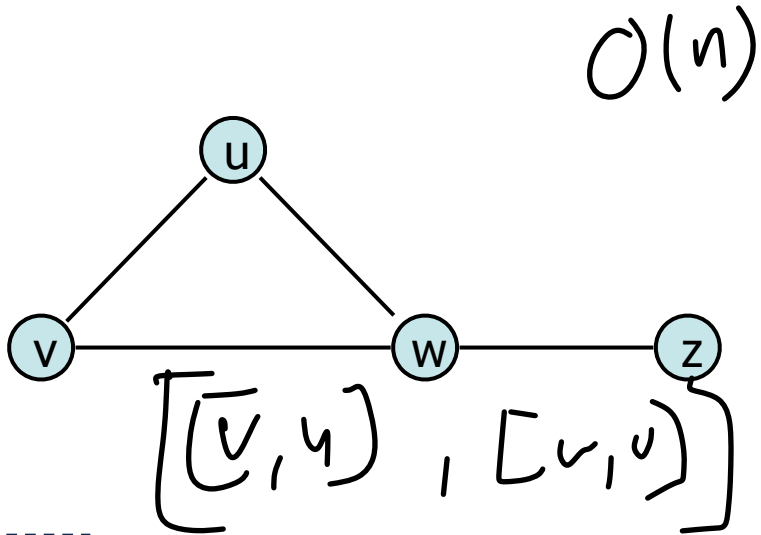
Dictionary.keys() gets list of keys

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

*if x in Dictionary.keys()
[]*

Graph Implementation: Adjacency Matrix



getEdges(v):

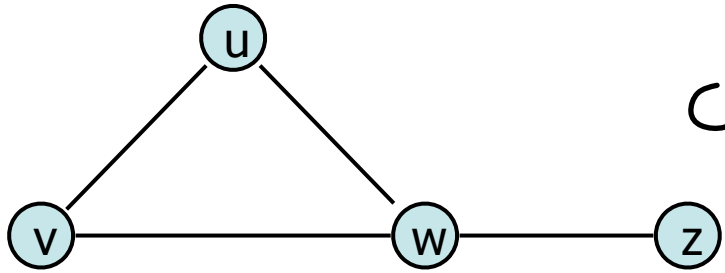
return list of edges
 $[v_1, v_2]$
 $[v_1, v_2], [v_3, v_4], \dots$

u	→	0
v	→	1
w	→	2
z	→	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

- 1) Look up label 'v' in vertex dictionary
 ↳ self. vertices[v] $O(1)^*$
- 2) Look up row or col in matrix
 ↳ self. edges[i] $O(1)^*$
- 3) Loop through all keys, looking up index
 ↳ The edge is $[v, v_2]$ call it v_2

Graph Implementation: Adjacency Matrix



$O(1)$ **areAdjacent(u, v):**

$O(1)$

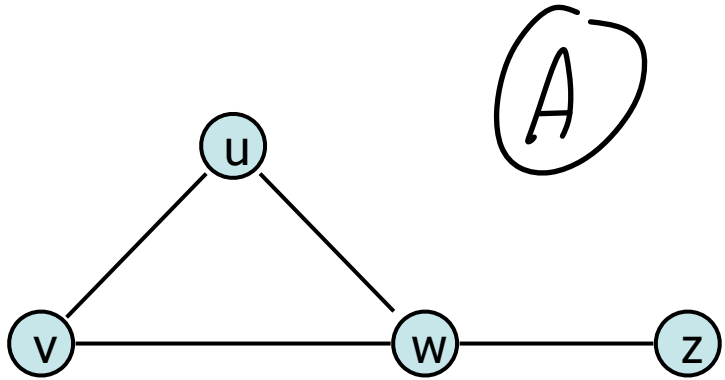
- 1) Look up u in vertex dictionary $\rightarrow i1$
- 2) Look up v in vertex dictionary $\rightarrow i2$

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

- 3) Look up matrix $[i1][i2]$
 - \hookrightarrow if 1, return True
 - \hookrightarrow if 0, return False

Graph Implementation: Adjacency Matrix



A

$O(N)$

$O(1)^*$

insertVertex(v):

1) Add vertex to vertex dictionary

↳ key / value pair
we know value

take dictionary and get length of keys

u	0
v	1
w	2
z	3

	u	v	w	z	A
u	0	1	1	0	0
v	1	0	1	0	0
w	1	1	0	1	0
z	0	0	1	0	0

2) Increase matrix size

↳ Append 0 to every list

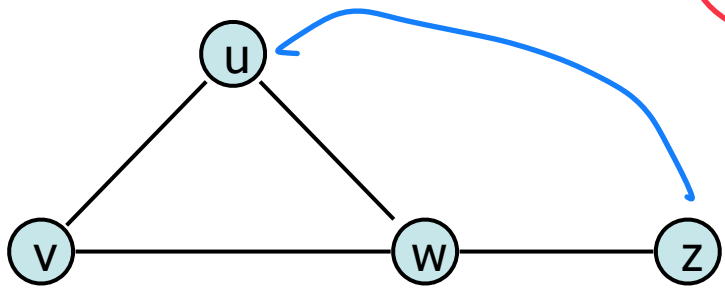
↳ Add a new $n+1$ list to matrix

n items + n items] \rightarrow list insert $O(1)^* \rightarrow O(n)$

$(A) / 4$

(A)

Graph Implementation: Adjacency Matrix



$O(1)$

insertEdge(u, v):

- 1) Look up v $(i2)$ $O(1)$
- 2) Look up u $(i1)$ $O(1)$
- 3) Change value $M[i1][i2] = 1$ $O(1)$
 $M[i2][i1] = 1$

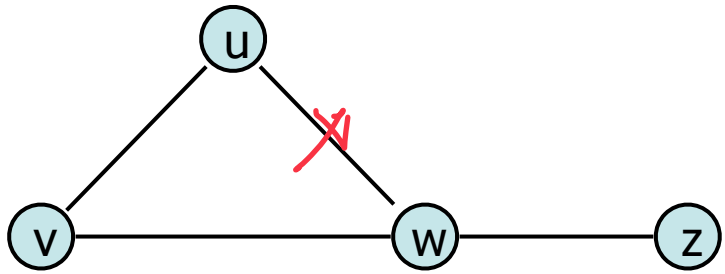
u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

$O(1)$ if no new vertex

-
- if new vertex $O(n)$
 - 1) insert vertex (u) $O(n)$
 - 2) insert vertex (v)
 - 3) insertEdge()
- $O(n)$

Graph Implementation: Adjacency Matrix



removeVertex(v):

↳ Messy! List removal $\times N$

u	0
v	1
w	2
x	3

	u	v	w	z
u	0	1	1 ⁰	0
v	1	0	1	0
w	1 ⁰	1	0	1
z	0	0	1	0

removeEdge(u, v):

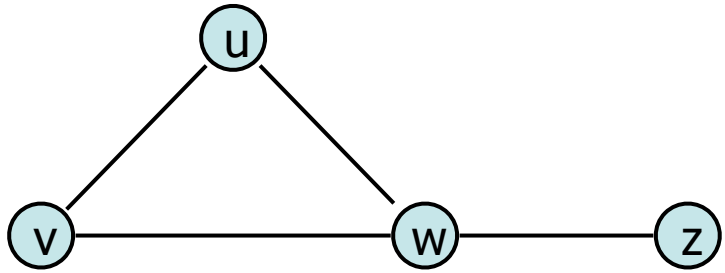
↳ Set $M[u][v] = 0$

$O(1)$

$M[v][u] = 0$

1D version: $X[u] = 0$

Graph Implementation: Adjacency Matrix



Pros: 1) Lookup is easy $\leftarrow O(1)$
2) Very good at changing edges

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

Cons: Adding ^{or removing} vertices is hard
is very inefficient

Its also very large
 N^2

Graph Implementations

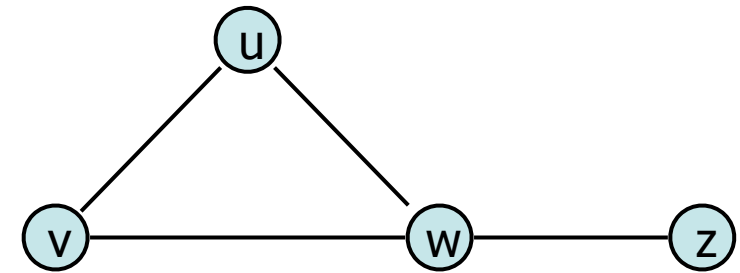
We want something...

Faster than an edge list

Less space than an adjacency matrix

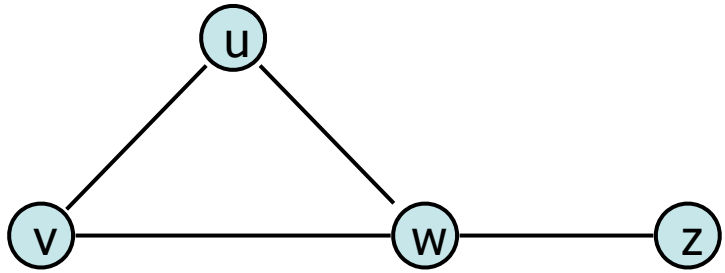
Particularly good at finding adjacent elements

+ better adding vertices to ~~the~~ graph



Graph Implementation: Edge List + ?

$$|V| = n, |E| = m$$



Vertex x

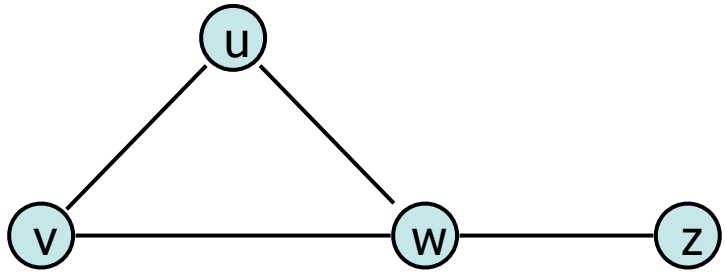
u
v
w
z

Edge List

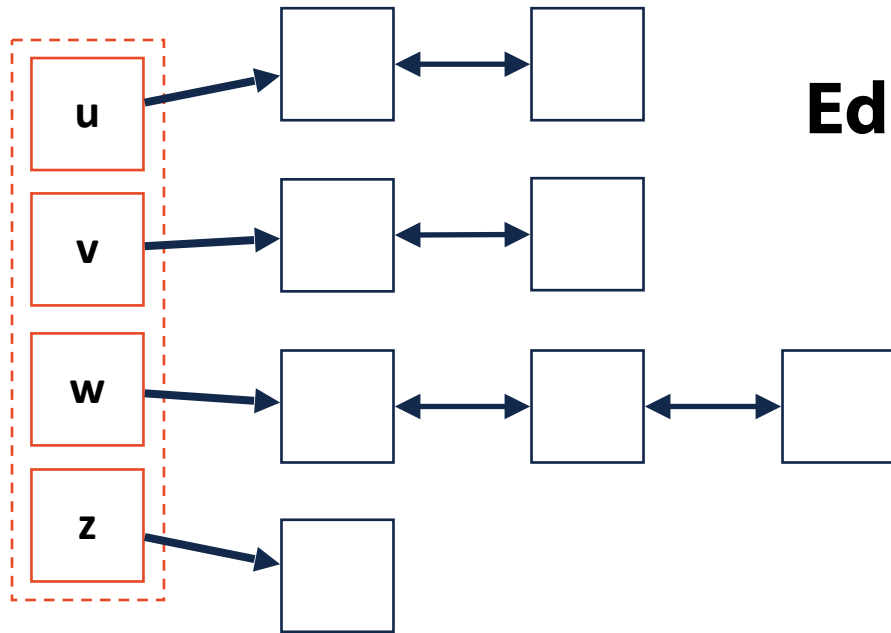
u	v
v	w
u	w
w	z

Why can't vertex
point to edges it has?

Adjacency List

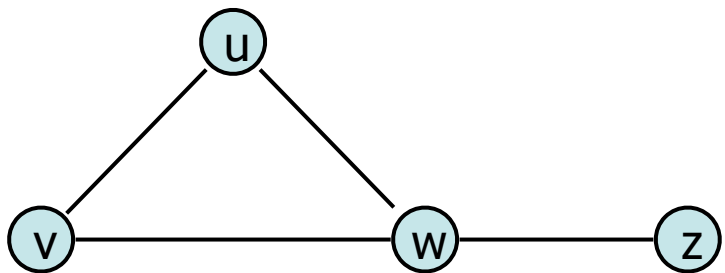


Vertex Storage:

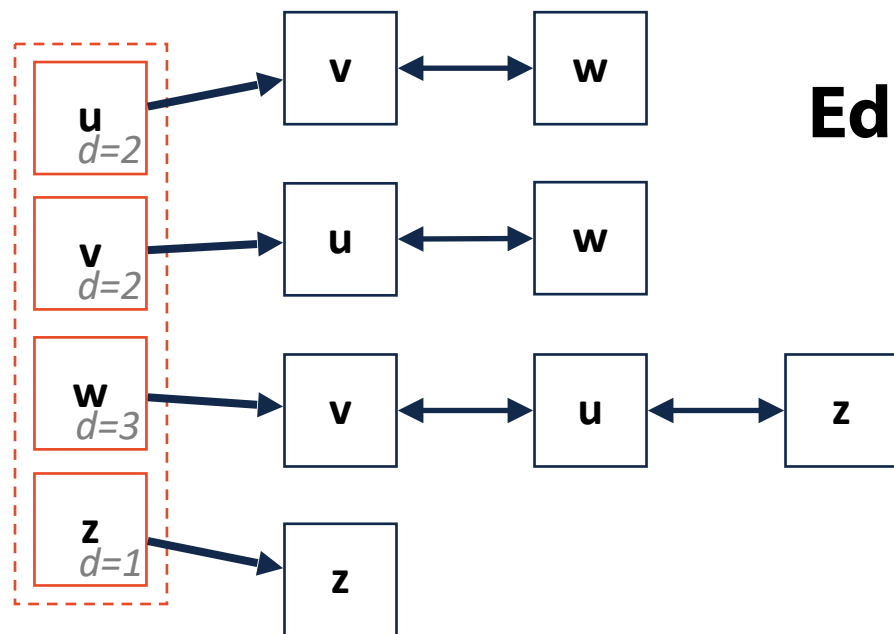


Edge Storage:

Adjacency List

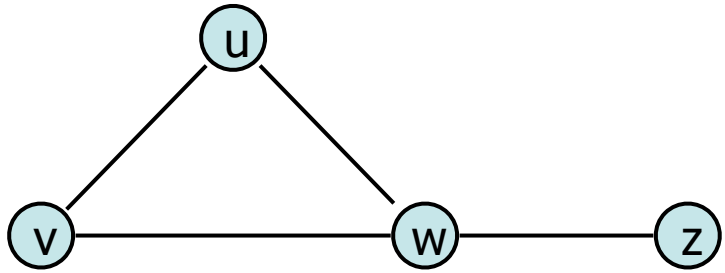


Vertex Storage:

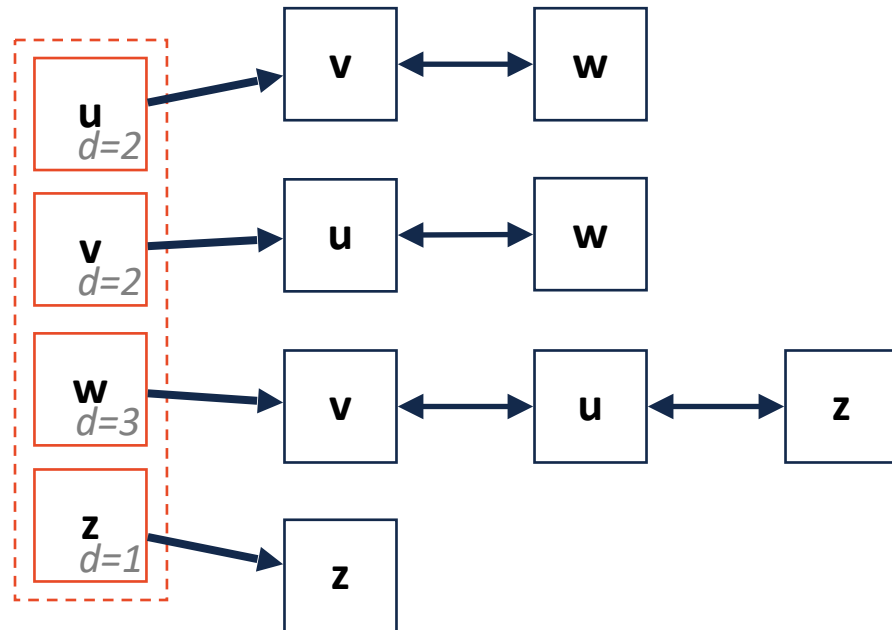


Edge Storage:

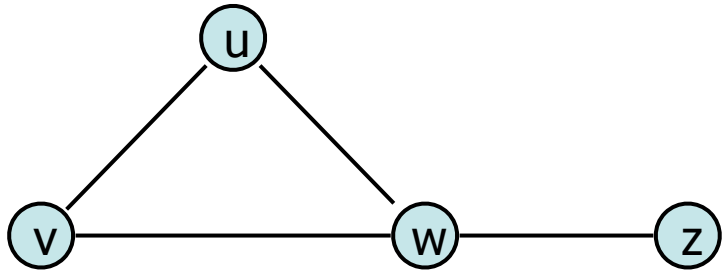
Adjacency List



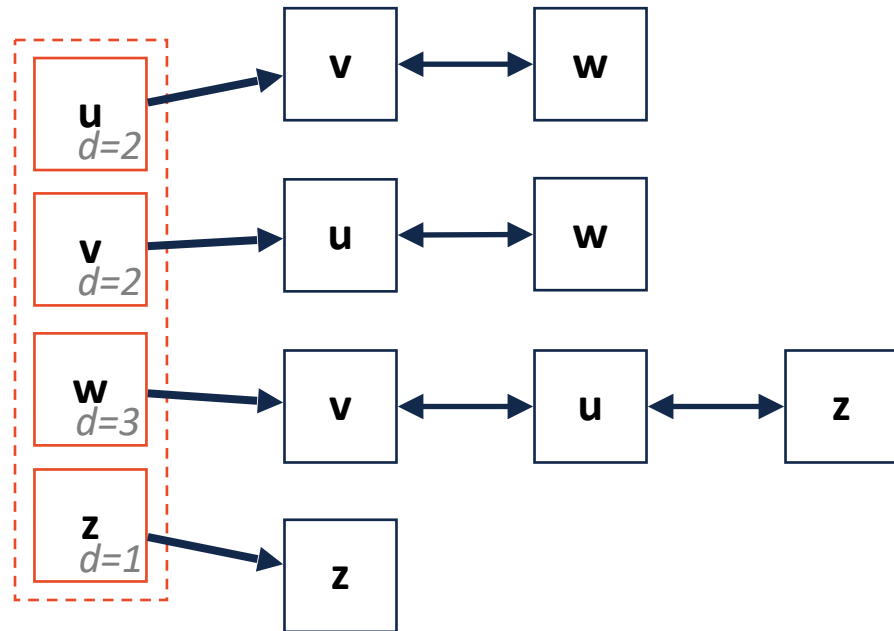
getVertices():



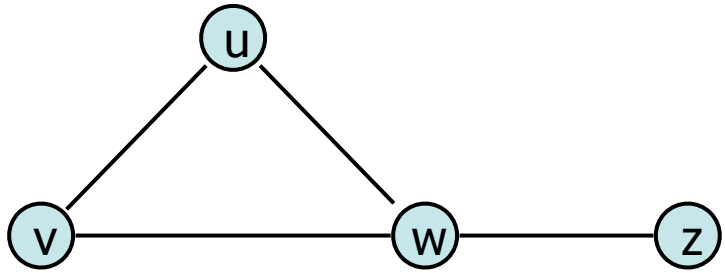
Adjacency List



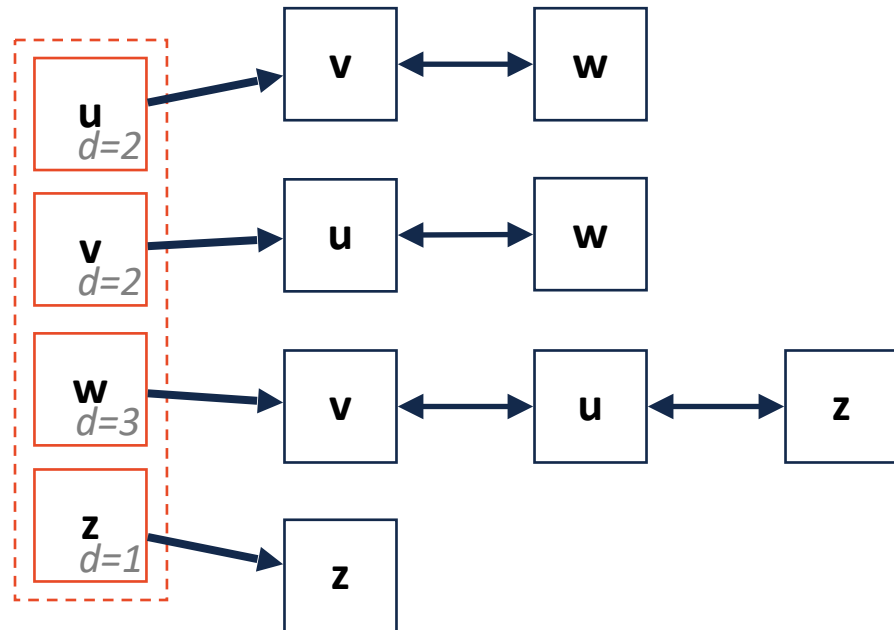
getEdges(v):



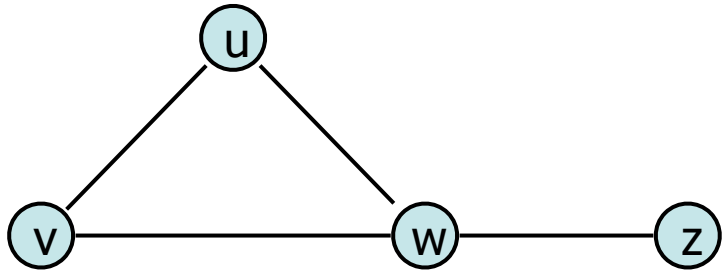
Adjacency List



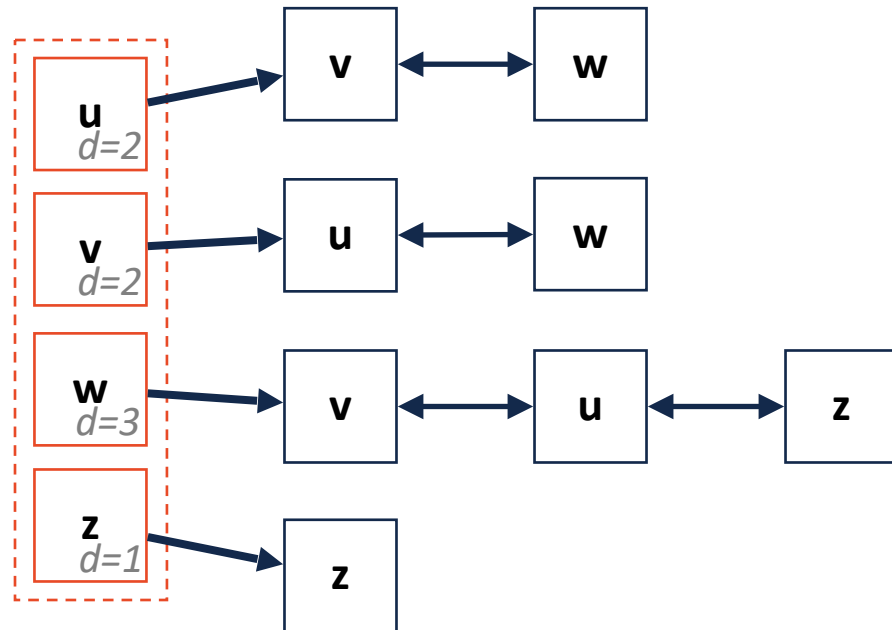
areAdjacent(u, v):



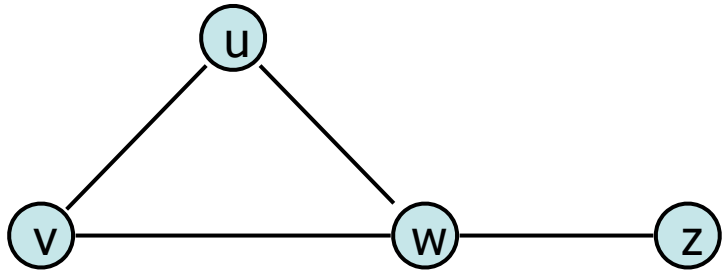
Adjacency List



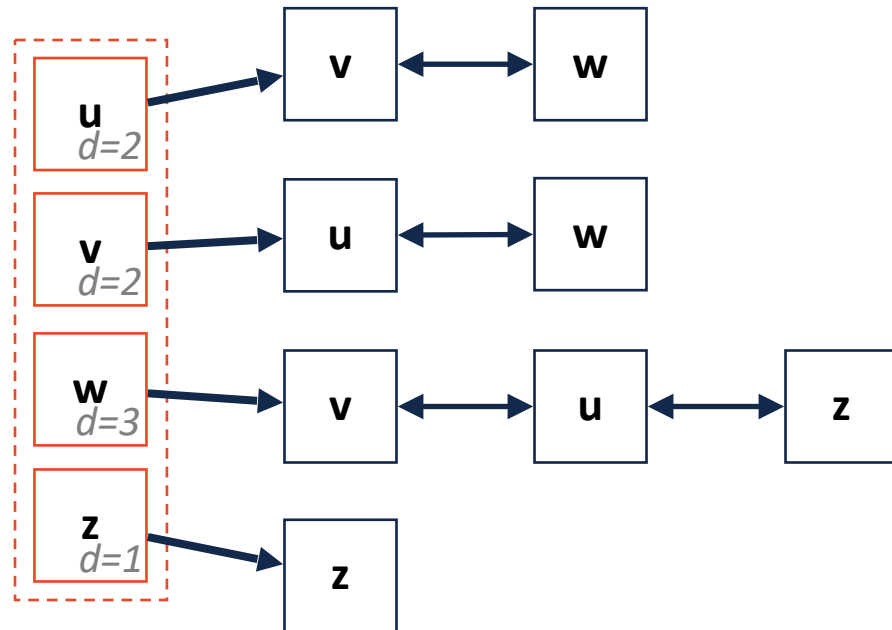
insertVertex(v):



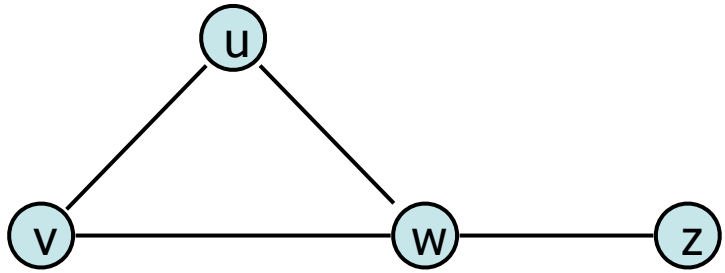
Adjacency List



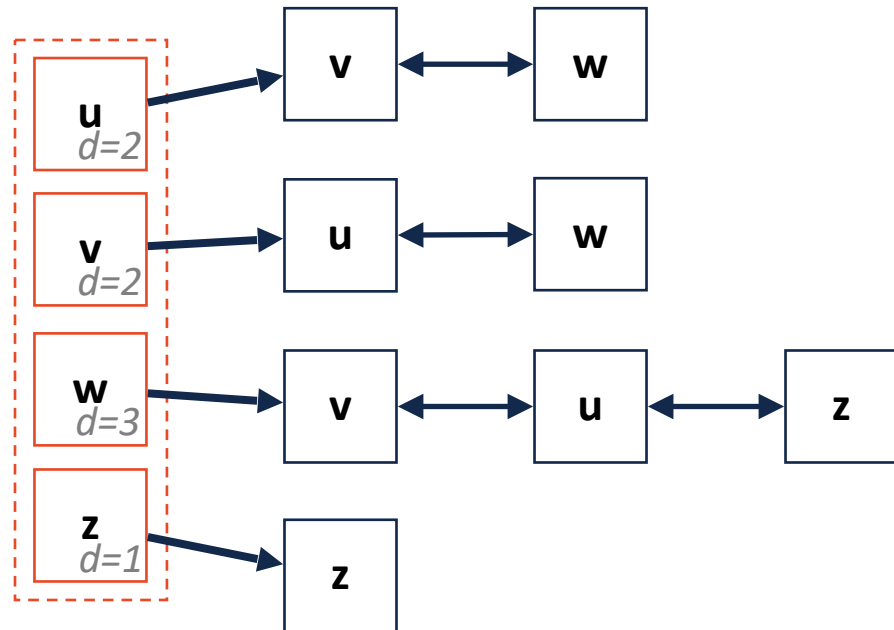
removeVertex(v):



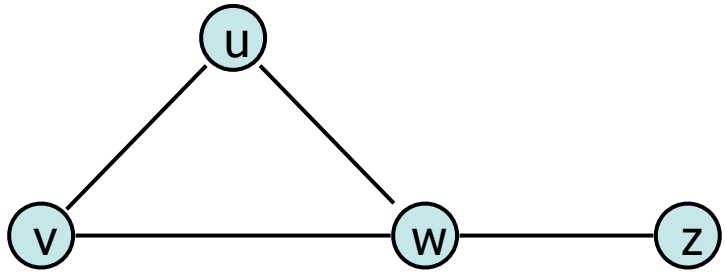
Adjacency List



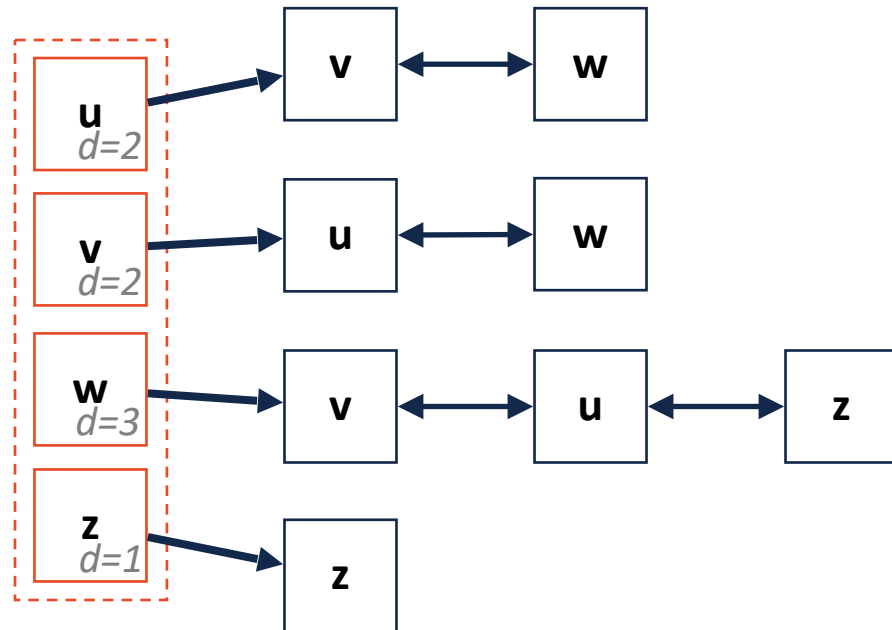
insertEdge(u, v):



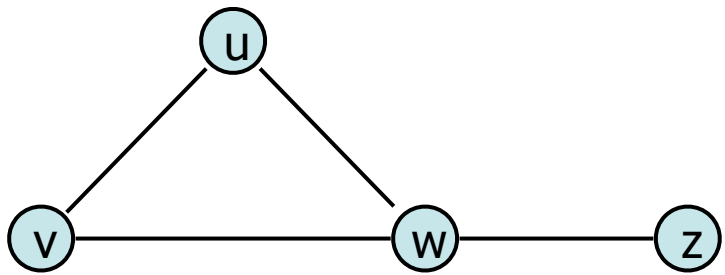
Adjacency List



removeEdge(u, v):

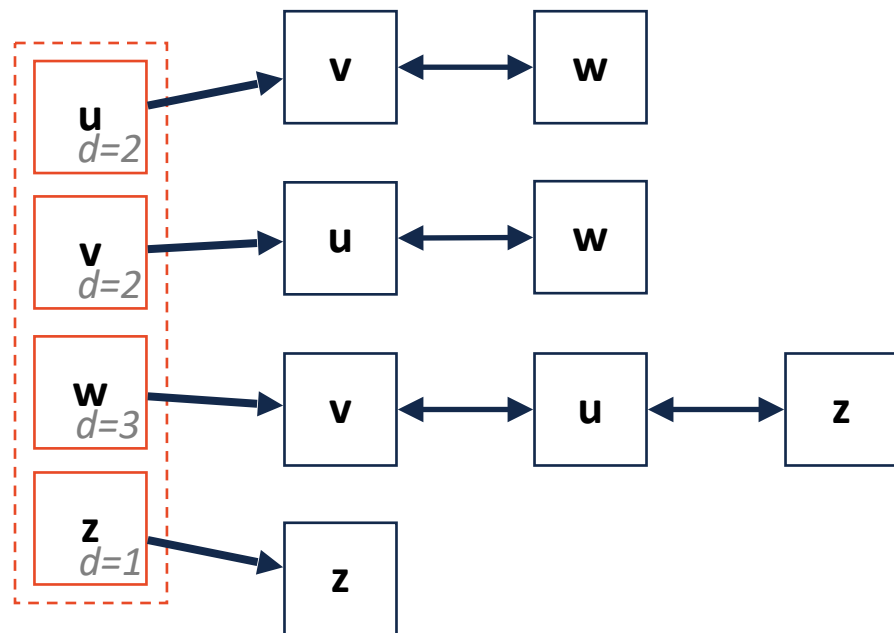


Adjacency List



Pros:

Cons:



$$|V| = n, |E| = m$$

Expressed as O(f)	Edge List	Adjacency Matrix	Adjacency List
Space			
insertVertex(v)			
removeVertex(v)			
insertEdge(u, v)			
removeEdge(u, v)			
getEdges(v)			
areAdjacent(u, v)			