

Algorithms and Data Structures for Data Science

Graph Implementations

CS 277

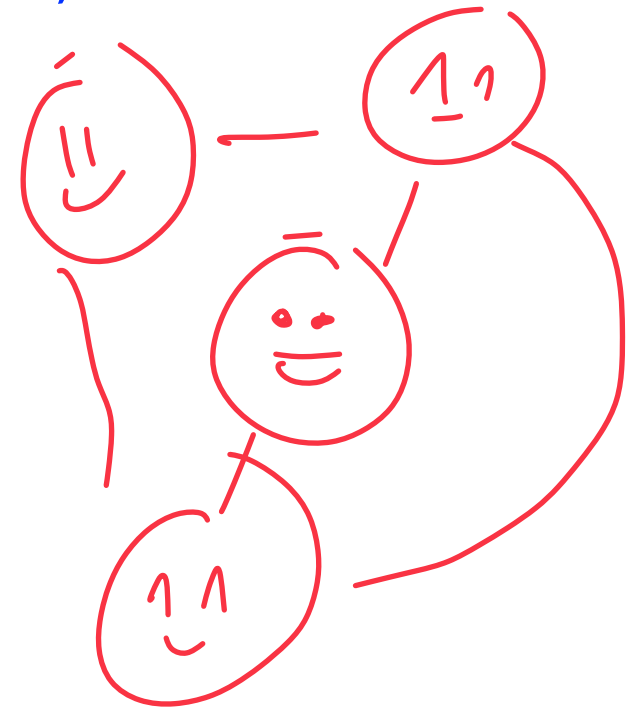
Brad Solomon

March 25, 2024



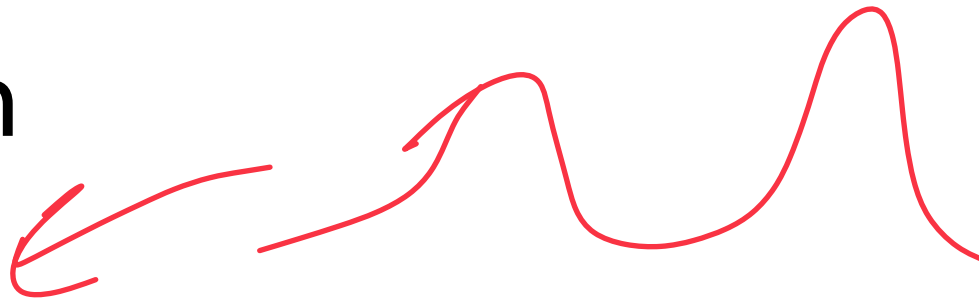
UNIVERSITY OF
ILLINOIS
URBANA - CHAMPAIGN

Department of Computer Science



Exam 2 Reflection

Average: 72%



The programming question was the most difficult

I also consider the programming question to be extremely fair

Consider going to office hours to go over exam

retake exam
retake exam

Learning Objectives

Practice implementing complex data structures (graphs)

Compare and contrast different implementations

Review Big O concepts in the context of graphs



Graph ADT

Constructor

Find: Need to be able to search for vertices, edges, and adjacency.

Insert: Need both a vertex and an edge insertion function

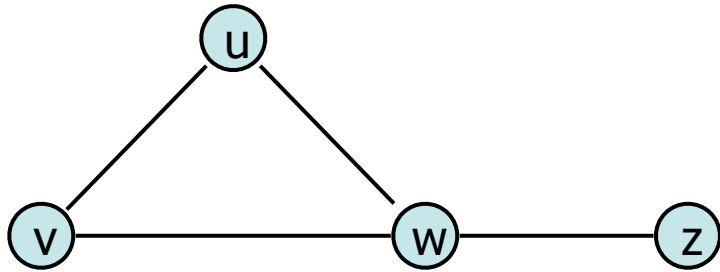
Remove: Need both a vertex and an edge removal function

Traversal: Need to be able to traversal a graph efficiently



Graph Implementation: Edge List $|V| = n, |E| = m$

The equivalent of an 'unordered' data structure



Vertex Storage:

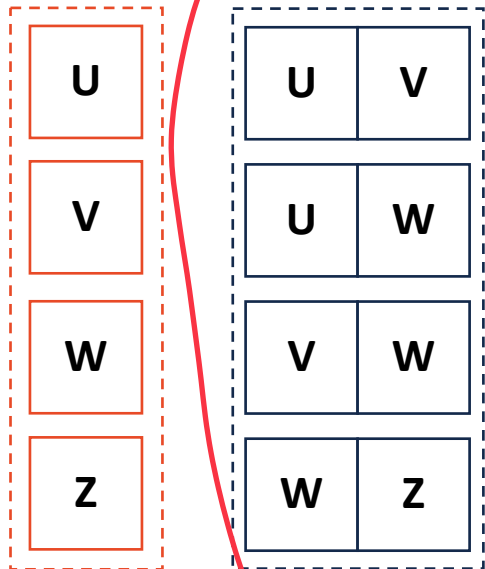
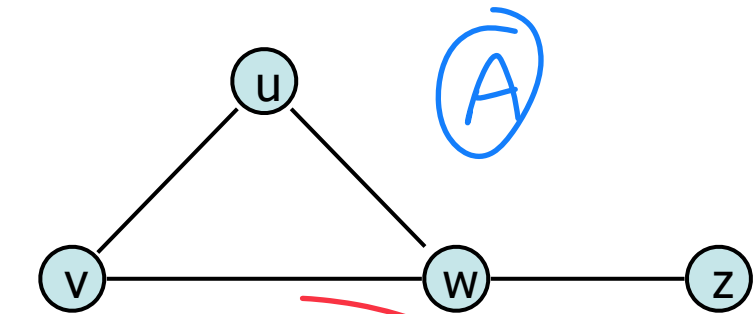
v	w
u	u

Edge Storage:

(list of)
↳ pairs of vertices

Graph Implementation: Edge List $|V| = n, |E| = m$

The equivalent of an 'unordered' data structure



Vertex Storage:

None

→ all vertices in edge list

A list of vertex labels

we will keep vertex list

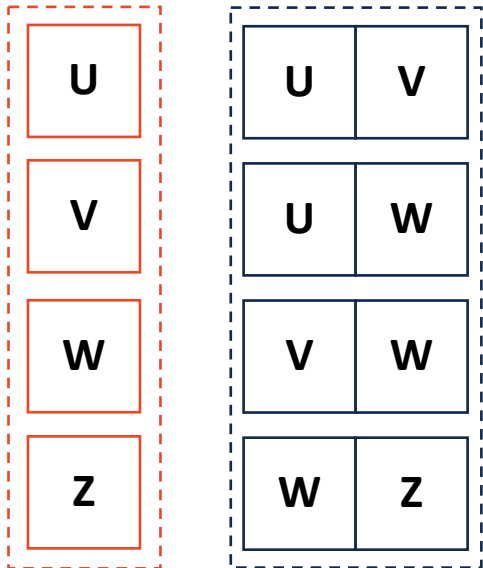
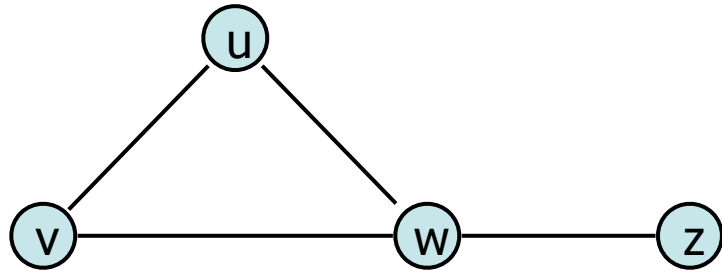
Edge Storage:

A list of paired vertex labels

Graph Implementation: Edge List

getVertices()

↳ return vertex list



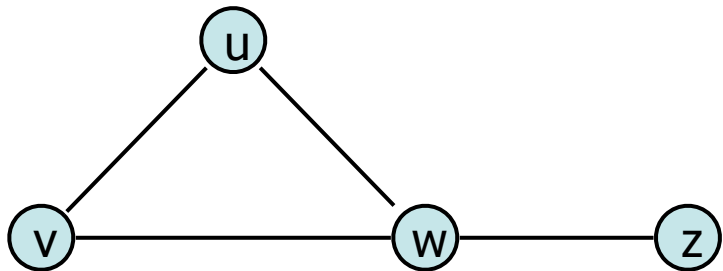
Graph Implementation: Edge List

$F_{in}(l)$

getVertices()

Return vertex list

~~Loop through edge list and build list of vertices~~



u	u	v
v	u	w
w	v	w
z	w	z

getEdges(w) → [(u,w), (v,w), (w,z)]

↳ loop through edge list

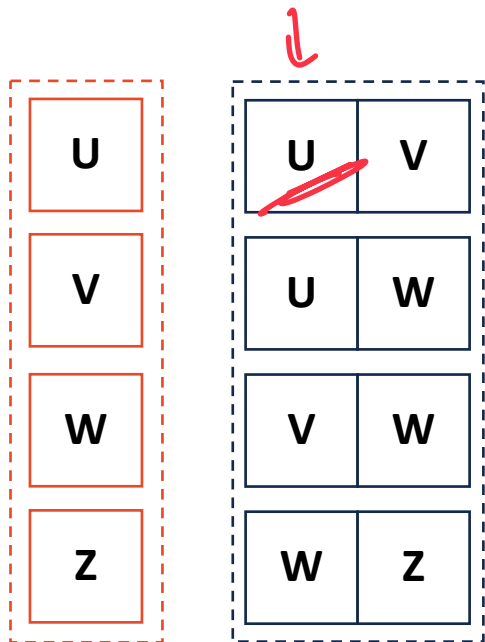
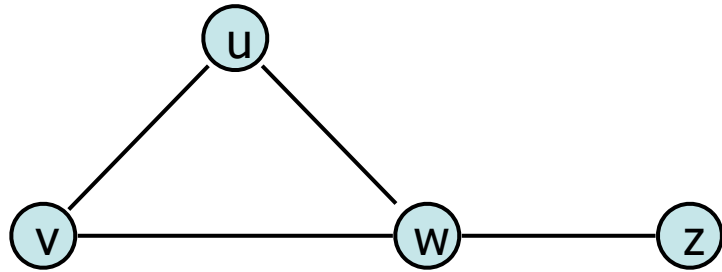
x = [(u,v), (u,w), (v,w), (w,z)]

x[0] → (u,v)

↳ x[0][0] ≠ 'w' & x[0][1] ≠ 'w'

Out. append[(u,w)] b/c x[1][1] = w

Graph Implementation: Edge List



getVertices()

Return vertex list

~~Loop through edge list and build list of vertices~~

getEdges(v)

Loop through edge list and build list of edges

areAdjacent(u, v)

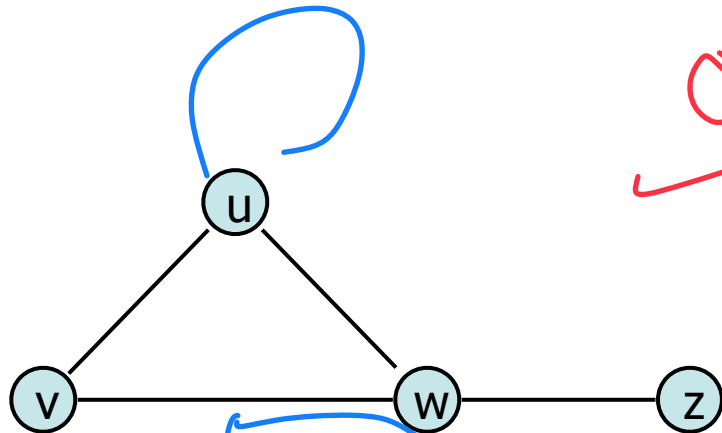
↳ getEdges(u) & getEdges(z)

if exists in both

Loop through EL and search for (u,z) or (z,u)

Graph Implementation: Edge List

Find



$O(1)$ **getVertices()**

Return vertex list

~~Loop through edge list and build list of vertices~~

getEdges(v)

Loop through edge list and build list of edges

areAdjacent(u, v)

Loop through edge list and find (u, v) or (v, u)

u	u	v
v	u	w
w	v	w
z	w	z

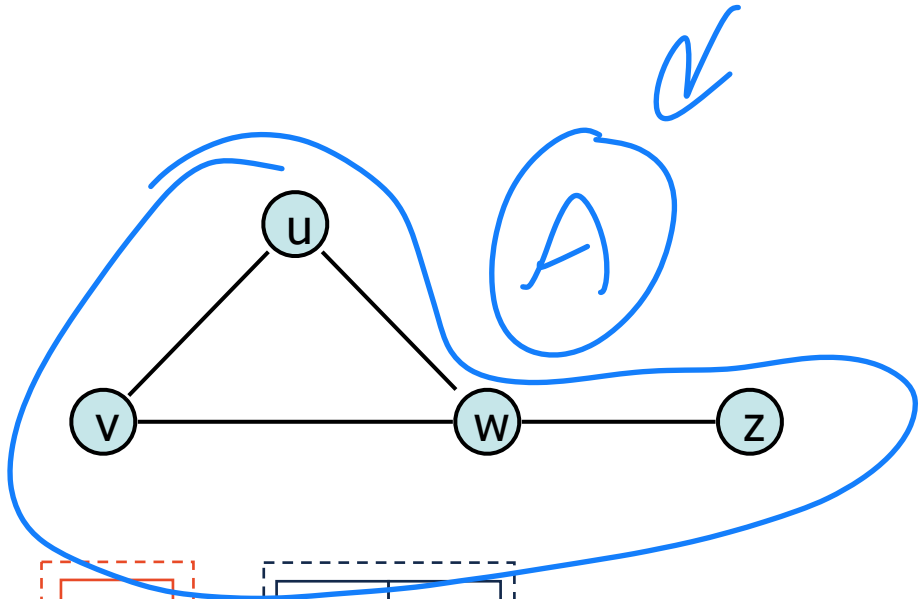
$O(m)$

Lets code this up!

$\hookrightarrow O(m^2)$

Graph Implementation: Edge List

Insert



insertVertex(v)

↳ list.append(v)
 ↑
 vertex list

u	u	v
v	u	w
w	v	w
z	w	z

A

insertEdge(u, v)

edge list . append (u, v)

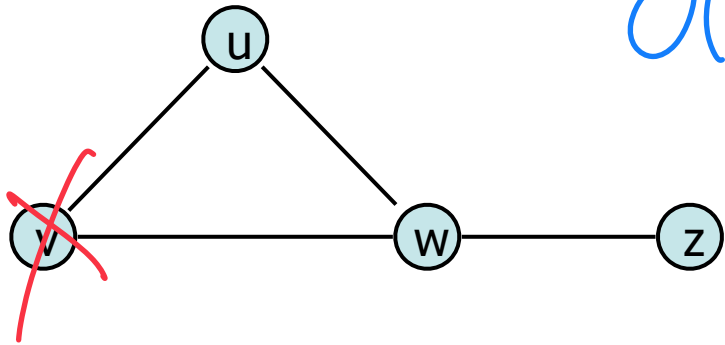
if u or v not in vertex list
 ↳ update (add) to vertex list

$O(1)$
 ↑
 as good

$O(1)$
 +
 $O(1)$

 $O(1)$

Graph Implementation: Edge List



removeVertex(v)

$O(n+m)$

↳ Search vertex list & remove $O(n)$
 (List Remove)

↳ Search edge list & remove

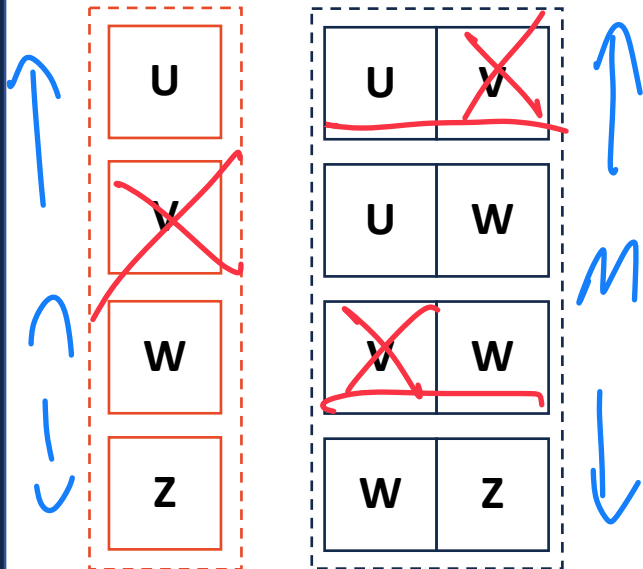
removeEdge(u, v)

↳ Search for loop through edge list

$O(m)$

(List Remove)

↳ maybe multiple times

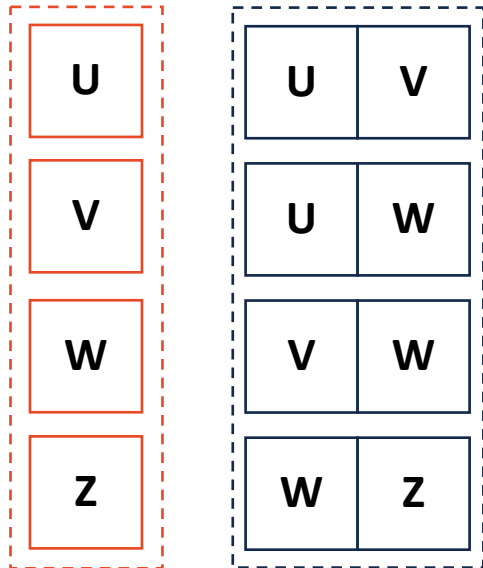
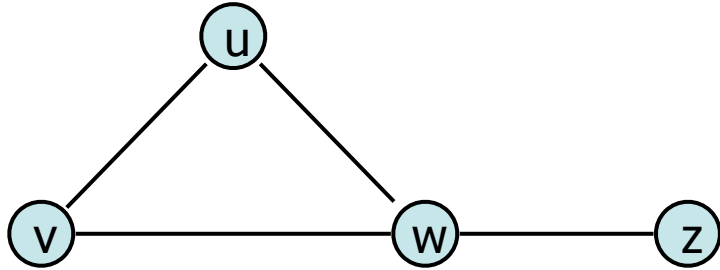




Graph Implementation: Edge List

Pros: Insert very efficient

Easy to implement

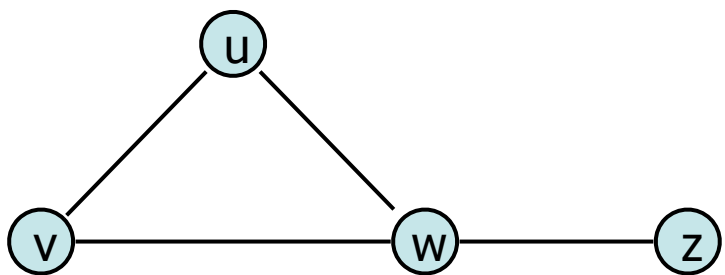


Cons:

↳ Doesn't track duplicates
(our version doesn't)

↳ Find / Remove is slow

Graph Implementation: Adjacency Matrix



Vertex Storage:

↳ as a dictionary / key - Vertex label / value

Draw arrow from label to row/col

Edge Storage:

↳ Store edges as a matrix (of bools)

↑ or ints

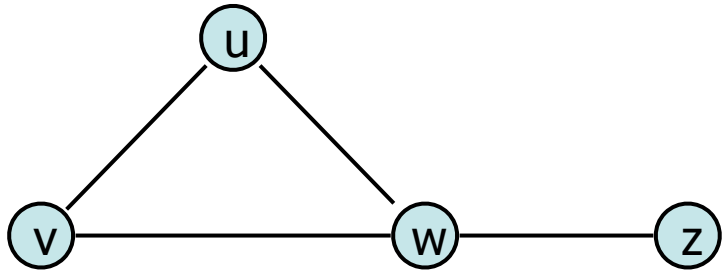
Vertex Dictionary

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	T	T	F
v		0	T	F
w	T	T	0	T
z	T	F	T	0

↑ Keys

Graph Implementation: Adjacency Matrix



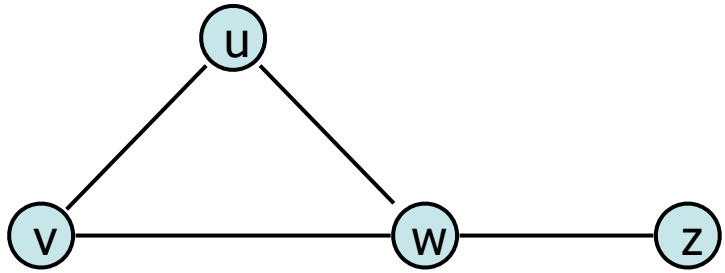
Vertex Storage:

Edge Storage:

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

Graph Implementation: Adjacency Matrix



getVertices():

↳ Dictionary, keys()

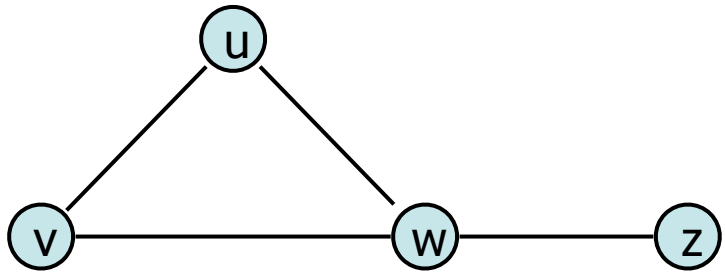
↳ unordered

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

Graph Implementation: Adjacency Matrix

$O(n^2)$



getEdges(v):

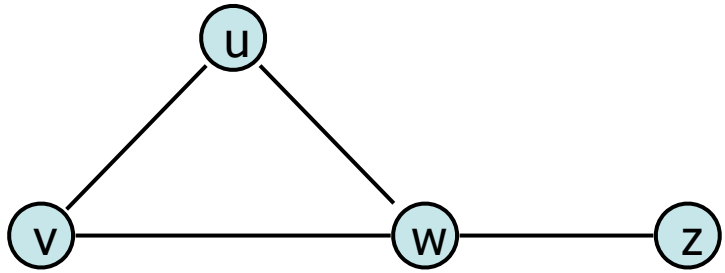
- 1) Look up value for "v" in dictionary
↳ that value is index in matrix
- 2) Look up row or col in matrix
- 3) Return list of edges where $(u,v) = 1$
 $x[u][v] = 1$

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

$\left[\begin{matrix} 1 \\ 1 \\ 1 \\ 1 \end{matrix} \right]$ or $x[i]$
 $x[0][i], x[1][i], \dots$

Graph Implementation: Adjacency Matrix

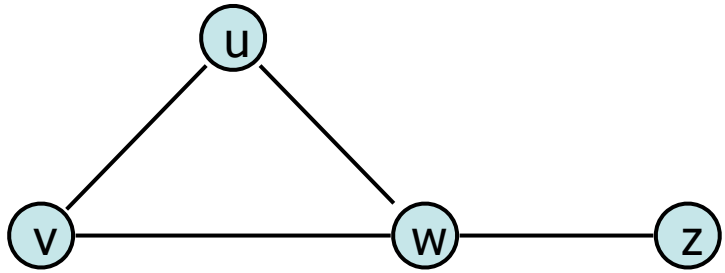


areAdjacent(u, v):

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

Graph Implementation: Adjacency Matrix

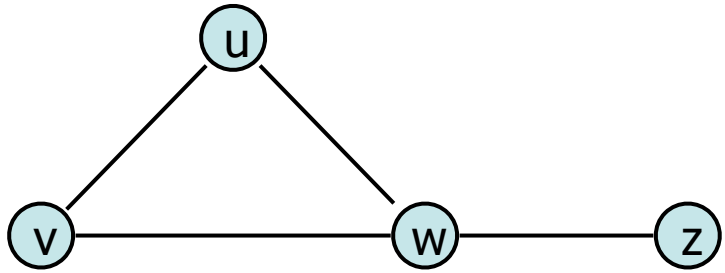


insertVertex(v):

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

Graph Implementation: Adjacency Matrix

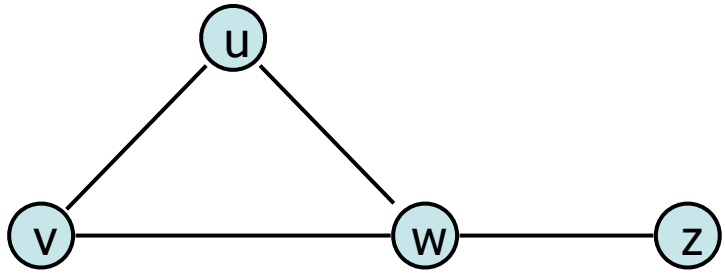


insertEdge(u, v):

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

Graph Implementation: Adjacency Matrix



removeVertex(v):

u	0
v	1
w	2
z	3

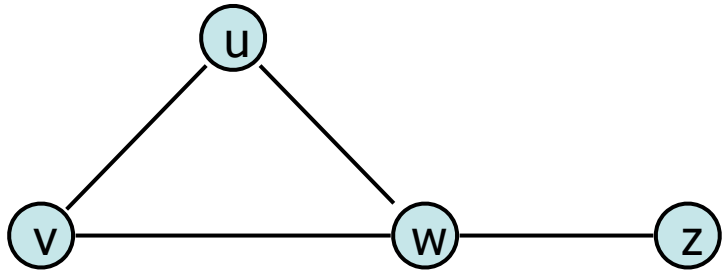
	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

removeEdge(u, v):

Graph Implementation: Adjacency Matrix



Pros:



Cons:

u	0
v	1
w	2
z	3

	u	v	w	z
u	0	1	1	0
v	1	0	1	0
w	1	1	0	1
z	0	0	1	0

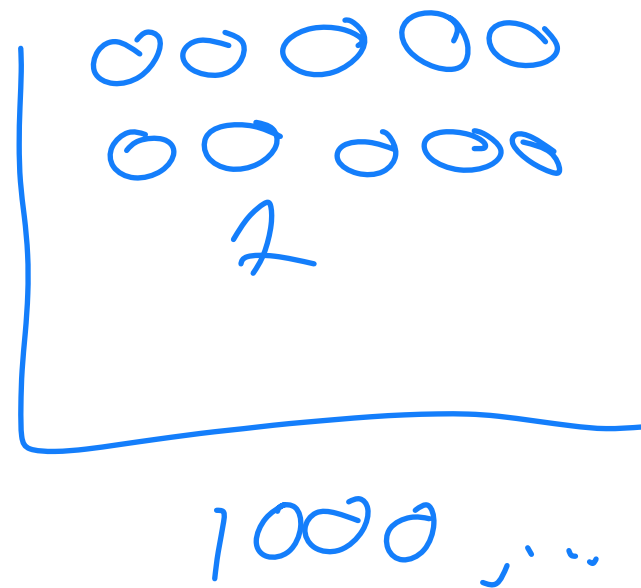
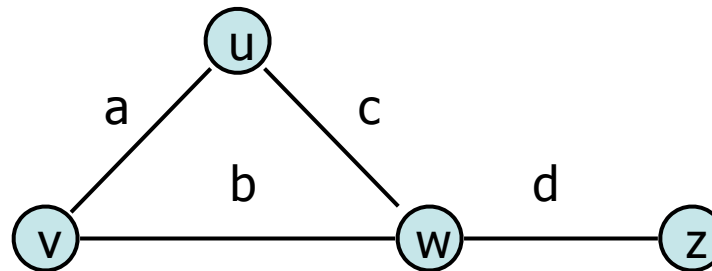
Graph Implementations

We want something...

Faster than an edge list

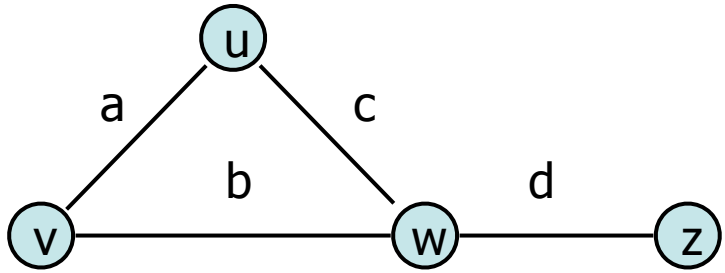
Less space than an adjacency matrix

Particularly good at finding adjacent elements



Graph Implementation: Edge List + ?

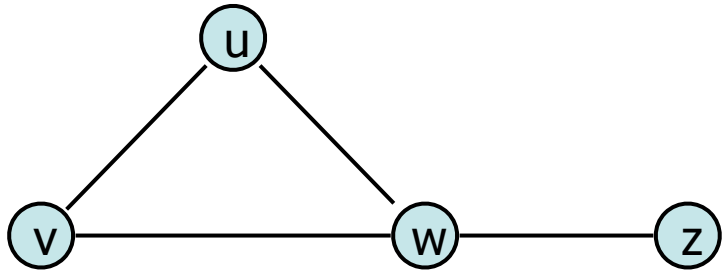
$$|V| = n, |E| = m$$



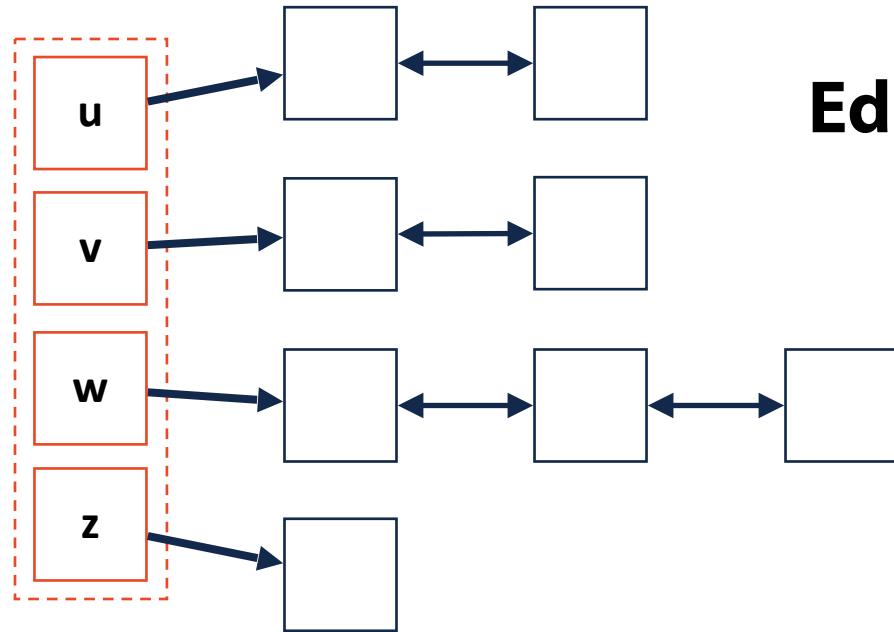
u
v
w
z

u	v	a
v	w	b
u	w	c
w	z	d

Adjacency List

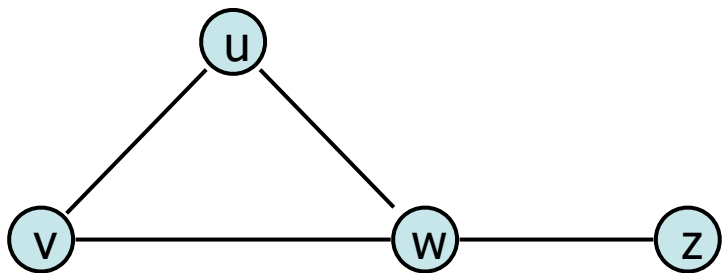


Vertex Storage:

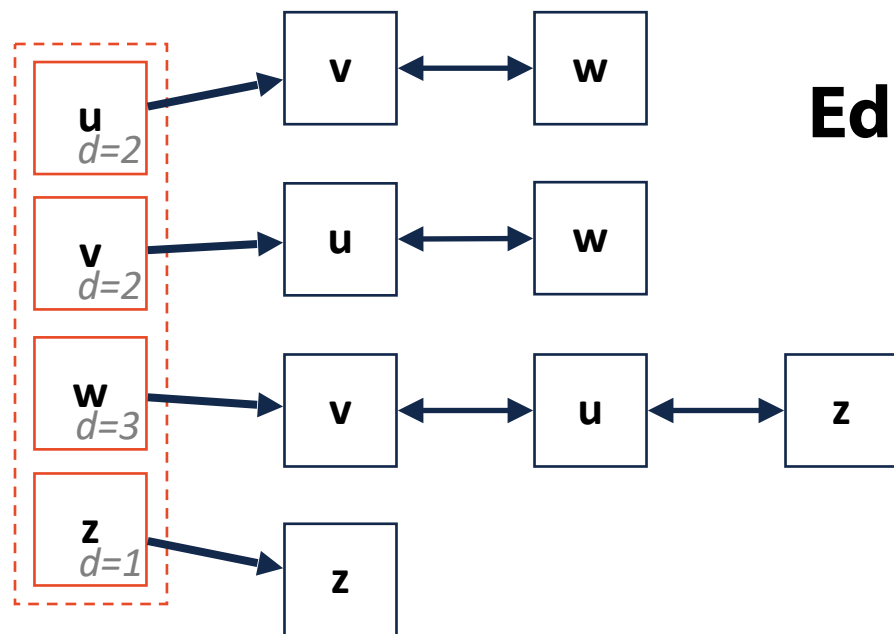


Edge Storage:

Adjacency List

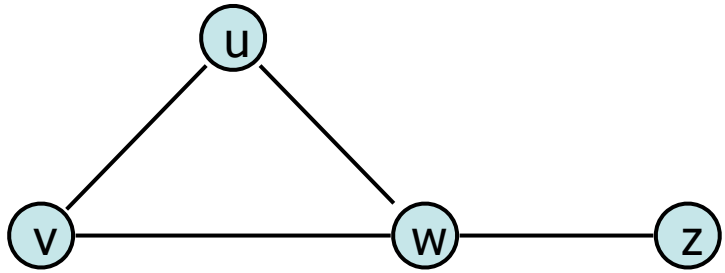


Vertex Storage:

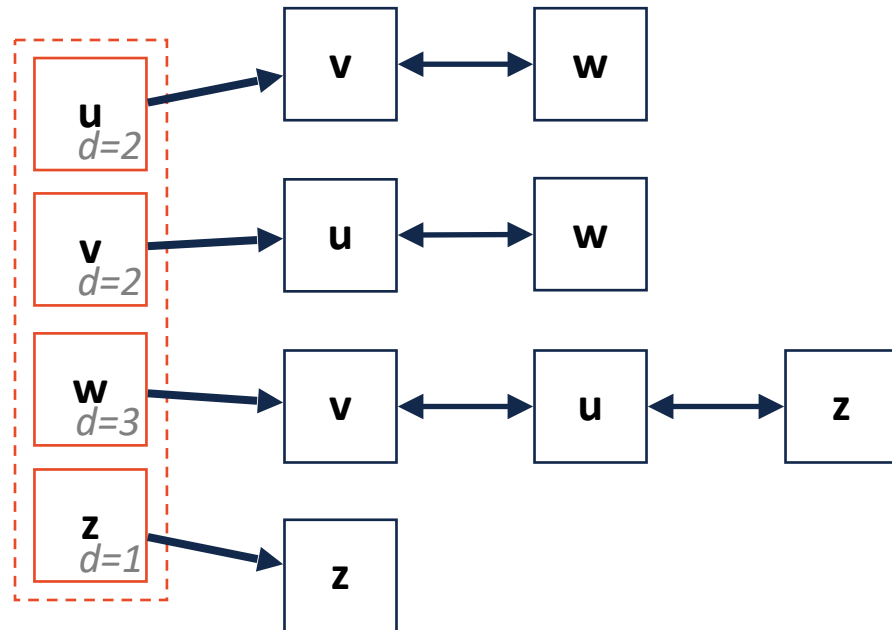


Edge Storage:

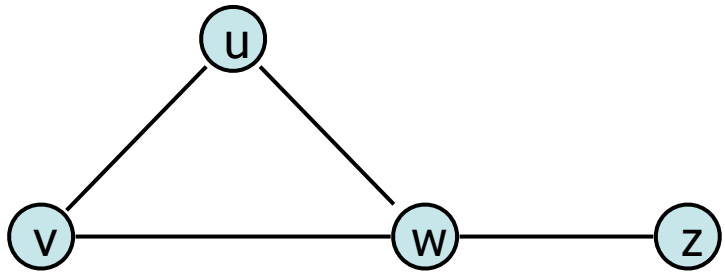
Adjacency List



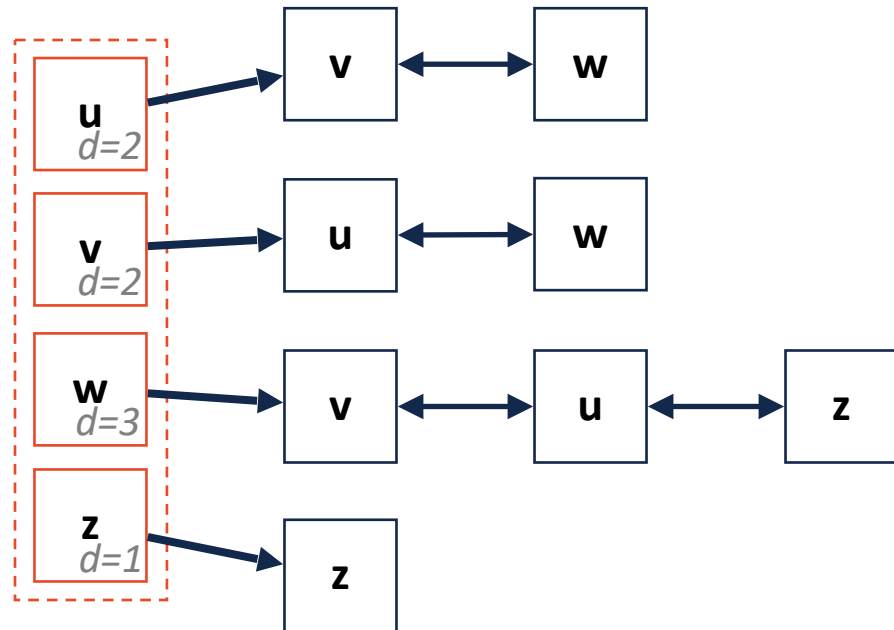
getVertices():



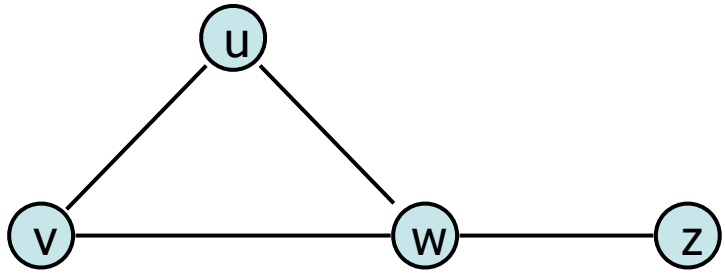
Adjacency List



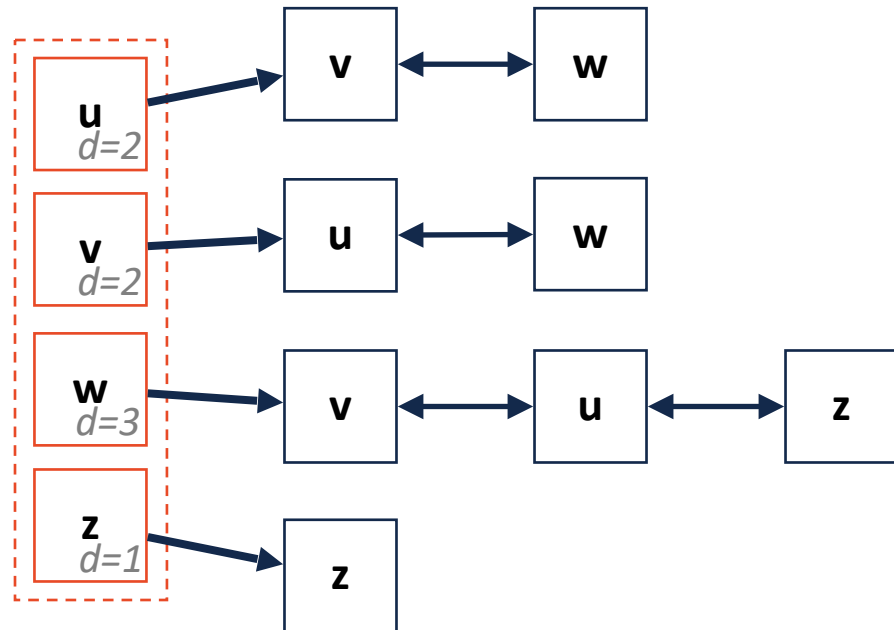
getEdges(v):



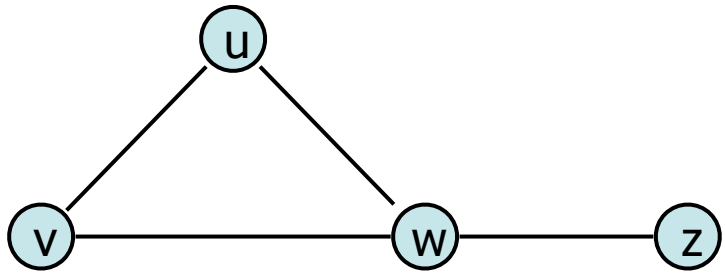
Adjacency List



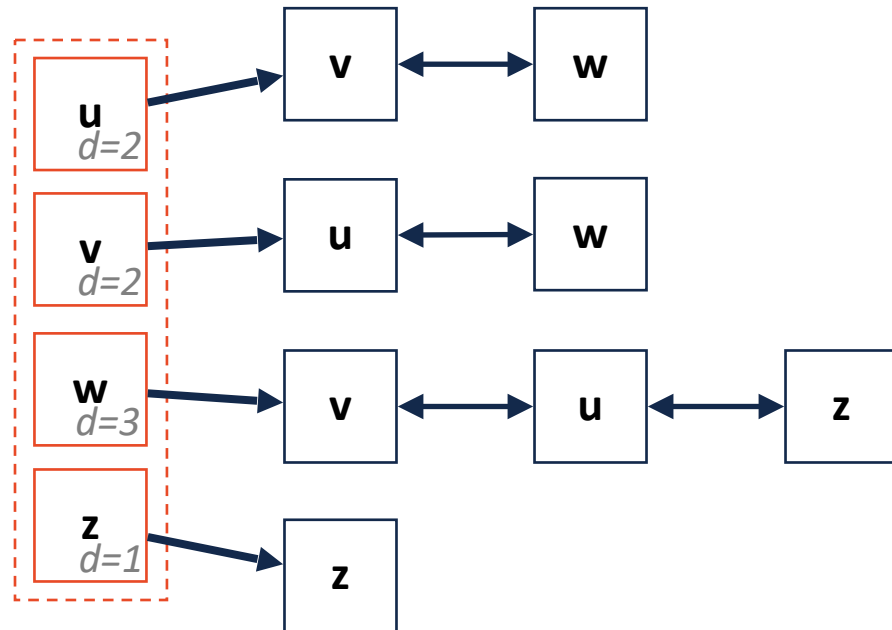
areAdjacent(u, v):



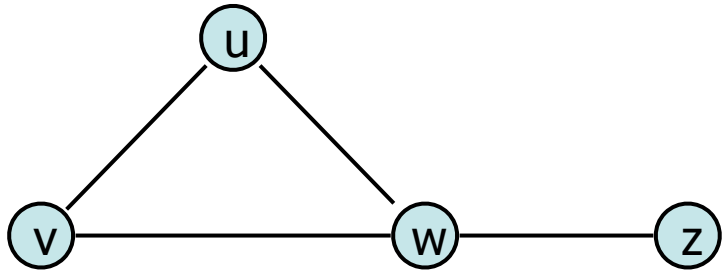
Adjacency List



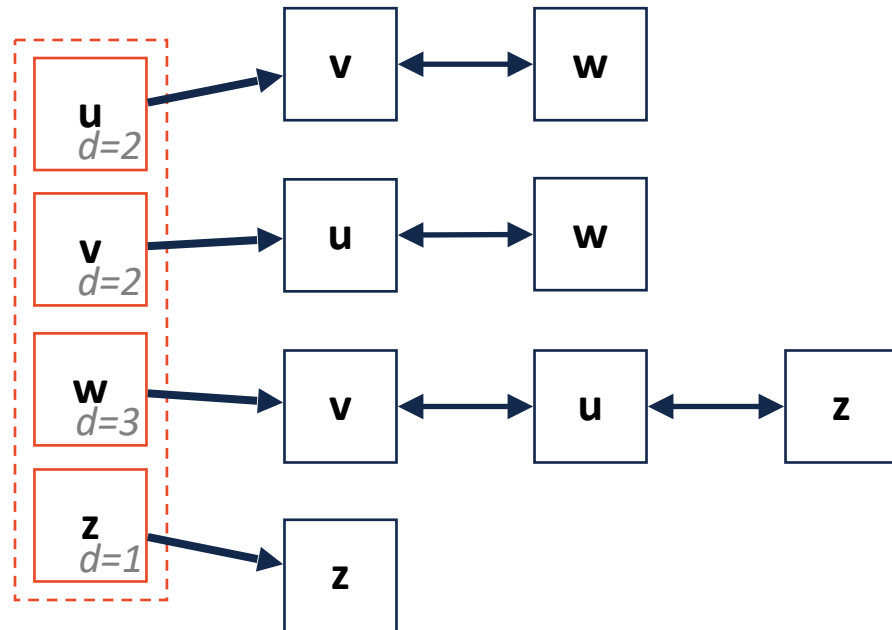
insertVertex(v):



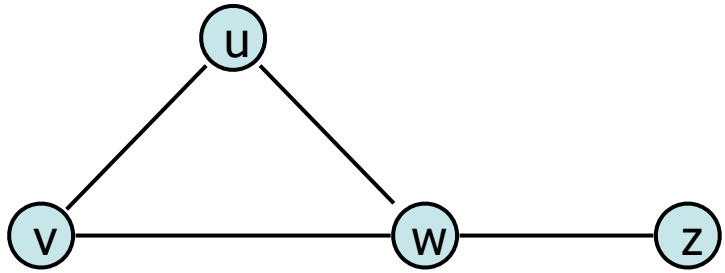
Adjacency List



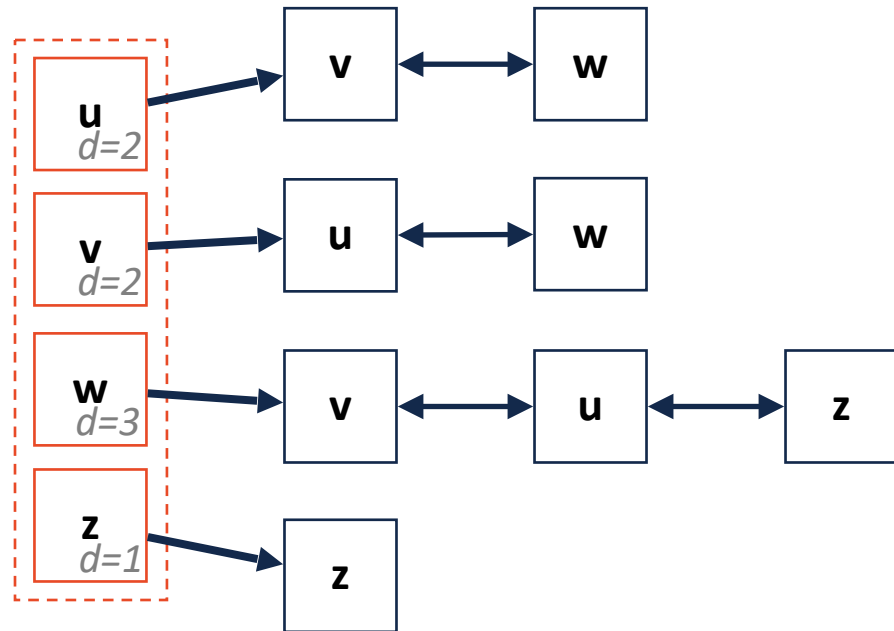
removeVertex(v):



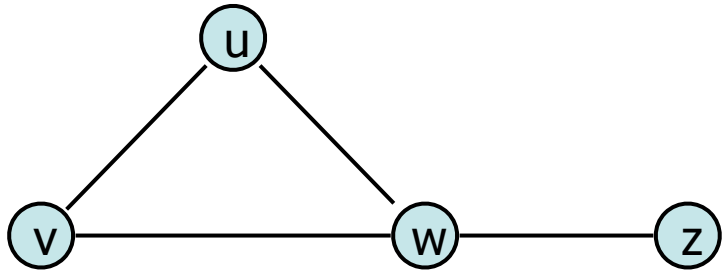
Adjacency List



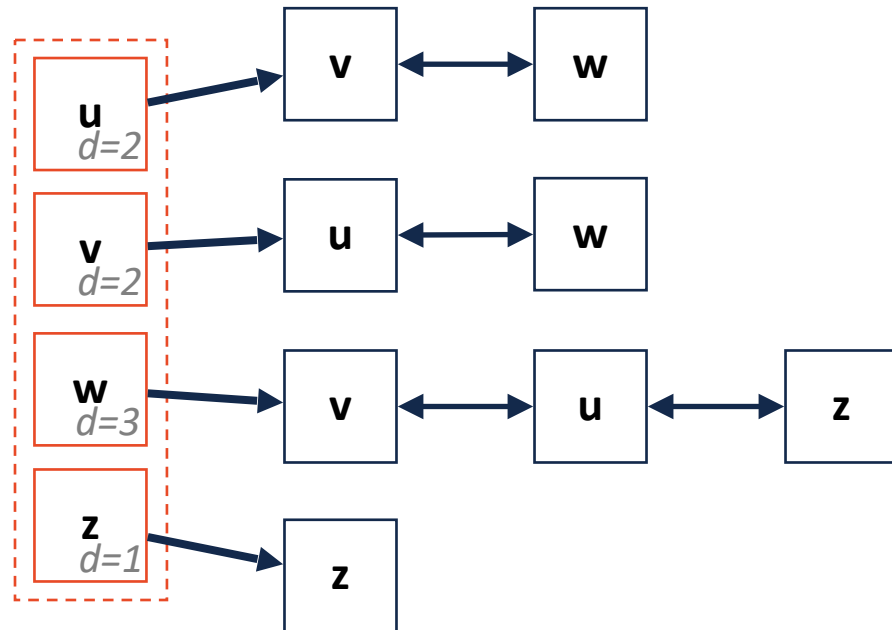
insertEdge(u, v):



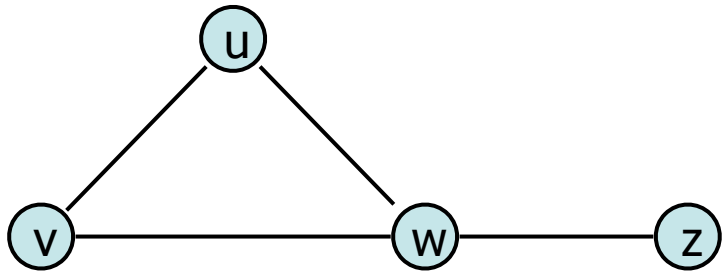
Adjacency List



removeEdge(u, v):

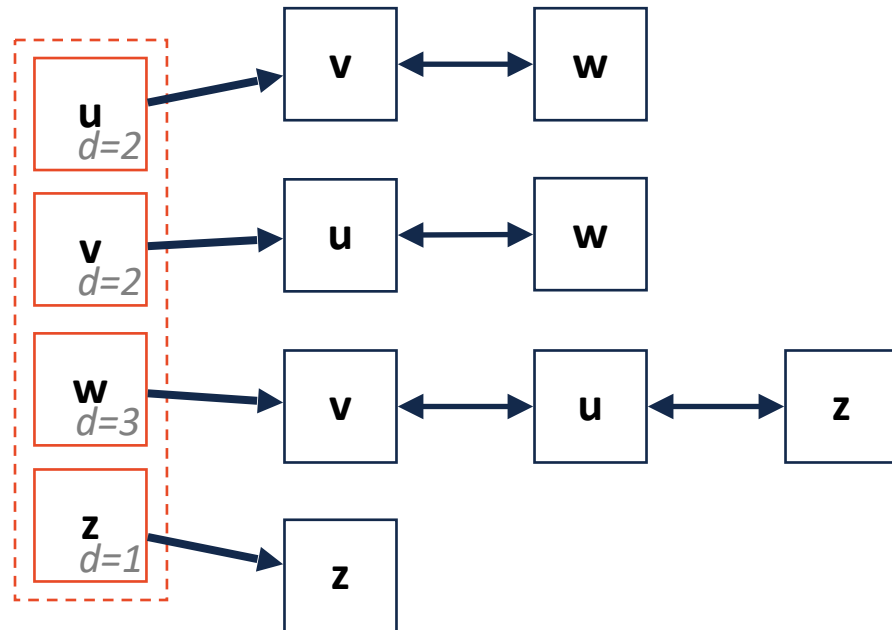


Adjacency List



Pros:

Cons:



$$|V| = n, |E| = m$$

Expressed as O(f)	Edge List	Adjacency Matrix	Adjacency List
Space			
insertVertex(v)			
removeVertex(v)			
insertEdge(u, v)			
removeEdge(u, v)			
getEdges(v)			
areAdjacent(u, v)			