

Algorithms and Data Structures for Data Science

Binary Search Tree

CS 277

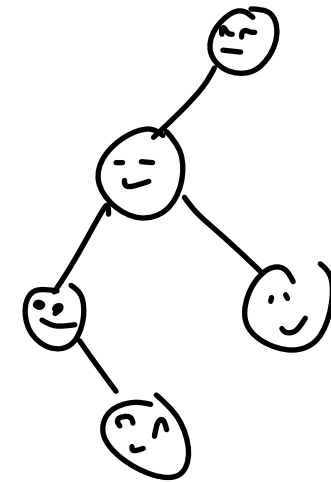
Brad Solomon

March 4, 2024



UNIVERSITY OF
ILLINOIS
URBANA - CHAMPAIGN

Department of Computer Science



Learning Objectives

Review understanding of Binary Trees

Introduce the dictionary ADT

Extend ADT to Binary Search Trees

Practice recursion in the context of trees

why is this useful?

Huffman Tree



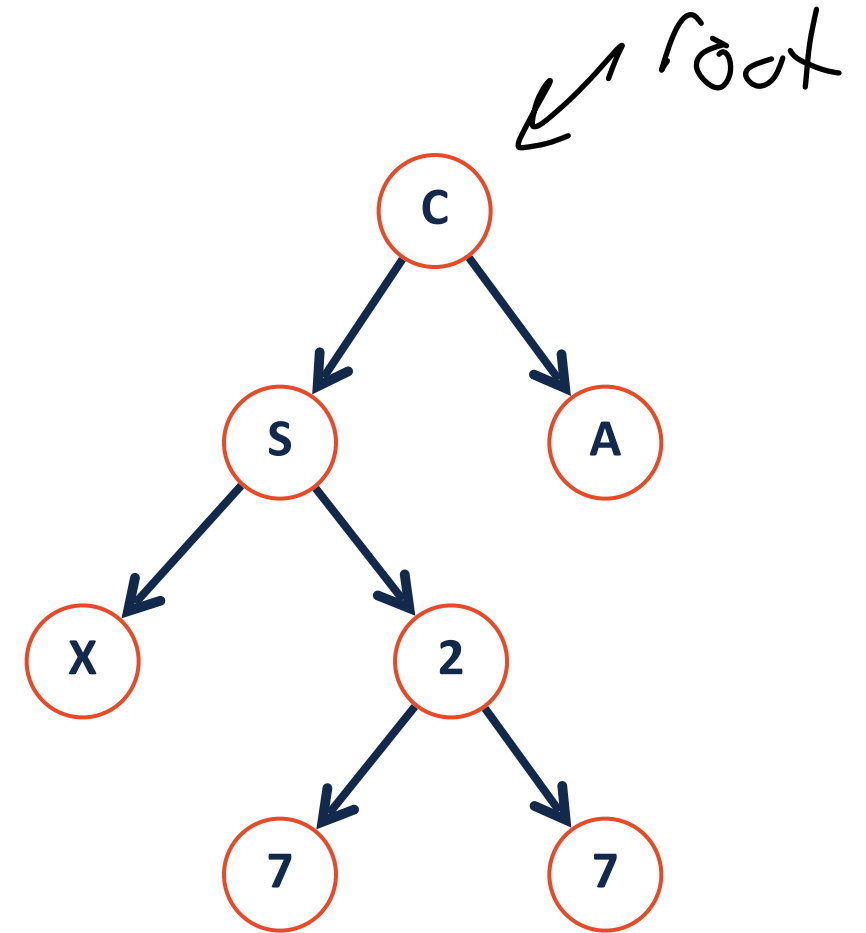
Binary Tree Recursion

A **binary tree** is a tree T such that:

$T = \text{None}$

or

$T = \text{treeNode}(\text{val}, T_L, T_R)$



```
1 class treeNode:
2     def __init__(self, val, left=None, right=None):
3         self.val = val
4         self.left = left
5         self.right = right
```

```
1 class binaryTree:
2     def __init__(self):
3         self.root = None
4
5
```

Tree ADT

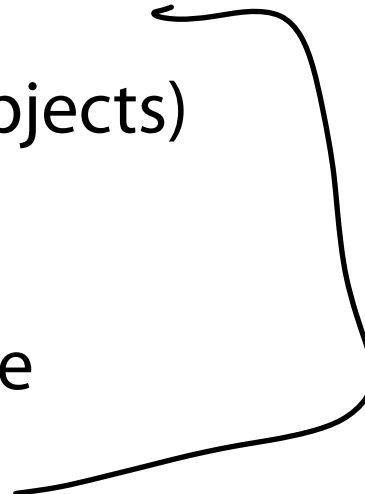
Constructor: Build a new (empty) tree

Insert: Add an object into tree

Remove: Remove a specific object from tree

Traverse: Visit every node in tree (all objects)

Search: Find a specific object in the tree



Binary Tree Traversal



Last class we implemented traversals using recursion, stacks, and queues.

What implementations led to a **depth first search traversal**?

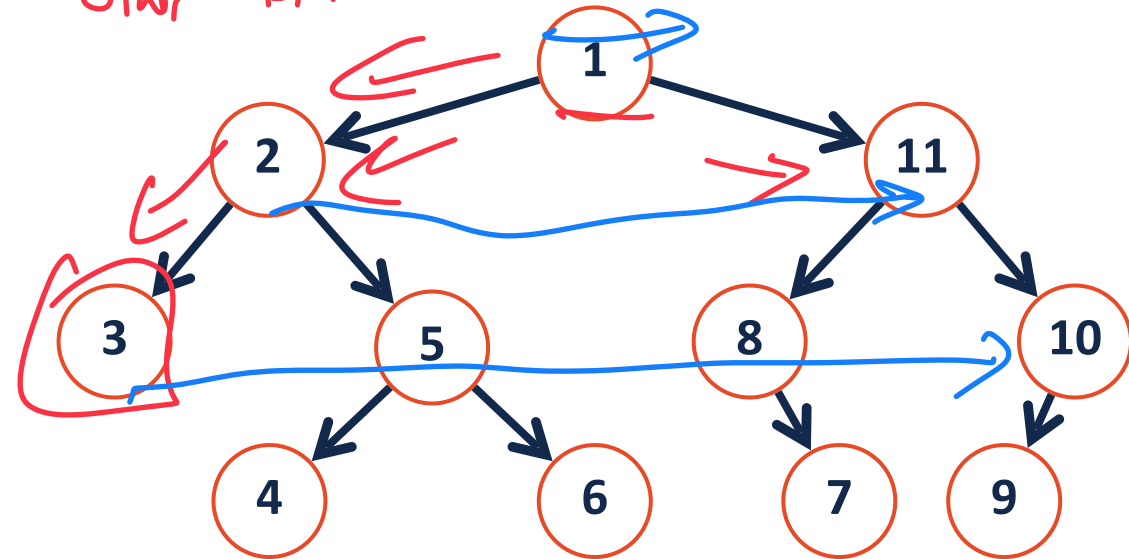
↳ pre, in, post order

↳ Descend down one branch as much as possible (to a leaf) before looking at other branches

↳ Stack

Which lead to **breadth first search**?

↳ clear a level before looking deeper



↳ Queue

1 2 11 3 5 8 10

Binary Tree Utility

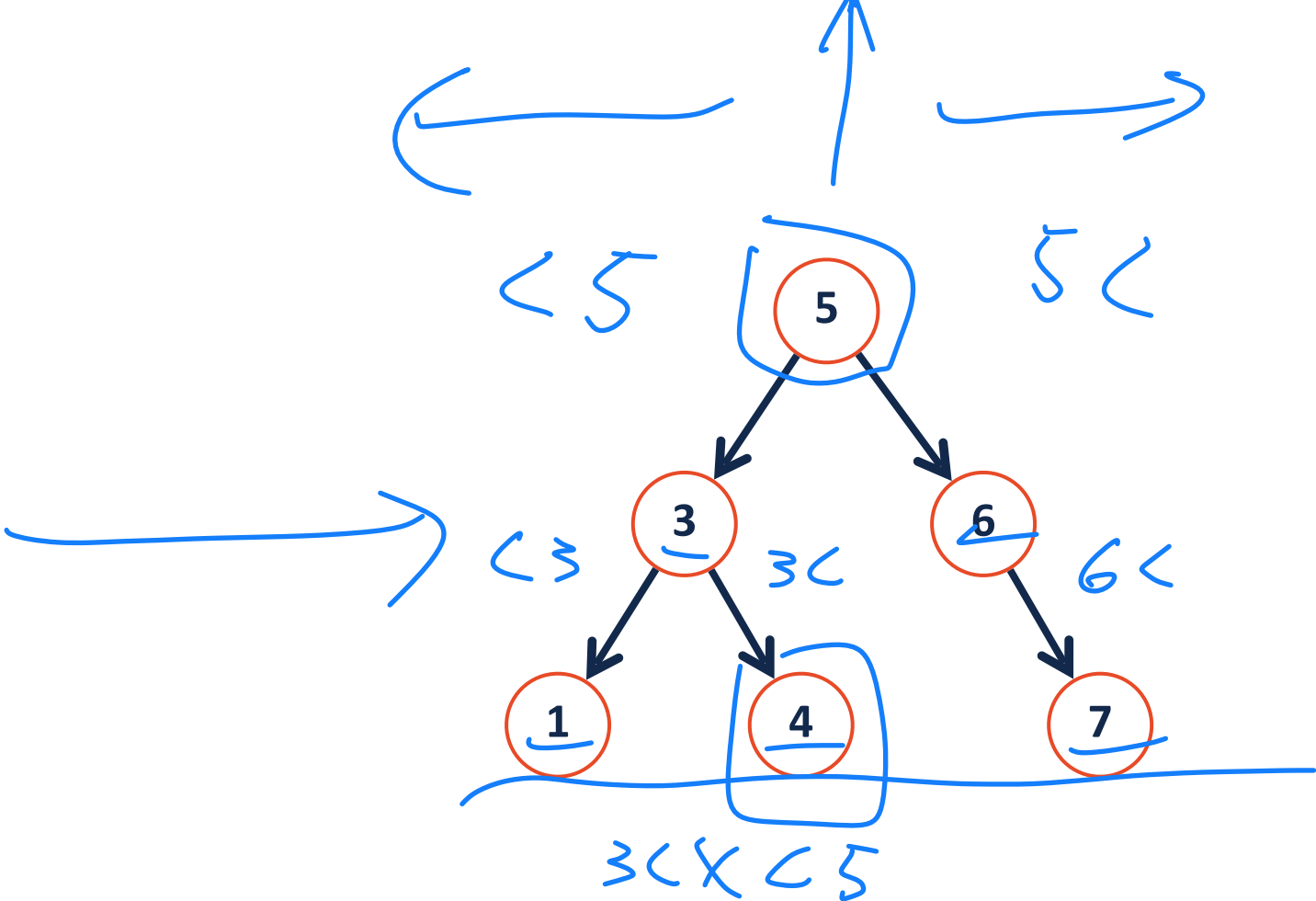
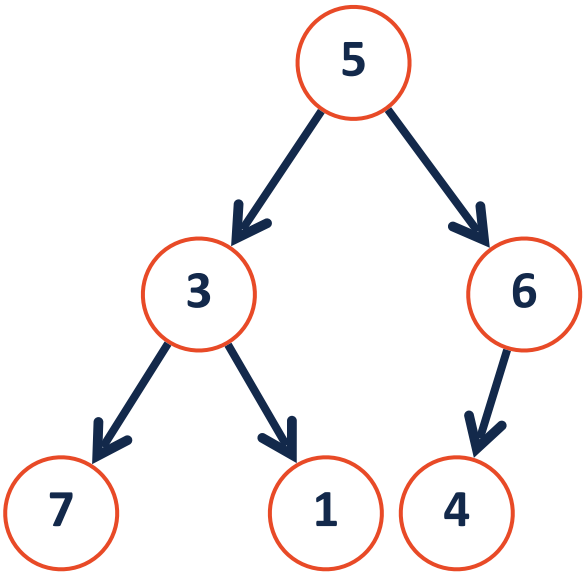
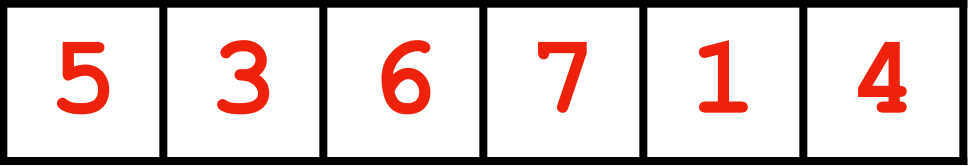
This week we will deep dive into useful implementations of binary trees

Binary Search Tree: An efficient implementation of a dictionary

Huffman Tree: A binary tree used to define an optimal text encoding

↳ 'Information theory'!!

Improved search on a binary tree



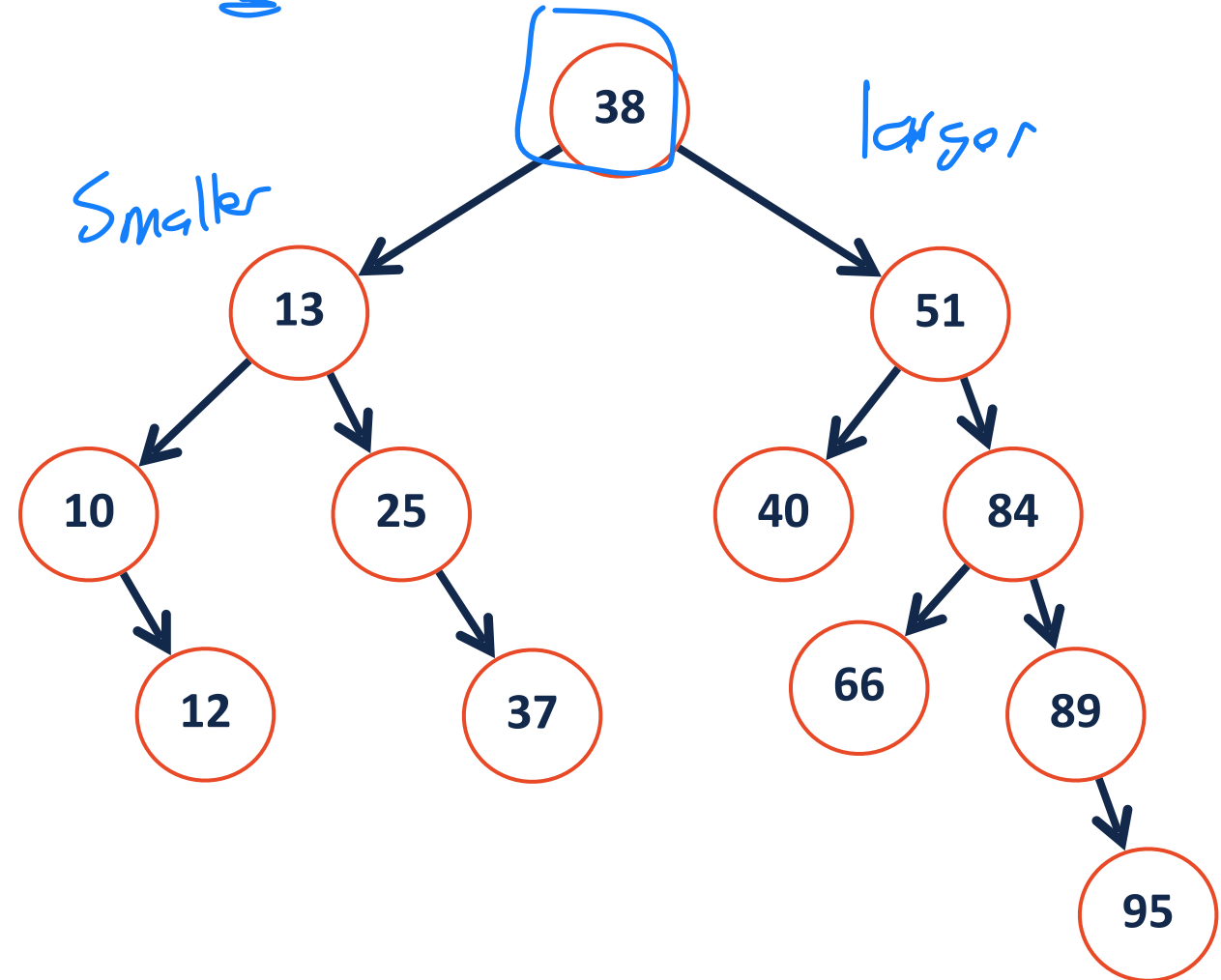
Binary Search Tree (BST)

A **BST** is a binary tree $T = \text{treeNode}(\text{val}, T_L, T_r)$ such that:

$\forall n \in T_L, n.\text{val} < T.\text{val}$

$\forall n \in T_R, n.\text{val} > T.\text{val}$

Added constraint!



Dictionary ADT

Data is often organized into key/value pairs:

Key → Value

Word → Definition

Course Number → Lecture/Lab Schedule

Node → Edges

Flight Number → Arrival Information

URL → HTML Page

Average Image Color → File Location of Image

Dictionary in Python

```

1 # The dictionary data structure
2 d = {}
3
4 # Change Value / Insert
5 d[key] = value
6 d[k2] = v2
7 d[key] = v3
8
9 # Remove value
10 d.pop(k2)
11
12 # Get Value
13 print(d[key])

```

$x = \{ (2, "Hi") \}$
 Bye
 X
 }

$d[key] = value$ $d[2] = "Hi"$

1) If key not in dictionary, (Insert)
 add k, v pair

2) If key is in dictionary,
 (change value) $d[2] = "Bye"$

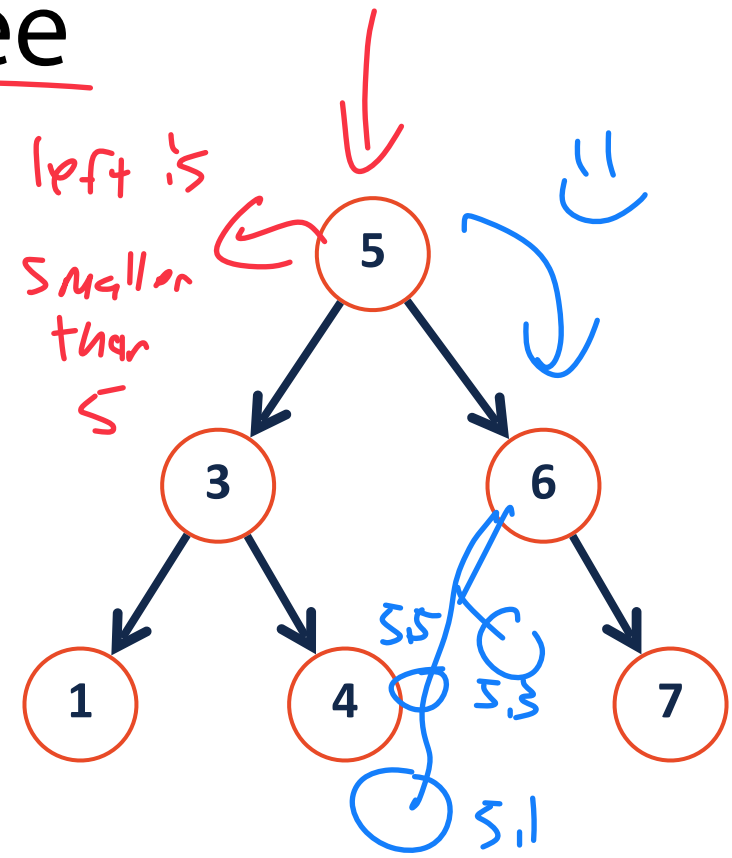
$x, pop(2)$ → will remove
 ↑
 key $(2, "Hi")$

Dictionary as a Binary Search Tree

```

1 class bstNode:
2     def __init__(self, key, val, left=None, right=None):
3         self.key = key
4         self.val = val
5         self.left = left
6         self.right = right
    
```

- val now has key + val



1) tree has structure! 😊
 ↳ More efficient (?)

2) type search has meaning
 ↳ Look up key to get value (value)

5	3	6	7	1	4
A	B	C	D	E	F

Binary Search Tree ADT — what changed?



Constructor: Build a new (empty) tree

Insert: Add an object into tree

old insert: Parent, direction, value

tree has structure! (Sorted)
↳ A new insert!

Remove: Remove a specific object from tree

→ A new remove

Traverse: Visit every node in tree (all objects)



Search: Find a specific object in the tree

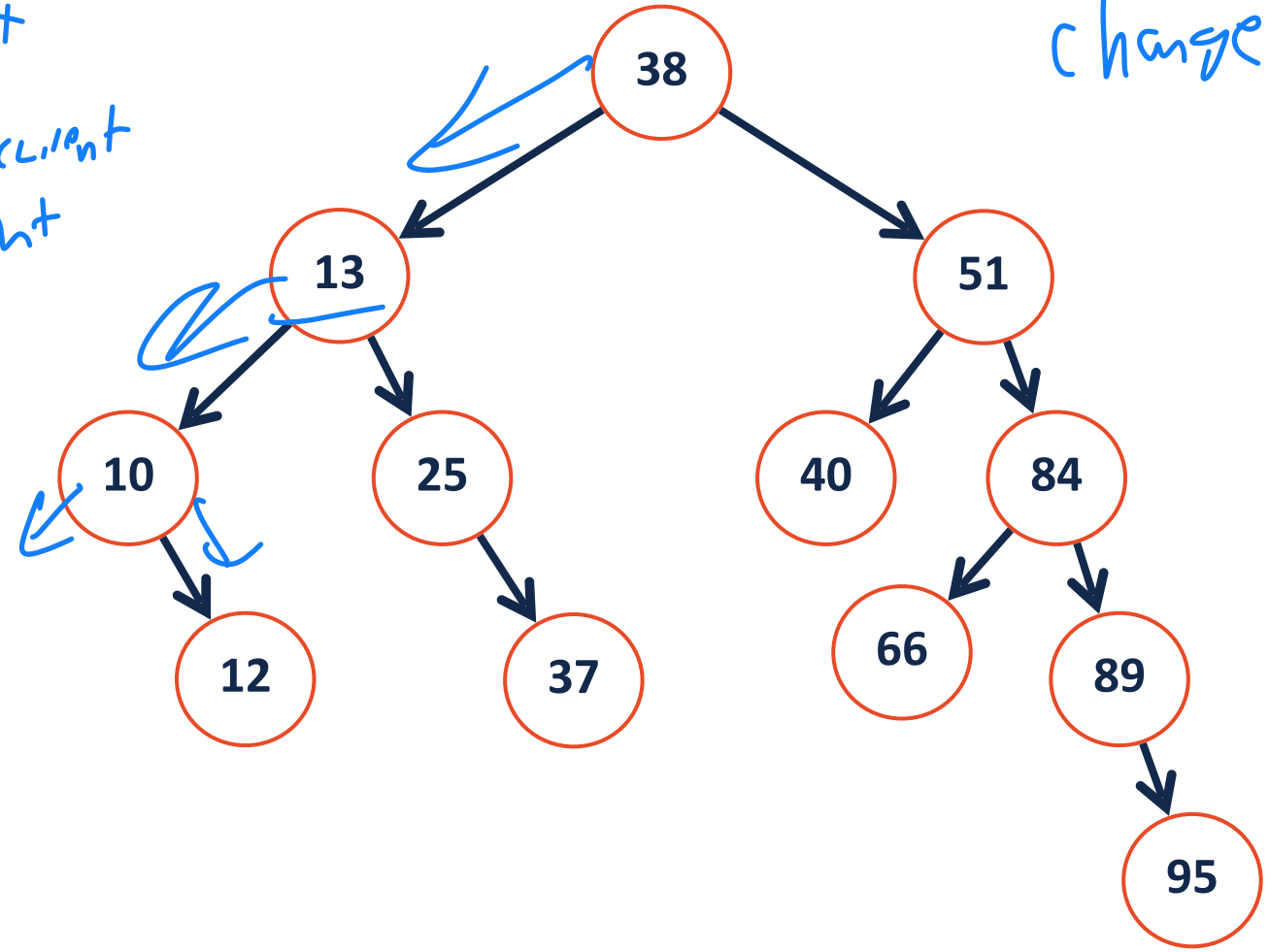
→ sorted order!

old search: Do a traversal

BST In-Order Traversal

No traversal doesn't change

- 1) Recurse left
- 2) Process mid/current
- 3) Recurse right



10 12 13 25 37 38 40 51 66 84 89 95

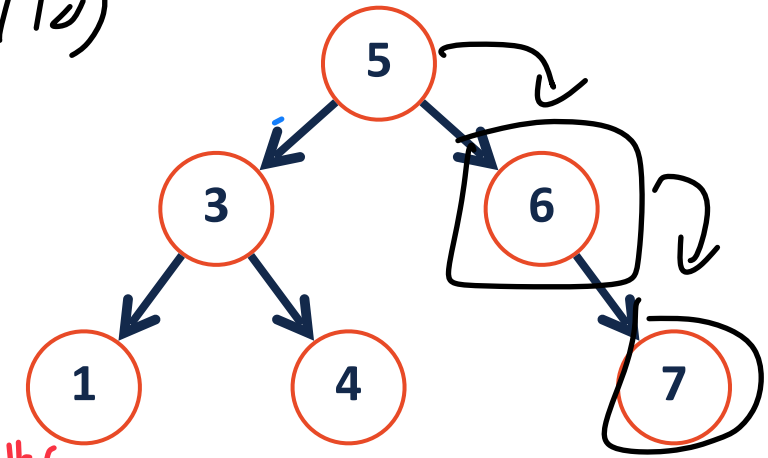
Insert(10)

BST Insert

Base Case:

Tree of size 0

→ make new bst Node
set to root

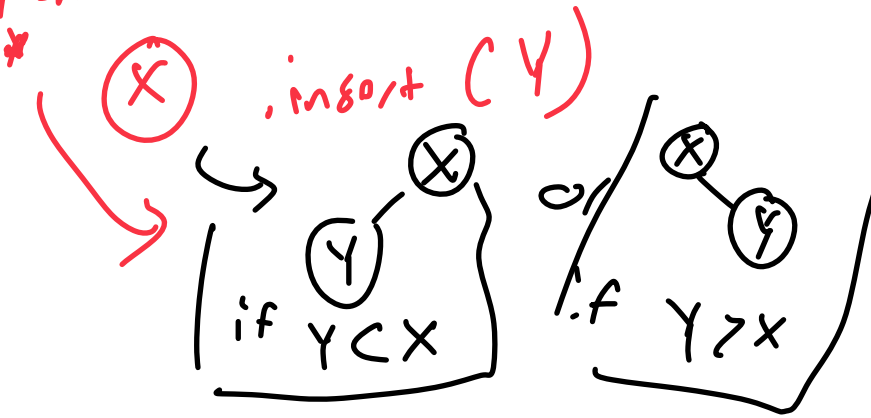


Tree of size 1

→ Check if insert larger or smaller
Set new Node as child

Recursive Step:

Check if key is larger or smaller than root node
↳ recurse insert to appropriate child

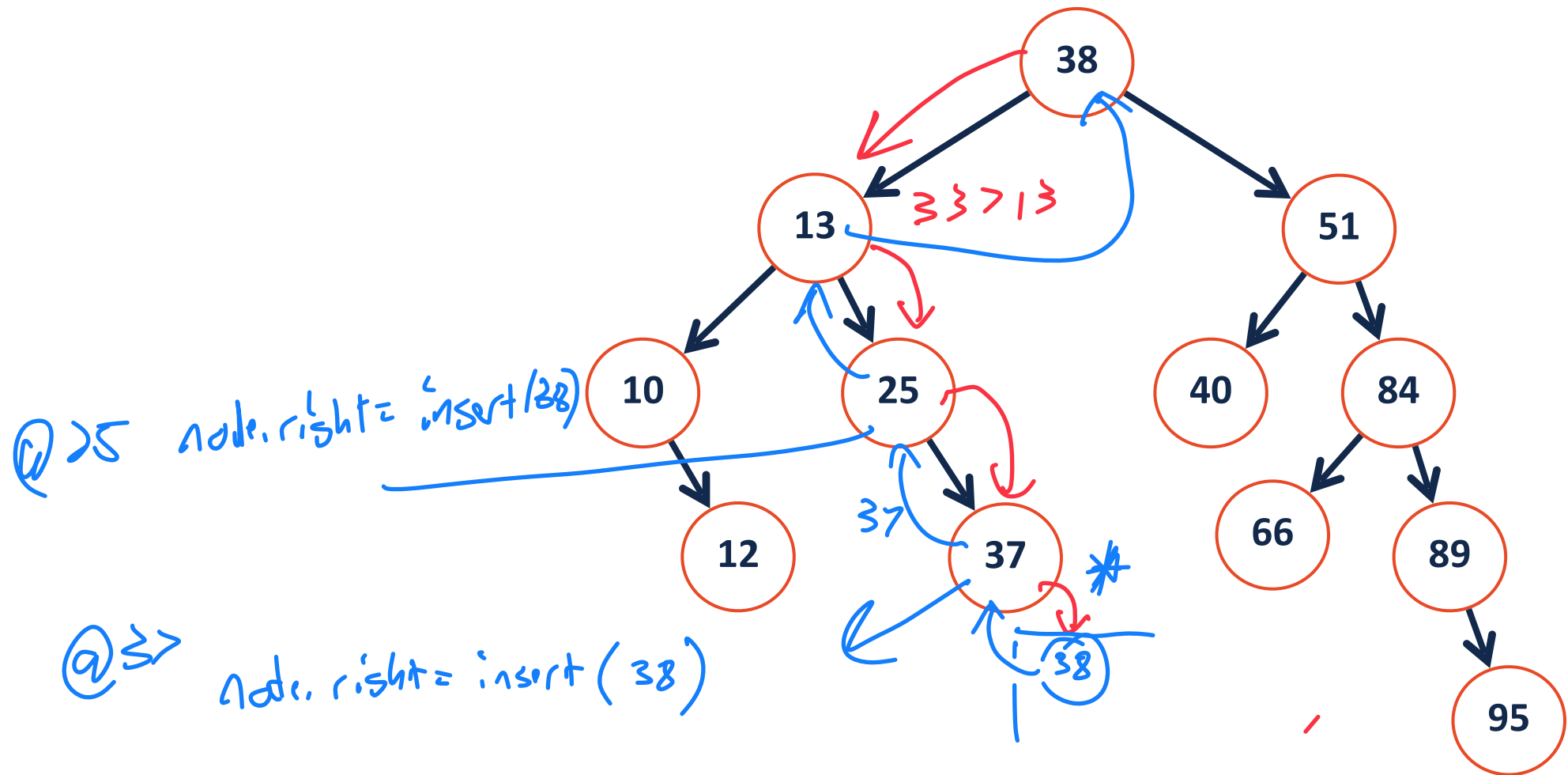


Combining: ??

BST Insert

insert(33)

$33 < 38$



BST Insert

Insert (2)

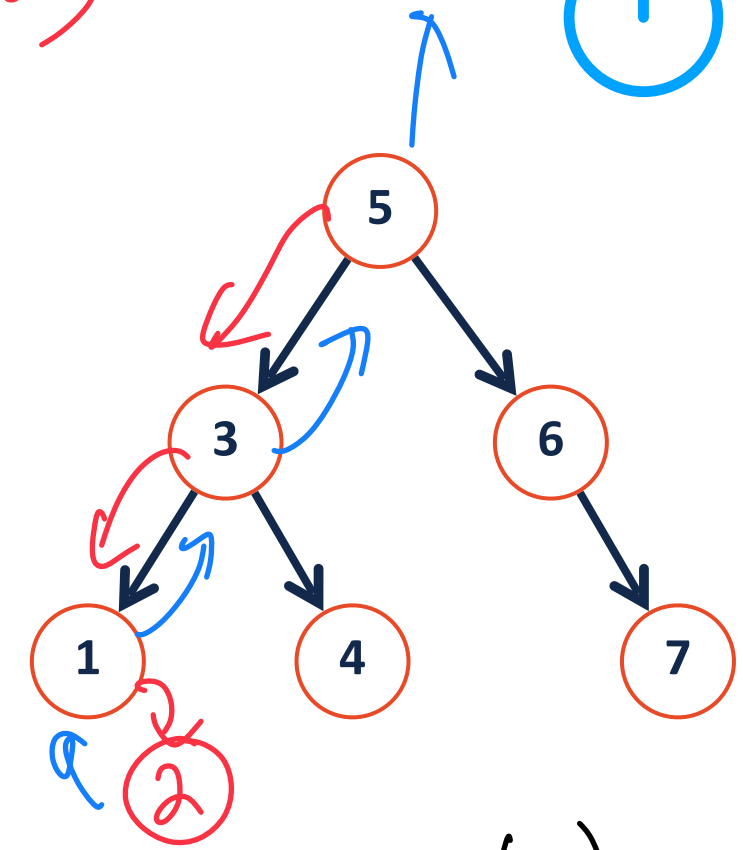


```
1 # inside class bst
2 def insert(self, key, val):
3     self.root = self.insert_helper(self.root, key, val)
4
5 def insert_helper(self, node, key, val):
6     if node == None:
7         return bstNode(key, val)
8
9     if key < node.key: # look left
10        node.left = self.insert_helper(node.left, key, val)
11    else: # look right
12        node.right = self.insert_helper(node.right, key, val)
13
14    return node
```

leaf

Base case
Recursion

combining

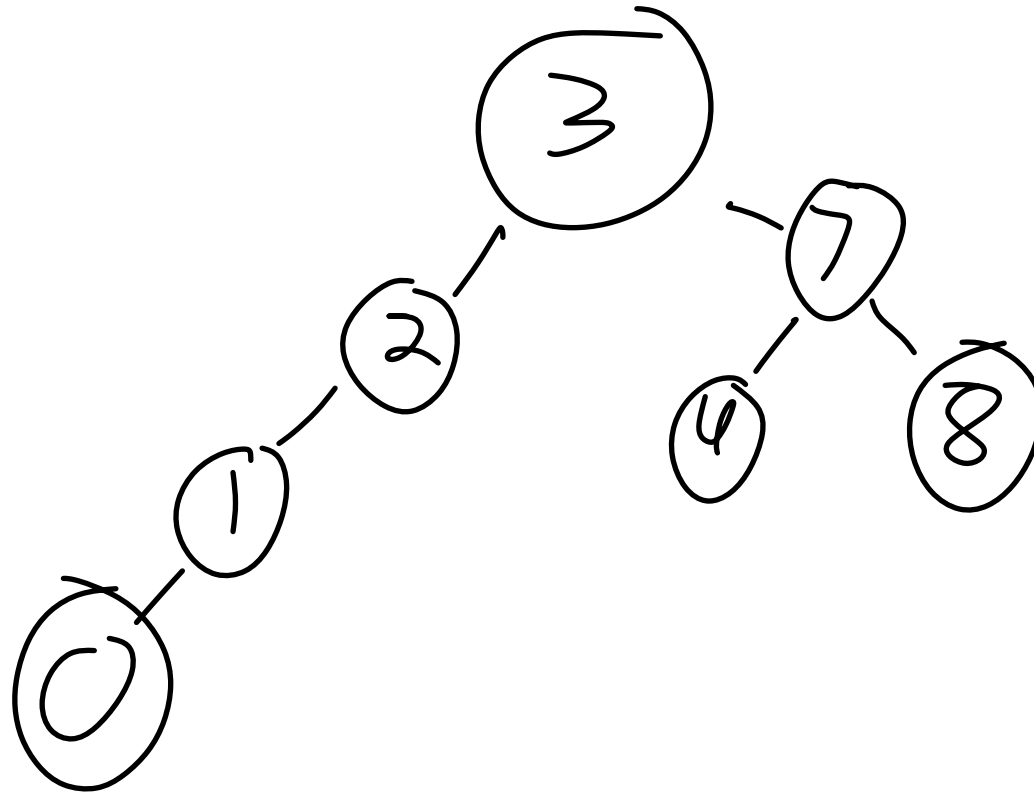


Big O : $O(n)$

b/c height can be n.
 $1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \rightarrow 6 \rightarrow 7$

BST Insert

What binary tree would be formed by inserting the following sequence of integers: [3, 7, 2, 1, 4, 8, 0]



BST Find

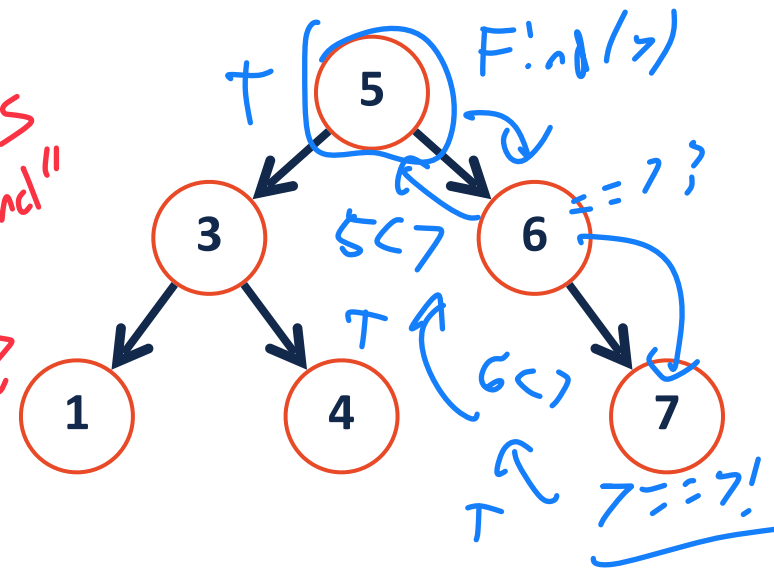
Base Case:

Tree is None / Empty

→ return

False
None
-1

What is
"Not found"
in context
of
problem?



Recursive Step:

If key being searched is smaller than Node.Key

↳ Recurse Left

else

↳ Recurse Right

before recursing,

check if Node.Key == key

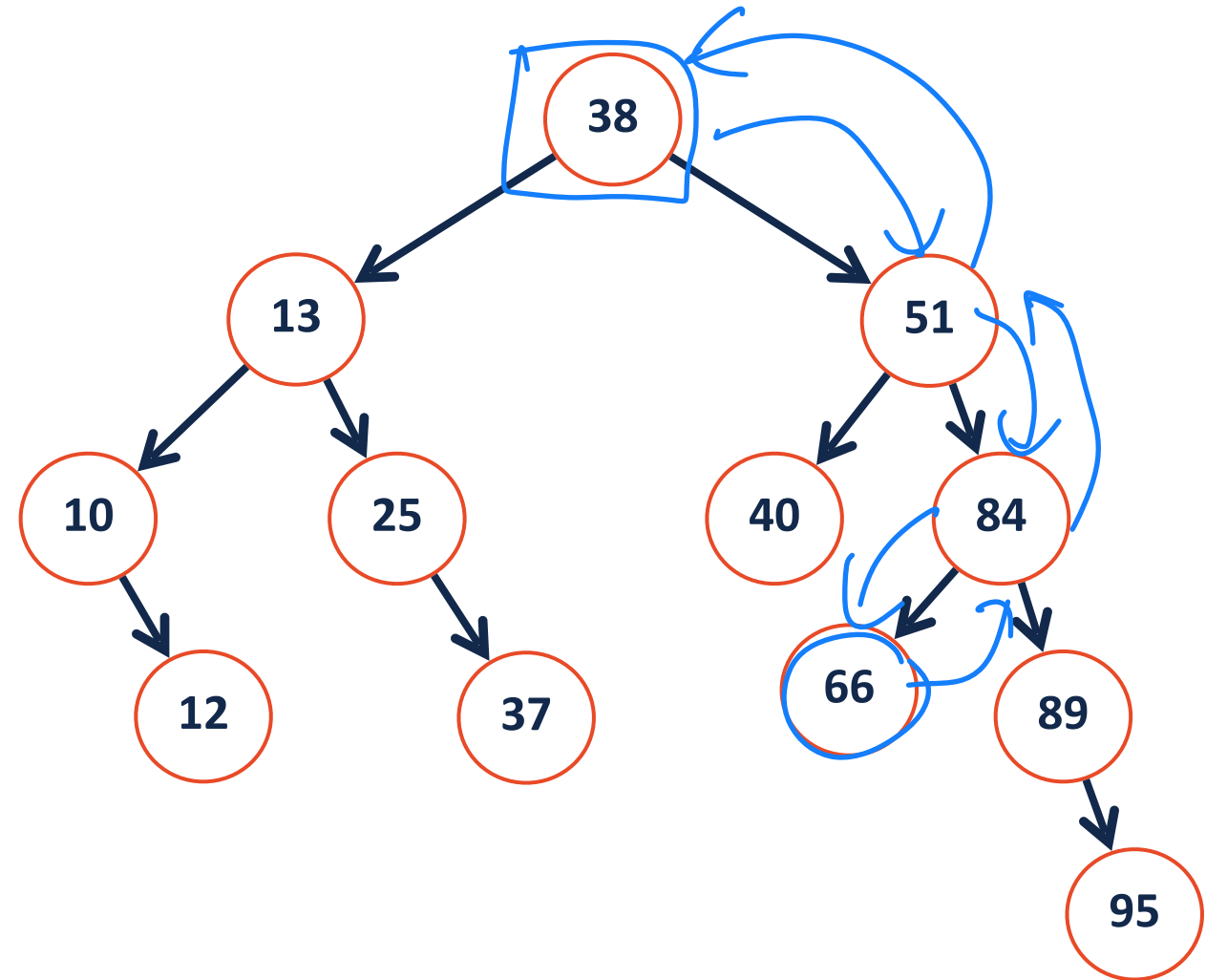
↳ return True!

Combining:

Return the value from recursion call

BST Find

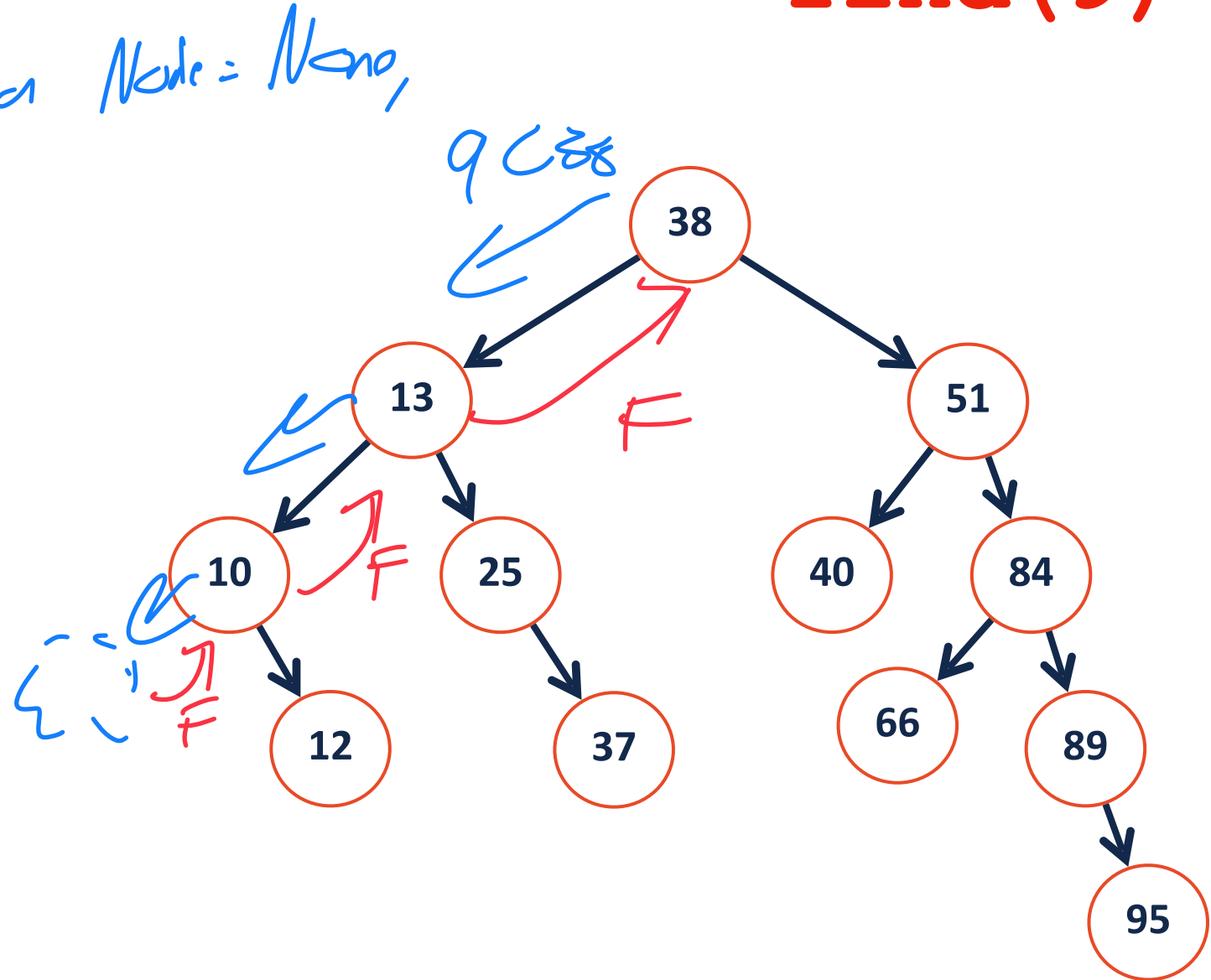
find(66)



BST Find

find(9)

If we ever reach a Node = None,
Value doesn't exist!
↳ False!



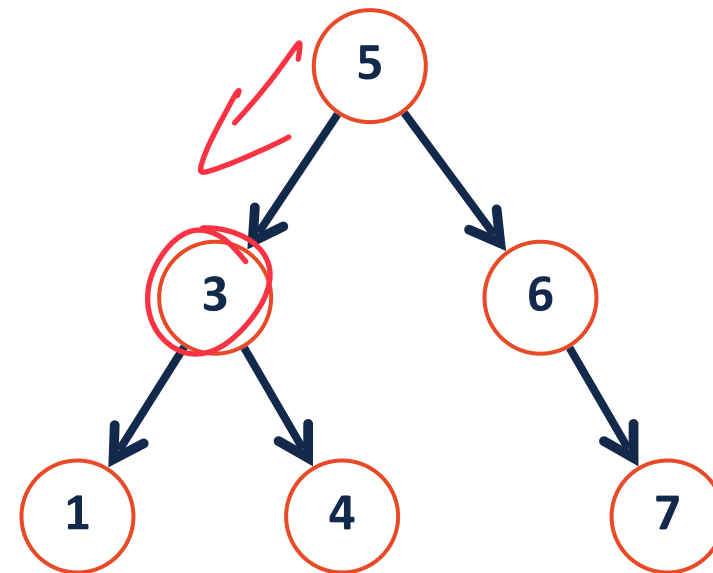
BST Find

My implementation

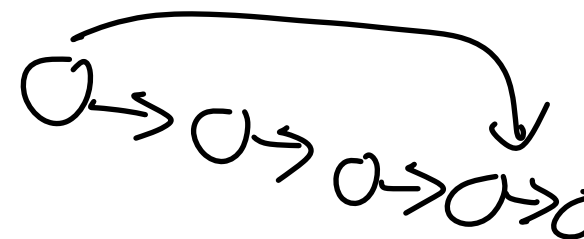
Find(3)



```
1 class bst:
2     def find(self, key):
3         n = self.find_helper(self.root, key)
4         if n:
5             return n.val
6         else:
7             return None
8
9     def find_helper(self, node, key):
10        nkey = node.key
11        > if nkey > key:
12            if node.left:
13                return self.find_helper(node.left, key)
14            else:
15                return None
16        C elif nkey < key:
17            if node.right:
18                return self.find_helper(node.right, key)
19            else:
20                return None
21        ← ← else:
22            return node
23
```



Big O: $O(n)$



BST Remove

remove (40)

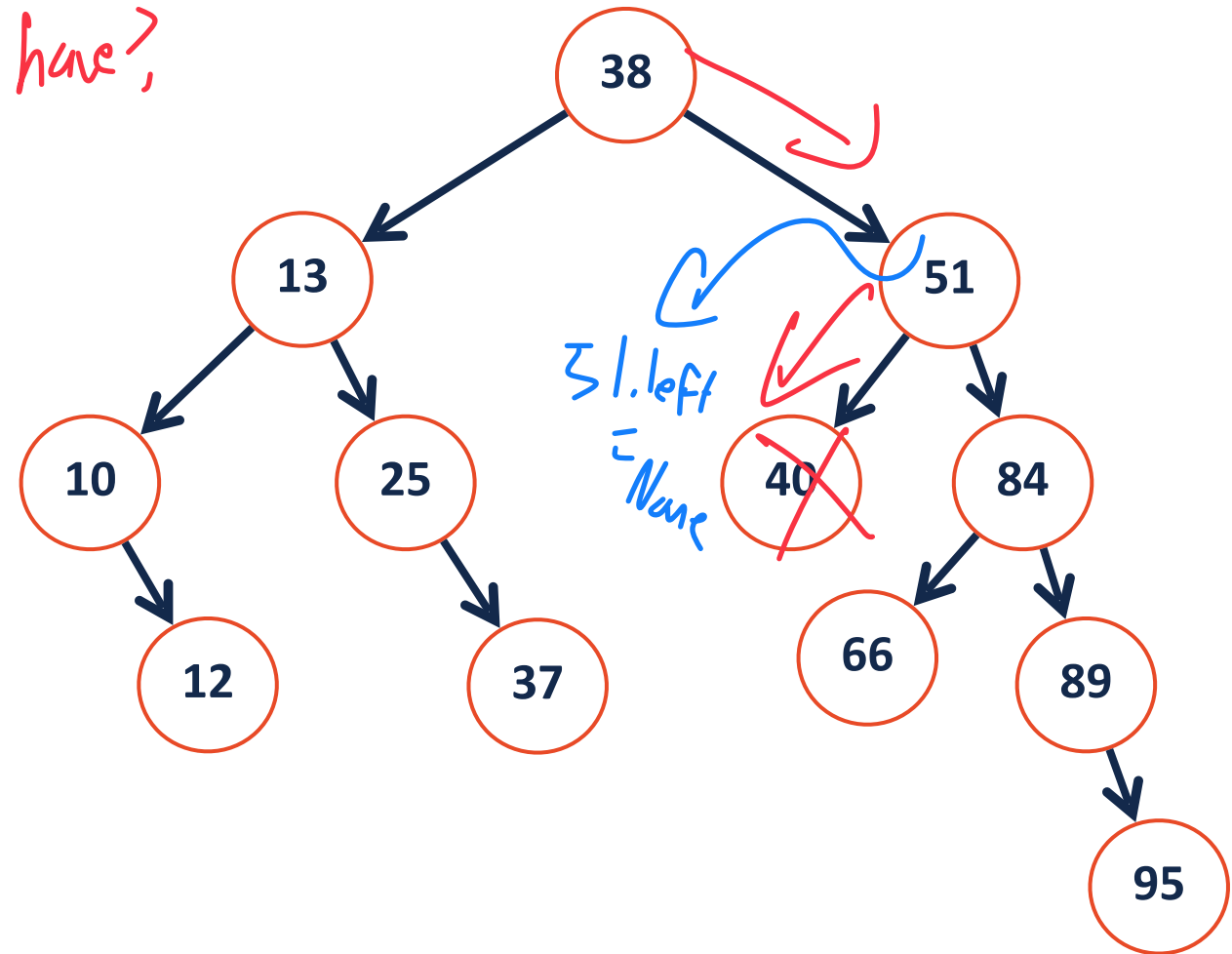
1) Find (40)

How many children does 40 have?

0 - child removal

↳ Set parent (get collect child)

↳ None



BST Remove

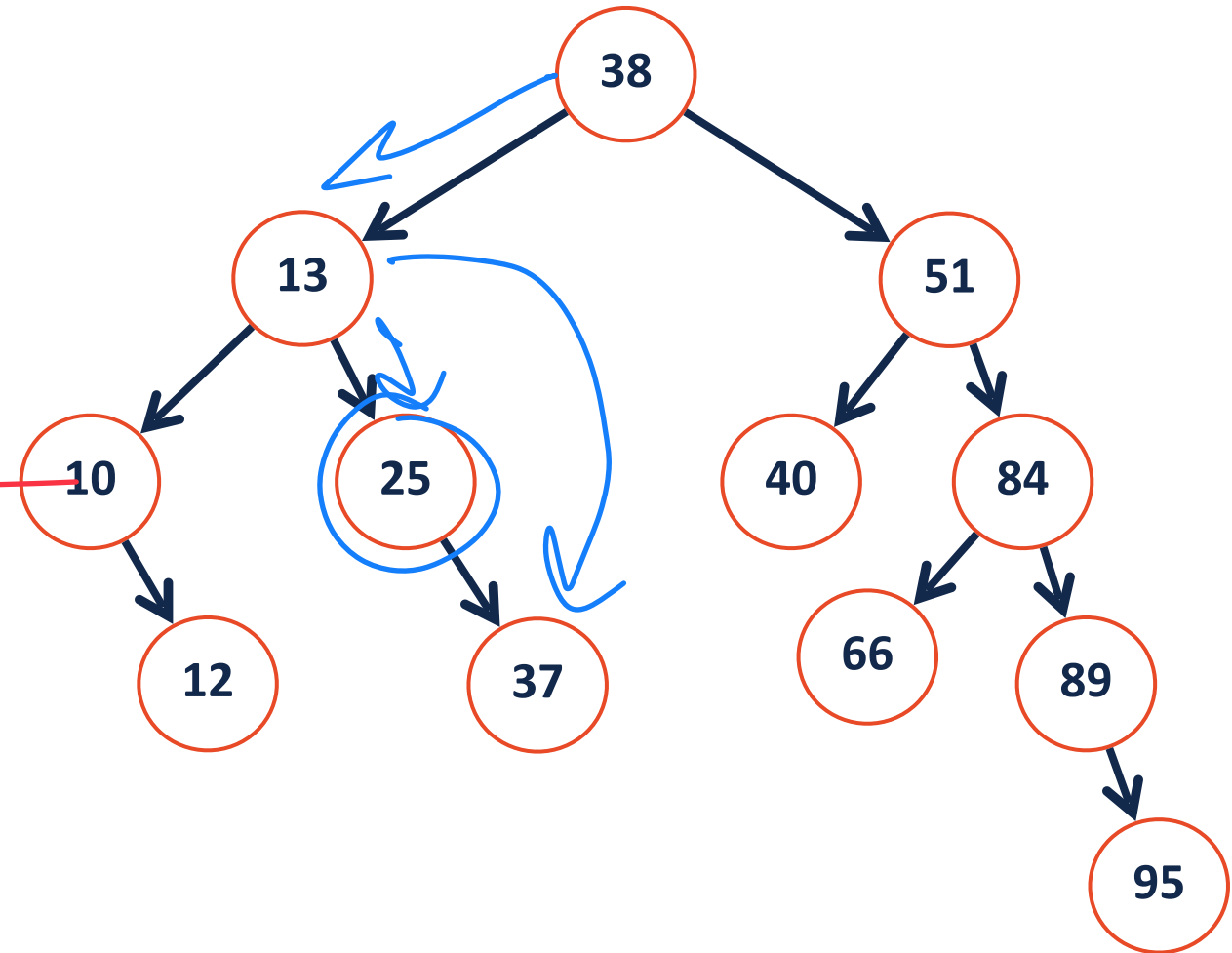
remove (25)

1) Find (25)

2) # children?

linked list removal *

parent direction = removed nodes child



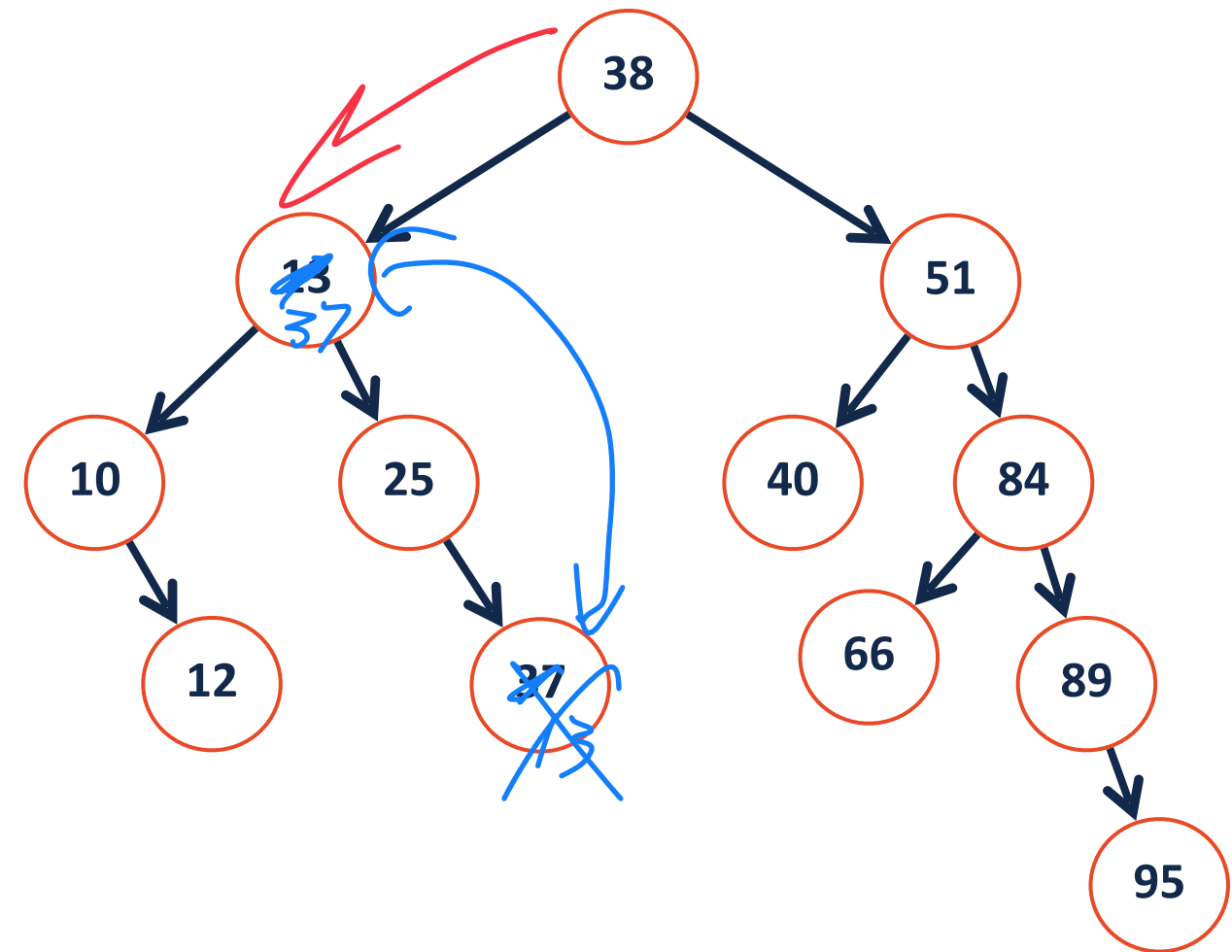
BST Remove

remove (13)

1) Find (13)

2) # children?
↳ 2 children

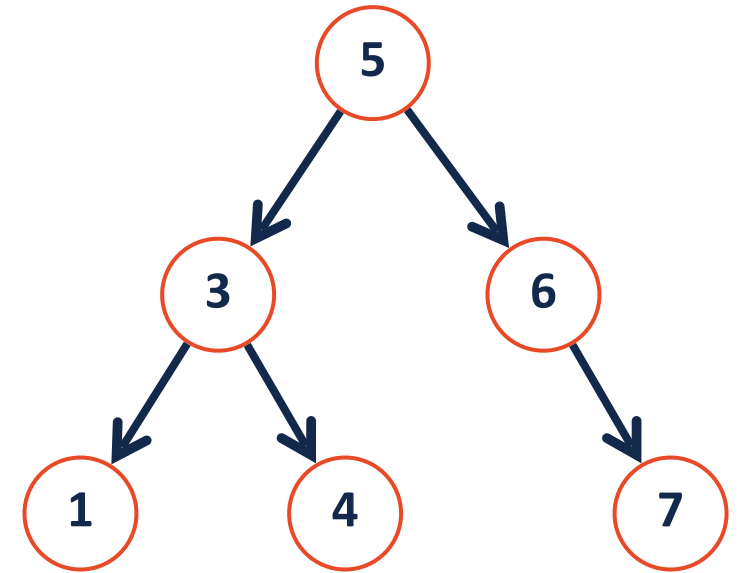
Before swap w/ leaf was ok!
But now... It can break
BSTree property



BST Remove

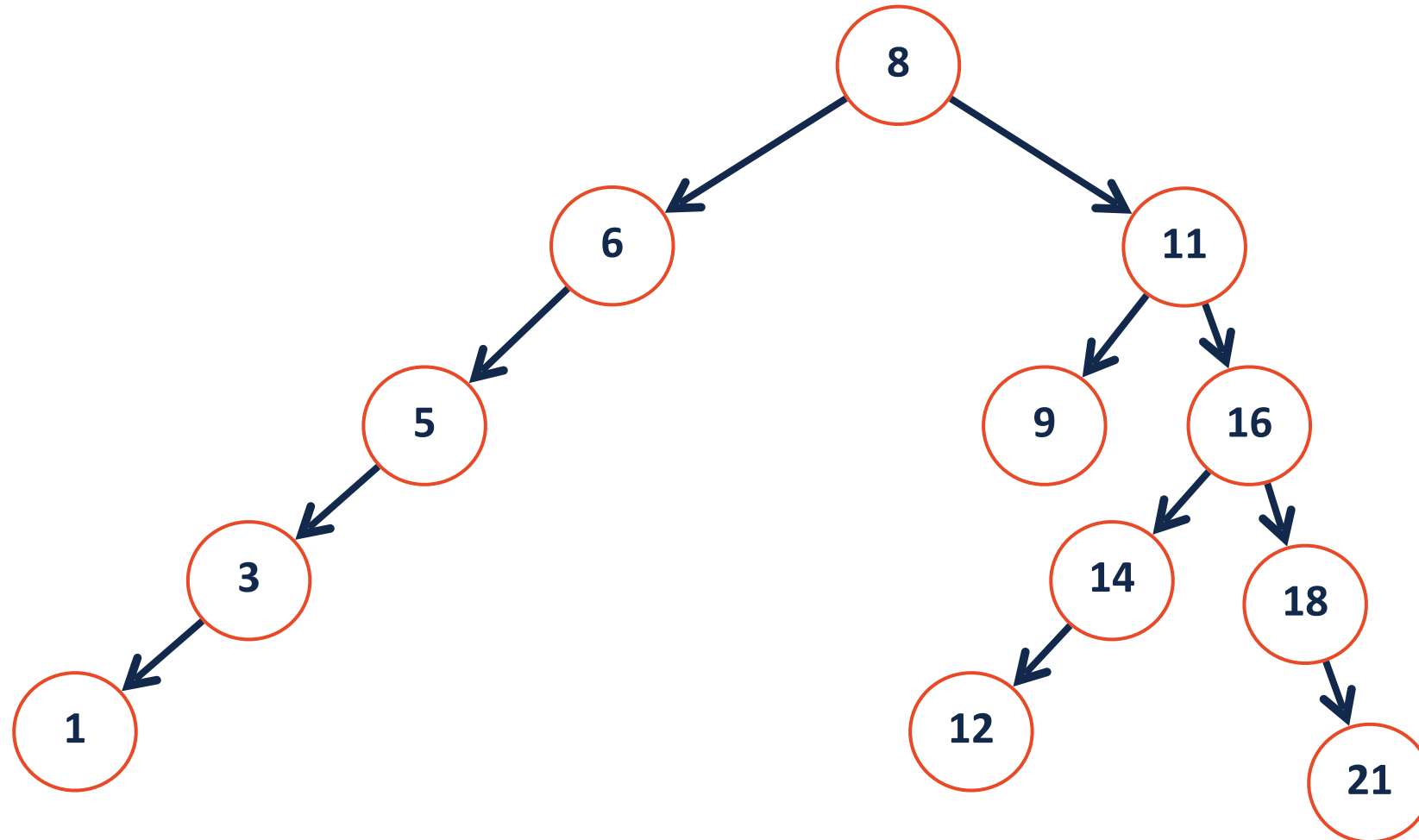


```
1 def remove(self, key):
2     self.root = self.remove_helper(self.root, key)
3
4 def remove_helper(self, node, key):
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25
```



BST Remove

What will the tree structure look like if we remove node 16 using IOS?

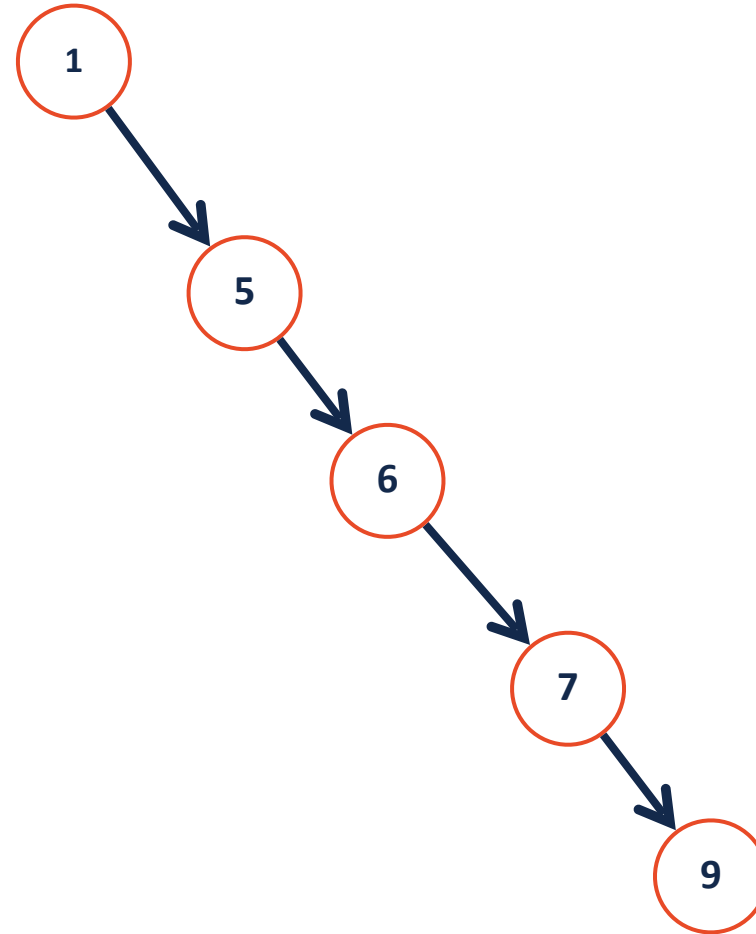
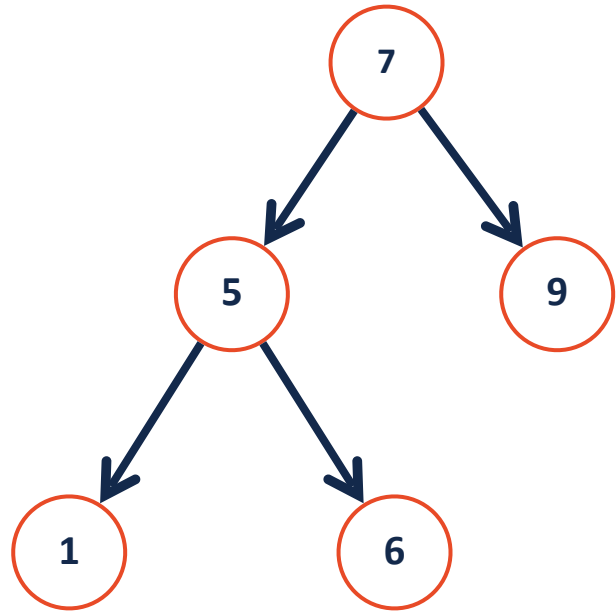


BST Analysis – Running Time



Operation	BST Worst Case
find	
insert	
delete	
traverse	

Limiting the height of a tree



Option A: Correcting bad insert order

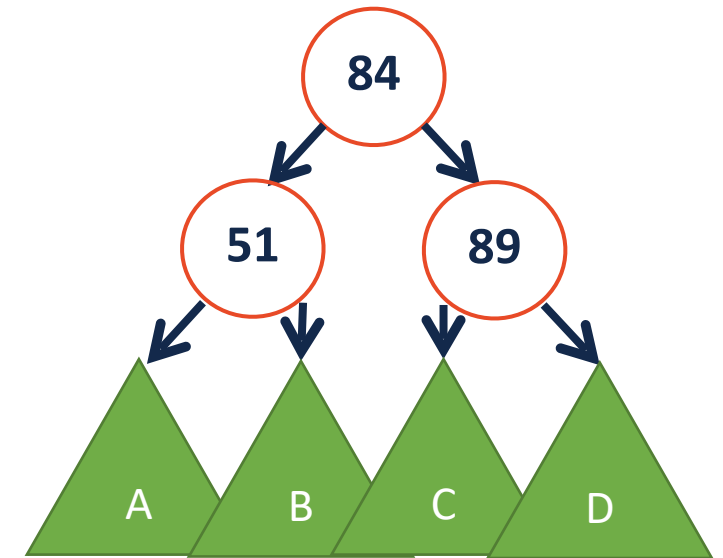
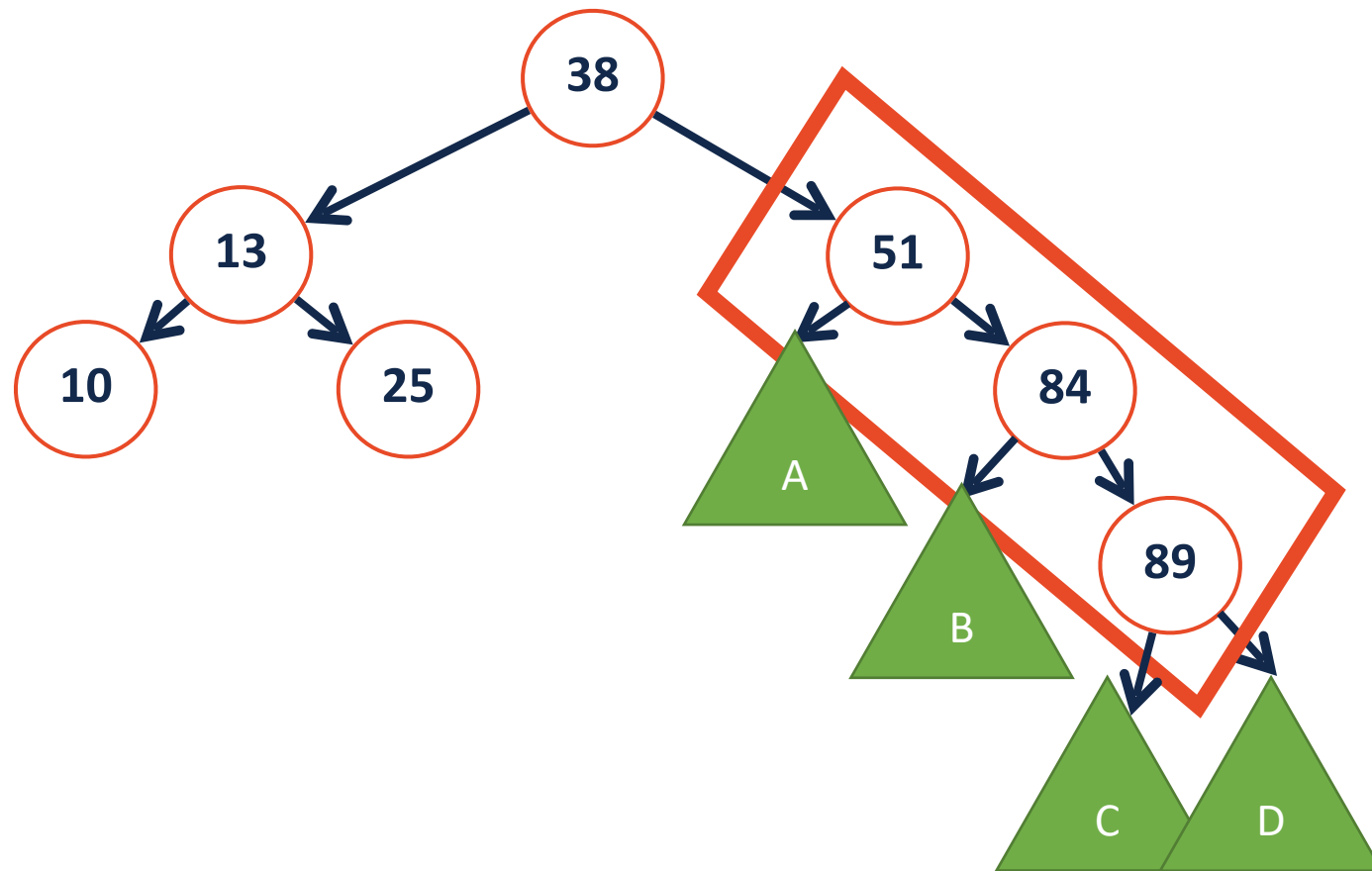
The height of a BST depends on the order in which the data was inserted

Insert Order: [1, 3, 2, 4, 5, 6, 7]

Insert Order: [4, 2, 3, 6, 7, 1, 5]

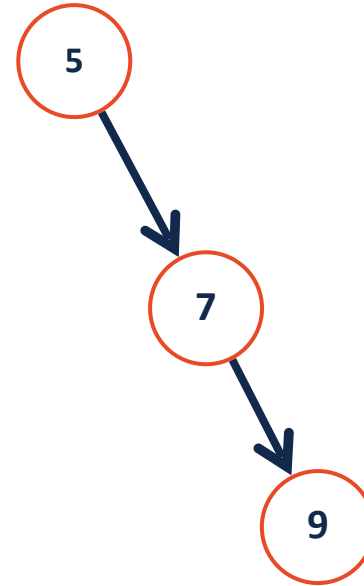
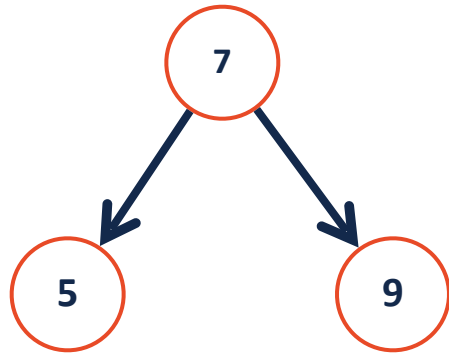
AVL-Tree: A self-balancing binary search tree

Rather than fixing an insertion order, just correct the tree as needed!



Height-Balanced Tree

What tree is better?



How would you describe this mathematically?