Algorithms and Data Structures for Data Science Graph Traversals

CS 277 Brad Solomon April 19, 2023



Department of Computer Science

Exam 3 Signups Available

April 24 — April 27

Very limited window for makeup exams (since end of semester is near)

Covers content from week 10 — 14

Learning Objectives

Re-introduce breadth first search traversal in a graph

Explore use cases for BFS traversal

Introduce depth first search traversal in a graph

Graph Traversals

There is no clear order in a graph (even less than a tree!)

How can we systematically go through a complex graph in the fewest steps?

Tree traversals won't work — lets compare:



- Rooted
- Acyclic
- Clear base cases ('doneness')



- Arbitrary starting point
- Can have cycles
- Must track visited nodes directly

Simple BFS Traversal 1) Create a queue and a visit list



2) Initialize both to contain our start

3) While queue not empty:

Remove front vertex of queue

Check if each edge has been seen before

Add unvisited edges to queue (and list)



Simple BFS Traversal 1) Create a queue and a visit list



2) Initialize both to contain our start

3) While queue not empty:

Remove front vertex of queue

Check if each edge has been seen before

Add unvisited edges to queue (and list)

What is my runtime?



Simple BFS Traversal



What is the shortest distance from **A** to **H**?

What is the shortest path from **A** to **H**?

What is the shortest path from **A** to **F**?

What is the shortest distance from **A** to **F**?

Simple BFS Traversal

What data structure is this?



Simple BFS Traversal



A **minimum spanning tree** is a tree formed by a subset of graph edges such that all vertices are connected with the smallest total possible edge weight

On an unweighted, undirected graph this MST can be built by tracking **discovery** edges during a BFS traversal



We call the remaining edges **cross** edges. What can I say about a graph with at least one **cross** edge?



If we modify our BFS traversal algorithm, we can track both **distances** and **discovery edges!**



Replace 'visited' list with a **distance** and a **previous**

When we add to queue, record **previous**.

When we process vertex from queue, record **distance**.

Queue

D

H

С

G

E

B

Vertex	Distance	Previous
A		
В		
C		
D		
E		
F		
G		
Н		

Replace 'visited' list with a **distance** and a **previous**

When we add to queue, record **previous**.

When we process vertex from queue, record **distance**.

A	
(D)	1
E F	
G	1

B

Distance	Previo
0	- (
1	Α
1	Α
1	Α
2	В
2	С
3	E
2	D
	Distance 0 1 1 1 1 2 2 3 2







Create a stack and a visit list
Initialize both to contain our start
While stack not empty:

Remove top item from stack (temporarily) If first time seeing top vertex, process it

If one or more adjacent edges unvisited, add the item back to stack before...

Add one unseen adjacent vertex to stack





Do our edge labels have meaning here?



Back Edge



DFS vs BFS	
DFS:	BFS:
Pros:	Pros:

Cons:

Cons:

