# Algorithms and Data Structures for Data Science Graph Implementations 2 

CS 277
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## This week only: Lab room and OH Changes

Friday April 14th: AE3's Celebration of Teaching in 1306 Everitt Our lab will be in 2101 Everitt instead!

Office Hour Changes: My OH will be Friday between 3:15 and 4:15 There will not be OH on Thursday April 13th!

## Lab Feedback

Still read them and appreciate feedback
lab_huffman needs work in the future in both presentation and content
lab_trees and lab_avl were both highly rated

## Learning Objectives

Review edge list and adjacency matrix graph implementations

Introduce adjacency list implementation

Discuss the strengths and weaknesses of each implementation

## Graphs

Given a roster of students for each class, build a graph which tracks whether there are at least three students in common between two classes What is a vertex?

What is an edge?

Are the edges directed or undirected?

Are the edges weighted or unweighted?

## Graph ADT

## Find

getVertices() - return the list of vertices in a graph
getEdges(v) - return the list of edges that touch the vertex $v$ areAdjacent( $u, v$ ) - returns a bool based on if an edge from $u$ to $v$ exists

## Insert

insertVertex(v) — adds a vertex to the graph
insertEdge( $u, v$ ) - adds an edge to the graph

## Remove

removeVertex(v) - removes a vertex from the graph
removeEdge( $u, v$ ) - removes an edge from the graph

## Graph Implementation: Edge List $|\mathrm{V}|=\mathrm{n},|\mathrm{E}|=\mathrm{m}$

The equivalent of an 'unordered' data structure


## Vertex Storage:

Not stored at all (recovered from edges) or
An unordered list of vertices

## Edge Storage:

An unordered list of edges (as tuples) [or equivalent]

## Graph Implementation: Edge List

 How would our data structure change if...

Edges are weighted:

## Graph Implementation: Edge List

 How would our data structure change if...

Edges are directed:

## Graph Implementation: Adjacency Matrix



## Vertex Storage:

Edge Storage:


## Graph Implementation: Adjacency Matrix



| U | 0 |  | U | v | W | z |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | u | 0 | 1 | 1 | 0 |
| V | 1 | v | 1 | 0 | 1 | 0 |
| W | 2 |  |  |  |  |  |
|  |  | W | 1 | 1 | 0 | 1 |
| Z | 3 | Z | 0 | 0 | 1 | 0 |

## Graph Implementation: Adjacency Matrix



| U | 0 |  | U | v | W | z |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | u | 0 | 1 | 1 | 0 |
| V | 1 | v | 1 | 0 | 1 | 0 |
| W | 2 |  |  |  |  |  |
|  |  | W | 1 | 1 | 0 | 1 |
| Z | 3 | Z | 0 | 0 | 1 | 0 |

## Graph Implementation: Adjacency Matrix


areAdjacent(u, v):

## Graph Implementation: Adjacency Matrix



| U | 0 |  | U | v | W | z |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | u | 0 | 1 | 1 | 0 |
| V | 1 | v | 1 | 0 | 1 | 0 |
| W | 2 |  |  |  |  |  |
|  |  | W | 1 | 1 | 0 | 1 |
| Z | 3 | Z | 0 | 0 | 1 | 0 |

## Graph Implementation: Adjacency Matrix



| U | 0 |  | U | v | W | z |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | u | 0 | 1 | 1 | 0 |
| V | 1 | v | 1 | 0 | 1 | 0 |
| W | 2 |  |  |  |  |  |
|  |  | W | 1 | 1 | 0 | 1 |
| Z | 3 | Z | 0 | 0 | 1 | 0 |

## Graph Implementation: Adjacency Matrix


removeVertex(v):

| $\mathbf{U}$ | $\mathbf{0}$ |
| :---: | :---: |
| $\mathbf{v}$ | $\mathbf{1}$ |
| $\mathbf{w}$ | $\mathbf{2}$ |
| $\mathbf{Z}$ | $\mathbf{3}$ |


|  | $u$ | $v$ | $w$ | $z$ |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{u}$ | 0 | 1 | 1 | 0 |
| $\mathbf{v}$ | 1 | 0 | 1 | 0 |
| $\mathbf{w}$ | 1 | 1 | 0 | 1 |
| $\mathbf{z}$ | 0 | 0 | 1 | 0 |

removeEdge(u, v):

## Graph Implementation: Adjacency Matrix

 Pros:

| U | 0 |  | u | v | w | z |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | U | 0 | 1 | 1 | 0 |
| V | 1 | v | 1 | 0 | 1 | 0 |
| W | 2 | w | 1 | 1 | 0 | 1 |
|  |  |  |  |  |  |  |
| Z | 3 | Z | 0 | 0 | 1 | 0 |

Cons:

## Graph Implementation: Adjacency Matrix

 How would our data structure change if...

Edges are directed:

## Graph Implementation: Adjacency Matrix

 How would our data structure change if...

Edges are weighted:

| $\mathbf{\| c \|}$ | $\mathbf{0}$ |
| :---: | :---: |
| $\mathbf{V}$ | $\mathbf{1}$ |
| $\mathbf{w}$ | $\mathbf{2}$ |
| $\mathbf{Z}$ | $\mathbf{3}$ |


|  | $\mathbf{u}$ | $\mathbf{v}$ | $\mathbf{w}$ | $\mathbf{z}$ |
| :---: | :---: | :---: | :---: | :---: |
| $\mathbf{u}$ | 0 | $\mathbf{1}$ | $\mathbf{1}$ | 0 |
| $\mathbf{v}$ | 1 | 0 | 1 | 0 |
| $\mathbf{w}$ | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ | $\mathbf{1}$ |
| $\mathbf{z}$ | $\mathbf{0}$ | $\mathbf{0}$ | $\mathbf{1}$ | $\mathbf{0}$ |

## Adjacency List



## Adjacency List



## Vertex Storage:

## Edge Storage:

## Adjacency List



## getVertices():

## Adjacency List



## getEdges(v):

## Adjacency List


areAdjacent(u, v):

## Adjacency List



## Adjacency List



## Adjacency List


insertEdge(u, v):

## Adjacency List



## removeEdge(u, v):

## Adjacency List

## Pros:



Cons:

## Adjacency List

How would our data structure change if...


## Edges are directed:

## Adjacency List

How would our data structure change if...


## Edges are weighted:

$|V|=n,|E|=m$

| Expressed as Off) | Edge List | Adjacency Matrix | Adjacency List |
| :---: | :---: | :---: | :---: |
| Space | n+m | $\mathrm{n}^{2}$ | n+m |
| insertVertex(v) | 1* | n* | 1* |
| removeVertex(v) | $\mathrm{m}^{* *}$ | n | deg(v) |
| insertEdge(u, v) | 1 | 1 | 1* |
| removeEdge(u, v) | m | 1 | $\begin{gathered} \min (\operatorname{deg}(u), \\ \operatorname{deg}(v)) \end{gathered}$ |
| getEdges(v) | m | n | deg(v) |
| areAdjacent(u, v) | m | 1 | $\begin{gathered} \min (\operatorname{deg}(u), \\ \operatorname{deg}(v)) \end{gathered}$ |

## Next week:Traversals

There is no clear order in a graph (even less than a tree!)
How can we systematically go through a complex graph in the fewest steps?

