Synchronization III

CS 241 Oct. 16, 2013

Review

- Mutex: A simple "lock"
 - pthread_mutex_lock()
 - pthread_mutex_unlock()

- Conditional Variable: "monitor" primitive
 - pthread_cond_wait()
 - pthread_cond_signal()
 - pthread_cond_broadcast()

void lock() {

}
void unlock() {

}

void wait() {

}
void post() {

}

Semaphore

 A semaphore is a "counting" mutex – sem_wait()

- sem_post()

Blocking Bounded Queue (v2)

void blocking_queue_push(queue_t *q, void *data) {

```
/* queue_push() adds the element to the queue;
   queue_push() is not thread-safe */
   queue_push(q, data);
```

Blocking Bounded Queue (v2)

void *blocking_queue_pop(queue_t *q) {

```
/* queue_pop() pops the top element;
   queue_pop() is not thread-safe */
void *data = queue pop(q);
```

Deadlock

```
void up() {
   pthread_mutex_lock(&mutex);
   ct++;
```

}

Four Conditions of Deadlock

In order to guarantee deadlock, four conditions <u>must</u> be true:

Dinning Philosophers Problem

- Five philosophers: Five silent philosophers sit around a table with a bowl of spaghetti.
- Five forks: A fork is placed between each pair of adjacent philosophers.

- **Two states**: Philosophers alternate between thinking and eating.
- **Condition**: To eat, a philosopher must have two forks: the fork to his right and the fork to his left.

Deadlock

