

### Heap Allocation

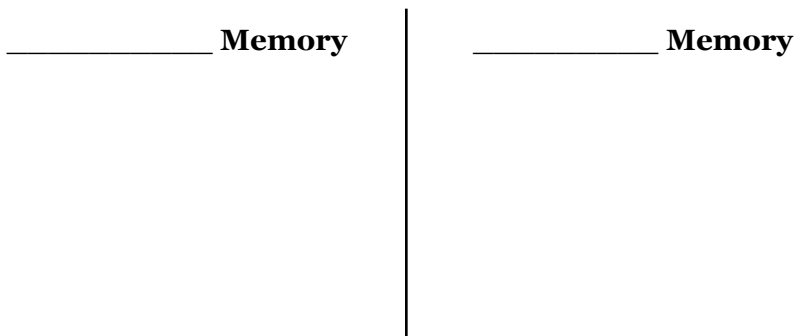
Up until now, we have arbitrarily placed memory with the process page table – however, all modern Operating Systems (OSes) organize the memory of a process in a predictable way:

06/memory-addr.c	
5	int val;
6	printf("&val: %p\n", &val);
7	
8	void *ptr = malloc(0x1000);
9	printf("&ptr: %p\n", &ptr);
10	printf(" ptr: %p\n", ptr);
11	
12	void *ptr2 = malloc(0x1000);
13	printf("&ptr2: %p\n", &ptr2);
14	printf(" ptr2: %p\n", ptr2);
15	
16	int arr[4096];
17	printf("&arr: %p\n", &arr);
18	
19	return 0;

Page Table:

....

As a programmer, we talk about these different regions of memory as different “types” of memory:



**Q1:** What if we access memory beyond the end of our heap? (Or any other region not allocated in our page table?)

### Efficient Use of Heap Memory

During the lifetime of a single process, we will allocate and free memory many times. Consider a simple program:

06/heap.c		
5	int *a = malloc(4096);	Heap v1: <small>(Without reuse after free)</small>
6	printf("a = %p\n", a);	
7	free(a);	Heap v2: <small>(With reuse after free)</small>
8		
9	int *b = malloc(4096);	Heap v1: <small>(Without reuse after free)</small>
10	printf("b = %p\n", b);	
11		Heap v2: <small>(With reuse after free)</small>
12	int *c = malloc(4096);	
13	printf("c = %p\n", c);	Heap v1: <small>(Without reuse after free)</small>
14		
15	int *d = malloc(4096);	Heap v2: <small>(With reuse after free)</small>
16	printf("d = %p\n", d);	
17		Heap v1: <small>(Without reuse after free)</small>
18	free(b);	
19	free(c);	Heap v2: <small>(With reuse after free)</small>
20		
21	int *e = malloc(5000);	Heap v1: <small>(Without reuse after free)</small>
22	printf("e = %p\n", e);	
23		Heap v2: <small>(With reuse after free)</small>
24	int *g = malloc(10);	
25	printf("g = %p\n", g);	Heap v1: <small>(Without reuse after free)</small>
26		
27	int *g = malloc(10);	Heap v2: <small>(With reuse after free)</small>
28	printf("g = %p\n", g);	

**Q2:** How much memory is used if we **do not** reuse memory?

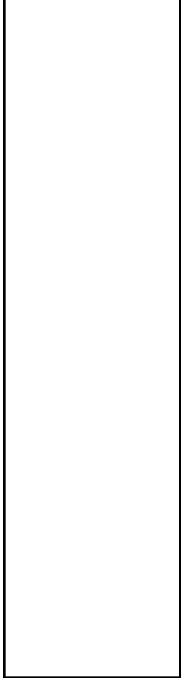
**Q3:** How much memory is used with **optimal** reuse of memory?

- What happens to our memory over time?
- When we have “holes” in our heap, how do we decide what hole to use?

## Data Structures for Heap Management

When we manage heap memory, we need to use memory to help us store memory:

- Overhead:
- Allocated Memory:

06/heap.c		
5	<code>int *a = malloc(4096);</code>	<b>Heap w/ Data Structures:</b> <i>(Without reuse after free)</i> 
6	<code>printf("a = %p\n", a);</code>	
7	<code>free(a);</code>	
8		
9	<code>int *b = malloc(4096);</code>	
10	<code>printf("b = %p\n", b);</code>	
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## Metadata-based Approach to Memory Storage

## Allocation Internals

Every process has a single heap starting point and a heap ending point in its virtual memory space that is provided by the Operating System.

- The initial heap size is: \_\_\_\_\_
  -
- A process grows/shrinks its heap using:  
`void *sbrk(intptr_t increment);`
- **MP3** (“malloc”) is released Friday and will have you build your own malloc, using the sbrk call, and require you to efficiently re-use memory just like the Linux kernel does!
  - EC Deadline:
  - Deadline #1:
  - Deadline #2:

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## Implementation Considerations

1. [Runtime]:
2. [Block Splitting]:
3. [Block Coalescing]: