

**Sample Programs:****04cr.c**

```

16 for (unsigned int c = 0; c < SIZE; c++) {
17     for (unsigned int r = 0; r < SIZE; r++) {
18         array[(r * SIZE) + c] = (r * SIZE) + c;
19     }
20 }

```

**04rc.c**

```

16 for (unsigned int r = 0; r < SIZE; r++) {
17     for (unsigned int c = 0; c < SIZE; c++) {
18         array[(r * SIZE) + c] = (r * SIZE) + c;
19     }
20 }

```

Running Times:     **04cr.c** (Program #1):**04rc.c** (Program #2):**Caching Strategies: Keeping Data Close**

**In working with memory in any computer system, we want to access it as quickly as possible.** However, space is extremely limited in the fastest memory, so we need strategies on what data to keep close. General Purpose Memory:

- CPU Registers:
- CPU Cache (i7-12700K, Released Q4'21):
- RAM:

**Key Idea: Locality of Reference****System Memory: Limited, Shared, and Simple**

- 1.
- 2.
- 3.

To help us to begin to organize this RAM, we divide the RAM up into chunks called \_\_\_\_\_.

On Linux, find the size of a page:

```
# getconf PAGESIZE
```

- On almost every modern system, a page is \_\_\_\_\_ KB.

**Virtual Memory:**

Modern systems provide an abstraction between the \_\_\_\_\_ and \_\_\_\_\_:

1. A \_\_\_\_\_ translates a \_\_\_\_\_ into a physical address.
2. Every memory address is made up of the \_\_\_\_\_ and the \_\_\_\_\_:
3. Virtual Memory is **NOT shared** between processes/apps:
4. **EVERY memory address** you have ever seen is a virtual memory address!

Let's explore a sequence of allocations using a page table:

P1 Page Table:	RAM	P2 Page Table:	P3 Page Table:

Allocation Sequence:

1. Process #1 (P1): `a = malloc(3 * 4096)`
2. Process #3 (P3): `b = malloc(5 * 4096)`
3. Process #1 (P1): `c = malloc(2 * 4096)`
4. Process #3 (P3) exits.
5. Process #2 (P2): `d = malloc(4 * 4096)`
6. Process #2 (P2): `e = malloc(5 * 4096)`
7. Process #1 (P1): `a = realloc(a, 5 * 4096)`

**With a virtual memory system:**

- Can we meet all of the allocation requests?
- Are we limited to just RAM?

### Advantages of a Virtual Memory System:

1.

05.c	
16	<code>printf(" Start of `array`: %p\n", array);</code>
17	<code>printf(" End of `array`: %p\n", &amp;(array[(SIZE * SIZE) -1]));</code>

2.

3.

### Simple Simulation of Page Tables with Disk Pages

RAM:	P1 Page Table:	Disk Pages:	
[0]:	[0]:	...	1: Load Program
[1]:	[1]:	./programCode (1/5)	2: Run PC, pg1: - malloc(4000)
[2]:	[2]:	./programCode (2/5)	3: Run PC, pg2: - malloc(10000)
[3]:	[3]:	./programCode (3/5)	- Open hiddenImage.png
	[4]:	./programCode (4/5)	- Read all of image
	[5]:	./programCode (5/5)	4: Run PC, pg3: - Access OG 4 KB
	[6]:		- Finish program
	[7]:		
	[8]:		
	[9]:		
	[10]:		
	[11]:	hiddenImage.png	
	[12]:	hiddenImage.png	
	[13]:	hiddenImage.png	
	[14]:		
	[15]:	...	

**Q1:** What is the range of possible file sizes for `hiddenImage.png`?

**Q2:** What is the range of possible file sizes for `./programCode`?

**Q3:** What is the size of the heap immediately before the program finishes?