CS 240 Week 5: File Formats and Memory

Computer Systems

CS 240, Spring 2021 - Week 5

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Virtual Memory - Limited Resources

In our memory hierarchy, we know memory pages are loaded into the processor cache -- nearly always a page at a time -- but what if we want more storage than our RAM?

RAM: [0]: [1]: [2]: [3]:	P1 Page Table: [0]: [1]: [2]: [3]: [4]: [5]: [6]: [7]: [8]: [9]: [10]: [11]: [12]: [13]: [14]: [15]:	Disk Pages:/programCode (1/5) ./programCode (2/5) ./programCode (3/5) ./programCode (4/5) ./programCode (5/5) ./programCode (5/5) ./programCode (5/5) ./programCode (5/5) ./programCode (5/5) ./programCode (5/5)	1: Load Program 2: Run PC1 - malloc(4000) 3. Run PC2: - malloc(10000) - Open hiddenImage.png - Read all of image 4: Run PC3 - Access OG 4 KB - Finish program
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...assume that this system has 4 KB pages.

Q1: What is the range of possible file sizes for hiddenImage.png?

Q2: What is the range of possible file sizes for ./programCode?

Q3: What is the size of the heap immediately before the program finished?

Page Eviction/Replacement Strategies:

When we need to remove a page from RAM and store it on disk, how do we decide which page to remove?

[1]:

Pag	де Ассе	esses:	17	33	40	17	43	8	99	33	99	17
R A												
M												

[2]:

Pag	де Ассе	esses:	17	33	40	17	43	8	99	33	99	17
R A												
M												

[3]:

Pag	де Ассе	esses:	17	33	40	17	43	8	99	33	99	17
_												
R A												
M												

[4]:

Pag	де Ассе	esses:	17 3	33	40	17	43	8 9	9 33	99	17
R A											
M											

[5]:

Pag	де Ассе	esses:	17 3	3 40	17	43	8 9	9 33	99	17
R										
M										

Remember: We count the number of **page faults**, or the number of times we needed to replace the contents of a page of RAM, as a metric to determine the performance of an eviction scheme.

To accomplish page eviction, every page table entry has many status bits:

- 1.
- 2.
- 3.
- 4.
- 5.

Segmentation Faults

RAM: [0]: [1]: [2]: [3]:	P1 Page Table: //programCode (PC1) //programCode (PC2) //programCode (PC3) //programCode (PC4) //programCode (PC5) Heap - 4 KiB Heap - 4 KiB Heap - 4 KiB	Disk Pages:	1: Load Program 2: Run PC1 - malloc(4000) 3. Run PC2: - malloc(10000) - Open hiddenImage.png - Read all of image 4: Run PC3 - Access OG 4 KB - Finish program
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Q: What happens if we access the address of the **bold** entry inside of our page table above?

We know the following:

- Every page table entry takes 4B of space
- Every address between 0 2⁶⁴ is addressable.
- Each 4 KB page needs its own page table entry.

Q: How big is our page table?

Q: How many entries can our page table have if we need to fit a page table into a single page?

Q: What can be to make this happen?