Synchronization, Dining Philosophers, and Deadlock

CS 240 - The University of Illinois

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Synchronization

Synchronization Technique #1

`pthread_mutex_init`: Creates a new lock in the “unlocked” state.

`pthread_mutex_lock(pthread_mutex_t *mutex)`:

- When `` `mutex` `` is unlocked, change the lock to the “locked” state and advance to the next line of code.
- When `` `mutex` `` is locked, this function **blocks** execution until the lock can be acquired.

`pthread_mutex_unlock`: Moves the lock to the “unlocked” state.

`pthread_mutex_destory`: Destroys the lock; frees memory.

09/count-with-lock.c

```
5  pthread_mutex_t lock;
6  int ct = 0;
7
8  void *thread_start(void *ptr) {
9    int countTo = *((int *)ptr);
10
11   int i;
12   for (i = 0; i < countTo; i++) {
13     pthread_mutex_lock(&lock);
14     ct = ct + 1;
15     pthread_mutex_unlock(&lock);
16   }
17
18   return NULL;
19 }
```

Synchronization Technique #2

`pthread_cond_init`: Create a new conditional variable.

`pthread_cond_wait(pthread_cond_t *cond, pthread_mutex_t *mutex)`: Performs two different synchronization actions:

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`pthread_cond_signal(pthread_cond_t *cond)`: Unblocks “at least one thread” that is blocked on **`cond`** (if any threads are blocked; otherwise an effective “NO OP”).

`pthread_cond_broadcast(pthread_cond_t *cond)`: Unblocks **ALL** threads blocked on **`cond`**.

`pthread_mutex_destory`: Destroys the lock; frees memory.

09/count-with-lock.c

```c
int things[THINGS_MAX];
int things_ct = 0;

void *producer(void *vptr) {
  while (1) {
    pthread_mutex_lock(&lock);

    // Cannot produce until there's space:
    while (things_ct >= THINGS_MAX) {
      pthread_cond_wait(&cond, &lock);
    }

    // Produce a thing:
    things[things_ct] = rand();
    printf("Produced [%d]: %d\n", things_ct, things[things_ct]);
    things_ct++;

    // Signal any waiting consumers:
    pthread_cond_broadcast(&cond);

    pthread_mutex_unlock(&lock);
  }
}
```

Synchronization Technique #3

`sem_init`: Creates a new semaphore with a specified “value”.

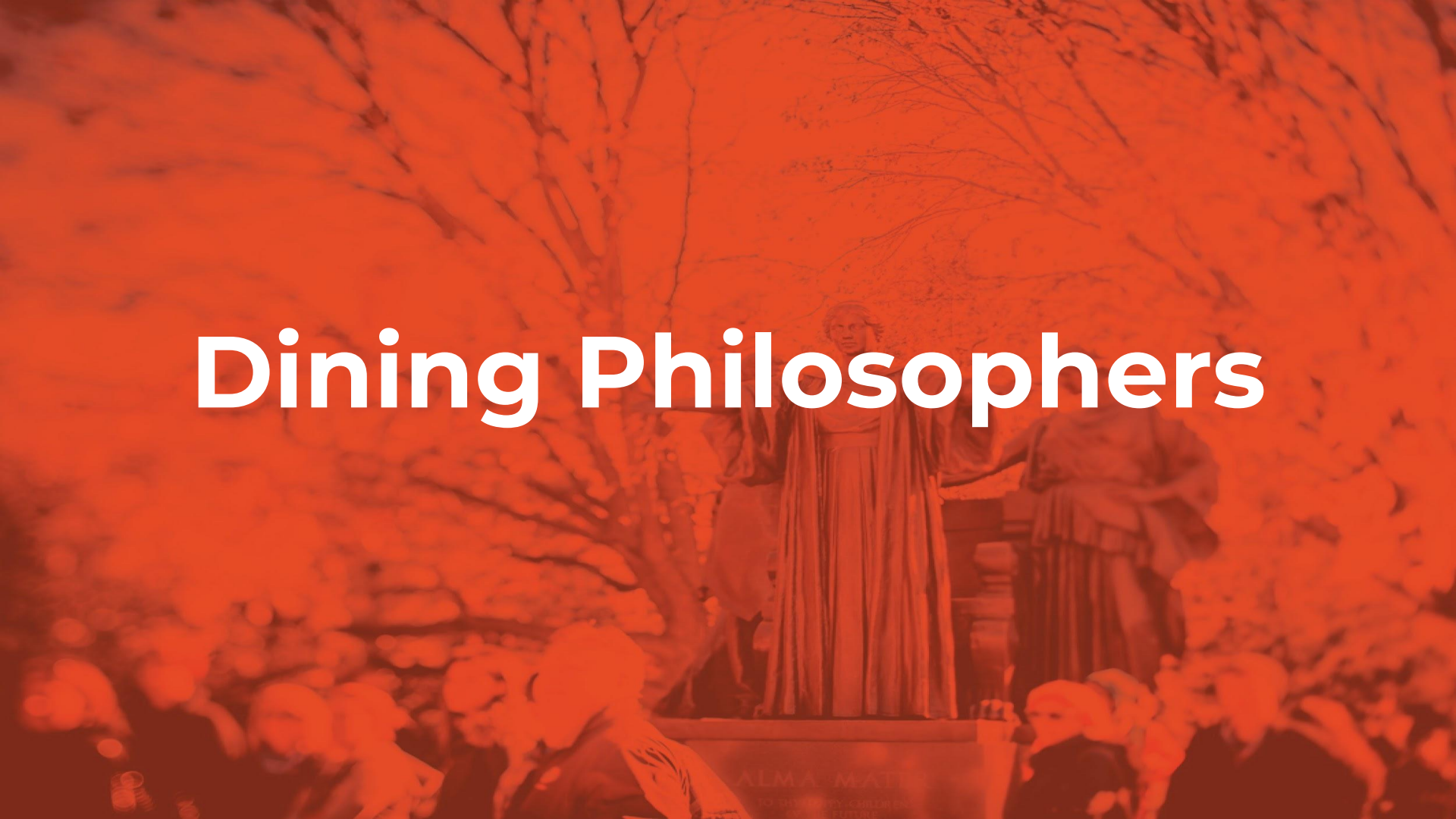
`sem_wait`: When the value is greater than zero, decreases the value and continues. Otherwise, **blocks** until the value is non-zero.

`sem_post`: Increments the value by one.

`sem_destroy`: Destroys the semaphore; frees memory.

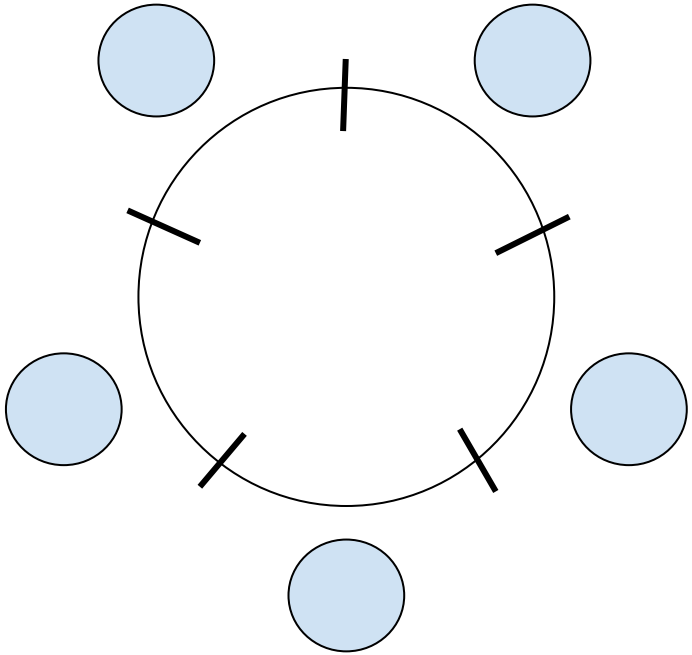
Critical Sections

Critical Section



Dining Philosophers

Dining Philosophers



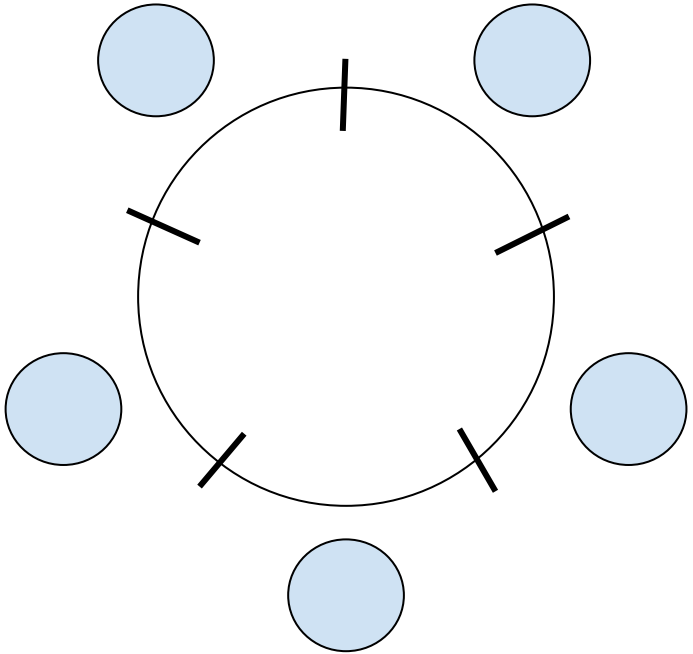
09/dining-philosopher.c

```
16    while (1) {
17      printf("%d is thinking...\n", id);
18
19      // Get left chopstick:
20      printf("%d is reaching for the left chopstick
              (chopstick=%d)...\n", id, left_chopstick_id);
21      pthread_mutex_lock(&locks[left_chopstick_id]);
22      printf("%d has the left chopstick
                (chopstick=%d).\n", id, left_chopstick_id);
23
24      // Get right chopstick:
25      printf("%d is reaching for the right chopstick
             (chopstick=%d)...\n", id, right_chopstick_id);
26      pthread_mutex_lock(&locks[right_chopstick_id]);
27      printf("%d has the right chopstick
                (chopstick=%d).\n", id, right_chopstick_id);
```

09/dining-philosopher.c

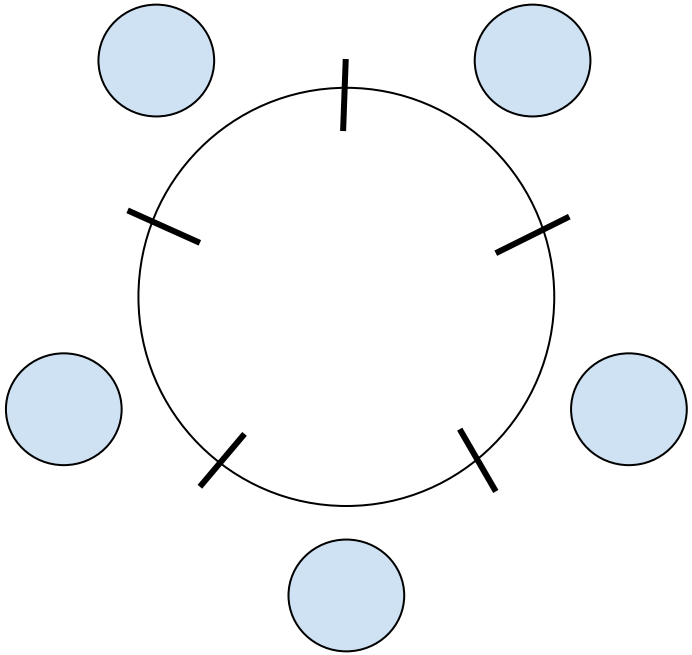
```
28
29      // Eat:
30      printf("%d is eating... 🍱🥢\n", id);
31
32      // Release chopsticks:
33      printf("%d is returning their chopsticks
                    (chopsticks: %d, %d)...\n", id,
                    left_chopstick_id, right_chopstick_id);
34      pthread_mutex_unlock(&locks[right_chopstick_id]);
35      pthread_mutex_unlock(&locks[left_chopstick_id]);
36    }
```

```
while (1) {
  think();
  lock_left(&mutex);
  lock_right(&mutex);
  eat();
  release_right(&mutex);
  release_left(&mutex);
}
```

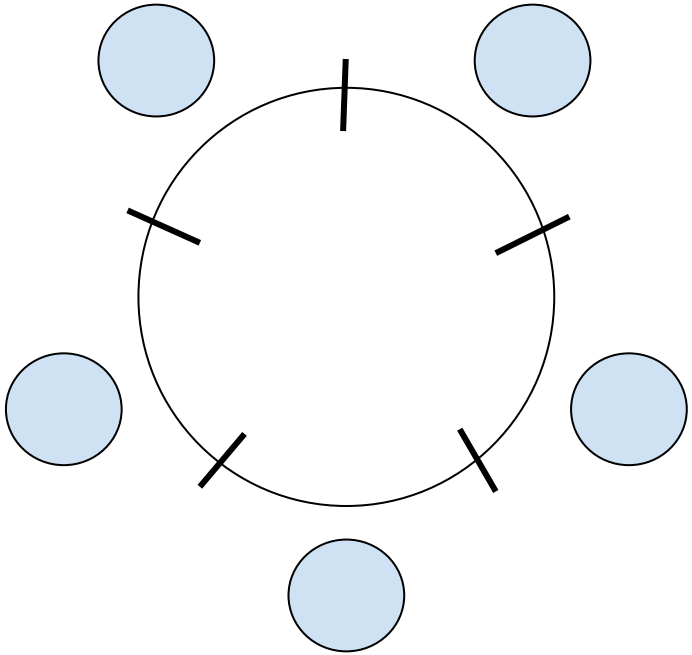


Four Necessary Conditions of Deadlock (“Hoffman Conditions”):

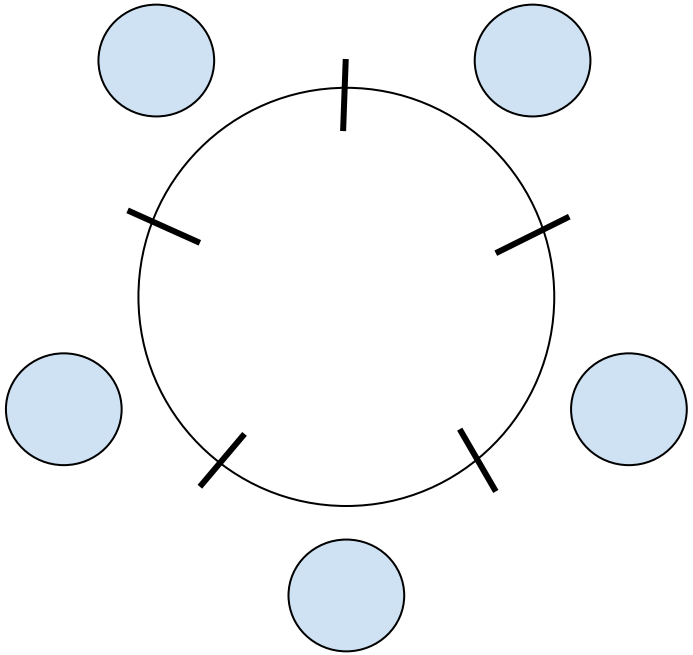
Mutual Exclusion



Circular Wait



Hold and Wait



No Preemption