

1 Problems

1. Understanding MIPS programs

```

flathead:
    addi $t0, $a2, 1
loop:
    bge $t0, $a1, exit
    mul $t1, $t0, 4
    add $t1, $t1, $a0
    lw  $t2, 0($t1)
    sub $t1, $t1, 4
    sw  $t2, 0($t1)
    addi $t0, $t0, 1
    j loop
exit:
    jr  $ra

```

- (a) Translate the `flathead` function above into a high-level language like C or Java.

Solution:

This function does stuff with `$a0`, `$a1` and `$a2`, which indicates that there are three arguments. Multiplying $(\$a2 + 1)$ by 4 suggests that `$a2` is an index, and adding that to `$a0` for a “lw” means `$a0` is an array of 32-bit values (most likely integers or floating-point numbers). The loop exits when $\$t0 \geq \$a1$, suggesting that `$a1` is some kind of maximum index. Since nothing is saved in `$v0` before the “jr”, there is no return value.

The problem says not to use gotos, so you should have translated the loop instructions into a higher-level “for” or “while” structure. You should also show array manipulations, instead of pointer arithmetic and dereferences.

```

void flathead(int a0[ ], int a1, int a2)
{
    int t0 = a2 + 1;
    while (t0 < a1) {
        a0[t0 - 1] = a0[t0];
        t0++;
    }
}

```

- (b) Describe briefly, in English, what this function does.

Solution:

This function takes an array `$a0` and shifts elements $(\$a2 + 1)$ through `$a1` into positions `$a2` through `$a1 - 1`. In other words it shifts those elements “downward” or “leftward”. You can also think of this as deleting element `$a2`.

2. Compute the result of the following expression and write it in hexadecimal notation:

`0xdeadbeef | (0x81 << 22)`

`0x81 = 1000 0001`, shifted left 22 bits gives you: `0010 0000 0100 0000 0000 0000 0000 0000`

Since in the above bits are only set in the 1st and 3rd hexadecimal characters, those are the only ones that change (recall that ORing with 0 preserves values).

```

    1101   E  1010   D    B    E    E    F
OR 0010 0000 0100 0000  0000 0000 0000 0000
= 1111   E  1110   D    B    E    E    F  = FEEDBEEF

```

3. C to MIPS translation

C Code provided for your reference:

```
int count(int A[], int n)
{
    int i;
    int num = 0;

    for(i = 0; i < n; i++){
        if (test(A[i]) == 1) {
            num++;
        }
    }
    return num;
}
```

Solution

The hard part here is preserving registers correctly. The `count` function acts as both a caller and callee. It has to save all the callee-saved registers that it uses and restore them before returning, but it is also responsible for preserving any caller-saved registers that it needs across the call to `test`. Regardless of whether you use the caller-saved registers, callee-saved registers or a combination of them, you'll have to push stuff onto the stack.

One example implementation is shown on the next page. It uses caller-saved registers for the temporary values `i` and `num`, so those must be saved and restored before and after the call to `test`. In addition the base address of the array and the array size, which are passed as arguments `$a0` and `$a1`, must also be saved. Finally, recall that the callee-saved register `$ra` has to be preserved because of the nested `jal`.

Translated MIPS code:

```
count:
    subu $sp, $sp, 20
    sw   $ra, 0($sp)      # Save $ra before jal
    li   $t0, 0           # $t0 = i
    li   $t1, 0           # $t1 = num
loop:
    bge  $t0, $a1, exit   # while (i < n)
    sw   $t0, 4($sp)      # Save caller-saved registers
    sw   $t1, 8($sp)      # that we need
    sw   $a0, 12($sp)
    sw   $a1, 16($sp)
    mul  $t2, $t0, 4       # We don't need to save $t2
    addu $t2, $a0, $t2
    lw   $a0, 0($t2)       # $a0 = A[i]
    jal  testfunc          # $v0 = testfunc(A[i])
    lw   $t0, 4($sp)      # Restore saved registers
    lw   $t1, 8($sp)
    lw   $a0, 12($sp)
    lw   $a1, 16($sp)
    beq  $v0, $0, next
    addi $t1, $t1, 1       # num++
next:
    addi $t0, $t0, 1       # i++
    j    loop
exit:
    move $v0, $t1          # return num
    lw   $ra, 0($sp)      # Restore $ra
    addiu $sp, $sp, 20
    jr   $ra
```