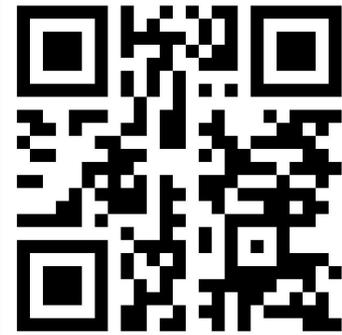


# Announcements

1. MP Mosaics EC due Today!
2. Exam 2 Grades In Process
3. Extra Credit Survey due Wednesday
  - a. 10 more people to hit threshold!



Extra Credit Survey!



Join Code: **225**

**Warm-Up Question:** What is the minimum number of nodes needed to make an AVL tree with height 3?



# AVL Tree Analysis

# Learning Objectives

---

1. Prove the upper bound on an AVL Tree's Height
2. Enumerate the runtime of Insert, Find, Remove, and Traverse for every Data Structure Seen so far



# AVL Tree Analysis

---

**We know:** insert, remove and find runs in: \_\_\_\_\_.

**We will show that:**  $h$  is \_\_\_\_\_.

Goal: Find the lowest upper bound, big O, on the maximum height of an AVL tree given the number of nodes.

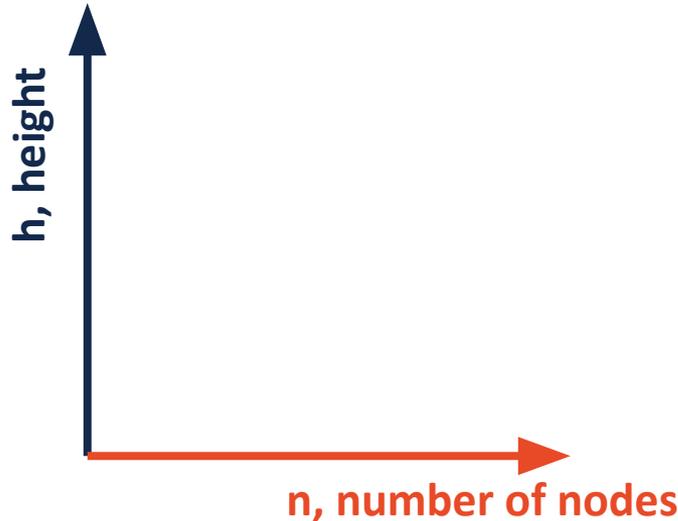


# Big O - Lowest Upper Bound

**Definition of big-O:** Given two functions  $f(n)$  and  $g(n)$ , we say that  $f(n)$  is  $O(g(n))$  if there exist constants  $c > 0$  and  $n_0 \geq 0$

such that  $f(n) \leq c \cdot g(n)$  for all  $n \geq n_0$ .

...or, with pictures:



## New Goal

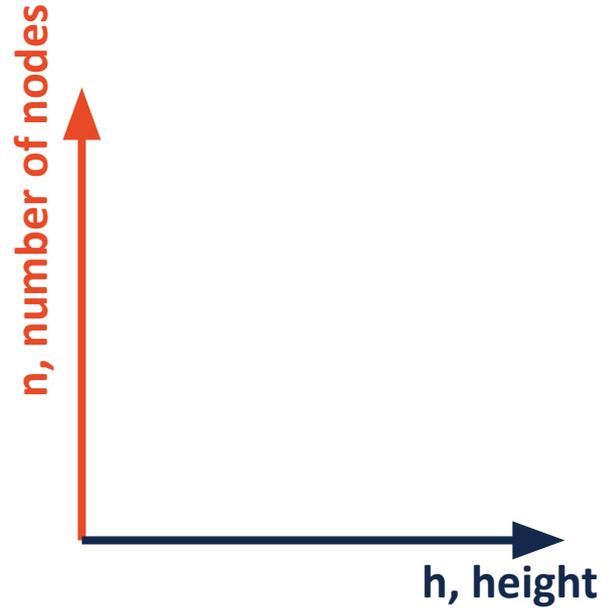
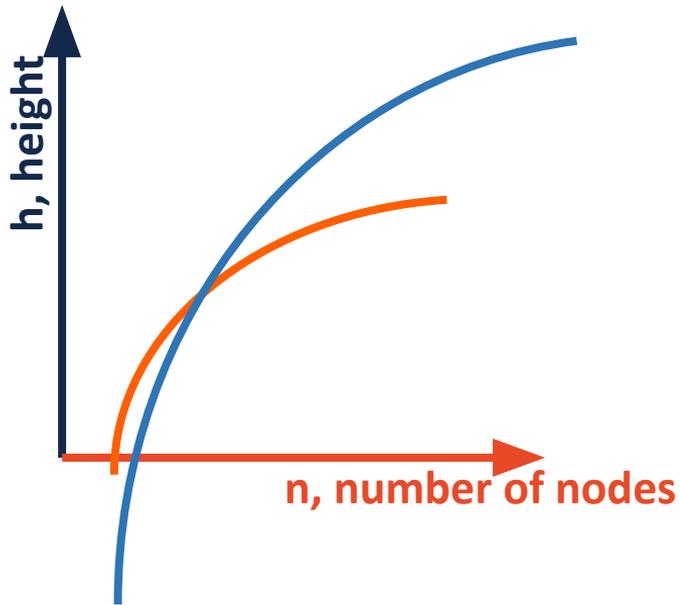
---

Goal: Find the lowest upper bound, big O, on the maximum height of an AVL tree given the number of nodes.

New Goal: Find the highest lower bound on the minimum number of nodes in an AVL tree given the height,  $N(h)$



# Invert and Find Highest Lower Bound



# Small Examples



$h = 0$

$h = 1$

$h = 2$

$h = 3$

$h = 4$



# Recurrence

---

$$N(h) = 1 + N(h - 1) + N(h - 2)$$



# Conclusion

---

$$h = O(\log(n))$$



# Summary

---

## **AVL Trees**

Max height:  $1.44 * \log(n)$

Rotations:

## **Red-Black Trees**

Max height:  $2 * \log(n)$

Constant number of rotations on insert, remove, and find



	Unsorted Array	Sorted Array	Unsorted List	Sorted List	Binary Tree	BST	AVL
Find							
Insert							
Remove							
Traverse							

**5:00**



	Unsorted Array	Sorted Array	Unsorted List	Sorted List	Binary Tree	BST	AVL
<b>Find</b>	$O(n)$	$O(\log n)$	$O(n)$	$O(n)$	$O(h)/O(n)$	$O(h)/O(n)$	$O(h)/O(\log n)$
<b>Insert</b>	$O(n)/O^*(1)$	$O(n)$	$O(1)$	$O(n)$	$O(h)/O(n)$	$O(h)/O(n)$	$O(h)/O(\log n)$
<b>Remove</b>	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(h)/O(n)$	$O(h)/O(n)$	$O(h)/O(\log n)$
<b>Traverse</b>	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$	$O(n)$

