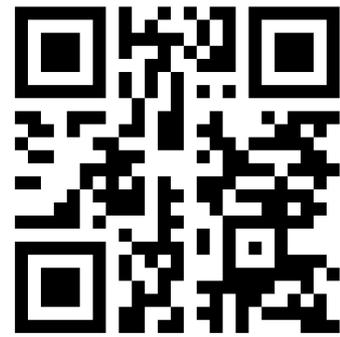


Announcements

1. MP1 Extra Credit due Tonight!

Exam 2 Review in Lab this week!

Prof. Solomon's OH are cancelled this week



Join Code: 225



Warm-Up Question: If you are searching through an infinite tree, should you use Breadth-First Search or Depth-First Search?



Iterative Deepening Search/Binary Search Trees

Learning Objectives

1. Understand Iterative Deepening
2. Know the runtime of iterative deepening on a binary tree
3. Understand the Dictionary ADT
4. Compare runtimes of operations on BST with binary trees



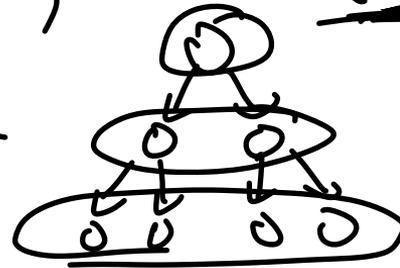
Tree Search

↳ Guarantees a solution
Breadth-First - in ∞ tree

Level Order Traversal, Queue, $O(n)$

Space Complexity

$O(2^h)$



Q: root
Q: left right

Depth-First -

Pre-Order Traversal, Stack, $O(n)$
Post-Order Traversal
In-Order Traversal

Time $O(n)$

Space $O(h)$

80% BFS
20% DFS

In an ∞ tree, can't guarantee a solution

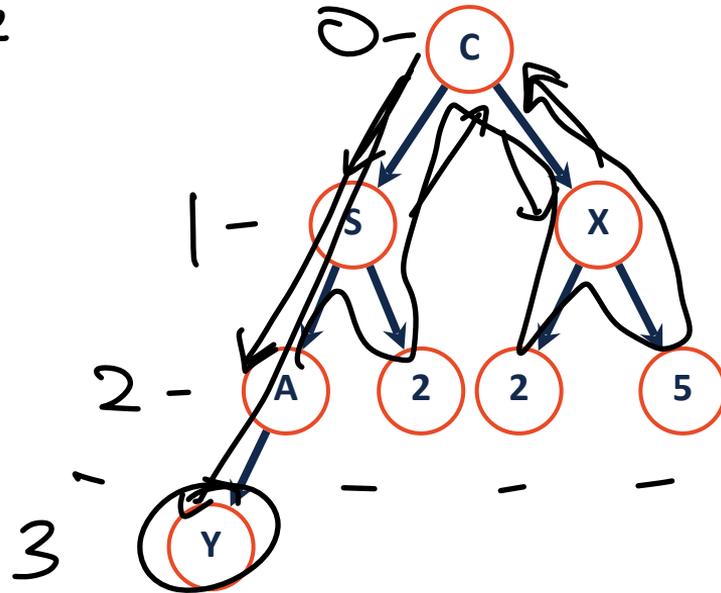


Iterative Deepening DFS

K - Current depth limit
 m - depth goal node

Goal is found is
 $K \geq m$

Ex. Goal $\rightarrow 4$
 $m \rightarrow 3$



Time
??

Space
 $O(h)$
 $O(m)$

$K=3$



Iterative Deepening Analysis

m - Goal depth

Level	0	1	2	...	m
Total Nodes	$2^0 = 1$	$1 + 2 = 3$ $2^0 + 2^1 = 3$	$1 + 2 + 4 = 7$ $2^0 + 2^1 + 2^2 = 7$...	

- 1. Guarantees solution
- 2. Saves space complexity

$\sum_{k=0}^m \sum_{j=0}^k 2^j$
 ↑ ↑
 # of times DFS is run For each level in DFS

$$\begin{aligned}
 &= \sum_{k=0}^m 2^{k+1} - 1 \\
 &= \sum_{k=0}^m 2^{k+1} - \sum_{k=0}^m 1 \\
 &= 2^{m+2} - 1 - m \\
 &\Rightarrow O(2^{m+2}) \Rightarrow O(4 \cdot 2^m) \Rightarrow O(n)
 \end{aligned}$$

Time Complexity

$O(n)$



Dictionary

Stores key -> value pairs

What applications of dictionaries can you think of?

Applications:

Words → Meanings / Definitions

Username → Passwords

Names → Numbers^{Phone}

UIN → Student Info

Flight No → Flight Info.

Array Compact Set of Integers → data at index

[0, 1, 2, 3...]

X[0, 2, 4]



Dictionary ADT

```
void Insert ( Key , Value )  
(Value, void) Remove ( Key )  
Value Find ( Key ) ←
```

Key musts be unique
Values dont need to be unique



```
1 #pragma once
2
3
4 class Dictionary {
5     public:
6         void insert(const K & key, V & value);
7         V remove(const K & key);
8         iterator find(const K & key) const;
9         iterator begin();
10        iterator end();
11        // ...
12
13    private:
14        // ...
15
16
17
18
19 };
```

What if my element
isn't in my dictionary?

What if you are trying to find a key that doesn't exist?

```
1 #pragma once
2
3
4 class Dictionary {
5     public:
6         void insert(const K & key, V & value);
7         V remove(const K & key);
8         V find(const K & key) const;
9         iterator begin();
10        iterator end();
11        // ...
12
13    private:
14        // ...
15        TreeNode * root_;
16
17
18
19};
```



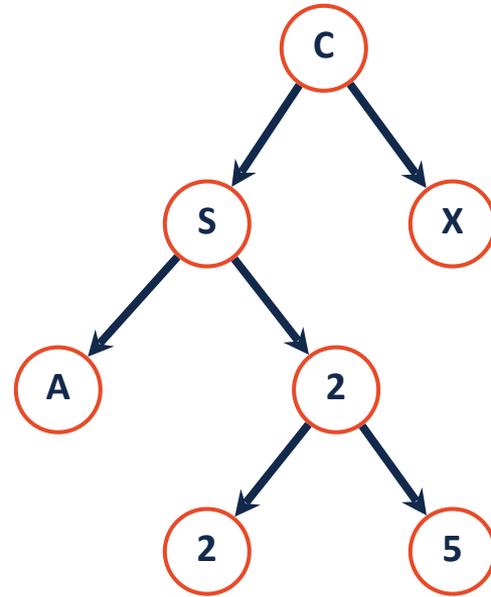
Binary Tree Definition

A binary tree T is either:

$$\underline{T = \emptyset}$$

OR

$$T = (\underline{r}, \underline{T}_L, \underline{T}_R)$$



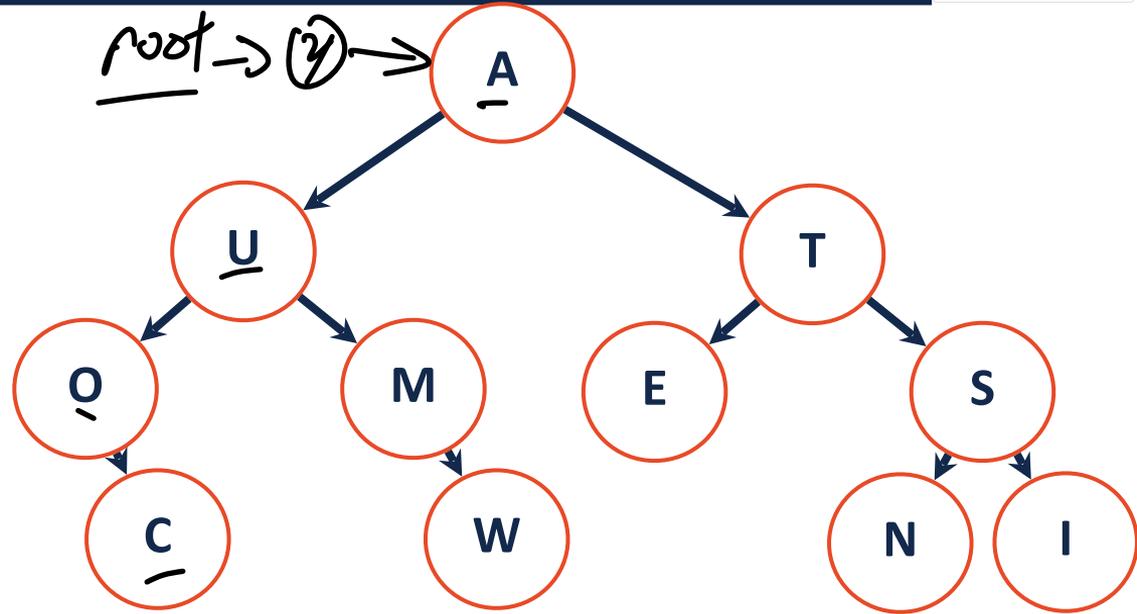
Searching through Binary Trees

1:00



Join Code: 225

Insert (key, value)
Remove (key)
value Find (key)



Function	Runtime
<u>Insert</u>	$O(1)$
Remove	$O(n)$
Find	$O(n)$



Binary Search Tree Definition

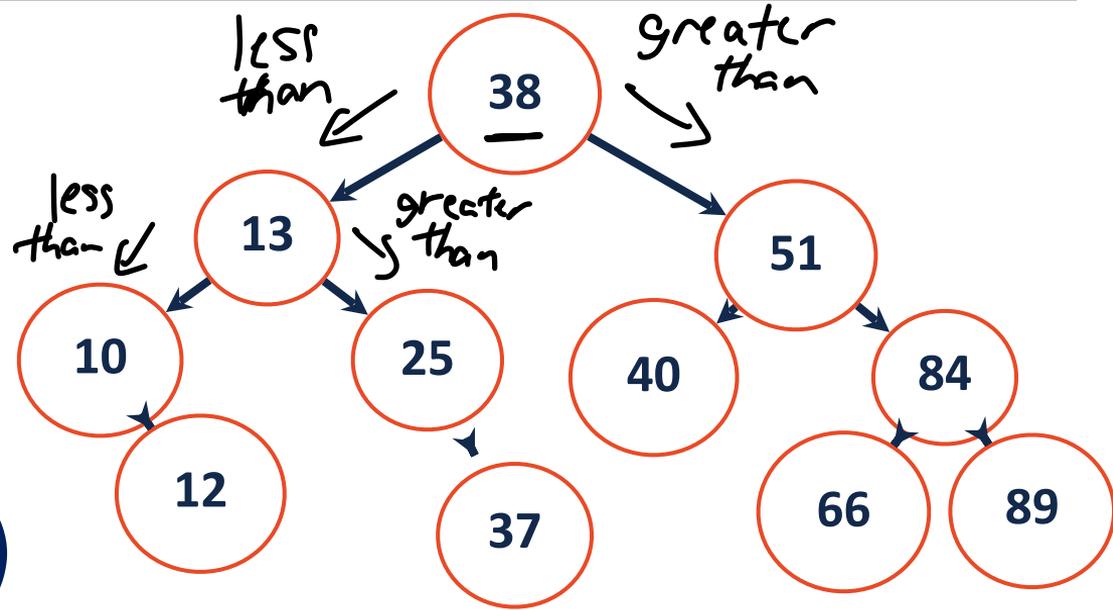
A binary search tree T is either:

$$\underline{T = \emptyset}$$

OR

$$\underline{T = (r, T_L, T_R)}$$

for $\forall x, x \in T_L, x < r$
 $\forall y, y \in T_R, y > r$



Binary Search Trees (BSTs)

How do my functions change given the structure of the BST?

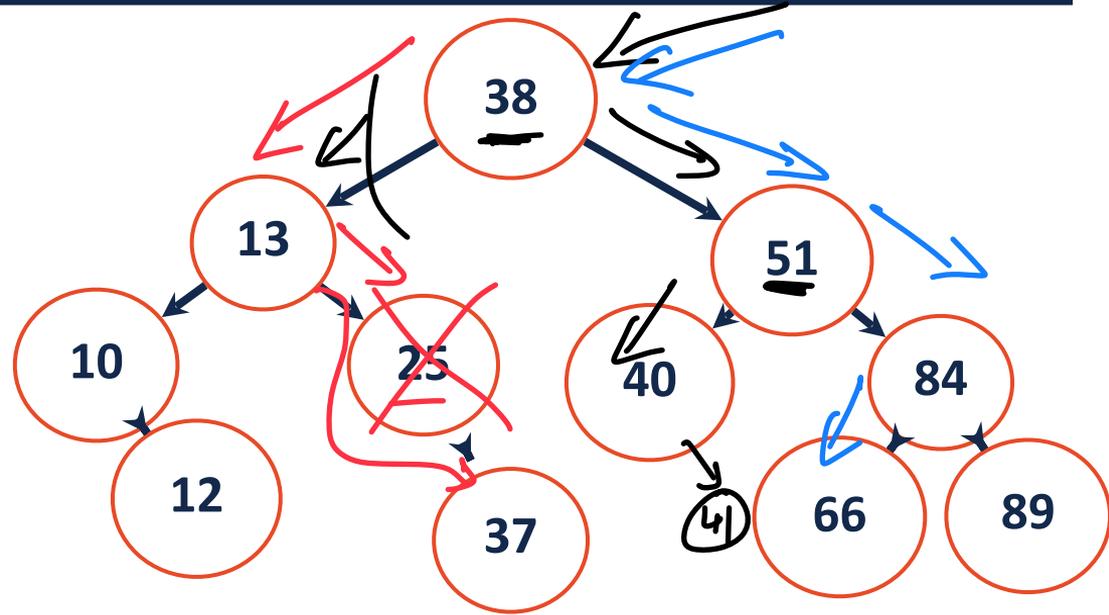
Insert: Insert 41

Find: Find(66)

-No traversal required

Removing: Remove (25)

Find (25)
Delete (25)

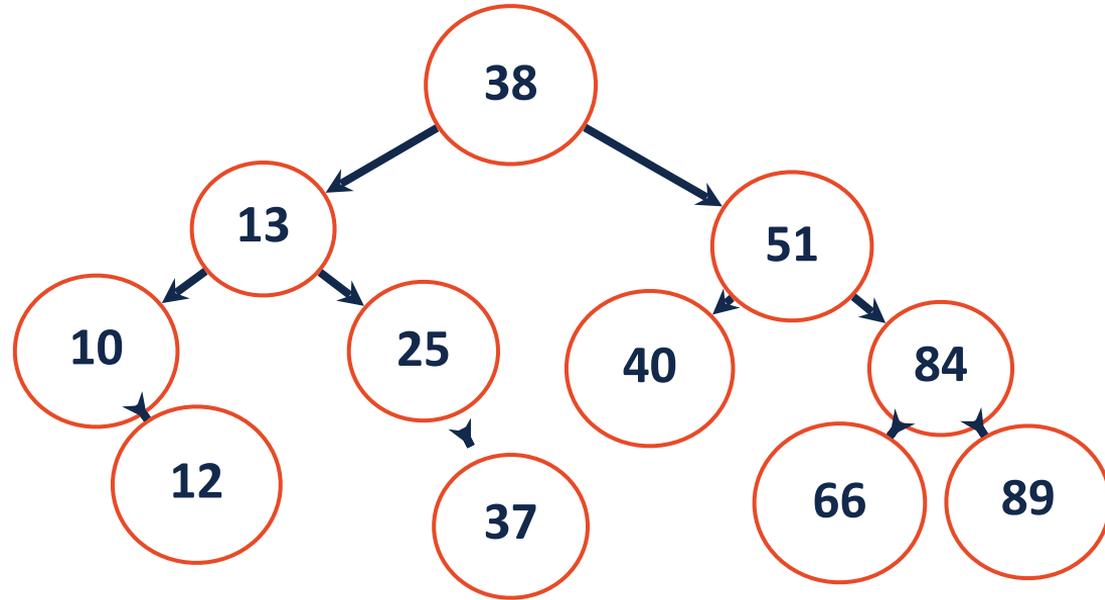


1:00



Join Code: 225

Insert (key, value)
Remove (key)
value Find (key)



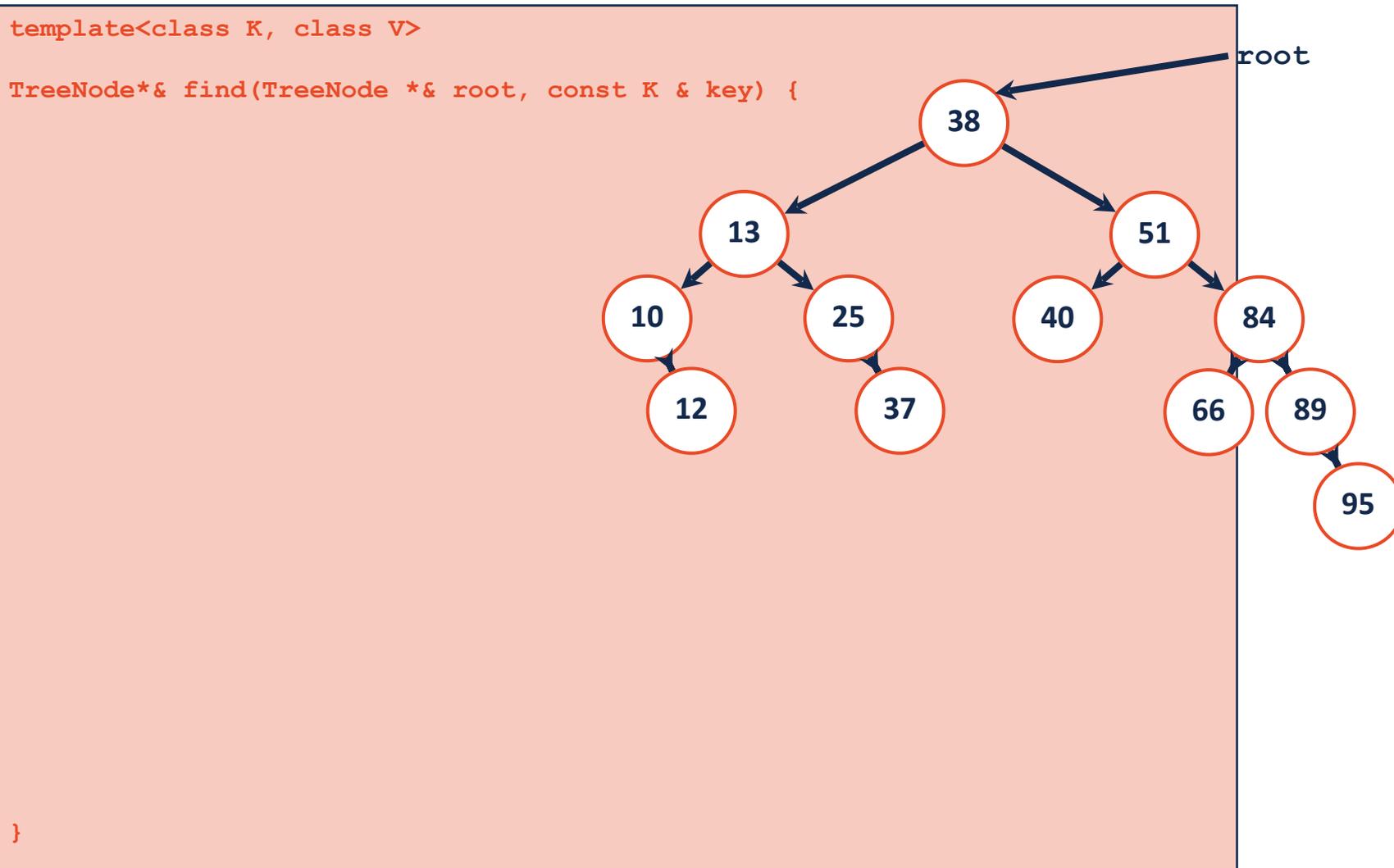
Function	Runtime
Insert	
Remove	
Find	



```
1 template<class K, class V>
```

```
2   TreeNode*& find(TreeNode *& root, const K & key) {
```

```
3  
4  
5  
6  
7  
8  
9  
10  
11  
12  
13  
14  
15  
16  
17  
18  
19  
20  
21  
22  
23  
24  
25  
26
```



```
1  template<class K, class V>
2  TreeNode*& find(TreeNode *& root, const K & key) {
3
4      if (root == NULL){
5          return root;
6      }
7
8      if (root->key == key){
9          return root;
10     }
11
12     if(key < root->key){
13         return find(root->left, key);
14     }
15
16     if(key > root->key){
17         return find(root->right, key);
18     }
19
20
21
22
23
24
25 }
26
```

