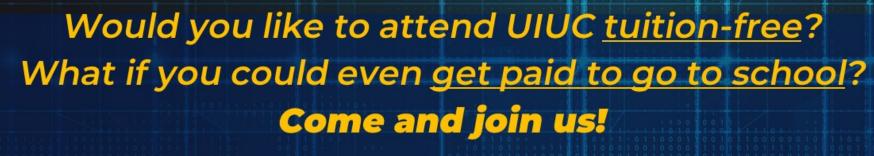
CS 225

Data Structures

February 24 – BBST Summary
G Carl Evans



Illinois Cyber Security Scholars Program
Info Session

Thursday, March 2, 2023 5:00 – 6:45 pm @ 2405 Siebel Center

PIZZA WILL BE PROVIDED!

To learn more, visit: https://publish.illinois.edu/cybersecurityscholars

Applicants must be US citizens or legal permanent residents.

Would you like to attend UIUC <u>tuition-free</u>?
What if you could even <u>get paid to go to school</u>?

Come and join us!

Illinois Cyber Security Scholars Program Info Session

Thursday, March 2, 2023

5:00-6:45 pm @ 2405 Siebel Center

PIZZA WILL BE PROVIDED!

This NSF-funded CyberCorps Scholarship for Service program provides scholarships to admitted, full-time UIUC students interested in cybersecurity.

Benefits include:

- Full tuition and fee waiver
- Stipend, allowance, and professional development funds
- Research experience with security experts
- Expertise in a growing professional field

Application Deadline:

March 31, 2023

To learn more, visit:

publish.illinois.edu/cybersecurityscholars

Applicants must be US citizens or legal permanent residents.

Summary of Balanced BST

Pros:

- Running Time:

- Improvement Over:

- Great for specific applications:

Summary of Balanced BST

Cons:

- Running Time:

- In-memory Requirement:

Q: Consider points in 1D: $p = \{p_1, p_2, ..., p_n\}$what points fall in [11, 42]?

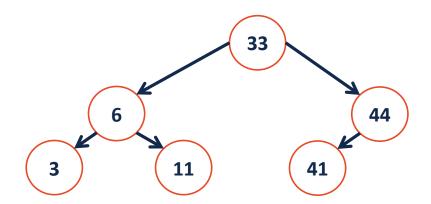
Tree construction:

Balanced BSTs are useful structures for range-based and nearest-neighbor searches.

Q: Consider points in 1D: $p = \{p_1, p_2, ..., p_n\}$what points fall in [11, 42]?



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C++ provides us a balanced BST as part of the standard library:

```
std::map<K, V> map;
```

```
V & std::map<K, V>::operator[]( const K & )
```

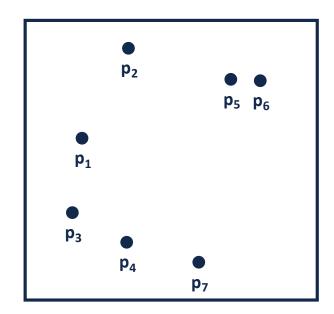
```
V & std::map<K, V>::operator[]( const K & )
std::map<K, V>::erase( const K & )
```

```
iterator std::map<K, V>::lower_bound( const K & );
iterator std::map<K, V>::upper_bound( const K & );
```

Consider points in 2D: $p = \{p_1, p_2, ..., p_n\}$.

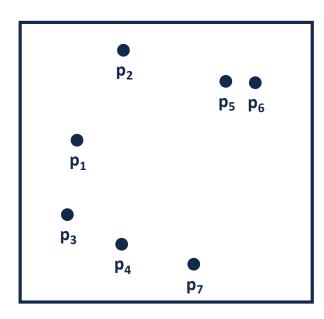
Q: What points are in the rectangle: $(x_1, y_1), (x_2, y_2)$]?

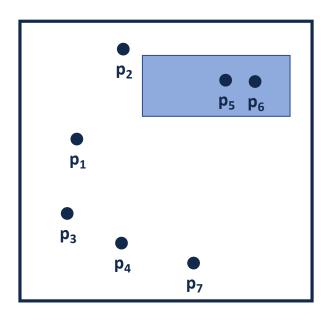
Q: What is the nearest point to (x_1, y_1) ?

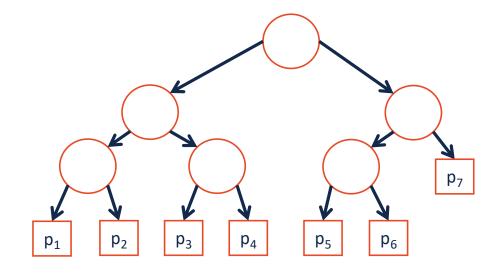


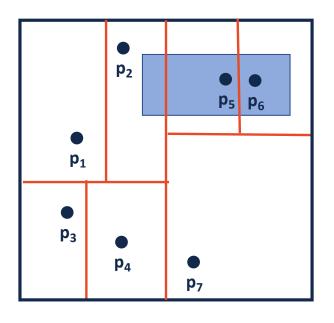
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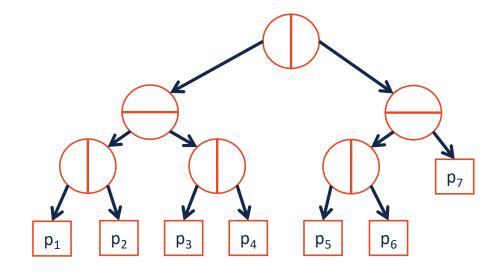
Tree construction:

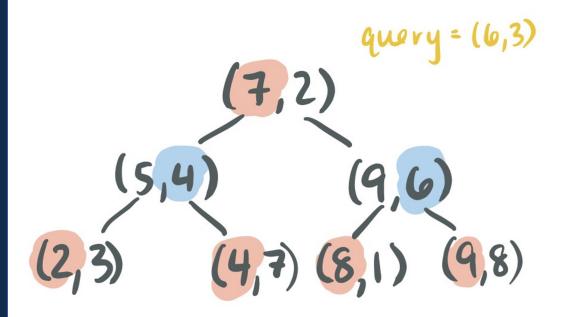


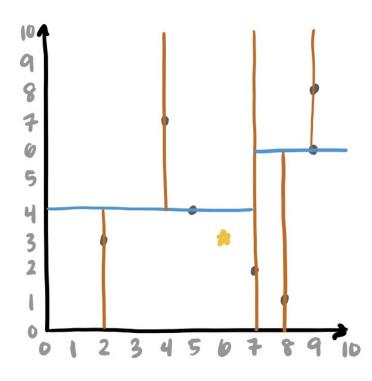


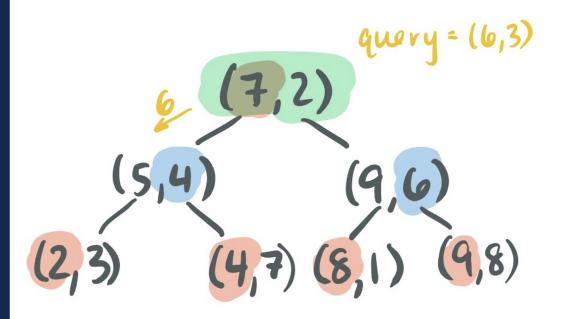


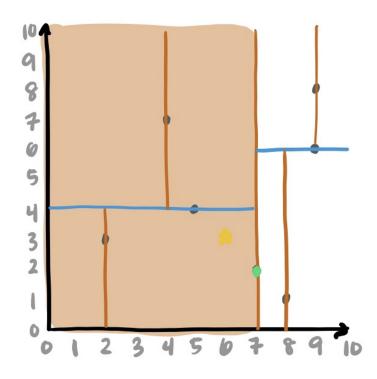


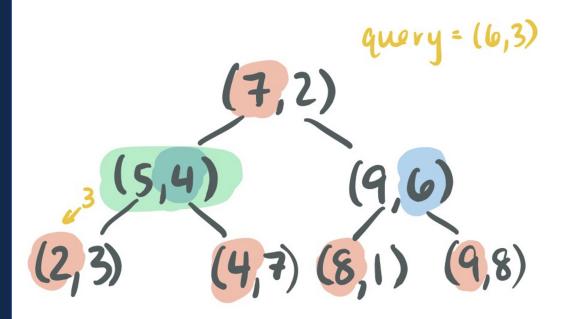


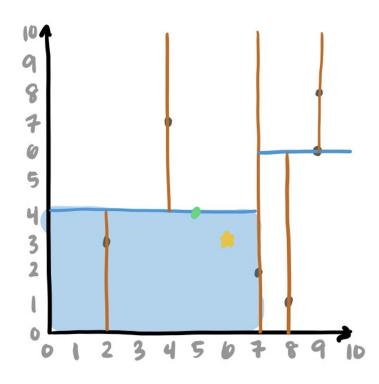


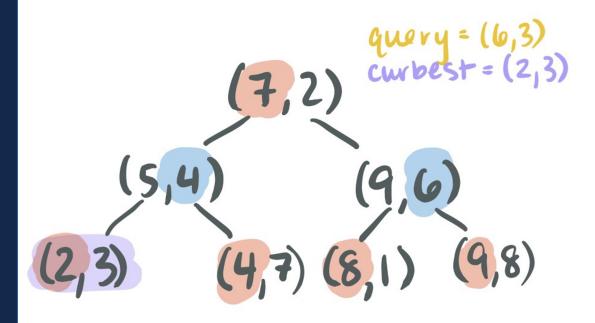


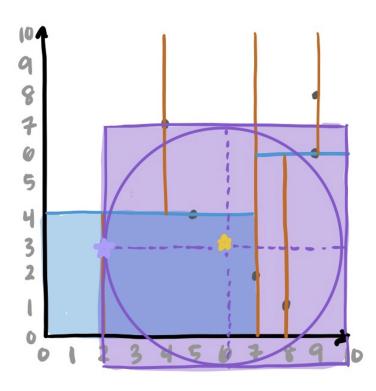




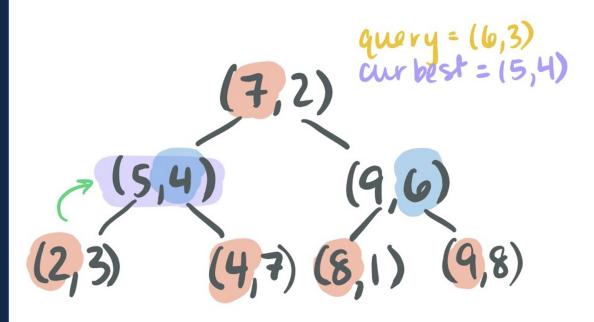


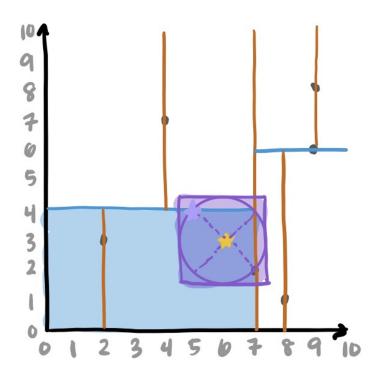


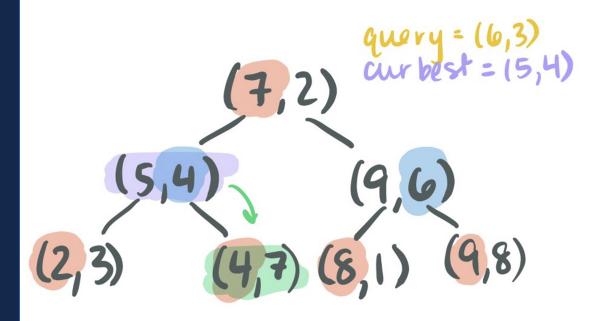


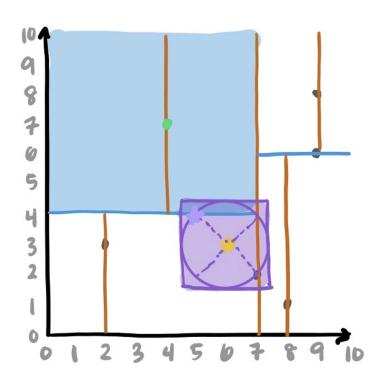


Backtracking: start recursing backwards -- store "best" possibility as you trace back

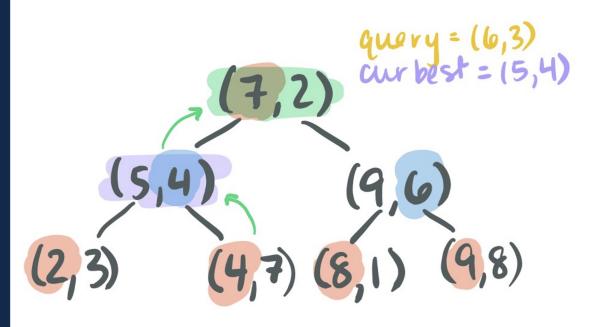


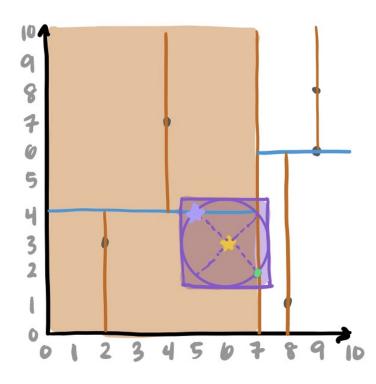


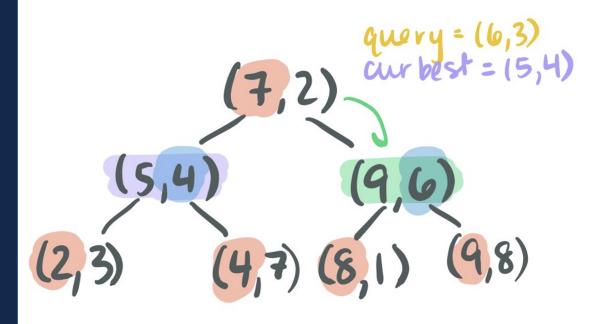


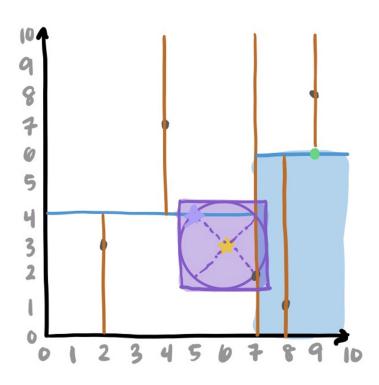


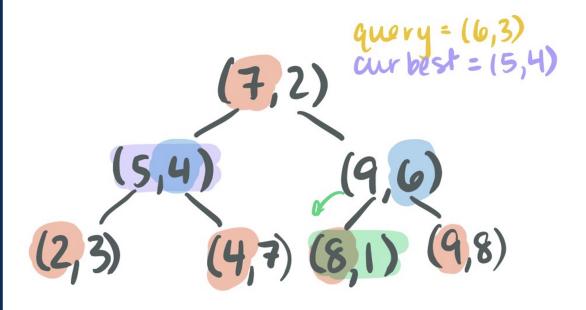
On ties, use smallerDimVal to determine which point remains curBest











9 8 7 0 5 4 3 2 1 0 0 1 2 3 4 5 6 7 8 9 1D

BEST: (5,4)

Every Data Structure So Far

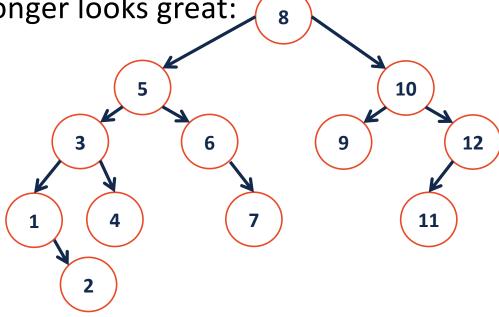
	Unsorted Array	Sorted Array	Unsorted List	Sorted List	Binary Tree	BST	AVL
Find							
Insert							
Remove							
Traverse							

B-Tree Motivation

In Big-O we have assumed uniform time for all operations, but this isn't always true.

However, seeking data from the cloud may take 40ms+.

...an O(lg(n)) AVL tree no longer looks great: 8



BTree Motivations

Knowing that we have large seek times for data, we want to: