# CS 225 

## Data Structures

February 22 - AVL Analysis

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AVL Tree Analysis
We know: insert, remove and find runs in: $\qquad$ .

We will argue that: $h$ is $\qquad$ .

## AVL Tree Analysis

## Definition of big-O:

...or, with pictures:


AVL Tree Analysis



AVL Tree Analysis



- The number of nodes in the tree, $\mathbf{f}^{-1}(\mathbf{h})$, will always be greater than $\mathbf{c} \times \mathbf{g}^{-1}(\mathbf{h})$ for all values where $\mathbf{n}>\mathbf{k}$.


## Plan of Action

Since our goal is to find the lower bound on $\mathbf{n}$ given $\mathbf{h}$, we can begin by defining a function given $\mathbf{h}$ which describes the smallest number of nodes in an AVL tree of height $\mathbf{h}$ :

Simplify the Recurrence $\mathbf{N}(\mathrm{h})=1+\mathrm{N}(\mathrm{h}-1)+\mathrm{N}(\mathrm{h}-2)$

## State a Theorem

Theorem: An AVL tree of height $h$ has at least $\qquad$ .

Proof:
I. Consider an AVL tree and let $\mathbf{h}$ denote its height.
II. Case: $\qquad$
$\qquad$ has at least $\qquad$ nodes.

## Prove a Theorem

III. Case:
$\qquad$ has at least $\qquad$ nodes.

## Prove a Theorem

## By an Inductive Hypothesis (IH):

We will show that:

An AVL tree of height $\qquad$ has at least $\qquad$ nodes.

## Prove a Theorem

V. Using a proof by induction, we have shown that:
...and inverting:

## Summary of Balanced BST

## Red-Black Trees

- Max height: 2 * $\lg (n)$
- Constant number of rotations on insert, remove, and find


## AVL Trees

- Max height: 1.44 * $\lg (\mathrm{n})$
- Rotations:


## Summary of Balanced BST

## Pros:

- Running Time:
- Improvement Over:
- Great for specific applications:


## Summary of Balanced BST

Cons:

- Running Time:
- In-memory Requirement:


## Range-based Searches

Q: Consider points in 1D: $p=\left\{p_{1}, p_{2}, \ldots, p_{n}\right\}$.
...what points fall in [11, 42]?

## Tree construction:

## Range-based Searches

Balanced BSTs are useful structures for range-based and nearest-neighbor searches.

Q: Consider points in 1D: $\mathbf{p}=\left\{\mathbf{p}_{1}, \mathbf{p}_{2}, \ldots, \mathbf{p}_{n}\right\}$.
...what points fall in [11, 42]?

Ex:


## Range-based Searches

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## Range-based Searches

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## Tree construction:

Range-based Searches


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## Range-based Searches

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Range-based Searches


Running Time


## Range-based Searches

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Ex:


## Red-Black Trees in C++

C++ provides us a balanced BST as part of the standard library: std::map<K, V>

V \& std::map<K, V>::operator[](const K \& )
iterator std::map<K, V>::lower_bound( const K \& )
iterator std::map<K, V>::upper_bound( const K \& )

## Every Data Structure So Far

|  | Unsorted Array | Sorted Array | Unsorted List | Sorted List | Binary Tree | BST | AVL |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Find |  |  |  |  |  |  |  |
| Insert |  |  |  |  |  |  |  |
| Remove |  |  |  |  |  |  |  |
| Traverse |  |  |  |  |  |  |  |

## CS 225 Final Project

Working with data and using graphs


## The Internet 2003

The OPTE Project (2003)
Map of the entire internet; nodes
are routers; edges are connections.



Conflict-Free Final Exam Scheduling Graph Unknown Source
Presented by Cinda Heeren, 2016



"Rush Hour" Solution
Unknown Source
Presented by Cinda Heeren, 2016

## Class Hierarchy At University of Illinois Urbana-Champaign <br> A. Mori, W. Fagen-Ulmschneider, C. Heeren <br> Graph of every course at UIUC; nodes are courses, edges are prerequisites <br> http://waf.cs.illinois.edu/discovery/class_hi erarchy_at_illinois/




MP Collaborations in CS 225
Unknown Source
Presented by Cinda Heeren, 2016

"Stanford Bunny"
Greg Turk and Mark Levoy (1994)

## B-Tree Motivation

In Big-O we have assumed uniform time for all operations, but this isn't always true.

However, seeking data from the cloud may take $40 \mathrm{~ms}+$.
...an $\mathrm{O}(\lg (\mathrm{n}))$ AVL tree no longer looks great:


## BTree Motivations

Knowing that we have large seek times for data, we want to:

