

Disjoint Sets

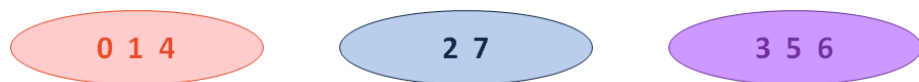
Let **R** be an equivalence relation. We represent R as disjoint sets

- Each element exists in exactly one set.
- Every set is an equitant representation.
 - Mathematically: $4 \in [0]_R \rightarrow 8 \in [0]_R$
 - Programmatically: `find(4) == find(8)`

Building Disjoint Sets:

- Maintain a collection $S = \{s_0, s_1, \dots, s_k\}$
- Each set has a representative member

```
void makeSet(const T & t);
void union(const T & k1, const T & k2);
T & find(const T & k);
```



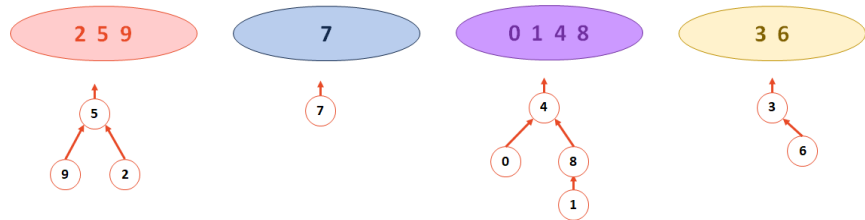
| | | | | | | | | | |
|-----|-----|-----|-----|-----|-----|-----|-----|-----|-----|
| | | | | | | | | | |
| [0] | [1] | [2] | [3] | [4] | [5] | [6] | [7] | [8] | [9] |

Operation: `find(k)`

Operation: `union(k1, k2)`

Implementation #2:

- Continue to use an array where the index is the key
- The value of the array is:
 - **-1**, if we have found the representative element
 - **The index of the parent**, if we haven't found the rep. element



| | | | | | | | | | |
|----------|----------|----------|-----------|-----------|-----------|----------|-----------|----------|----------|
| 4 | 8 | 5 | -1 | -1 | -1 | 3 | -1 | 4 | 5 |
| [0] | [1] | [2] | [3] | [4] | [5] | [6] | [7] | [8] | [9] |

Implementation – DisjointSets::find

```
DisjointSets.cpp (partial)
1 int DisjointSets::find(int i) {
2   if ( s[i] < 0 ) { return i; }
3   else { return _find( s[i] ); }
4 }
```

What is the running time of `find`?

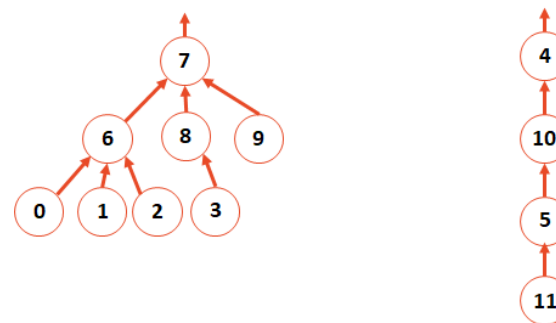
What is the ideal UpTree?

Implementation – DisjointSets::union

```
DisjointSets.cpp (partial)
1 void DisjointSets::union(int r1, int r2) {
2
3
4 }
```

How do we want to union the two UpTrees?

Building a Smart Union Function



The implementation of this visual model is the following:

| | | | | | | | | | | | |
|----------|----------|----------|----------|-----------|-----------|----------|-----------|----------|----------|----------|----------|
| 6 | 6 | 6 | 8 | -1 | 10 | 7 | -1 | 7 | 7 | 4 | 5 |
| [0] | [1] | [2] | [3] | [4] | [5] | [6] | [7] | [8] | [9] | [10] | [11] |

What are possible strategies to employ when building a “smart union”?

Smart Union Strategy #1: _____

Idea: Keep the height of the tree as small as possible!

Metadata at Root:

After `union(4, 7)`:

| | | | | | | | | | | | |
|----------|----------|----------|----------|-----|-----------|----------|-----|----------|----------|----------|----------|
| 6 | 6 | 6 | 8 | | 10 | 7 | | 7 | 7 | 4 | 5 |
| [0] | [1] | [2] | [3] | [4] | [5] | [6] | [7] | [8] | [9] | [10] | [11] |

Smart Union Strategy #2: _____

Idea: Minimize the number of nodes that increase in height.
(Observe that the tree we union have all their nodes gain in height.)

Metadata at Root:

After `union(4, 7)`:

| | | | | | | | | | | | |
|----------|----------|----------|----------|-----|-----------|----------|-----|----------|----------|----------|----------|
| 6 | 6 | 6 | 8 | | 10 | 7 | | 7 | 7 | 4 | 5 |
| [0] | [1] | [2] | [3] | [4] | [5] | [6] | [7] | [8] | [9] | [10] | [11] |

Smart Union Implementation:

```

DisjointSets.cpp (partial)
1 void DisjointSets::unionBySize(int root1, int root2) {
2     int newSize = arr_[root1] + arr_[root2];
3
4     if ( arr_[root1] < arr_[root2] ) {
5         arr_[root2] = root1; arr_[root1] = newSize;
6     } else {
7         arr_[root1] = root2; arr_[root2] = newSize;
8     }
9 }

```

How do we improve this?

```

DisjointSets.cpp (partial)
1 int DisjointSets::find(int i) {
2     if ( arr_[i] < 0 ) { return i; }
3     else { return _find( arr_[i] ); }
4 }

```

```

DisjointSets.cpp (partial)
1 void DisjointSets::unionBySize(int root1, int root2) {
2     int newSize = arr_[root1] + arr_[root2];
3
4     // If arr_[root1] is less than (more negative), it is the
5     // larger set; we union the smaller set, root2, with root1.
6     if ( arr_[root1] < arr_[root2] ) {
7         arr_[root2] = root1;
8         arr_[root1] = newSize;
9     }
10    // Otherwise, do the opposite:
11    else {
12        arr_[root1] = root2;
13        arr_[root2] = newSize;
14    }
15 }

```

Running Time:

- Worst case running time of `find(k)`:
- Worst case running time of `union(r1, r2)`, given roots:
- New function: “Iterated Log”:
 $\log^*(n) :=$
- Overall running time:
 - A total of **m** union/find operation runs in:

| CS 225 – Things To Be Doing: |
|---|
| <ol style="list-style-type: none"> 1. mp_mosaics due today. 2. Exam on Friday practice on PrairieLearn now 3. Daily POTDs are ongoing! |