CS 225

**Data Structures** 

March 13 – BTree Analysis

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The height of the BTree determines maximum number of \_\_\_\_\_ possible in search data.

...and the height of the structure is: \_\_\_\_\_.

**Therefore:** The number of seeks is no more than \_\_\_\_\_\_.

...suppose we want to prove this!

In our AVL Analysis, we saw finding an upper bound on the height (given **n**) is the same as finding a lower bound on the nodes (given **h**).

We want to find a relationship for BTrees between the number of keys (n) and the height (h).

#### **Strategy:**

We will first count the number of nodes, level by level.

Then, we will add the minimum number of keys per node (n).

The minimum number of nodes will tell us the largest possible height (h), allowing us to find an upper-bound on height.

The minimum number of **nodes** for a BTree of order m **at each level**:

root:

level 1:

level 2:

level 3:

• • •

level h:

The total number of nodes is the sum of all of the levels:

The total number of keys:

The smallest total number of keys is:

So an inequality about **n**, the total number of keys:

Solving for **h**, since **h** is the number of seek operations:

Given m=101, a tree of height h=4 has:

Minimum Keys:

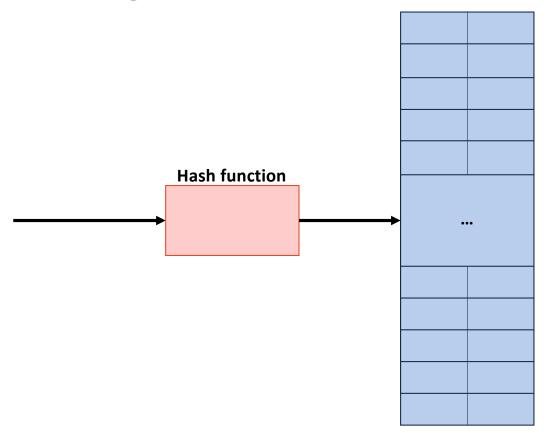
Maximum Keys:

#### **Goals:**

We want to define a **keyspace**, a (mathematical) description of the keys for a set of data.

...use a function to map the **keyspace** into a small set of integers.

Locker Number	Name
103	
92	
330	
46	
124	



## A Hash Table based Dictionary

#### **Client Code:**

```
Dictionary<KeyType, ValueType> d;
d[k] = v;
```

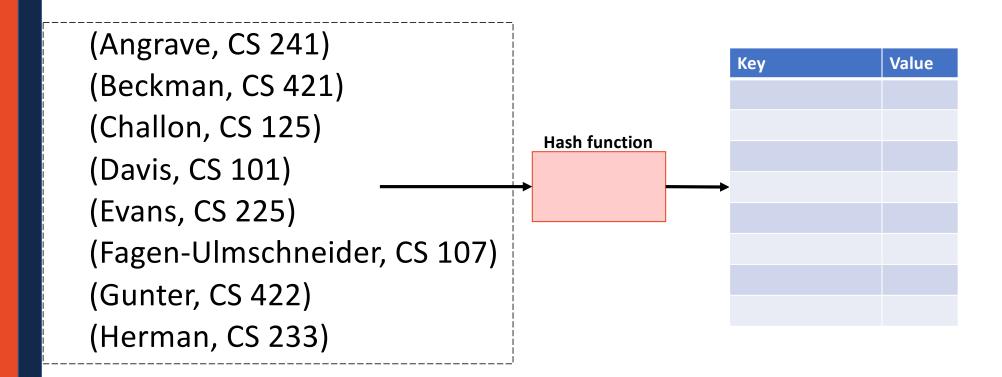
#### A **Hash Table** consists of three things:

1.

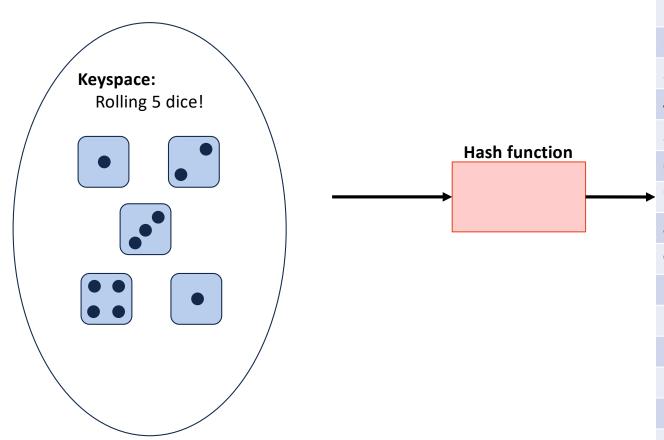
2.

3.

#### A Perfect Hash Function



#### A Perfect Hash Function



	Key	Value
	0	
	1	
	2	
	3	
	4	
	5	
	6	
<b>→</b>	7	
	8	
	9	
	10	
	11	
	12	
	13	
	14	
	15	