



CS 225

Data Structures

March 2 – BST Balance

G Carl Evans



BST Analysis

Therefore, for all BST:

Lower bound: $h \geq O(\lg(n))$

Upper bound: $h \leq O(n)$



BST Analysis

The height of a BST depends on the order in which the data is inserted into it.

ex: 1 3 2 4 5 7 6

vs.

4 2 3 6 7 1 5

Q: How many different ways are there to insert keys into a BST?

Q: What is the average height of all the arrangements?



BST Analysis

Q: How many different ways are there to insert keys into a BST?

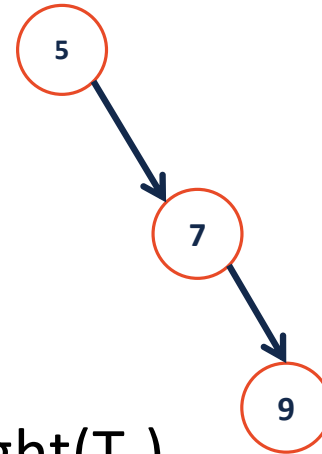
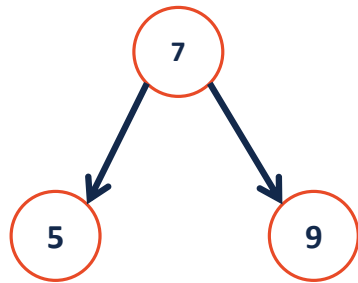
Q: What is the average height of all the arrangements?

BST Analysis – Running Time

Operation	BST Average case	BST Worst case	Sorted array	Sorted List
find				
insert				
delete				
traverse				

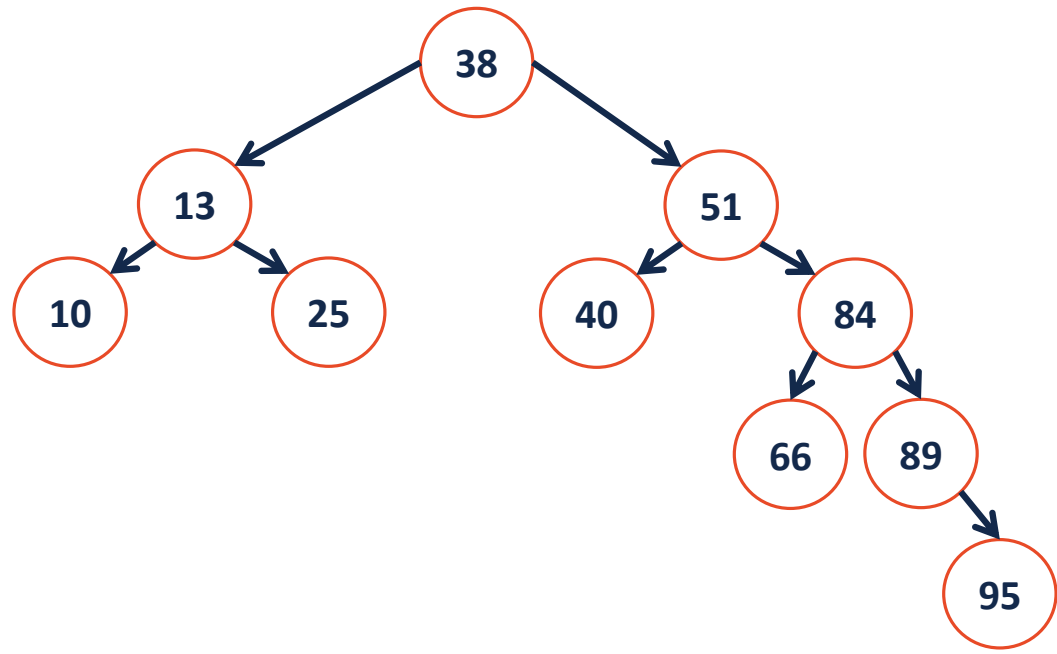
Height-Balanced Tree

What tree makes you happier?



Height balance: $b = \text{height}(T_R) - \text{height}(T_L)$

A tree is height balanced if:



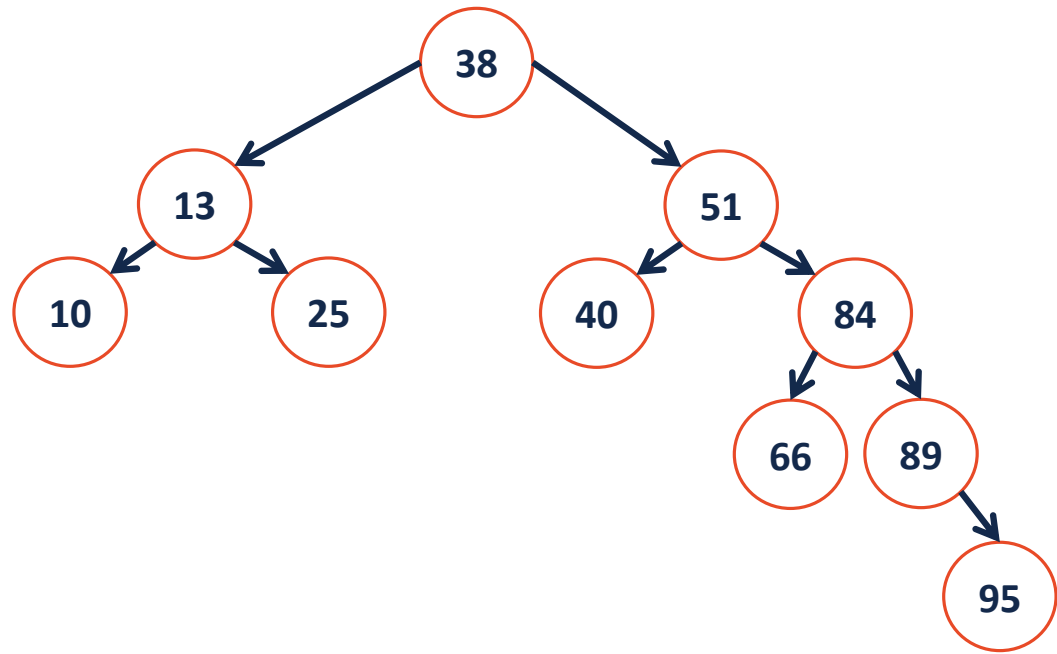


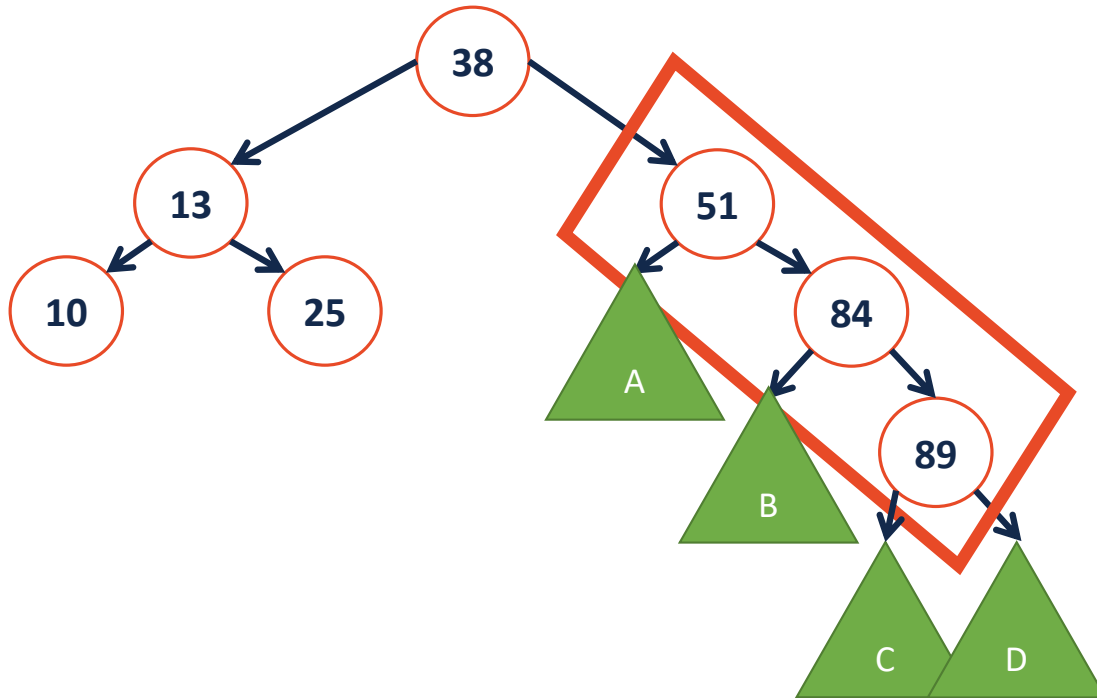
BST Rotation

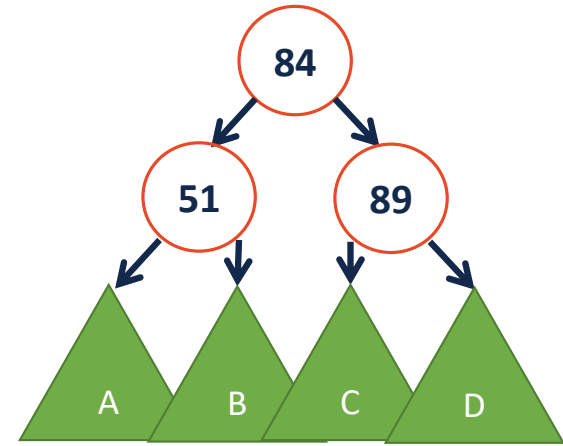
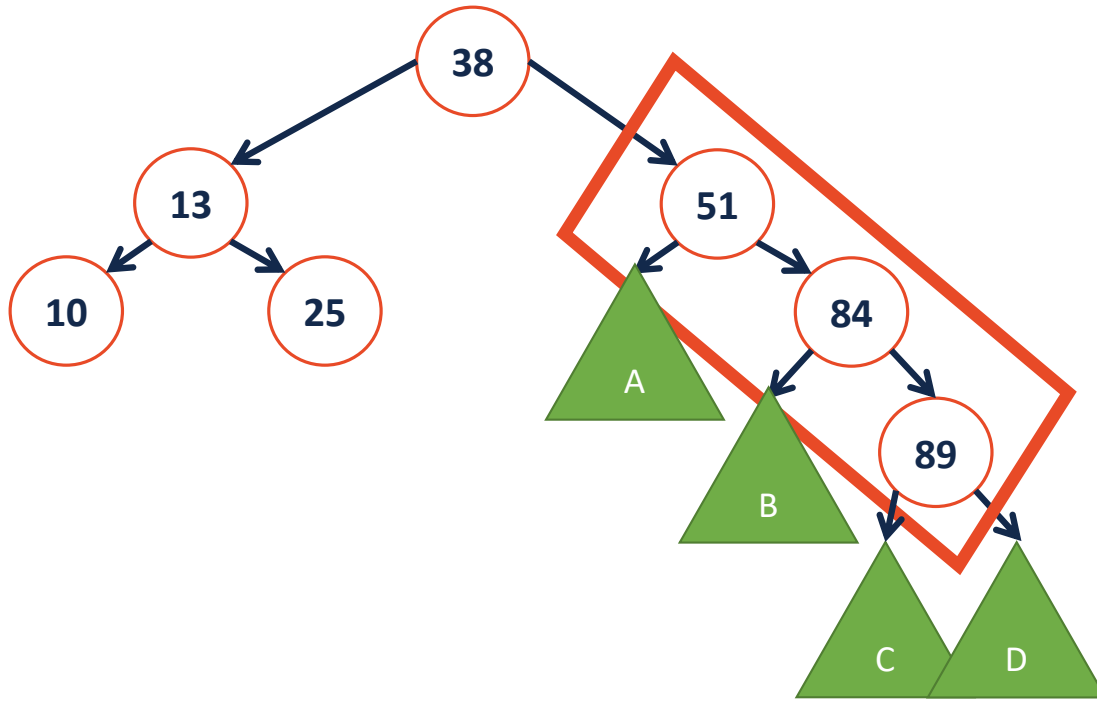
We will perform a rotation that maintains two properties:

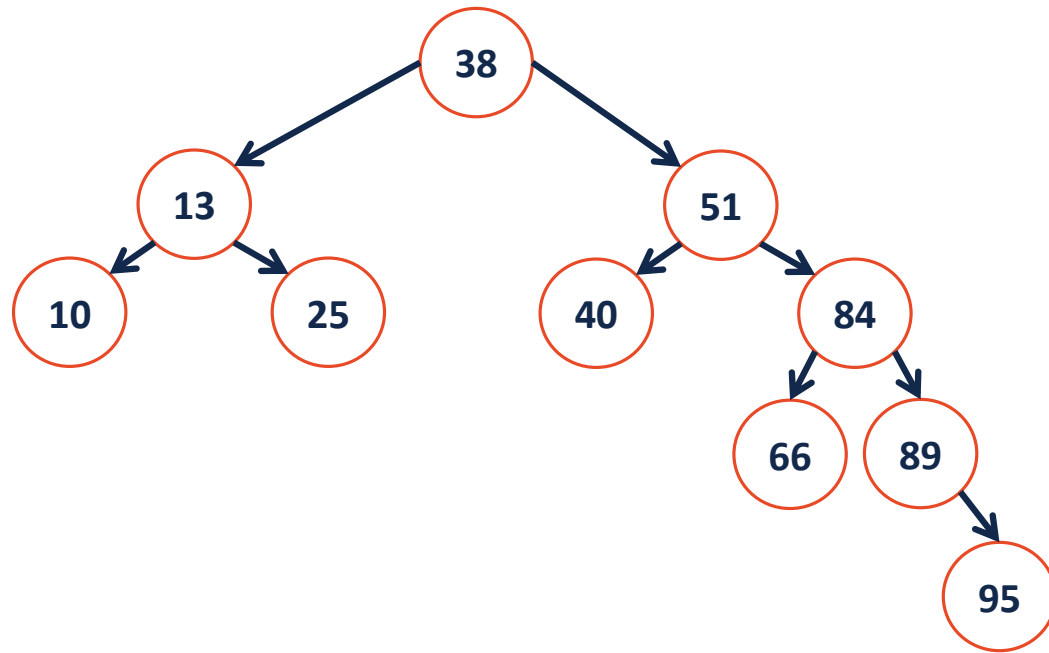
1.

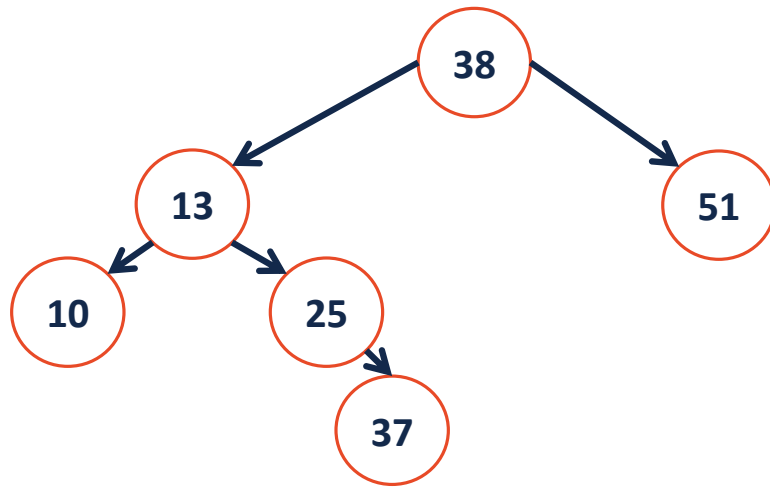
2.

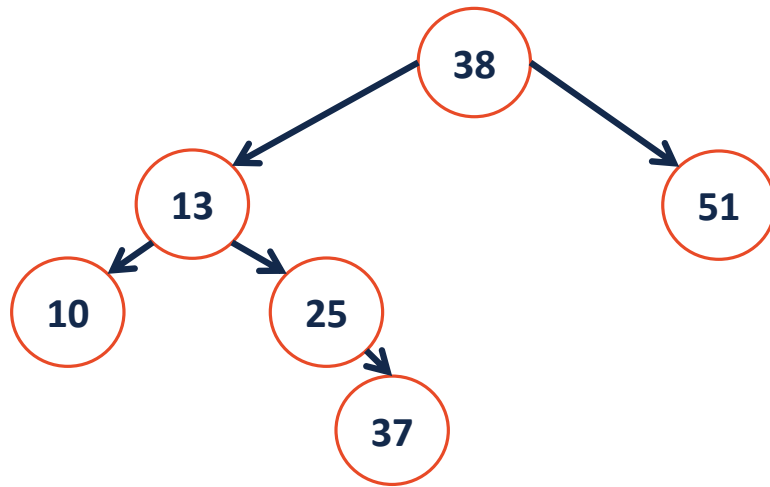














BST Rotation Summary

- Four kinds of rotations (L, R, LR, RL)
- All rotations are local (subtrees are not impacted)
- All rotations are constant time: $O(1)$
- BST property maintained

GOAL:

We call these trees:



AVL Trees

Three issues for consideration:

- Rotations
- Maintaining Height
- Detecting Imbalance