

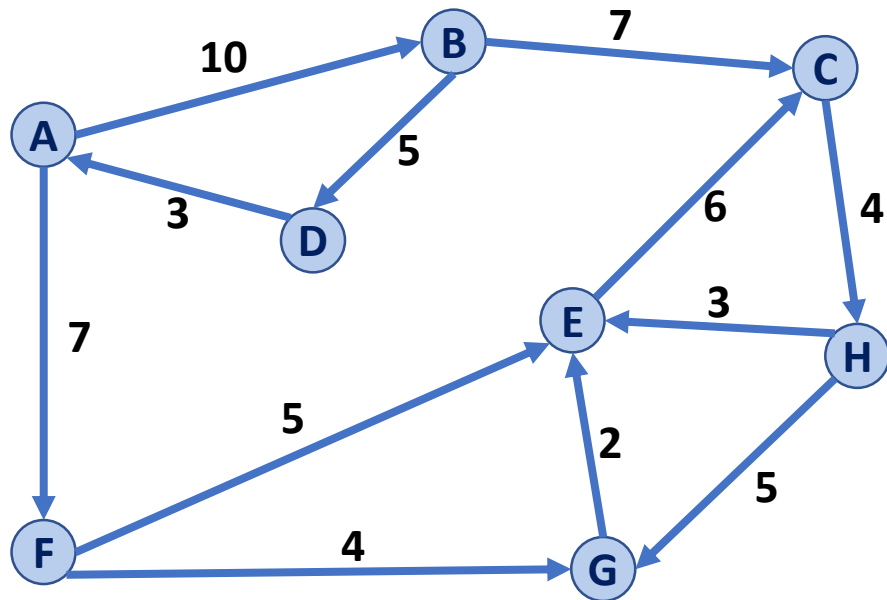
CS 225

Data Structures

April 27 – Graph Problems + End of Semester

Wade Fagen-Ulmschneider

Dijkstra's Algorithm (SSSP)



```
DijkstraSSSP(G, s):
```

```
6  foreach (Vertex v : G):
```

```
7      d[v] = +inf
```

```
8      p[v] = NULL
```

```
9      d[s] = 0
```

```
10
```

```
11  PriorityQueue Q // min distance, defined by d[v]
```

```
12  Q.buildHeap(G.vertices())
```

```
13  Graph T          // "labeled set"
```

```
14
```

```
15  repeat n times:
```

```
16      Vertex u = Q.removeMin()
```

```
17      T.add(u)
```

```
18      foreach (Vertex v : neighbors of u not in T):
```

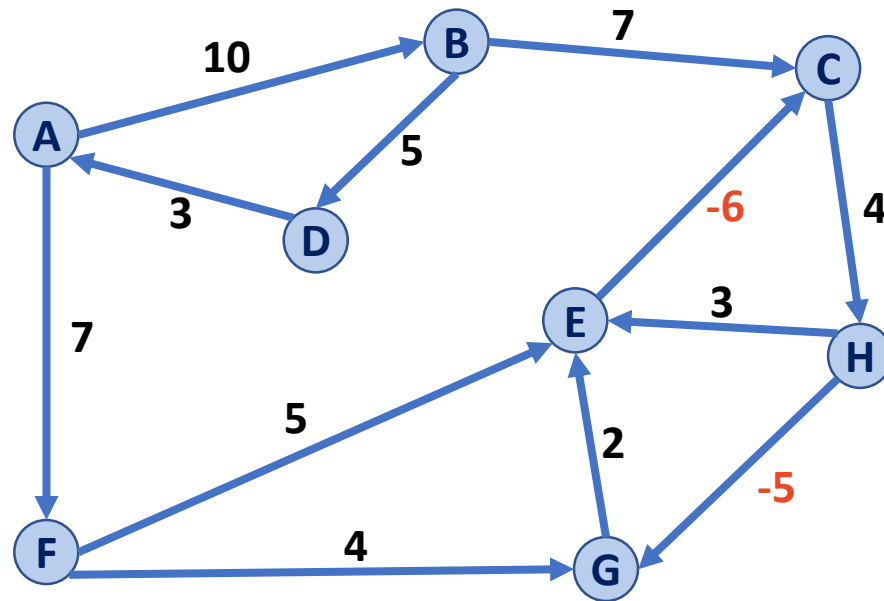
```
19          if cost(u, v) + d[u] < d[v]:
```

```
20              d[v] = cost(u, v) + d[u]
```

```
21              p[v] = u
```

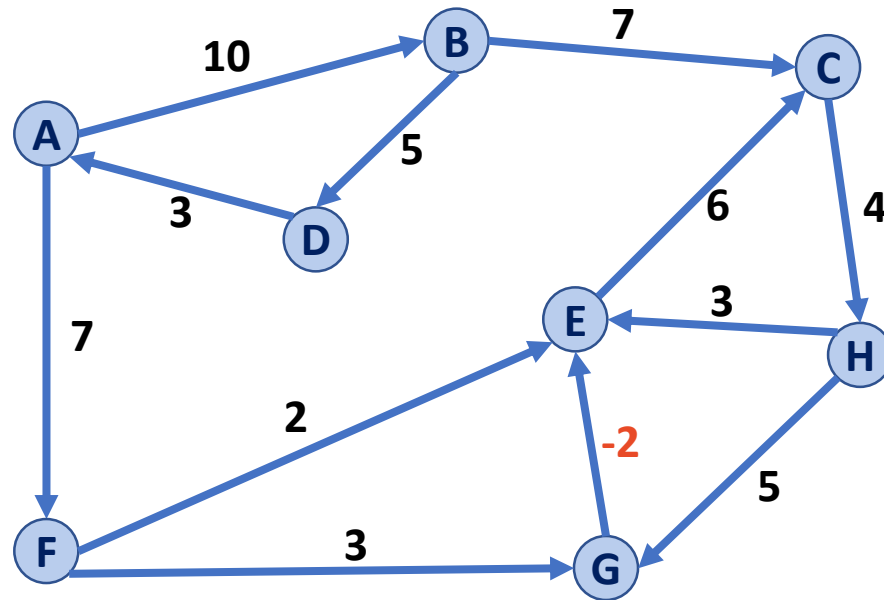
Dijkstra's Algorithm (SSSP)

Q: How does Dijkstra handle negative weight cycles?



Dijkstra's Algorithm (SSSP)

Q: How does Dijkstra handle negative weight edges, without a negative weight cycle?



Dijkstra's Algorithm (SSSP)

What is Dijkstra's running time?

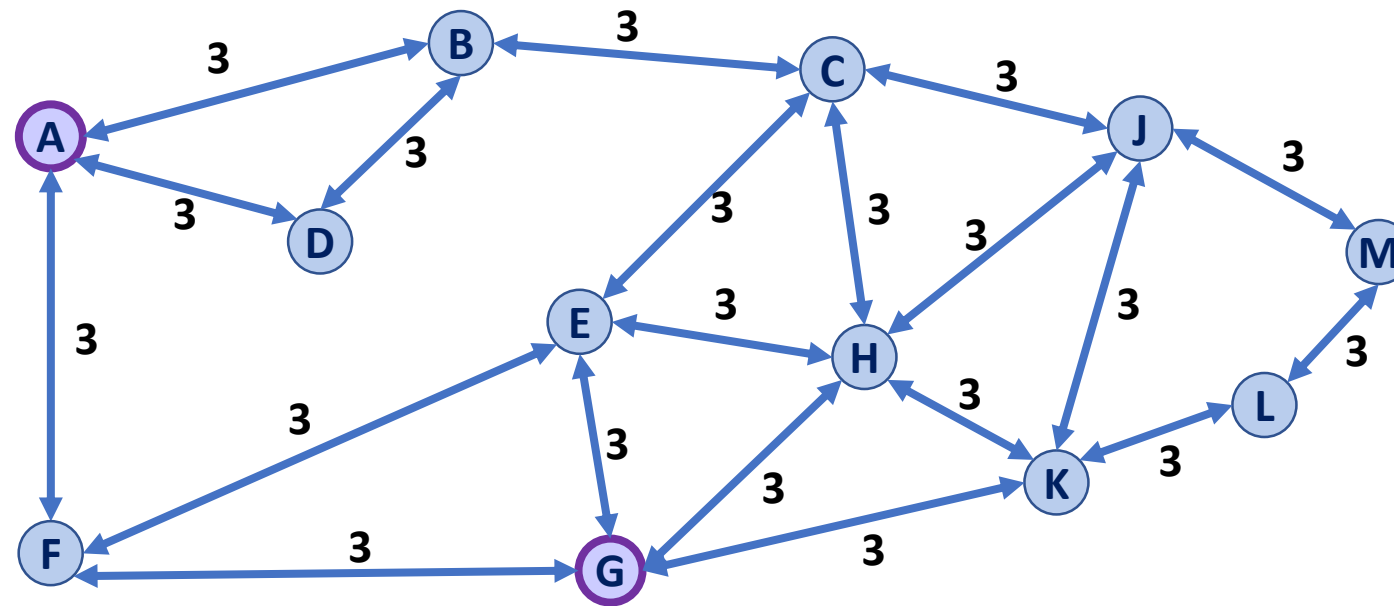
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12  Q.buildHeap(G.vertices())
13  Graph T          // "labeled set"
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15  repeat n times:
16      Vertex u = Q.removeMin()
17      T.add(u)
18      foreach (Vertex v : neighbors of u not in T):
19          if cost(u, v) + d[u] < d[v]:
20              d[v] = cost(u, v) + d[u]
21              p[v] = m
22
23  return T
```

Landmark Path Problem

Suppose you want to travel from **A** to **G**.

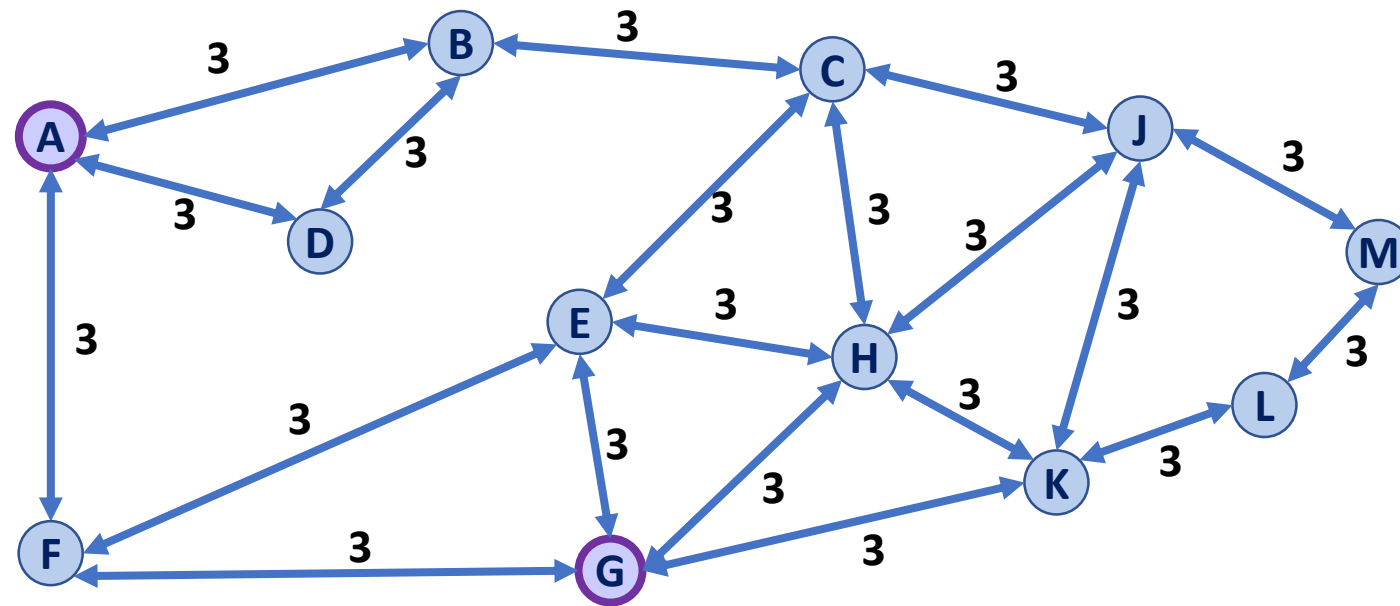
Q1: What is the shortest path from **A** to **G**?



Landmark Path Problem

Suppose you want to travel from **A** to **G**.

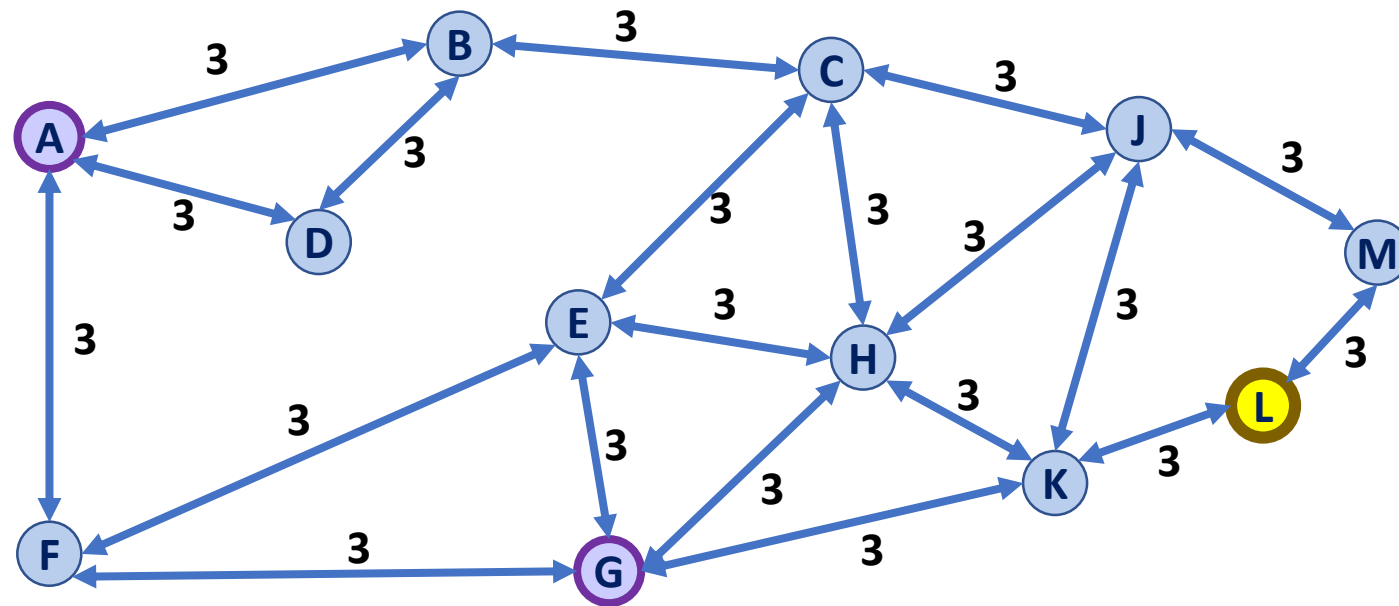
Q2: What is the fastest algorithm to use to find the shortest path?



Landmark Path Problem

In your journey between **A** and **G**, you also want to visit the landmark **L**.

Q3: What is the shortest path from **A** to **G** that visits **L**?

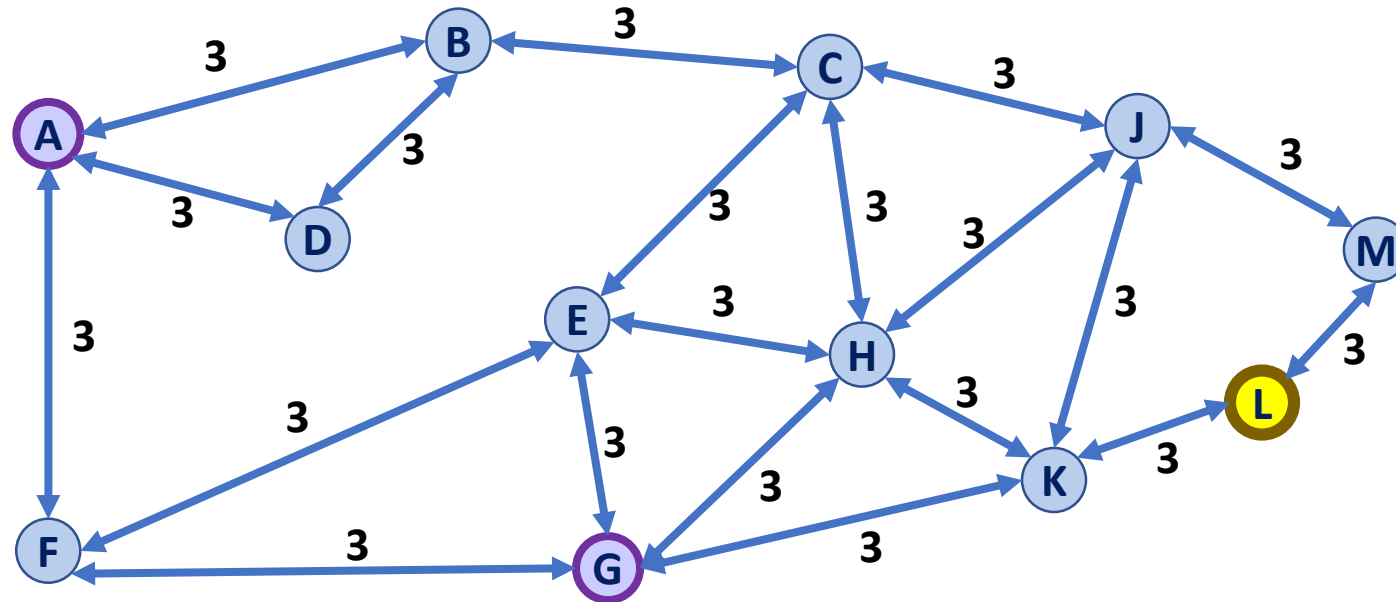


Landmark Path Problem

In your journey between **A** and **G**, you also want to visit the landmark **L**.

Q4: What is the fastest algorithm to find this path?

Q5: What are the specific call(s) to this algorithm?



End of Semester Logistics

CS 225 Final Exam

- The final exam begins on Thursday, May 3rd
- The final exam is a **3 hour CBTF exam**, is a **cumulative exam**, and has the format of a **combined theory + programming exam**
- The last office hours is Wednesday, May 2nd
- *We'll use lecture on Wednesday, May 2nd as a final exam review!*

End of Semester Logistics

“Pre-Final” Grade Dump

- I believe there’s only a few remaining issues left with grading; I’ll be starting to wrap these up myself over the weekend:
 - +EC from creative components
 - Working on recovering repos that were force deleted
- As soon as possible after MP7’s deadline, we’ll provide a “Pre-Final” grade in Compass that incorporates everything except the final exam into your CS 225 grade.

End of Semester Logistics

End of Semester Grade Review

- Excel sheet will be provided once final grades are posted.
- Must submit an Excel sheet for this review.

End of Semester Logistics

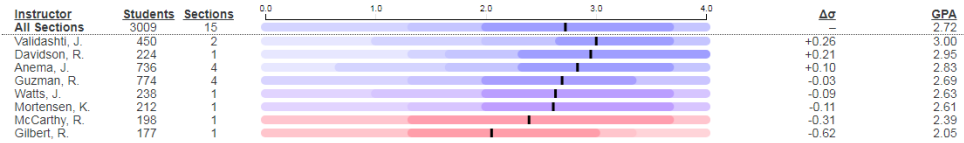
You're Awesome -- +1 To Your Skills!

My Passion: Data Discovery

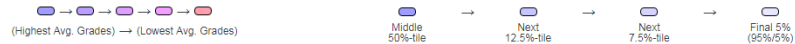
Diversity at Illinois:

GPAs at Illinois:

MATH 221: Calculus I

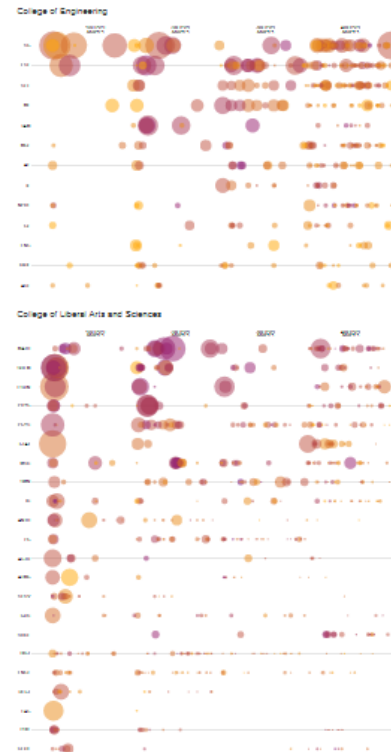
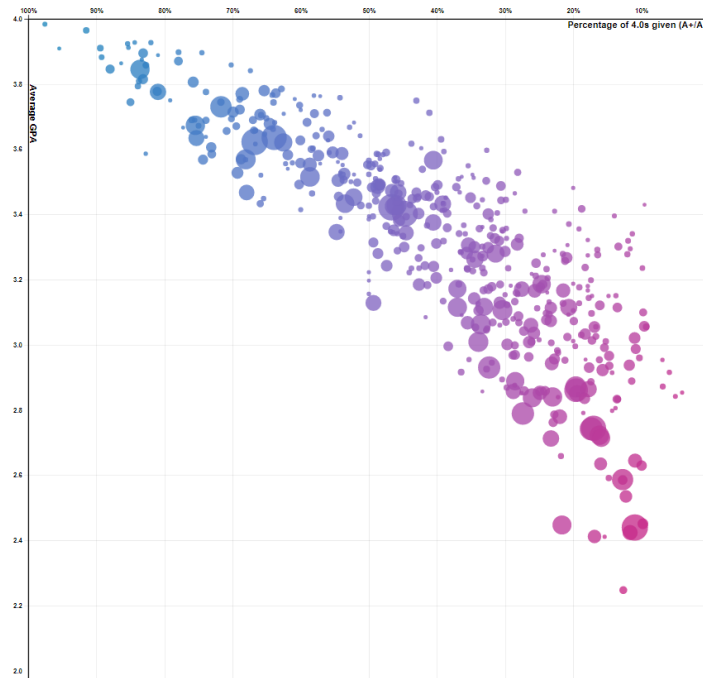


Legend



Instructors with average grades significantly lower than the average grade for a course have increasing red hues.

The darkest shading shows the median grades in a course, with each lighter showing grades further from the median.



Department of Computer Science

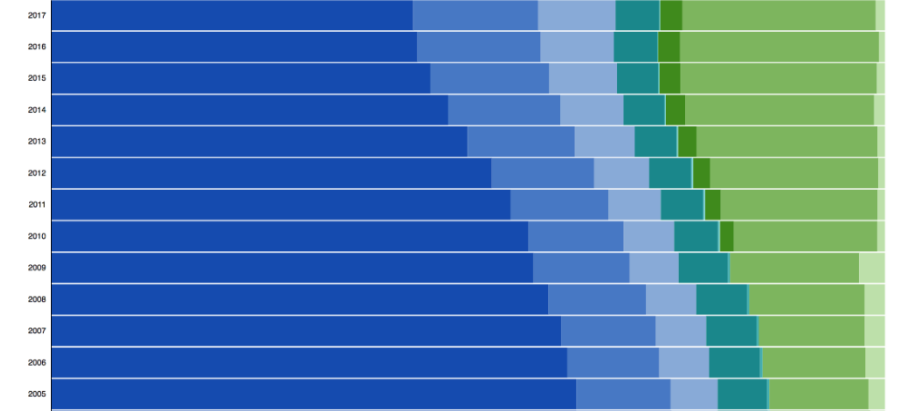


Legend

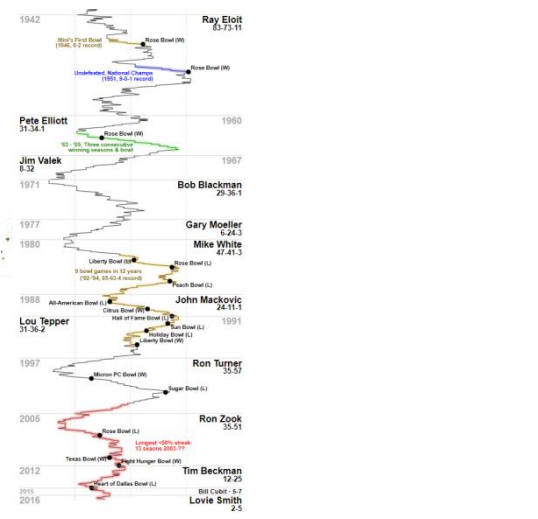
- White
- Asian American
- Hispanic/Latino
- Black
- Native American
- Hawaiian/Pacific Islander
- Multiracial
- International
- Unknown

College of Engineering in

- Aerospace Engineering
- Agriculture & Biological Eng.
- Biomedical Eng.
- Bioengineering
- Bioinformatics
- Bioinstrumentation
- Civil Engineering
- Computer Engineering



And others:



Floyd-Warshall Algorithm

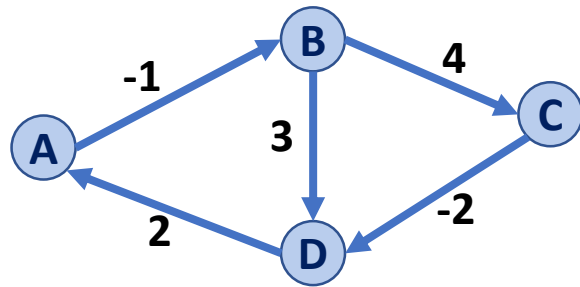
Floyd-Warshall's Algorithm is an alternative to Dijkstra in the presence of **negative-weight edges** (not **negative weight cycles**).

```
FloydWarshall(G) :
6   Let d be a adj. matrix initialized to +inf
7   foreach (Vertex v : G) :
8       d[v][v] = 0
9   foreach (Edge (u, v) : G) :
10      d[u][v] = cost(u, v)
11
12  foreach (Vertex u : G) :
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15              if d[u, v] > d[u, w] + d[w, v] :
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Floyd-Warshall Algorithm

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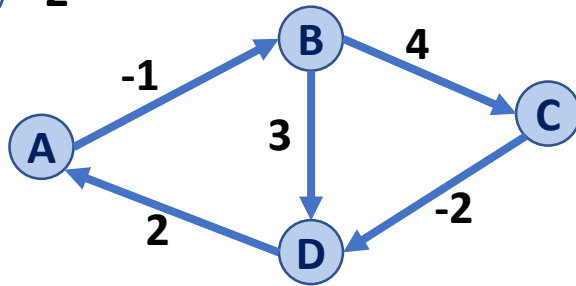
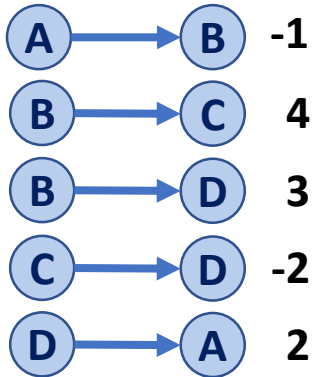
	A	B	C	D
A				
B				
C				
D				

Floyd-Warshall Algorithm

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12  foreach (Vertex u : G):  
13      foreach (Vertex v : G):  
14          foreach (Vertex w : G):  
15              if  $d[u, v] > d[u, w] + d[w, v]$ :  
16                   $d[u, v] = d[u, w] + d[w, v]$ 
```

	A	B	C	D
A	0	-1		
B		0	4	3
C			0	-2
D	2			0

Initially:



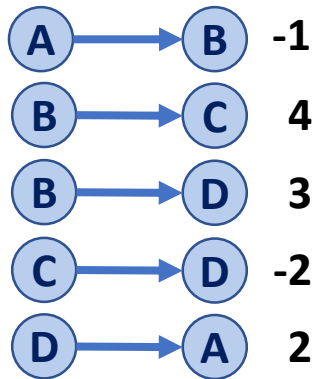
Floyd-Warshall Algorithm

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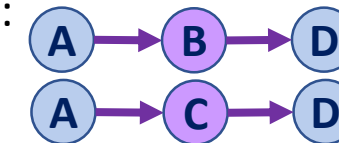
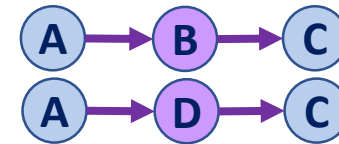
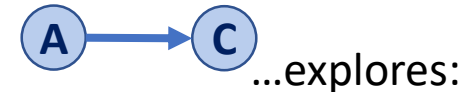
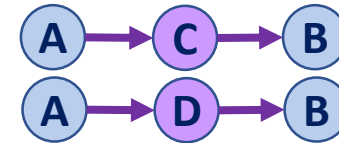
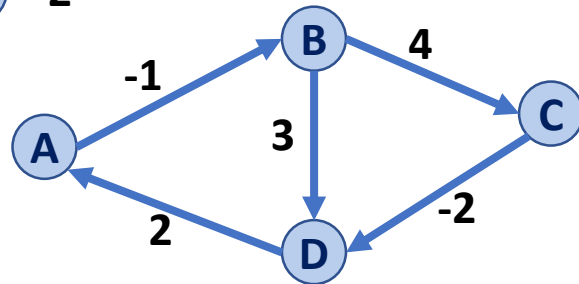
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```

Initially:



Let u = A; v and w explores for better paths:



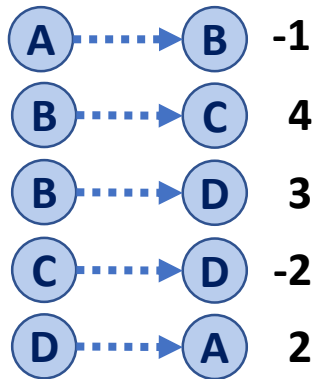
Floyd-Warshall Algorithm

	A	B	C	D
A	0	-1	2	1
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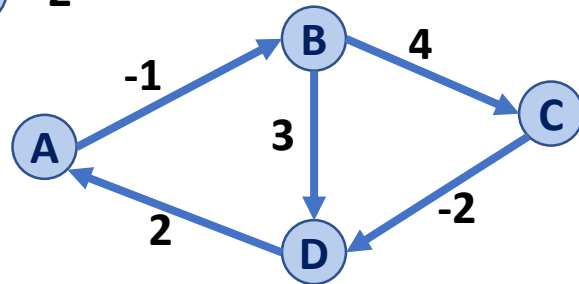
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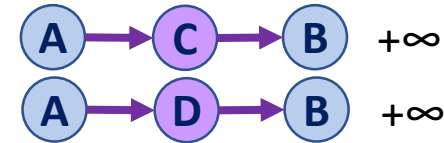
Initially:



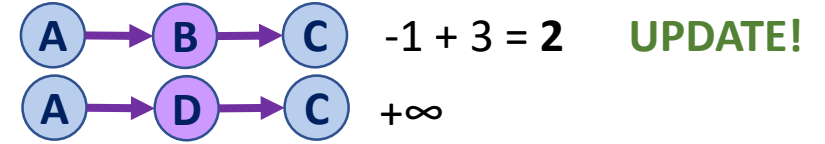
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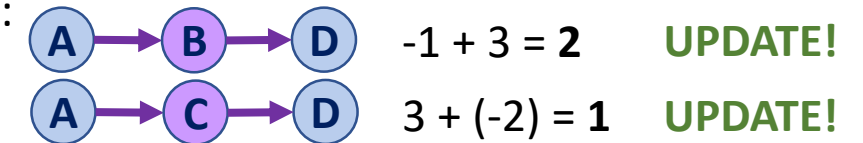
...explores:



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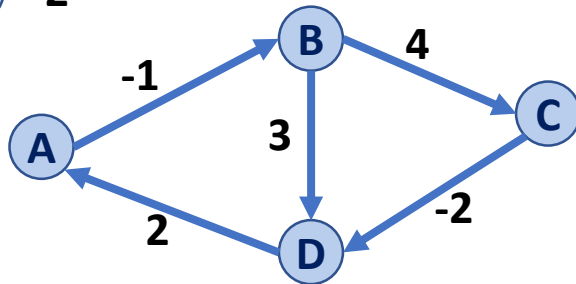
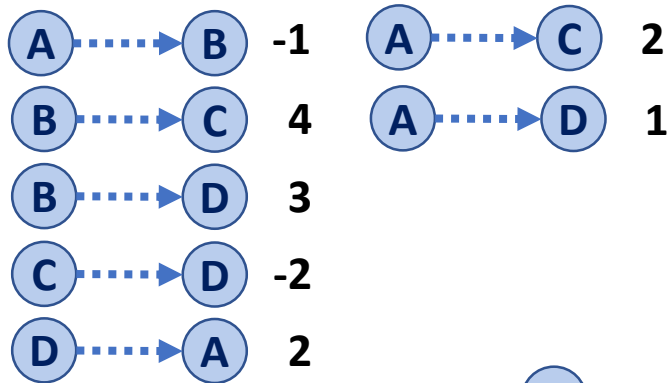
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Floyd-Warshall Algorithm

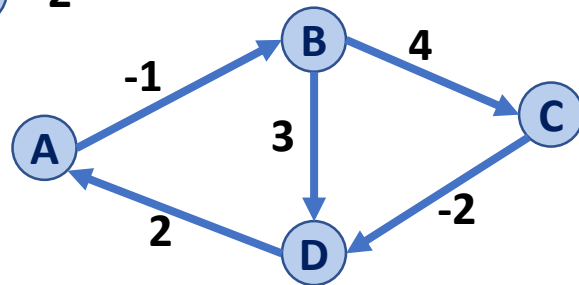
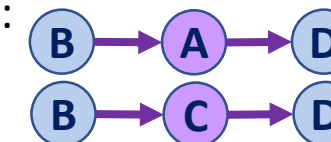
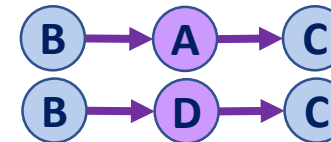
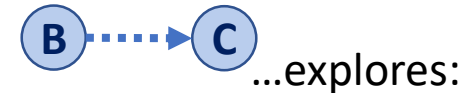
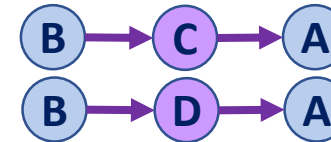
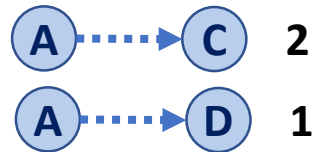
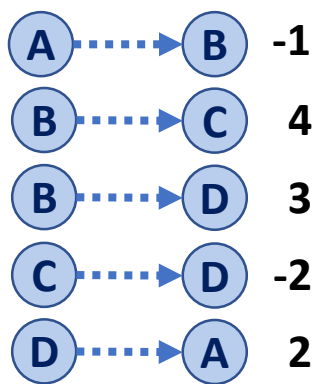
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C			0	-2
D	2			0

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Let u = B; v and w explores for better paths:



Floyd-Warshall Algorithm

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