CS 225

Data Structures

March 30 – Heap Operations

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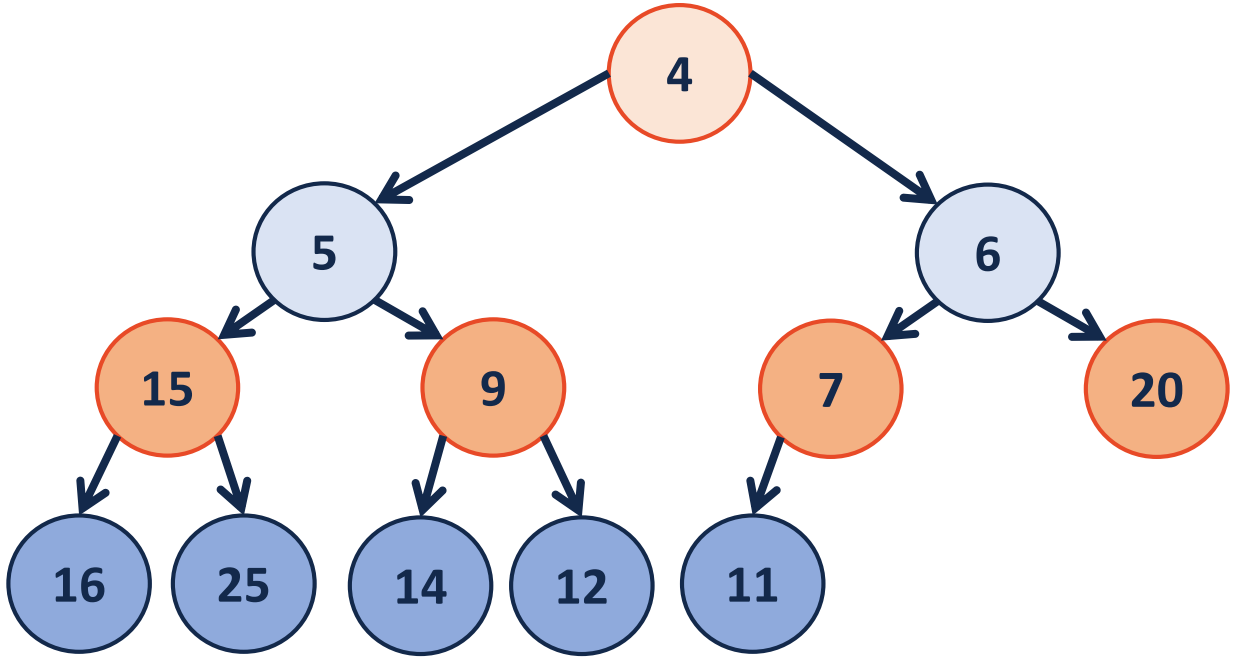
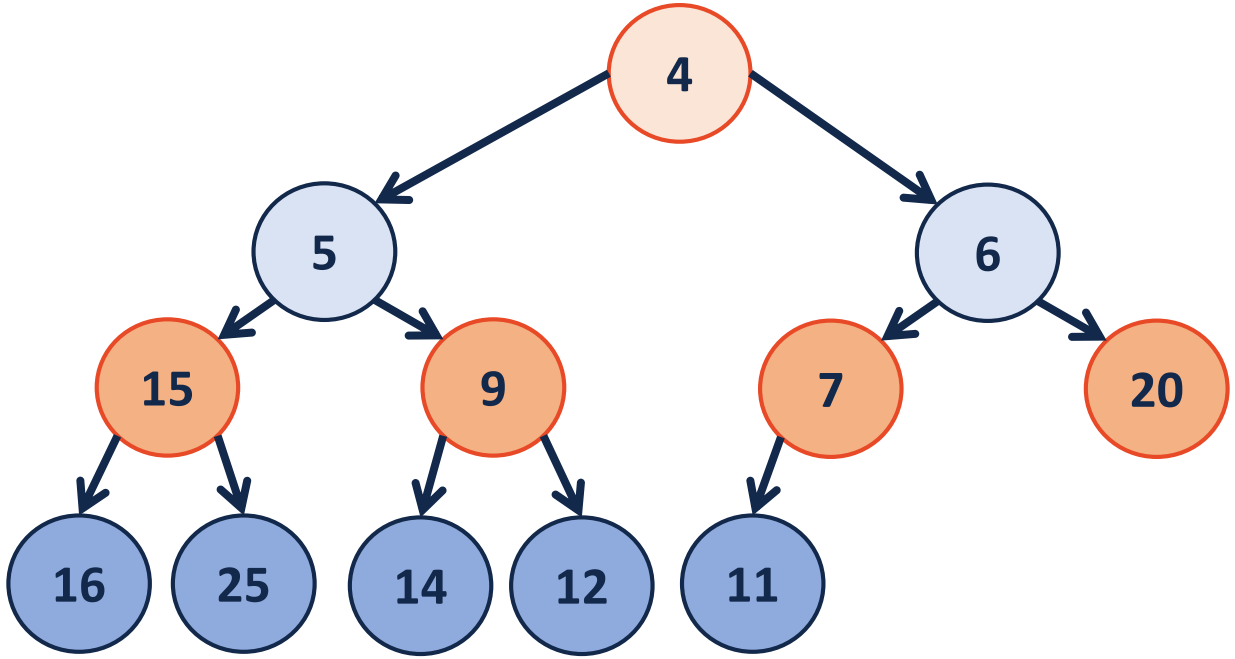
Heap



	4	5	6	15	9	7	20	16	25	14	12	11		

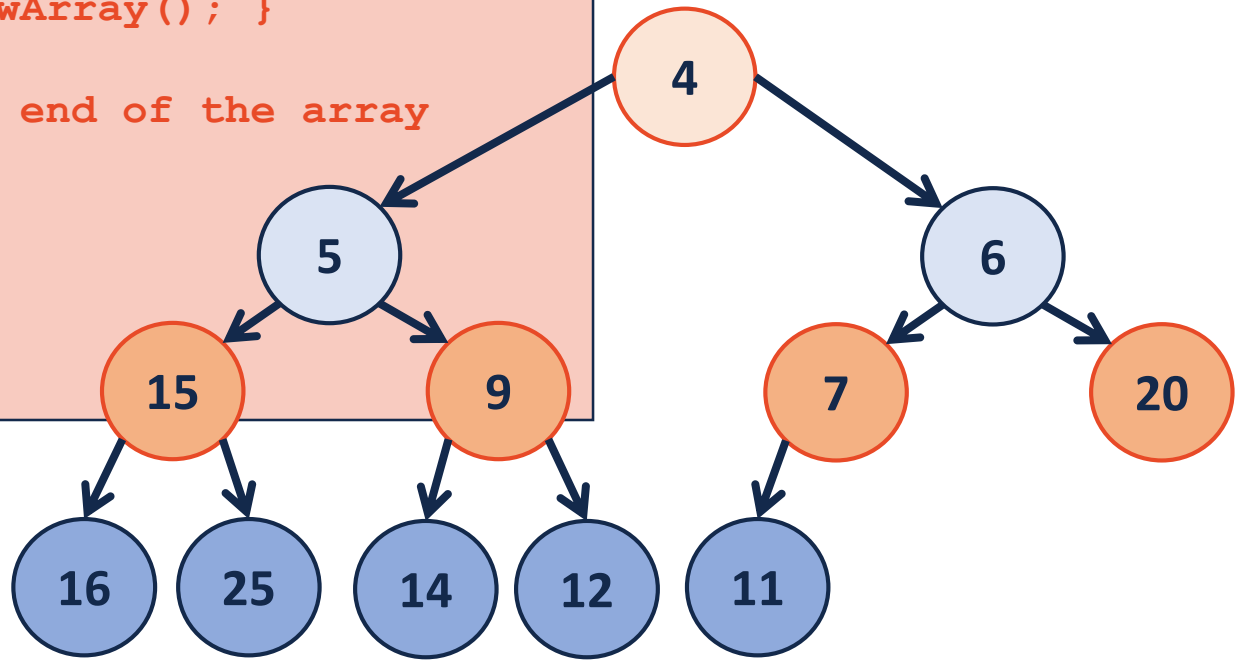
insert



	4	5	6	15	9	7	20	16	25	14	12	11		

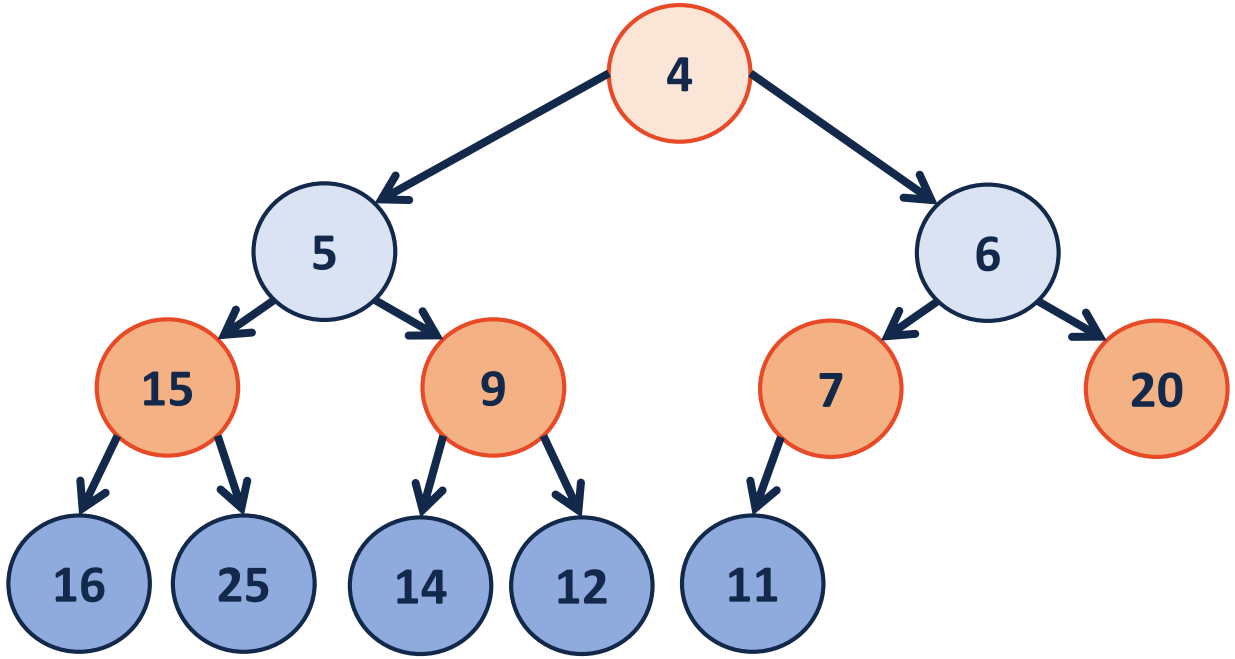
insert

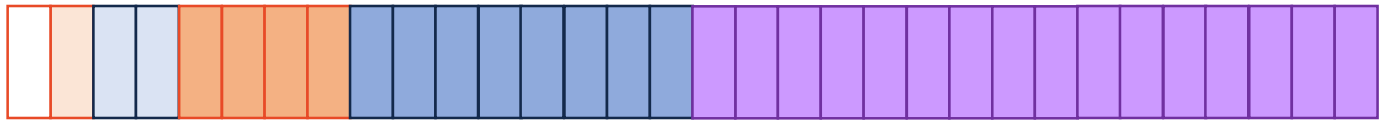
```
1  template <class T>
2  void Heap<T>::_insert(const T & key) {
3    // Check to ensure there's space to insert an element
4    // ...if not, grow the array
5    if ( size_ == capacity_ ) { _growArray(); }
6
7    // Insert the new element at the end of the array
8    item_[++size] = key;
9
10   // Restore the heap property
11   _heapifyUp(size);
12 }
```



	4	5	6	15	9	7	20	16	25	14	12	11			

growArray





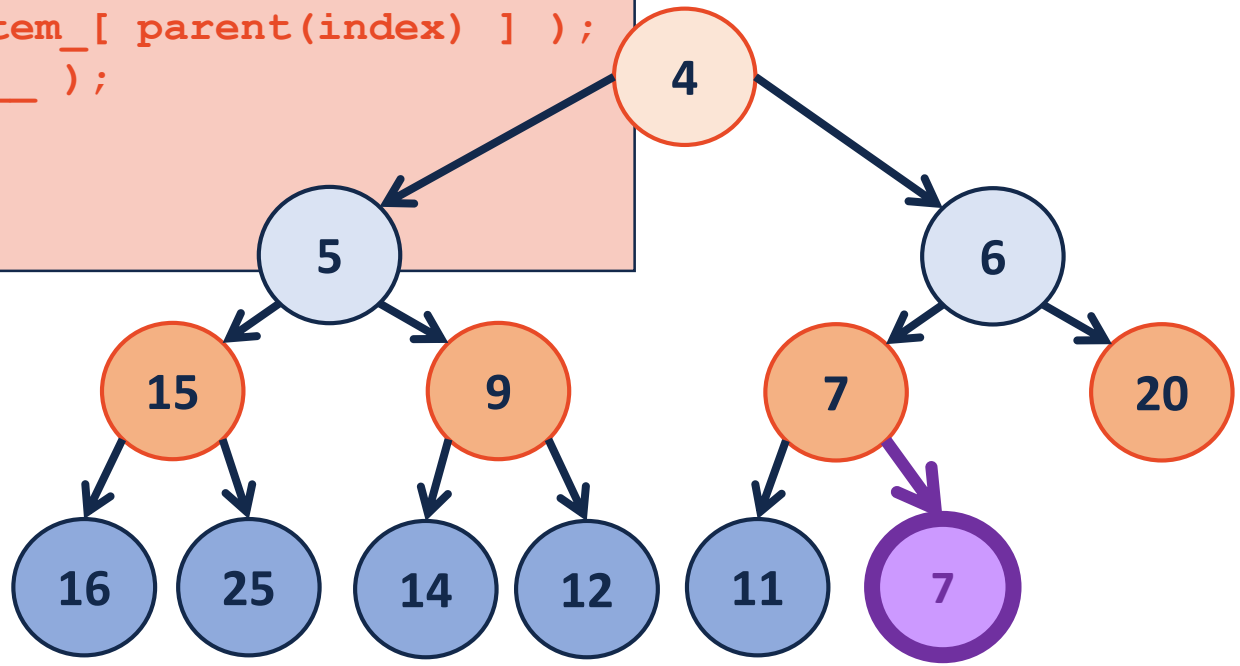
insert - heapifyUp

```
1  template <class T>
2  void Heap<T>::_insert(const T & key) {
3    // Check to ensure there's space to insert an element
4    // ...if not, grow the array
5    if ( size_ == capacity_ ) { _growArray(); }
6
7    // Insert the new element at the end of the array
8    item_[++size] = key;
9
10   // Restore the heap property
11   _heapifyUp(size);
12 }
```

```
1  template <class T>
2  void Heap<T>::_heapifyUp( ____________________ ) {
3    if ( index > __________ ) {
4      if ( item_[index] < item_[ parent(index) ] ) {
5         std::swap( item_[index], item_[ parent(index) ] );
6         _heapifyUp( ___________________ );
7      }
8    }
9  }
```

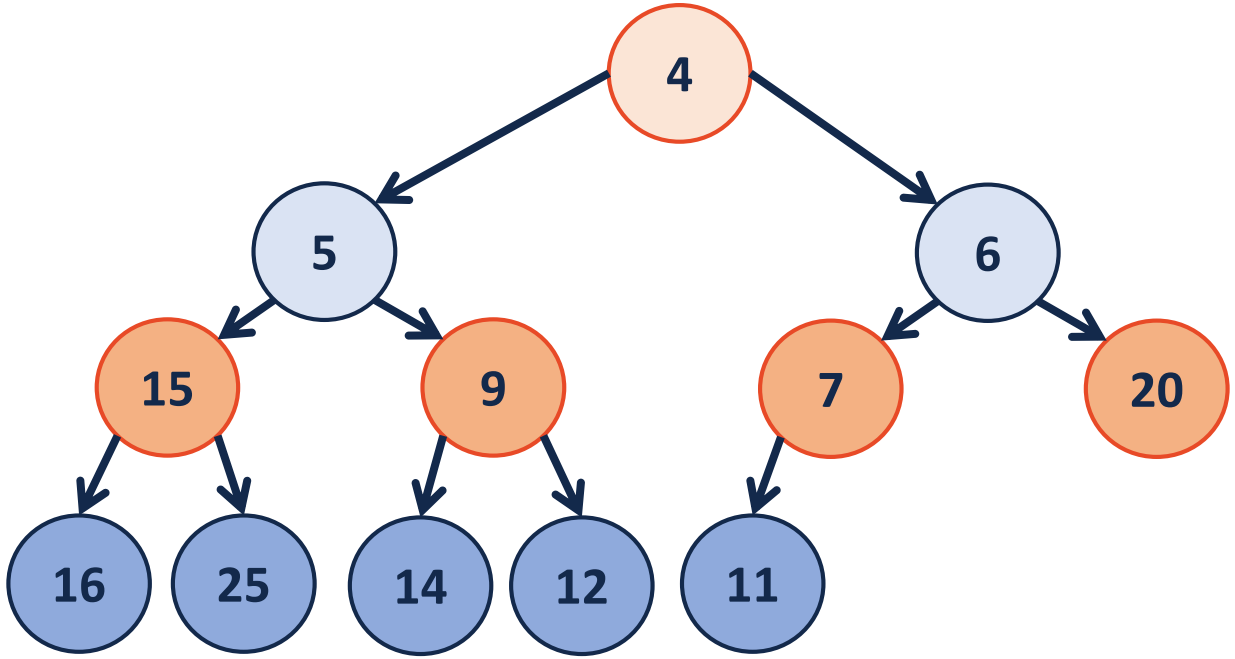
heapifyUp

```
1  template <class T>
2  void Heap<T>::_heapifyUp( ____________________ ) {
3    if ( index > _________ ) {
4      if ( item_[index] < item_[ parent(index) ] ) {
5        std::swap( item_[index], item_[ parent(index) ] );
6        _heapifyUp( ________________ );
7      }
8    }
9  }
```



	4	5	6	15	9	7	20	16	25	14	12	11			

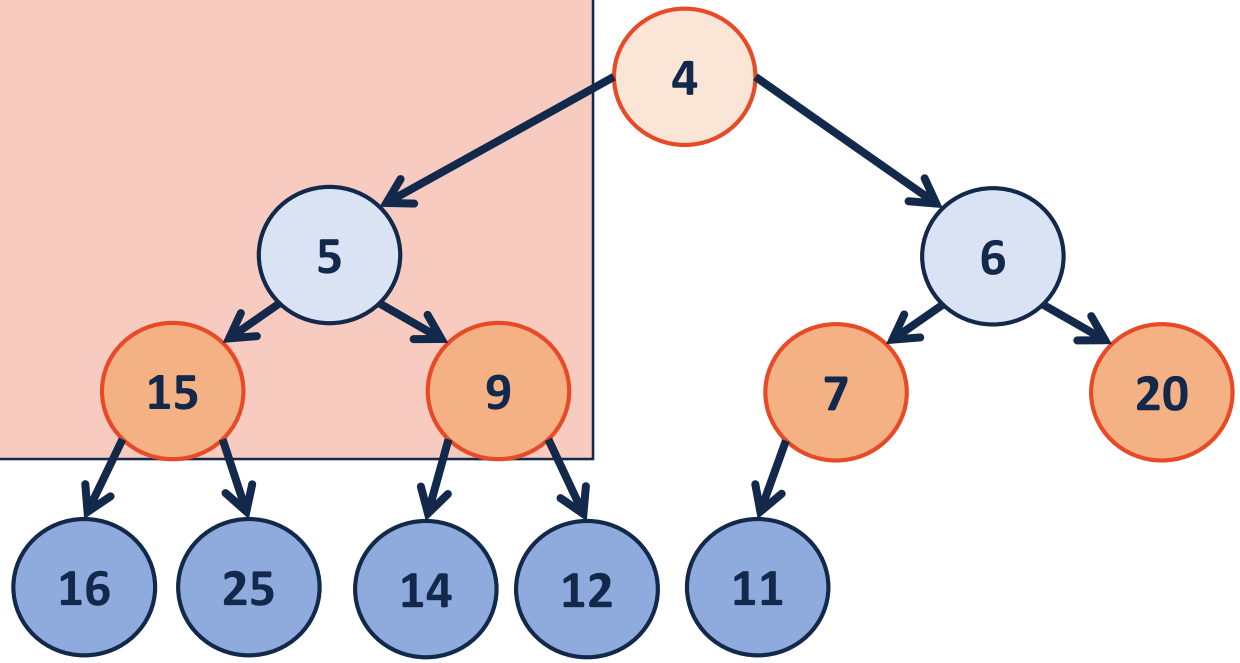
removeMin



	4	5	6	15	9	7	20	16	25	14	12	11			

removeMin

```
 1  template <class T>
 2  void Heap<T>::_removeMin() {
 3    // Swap with the last value
 4    T minValue = item_[1];
 5    item_[1] = item_[size_];
 6    size--;
 7
 8    // Restore the heap property
 9    heapifyDown();
10
11    // Return the minimum value
12    return minValue;
13  }
```



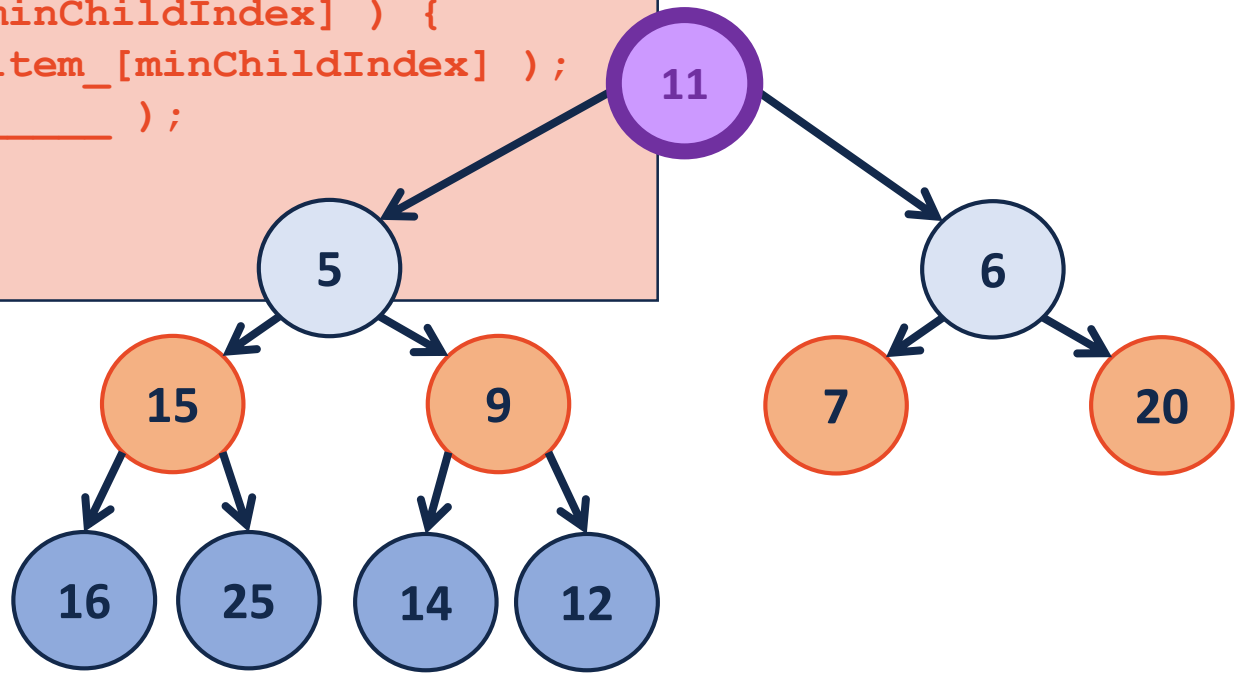
removeMin - heapifyDown

```
 1  template <class T>
 2  void Heap<T>::_removeMin() {
 3    // Swap with the last value
 4    T minValue = item_[1];
 5    item_[1] = item_[size_];
 6    size--;
 7
 8    // Restore the heap property
 9    _heapifyDown();
10
11    // Return the minimum value
12    return minValue;
13  }
```

```
 1  template <class T>
 2  void Heap<T>::_heapifyDown(int index) {
 3    if ( !_isLeaf(index) ) {
 4      T minChildIndex = _minChild(index);
 5      if ( item_[index] ___ item_[minChildIndex] ) {
 6         std::swap( item_[index], item_[minChildIndex] );
 7         _heapifyDown( ________________ );
 8      }
 9    }
10  }
```

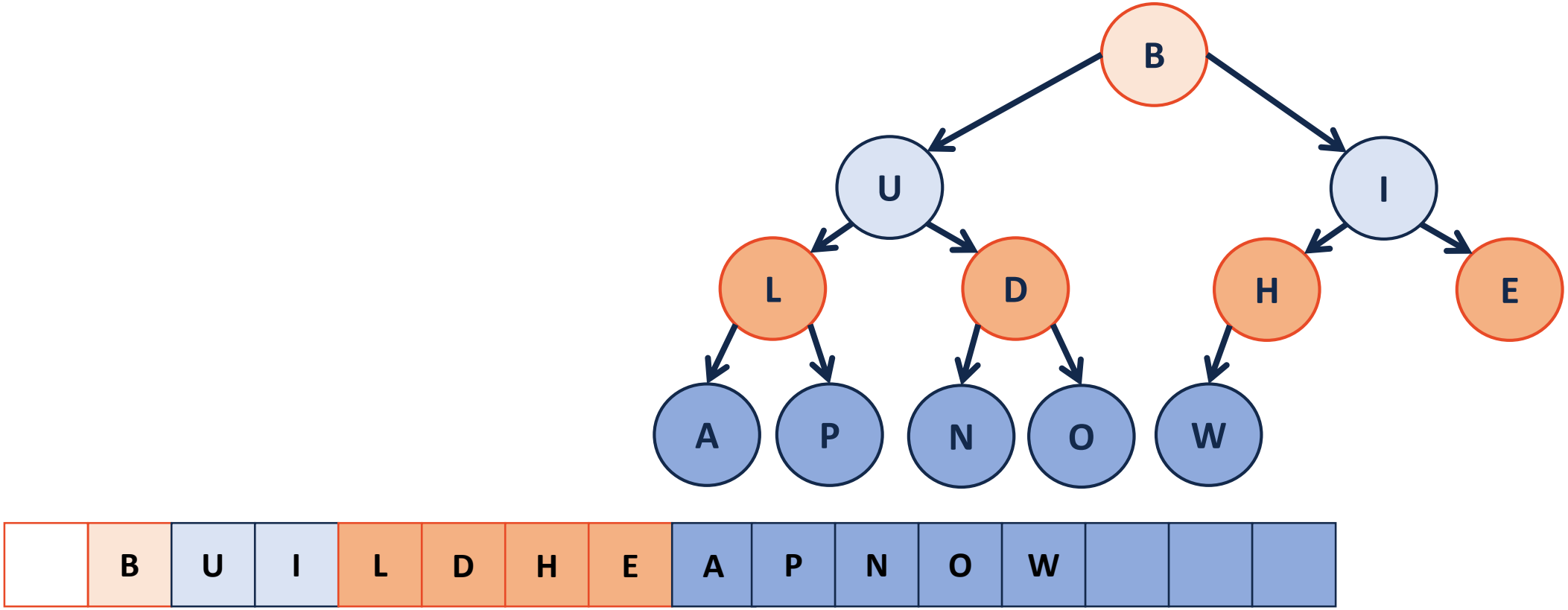
removeMin

```
1  template <class T>
2  void Heap<T>::_heapifyDown(int index) {
3    if ( !_isLeaf(index) ) {
4      T minChildIndex = _minChild(index);
5      if ( item_[index] ___ item_[minChildIndex] ) {
6         std::swap( item_[index], item_[minChildIndex] );
7         _heapifyDown( _________________ );
8      }
9    }
10 }
```

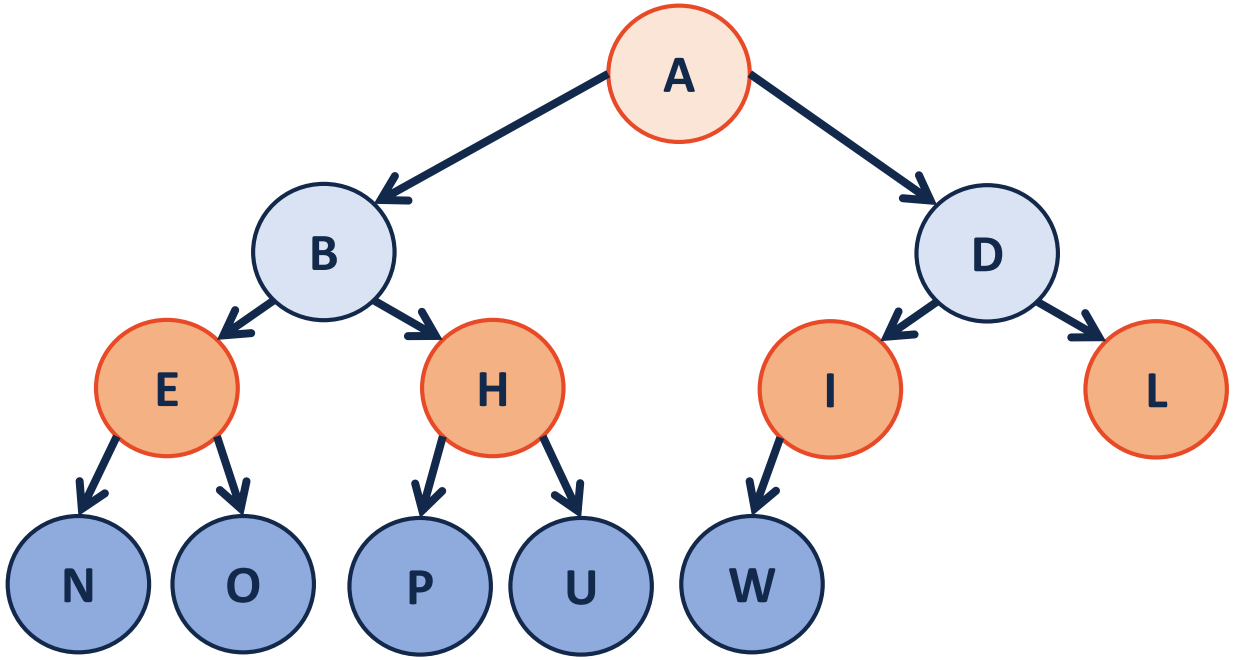


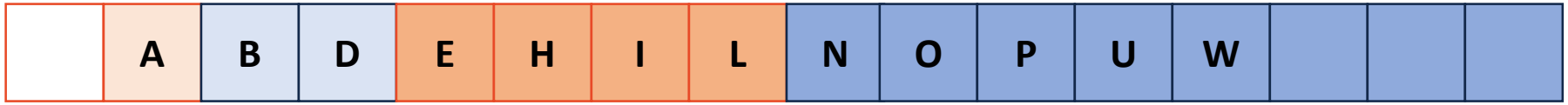
	4	5	6	15	9	7	20	16	25	14	12	11			

buildHeap

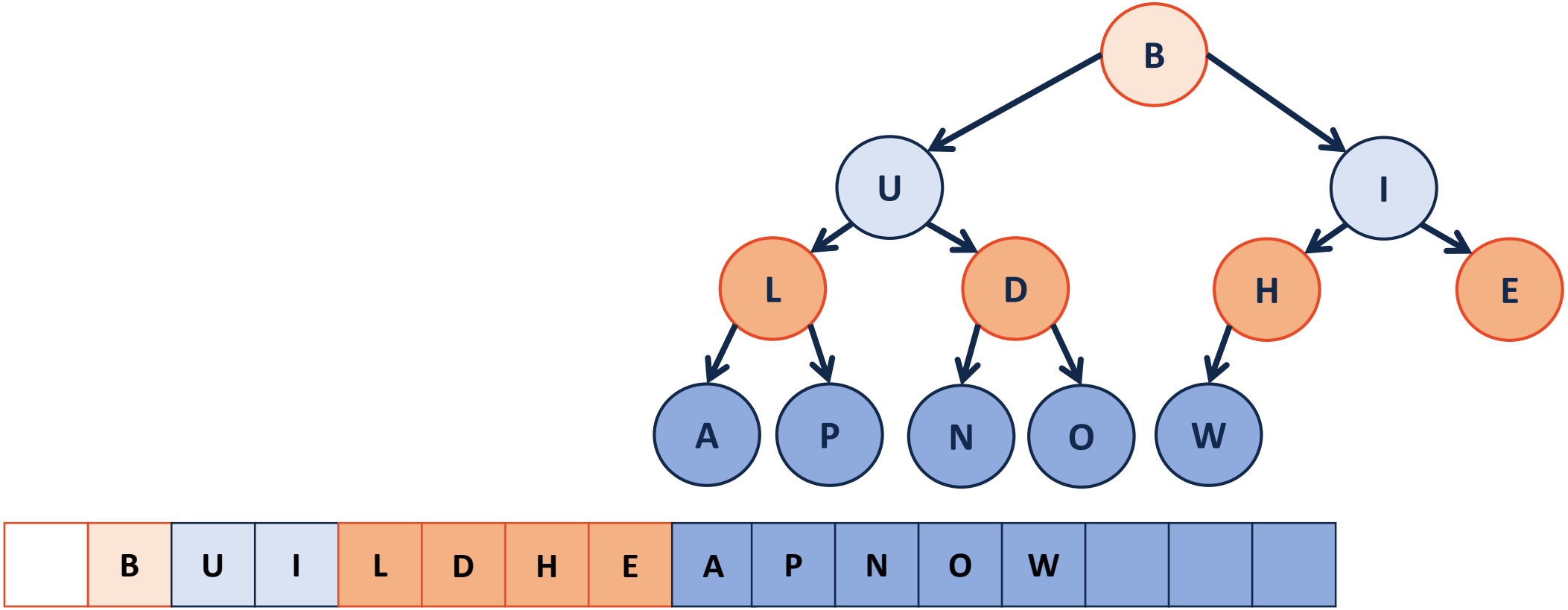


buildHeap – sorted array

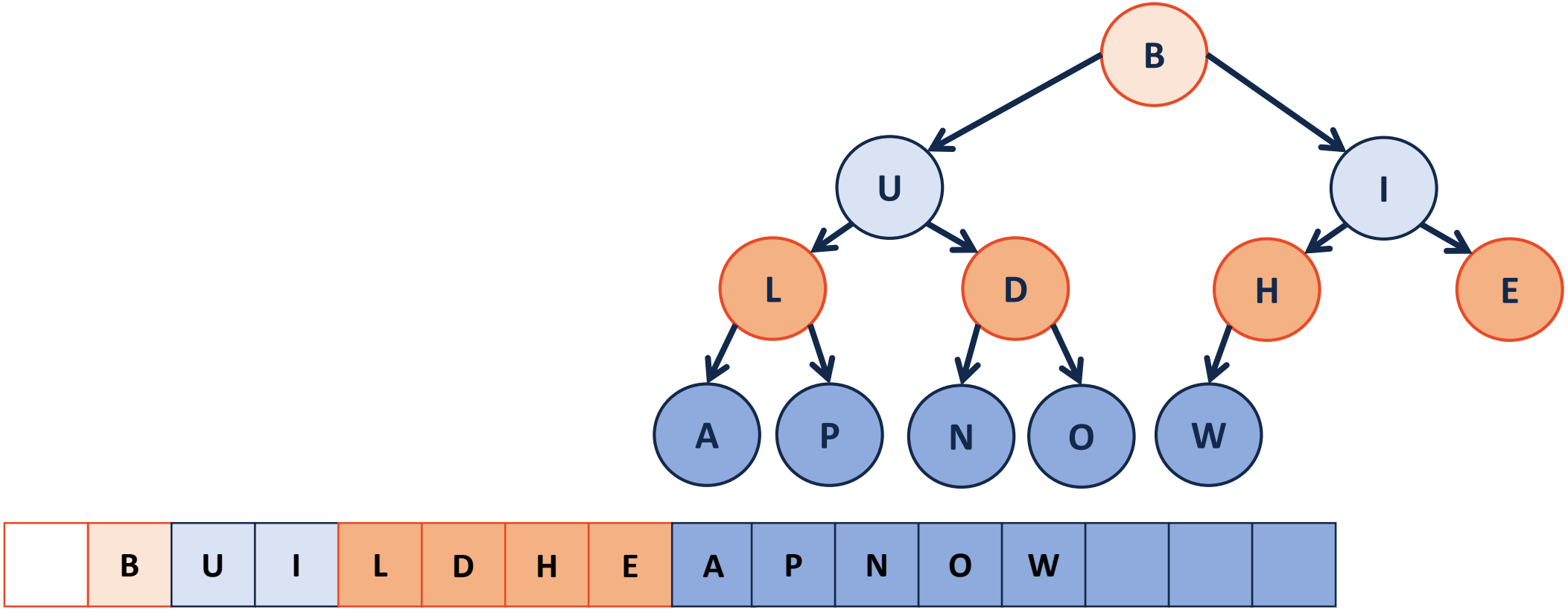




buildHeap - heapifyUp

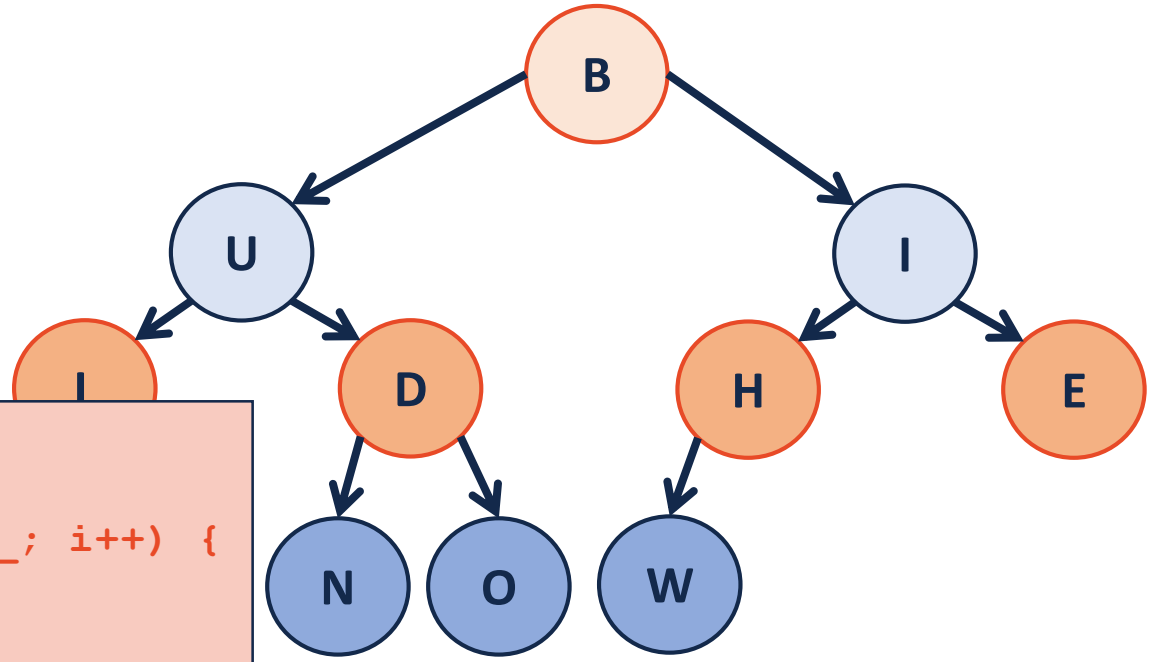


buildHeap - heapifyDown



buildHeap

1. Sort the array – it's a heap!



2.

```
1  template <class T>
2  void Heap<T>::buildHeap() {
3    for (unsigned i = 2; i <= size_; i++) {
4      heapifyUp(i);
5    }
6  }
```

3.

```
1  template <class T>
2  void Heap<T>::buildHeap() {
3    for (unsigned i = parent(size); i > 0; i--) {
4      heapifyDown(i);
5    }
6  }
```

	B	U	I	L	D	H	E	A	P	N	O	W			

Proving buildHeap Running Time

Theorem: The running time of buildHeap on array of size **n** is: _____.

Strategy:

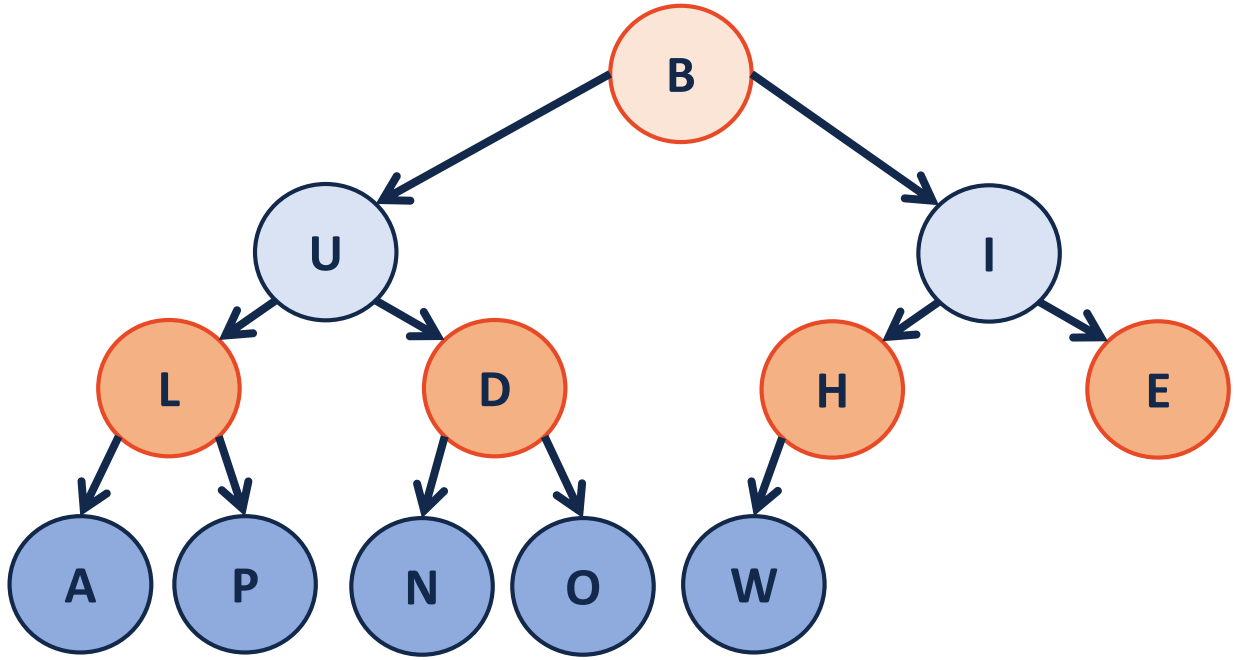
-

-

-

Proving buildHeap Running Time

S(h): Sum of the heights of all nodes in a complete tree of height **h**.



S(0) =

S(1) =

S(h) =

Proving buildHeap Running Time

Proof the recurrence:

Base Case:

General Case:

Proving buildHeap Running Time

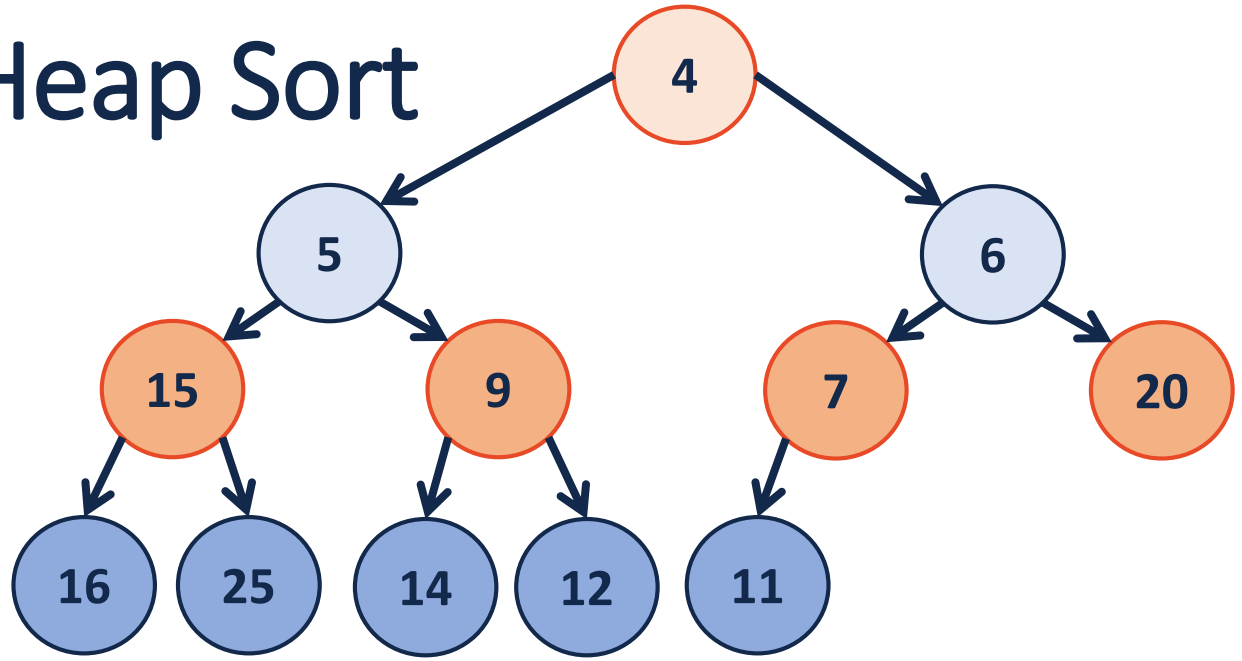
From S(h) to RunningTime(n):

S(h):

Since $h \leq \lg(n)$:

RunningTime(n) ≤

Heap Sort



1.

2.

3.

Running Time?

Why do we care about another sort?

A(nother) throwback to CS 173…

Let **R** be an equivalence relation on *us* where **(s, t) ∈ R** if **s** and **t** have the same favorite among:

{ ___, ___, _____, ___, _____, }