



# CS 225

## Data Structures

*Feb. 28 – AVL Trees*

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# Course Logistics Update

**CBTF exams will go on as-scheduled:**

- Theory Exam 2 is ongoing
- Sample Exam available on PL

**MPs and Lab assignments will be due on schedule:**

- MP4 is released; due March 12, 2018
- lab\_huffman is released later today

**My office hours are cancelled today.**

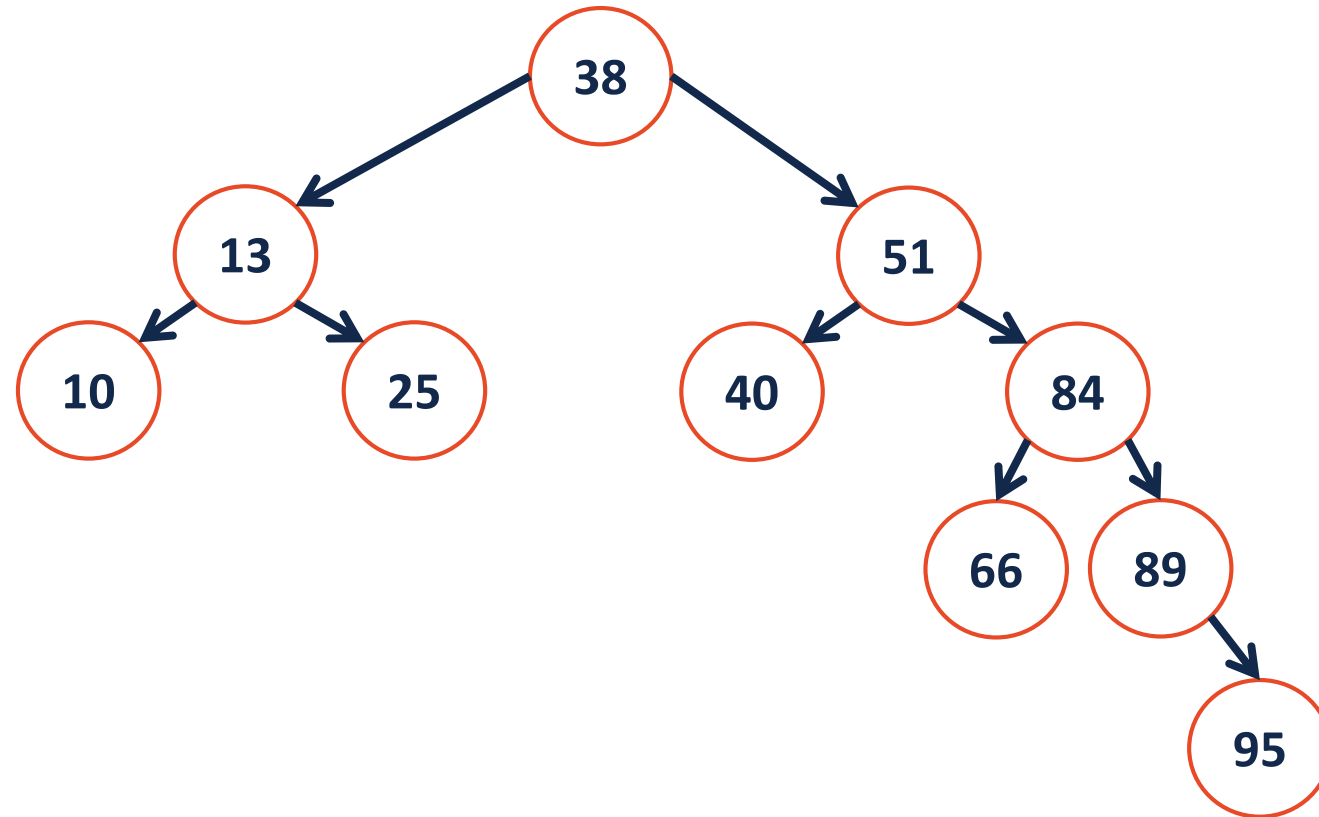
# Lab Sections

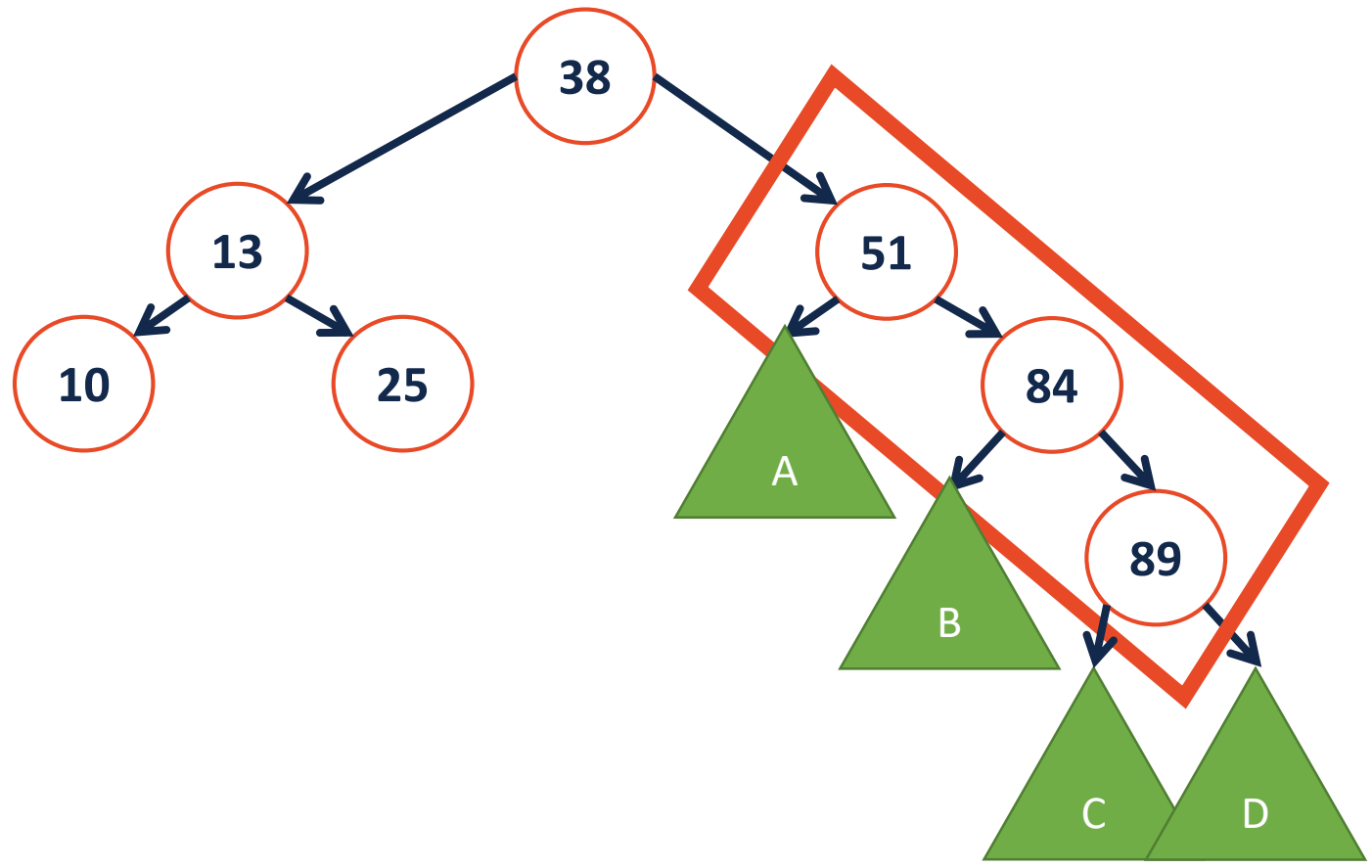
All lab sections are not meeting this week.

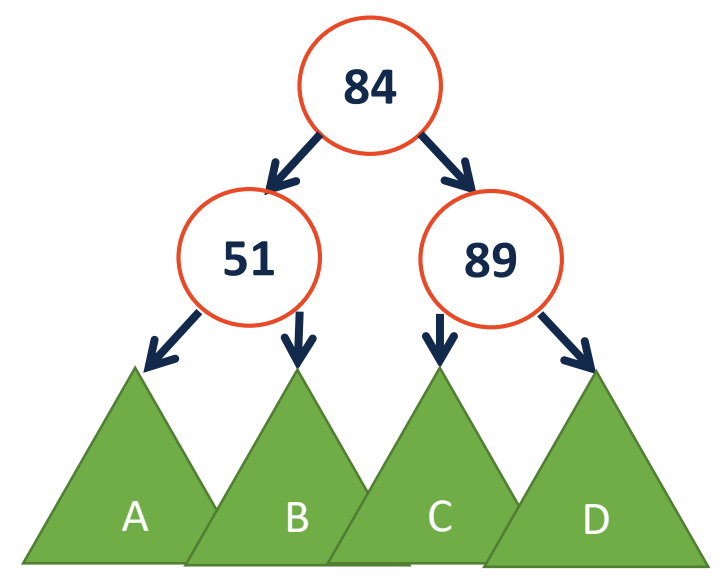
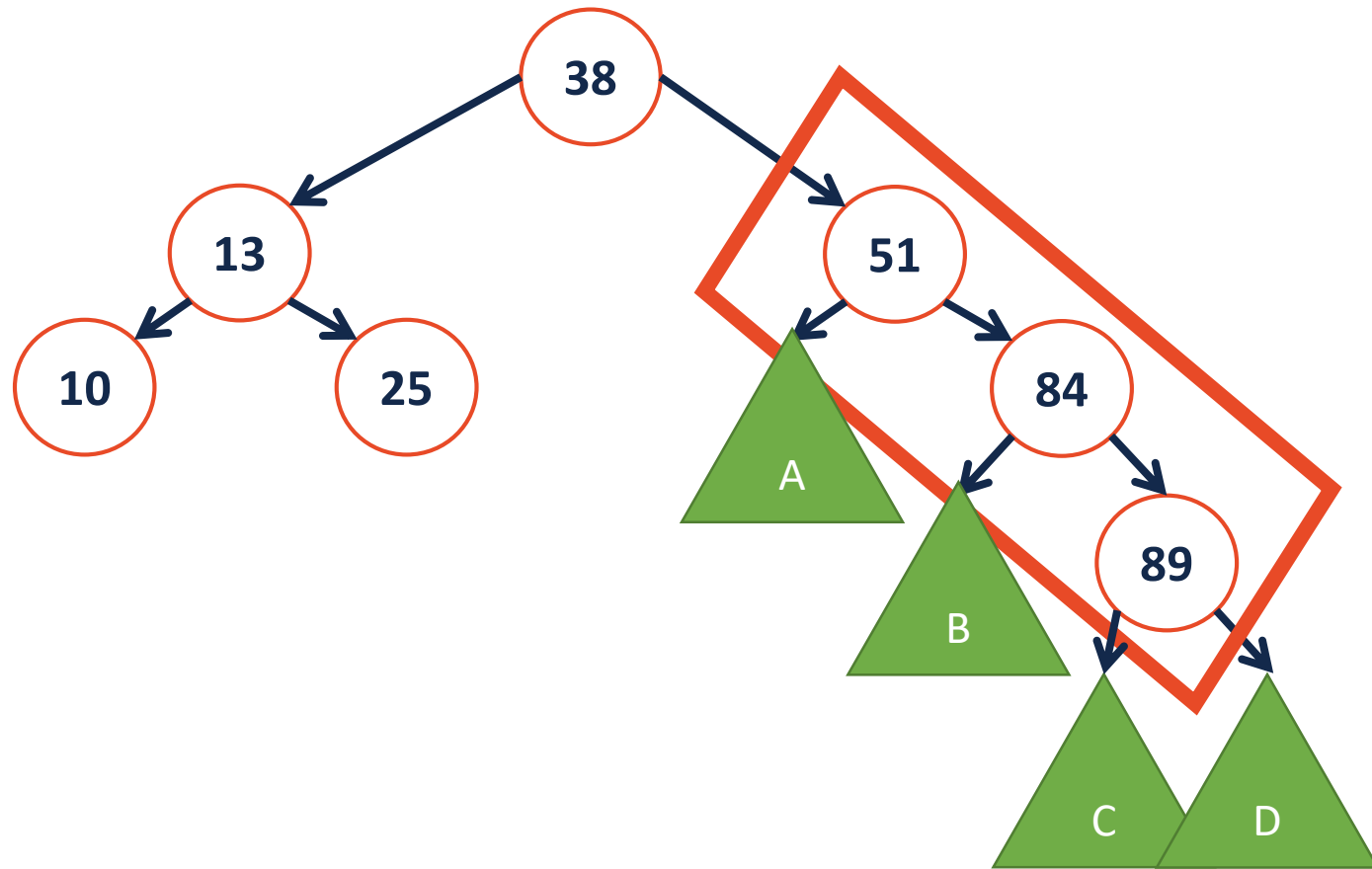
Instead, all CAs and non-striking TAs will hold open office hours (using the regular queue, held in the basement):

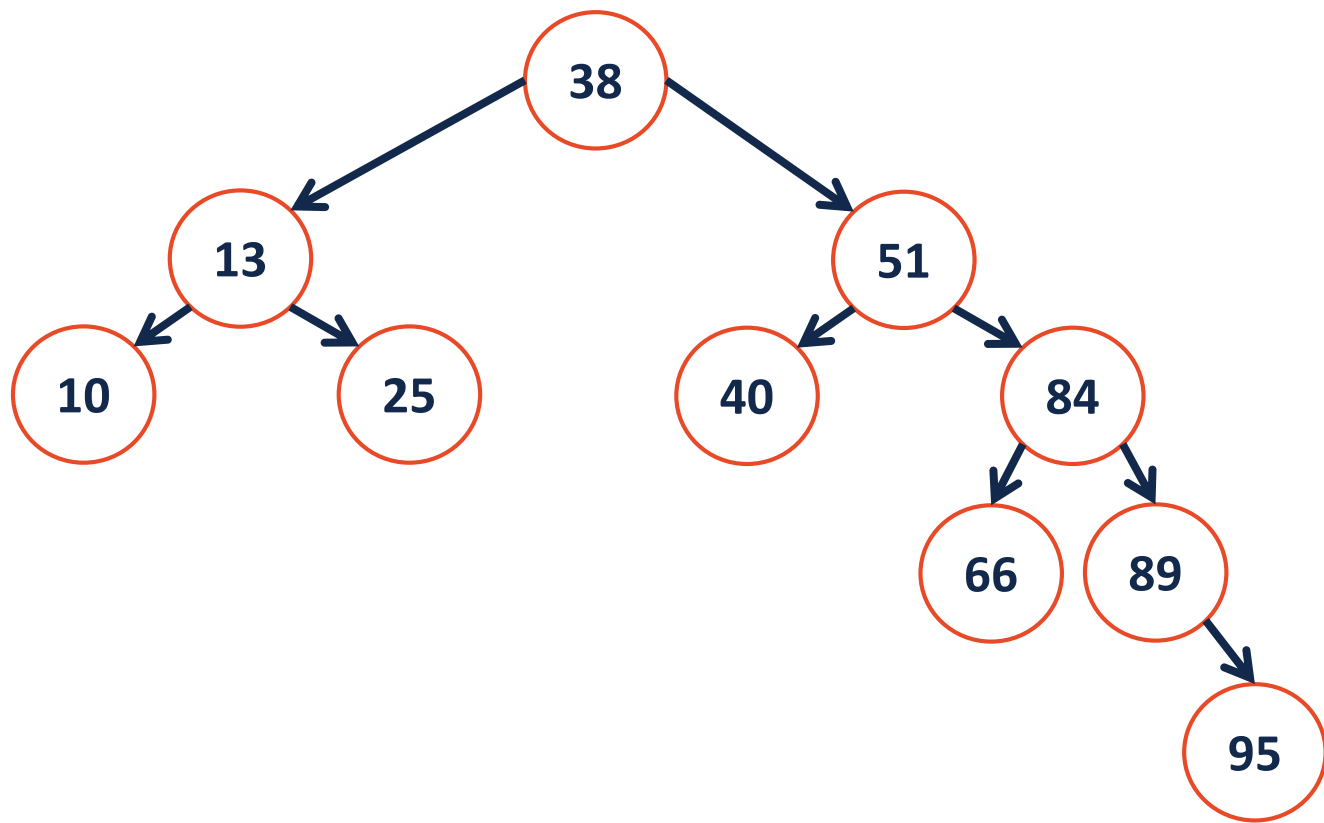
- Feel free to use the room to work with your peers on the lab. Staff will be available in open office hours in the basement of Siebel.
- *An intro video on Huffman trees will be provided.*

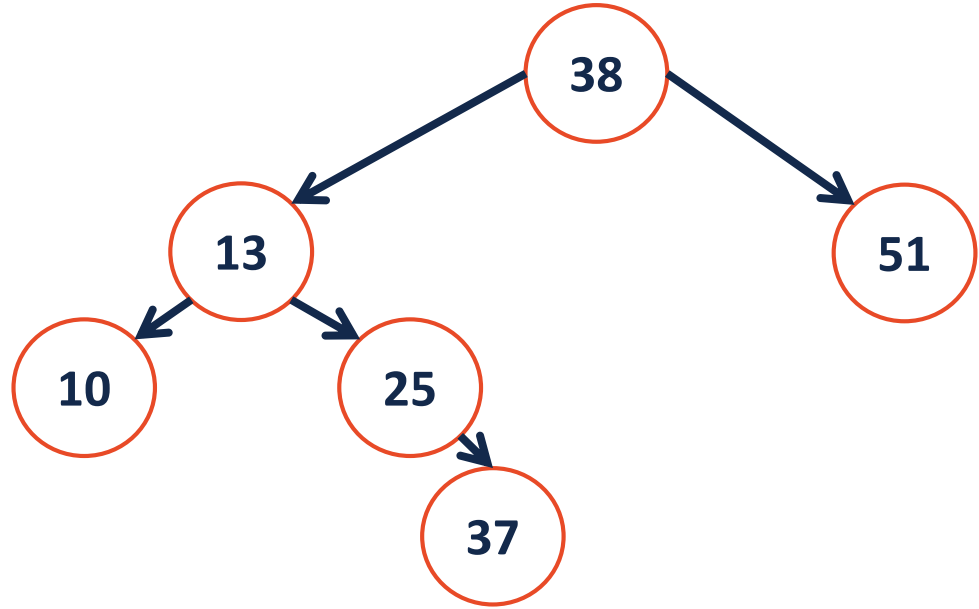
# Left Rotation



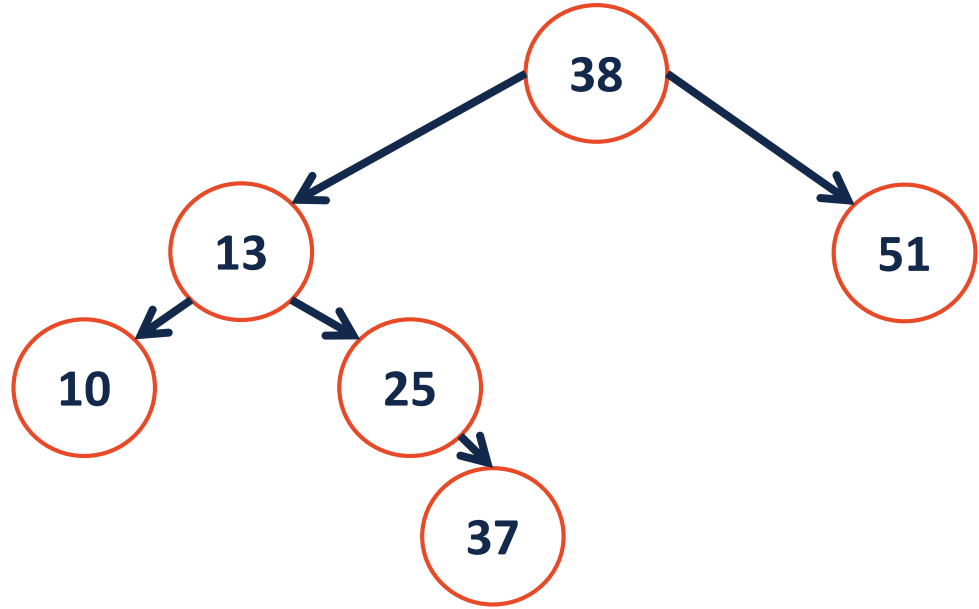












# BST Rotation Summary

- Four kinds of rotations (L, R, LR, RL)
- All rotations are local (subtrees are not impacted)
- All rotations are constant time:  $O(1)$
- BST property maintained

**GOAL:**

We call these trees:

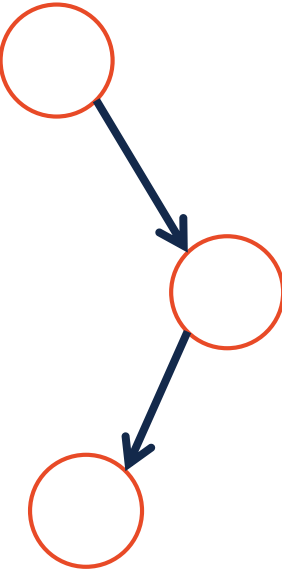
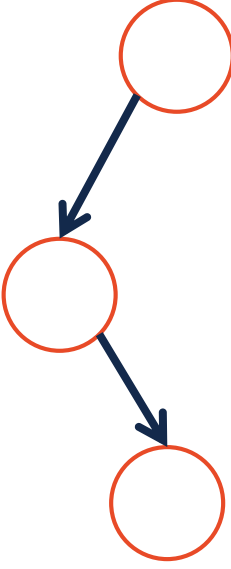
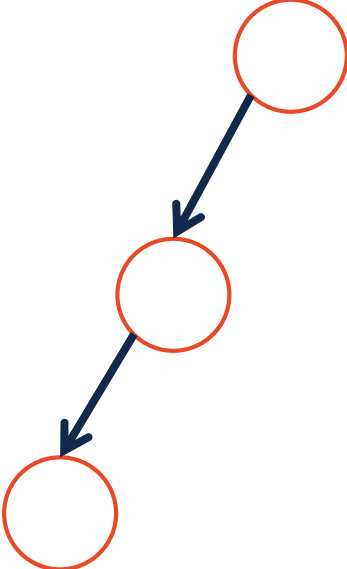
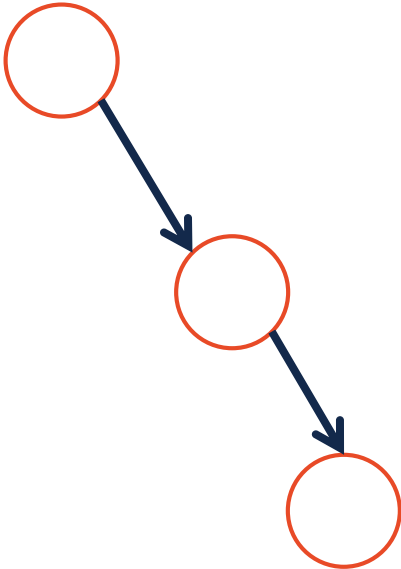
# AVL Trees

Three issues for consideration:

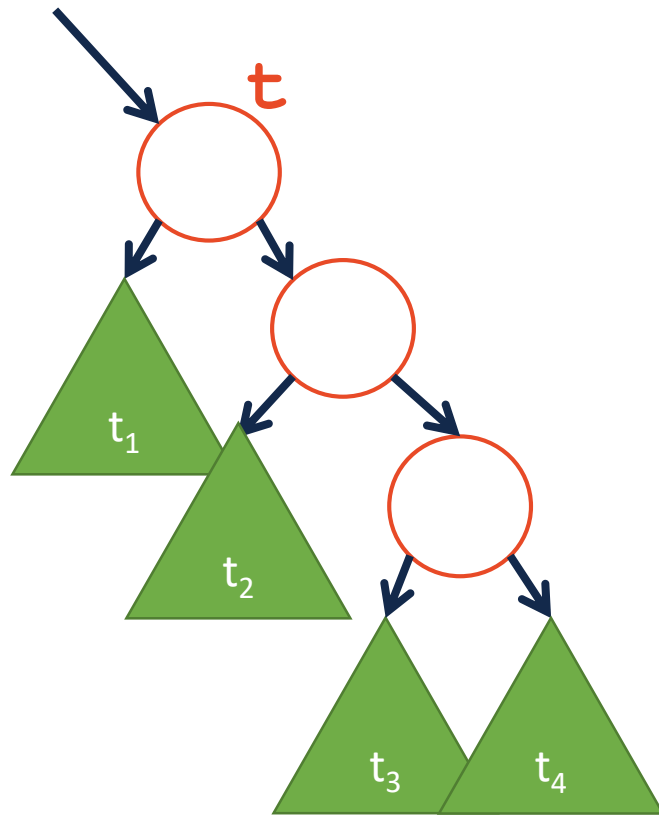
- Rotations
- Maintaining Height
- Detecting Imbalance

# AVL Tree Rotations

Four templates for rotations:



# Finding the Rotation

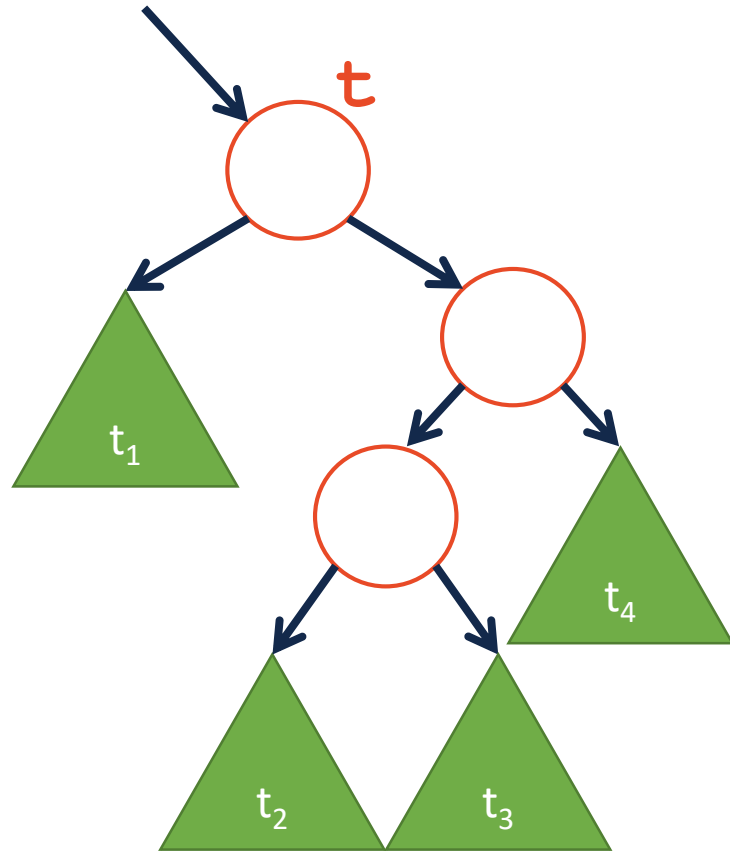


## Theorem:

If an insertion occurred in subtrees  $t_3$  or  $t_4$  and a subtree was detected at  $t$ , then a \_\_\_\_\_ rotation about  $t$  restores the balance of the tree.

We gauge this by noting the balance factor of **t-right** is \_\_\_\_\_.

# Finding the Rotation



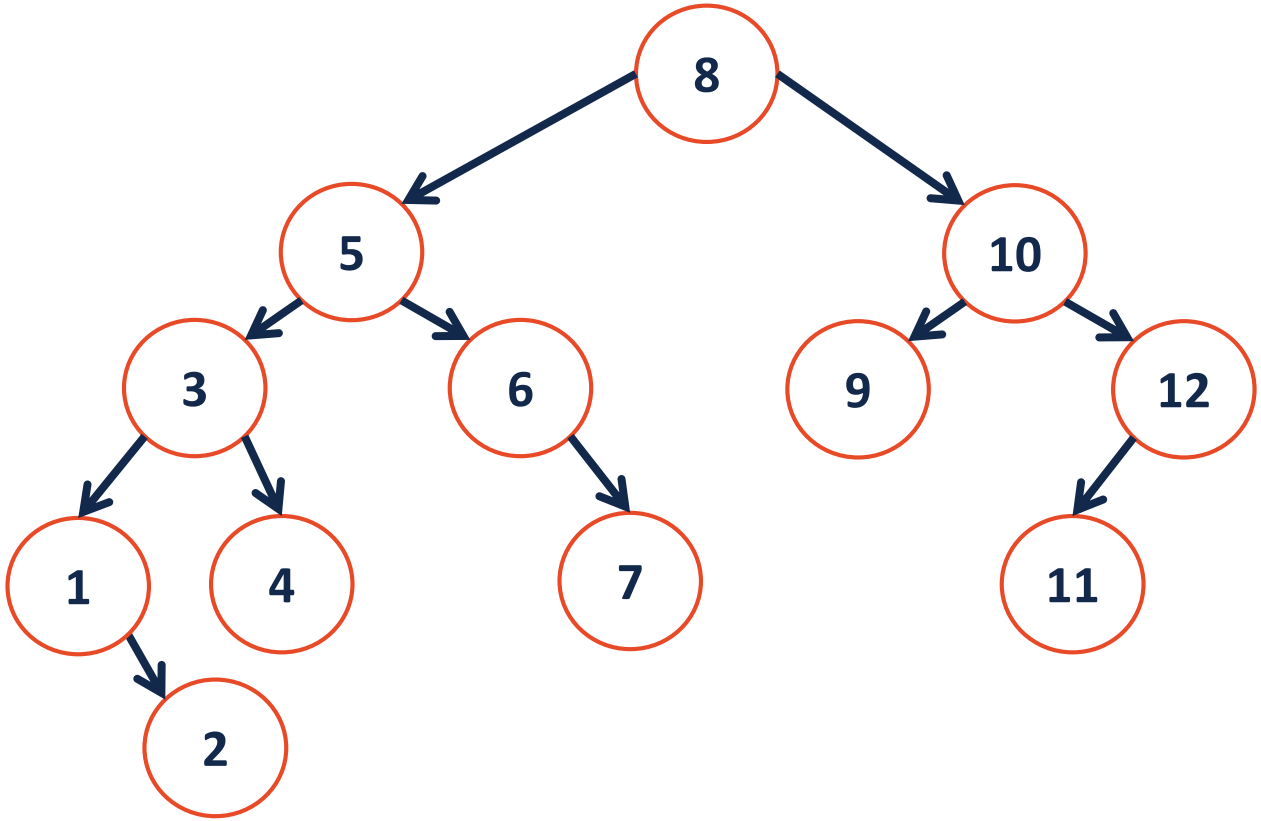
## Theorem:

If an insertion occurred in subtrees  $t_2$  or  $t_3$  and a subtree was detected at  $t$ , then a \_\_\_\_\_ rotation about  $t$  restores the balance of the tree.

We gauge this by noting the balance factor of  **$t \rightarrow \text{right}$**  is \_\_\_\_\_.

# Insertion into an AVL Tree

```
1 struct TreeNode {  
2     T key;  
3     unsigned height;  
4     TreeNode *left;  
5     TreeNode *right;  
6 };
```

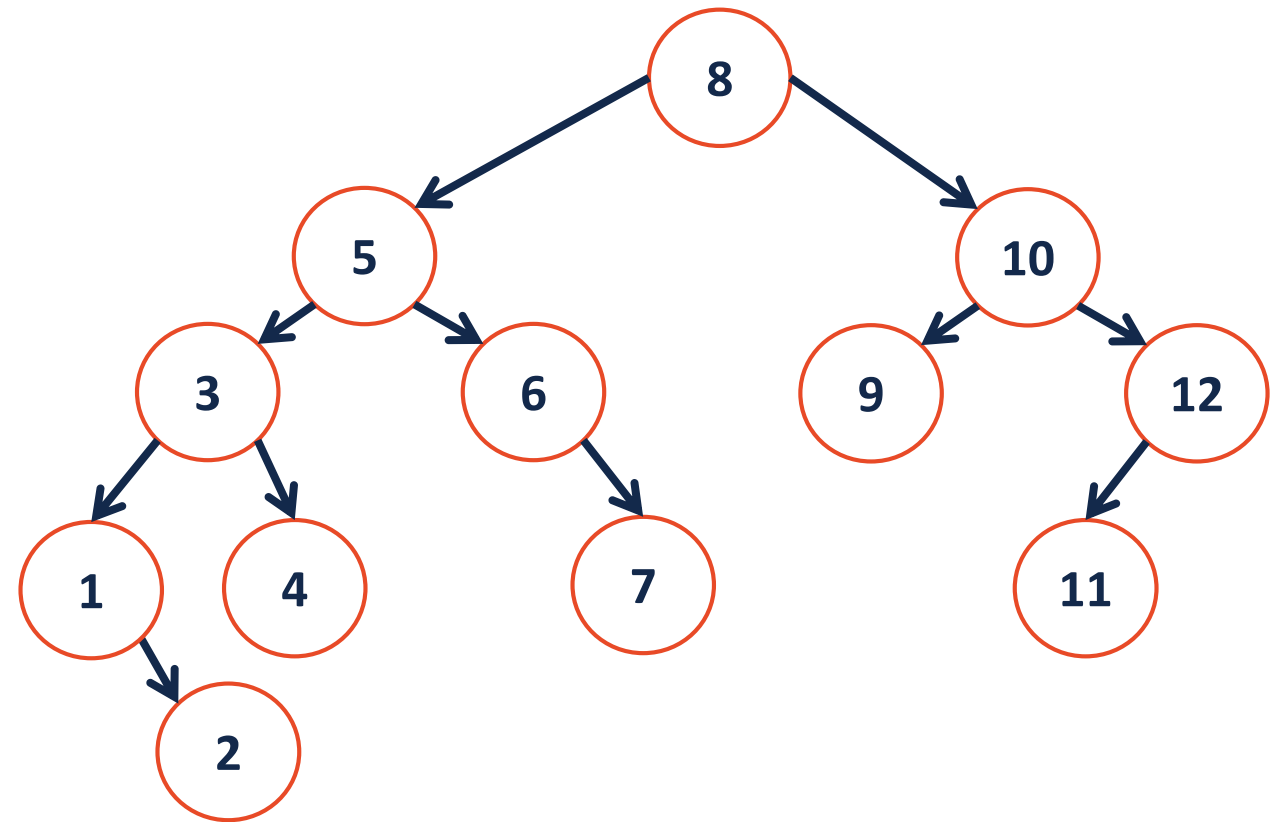


# Insertion into an AVL Tree

## Insert (pseudo code):

- 1: Insert at proper place
- 2: Check for imbalance
- 3: Rotate, if necessary
- 4: Update height

```
1 struct TreeNode {  
2     T key;  
3     unsigned height;  
4     TreeNode *left;  
5     TreeNode *right;  
6 };
```

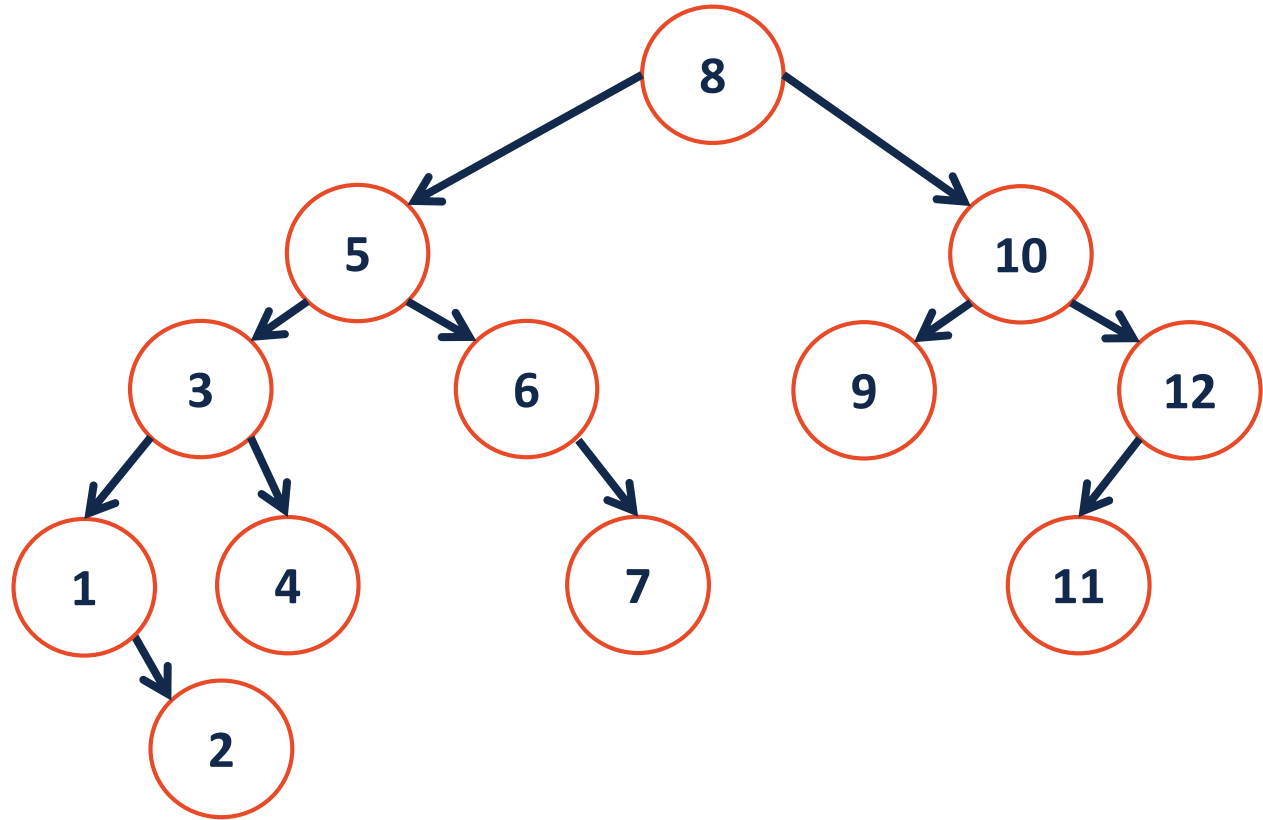




```
1  template <class T> void AVLTree<T>::_insert(const T & x, treeNode<T> * & t ) {
2      if( t == NULL ) {
3          t = new TreeNode<T>( x, 0, NULL, NULL);
4      }
5
6      else if( x < t->key ) {
7          _insert( x, t->left );
8          int balance = height(t->right) - height(t->left);
9          int leftBalance = height(t->left->right) - height(t->left->left);
10         if ( balance == -2 ) {
11             if ( leftBalance == -1 ) { rotate_____ ( t ); }
12             else { rotate_____ ( t ); }
13         }
14     }
15
16     else if( x > t->key ) {
17         _insert( x, t->right );
18         int balance = height(t->right) - height(t->left);
19         int rightBalance = height(t->right->right) - height(t->right->left);
20         if( balance == 2 ) {
21             if( rightBalance == 1 ) { rotate_____ ( t ); }
22             else { rotate_____ ( t ); }
23         }
24     }
25
26     t->height = 1 + max(height(t->left), height(t->right));
27 }
```

# Height-Balanced Tree

Height balance:  $b = \text{height}(T_R) - \text{height}(T_L)$



# AVL Tree Analysis

**We know:** insert, remove and find runs in: \_\_\_\_\_.

**We will argue that:**  $h =$  \_\_\_\_\_.

# AVL Tree Analysis

Definition of big-O:

...or, with pictures:

