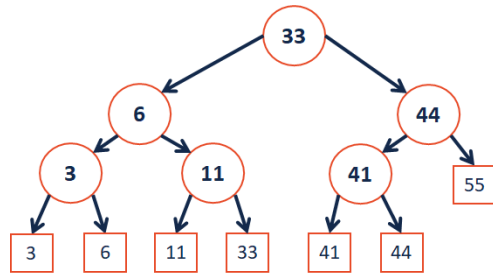


**Range-based Searches:**

Q: Consider points in 1D:  $p = \{p_1, p_2, \dots, p_n\}$ .  
...what points fall in  $[11, 42]$ ?



**Tree Construction:**

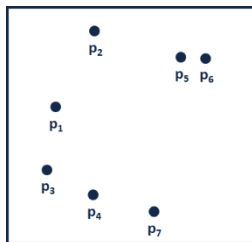


**Range-based Searches:**

**Running Time:**

**Extending to k-dimensions:**

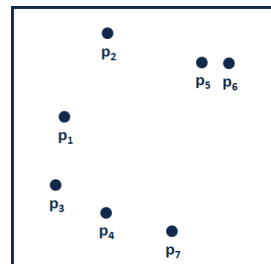
Consider points in 2D:  $p = \{p_1, p_2, \dots, p_n\}$ :



...what points are inside a range (rectangle)?  
...what is the nearest point to a query point  $q$ ?

**kd-Tree Motivation:**

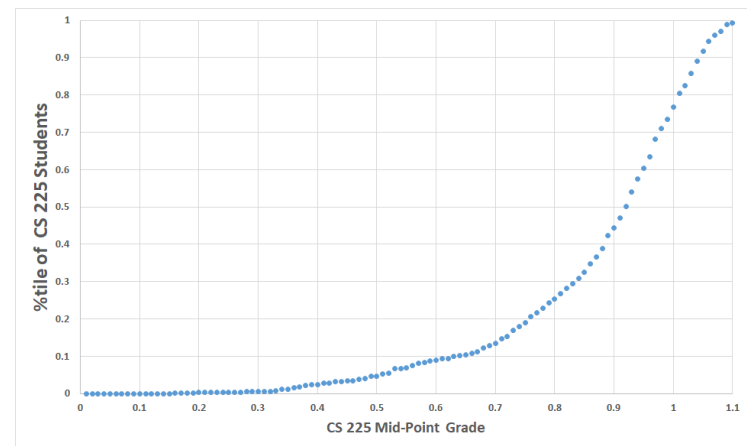
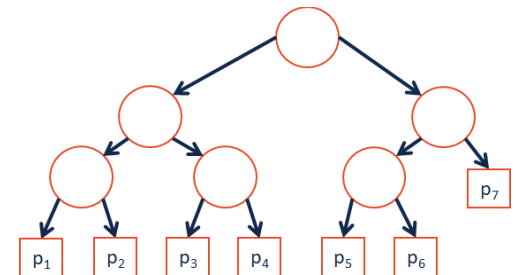
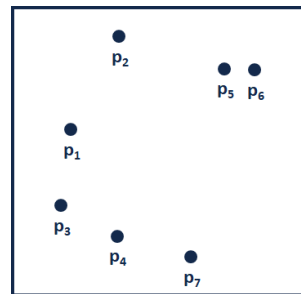
First, let's try and divide our space up:



**kd-Tree Construction:**

How many dimensions exist in our input space?

How do we want to "order" our dimensions?



## Motivation

Can we always fit our data in main memory?

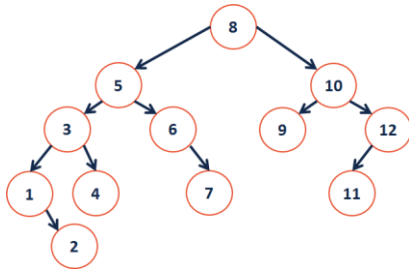
Where else do we keep our data?

-

-

vs. CPU: 3 GHz == 3m ops / \_\_\_\_\_ \* \_\_\_\_\_ cores

## AVL Operations on Disk:



How deep do AVL trees get?

## BTree Motivations

Knowing that we have long seek times for data, we want to build a data structure with two (related) properties:

1.

2.

---

---

## BTree<sub>m</sub>



**Goal:** Build a tree that uses \_\_\_\_\_ /node!  
...optimize the algorithm for your platform!

A **BTree of order m** is an m-way tree where:

1. All keys within a node are ordered.
2. All leaves contain no more than **m-1** nodes.

---

## BTree Insert, using m=5

...when a BTree node reaches **m** keys:

### CS 225 – Things To Be Doing:

1. Programming Exam B starts next Tuesday (March 13<sup>th</sup>)
2. MP4 extra credit ongoing (final deadline March 12<sup>th</sup>)
3. lab\_avl released this week; lab sections **are** being held this week!
4. Daily POTDs are ongoing!