

starter slides
from Wednesday
(if following along using PDFs)
Queues +

Data Structures

Iterators & Exam 1 Review

CS 225

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Announcements → *Schedule*

Reminder: Labs are due on Sunday!

MP is due on Monday (late day Tuesday for 93%)

Next MP (MP_lists, focus on iterators and lists) out Monday



Exam 1 (9/16 — 9/18)

Autograded MC and one coding question

Manually graded short answer prompt

out since Wednesday
Practice exam ~~will be released~~ on PL

Topics covered can be found on website

Register now

<https://courses.engr.illinois.edu/cs225/exams/>

Learning Objectives

Finish discussing ^{Introduces} queue implementation details

Discuss the importance of iterators

Review content seen on exam 1 (and how to prepare)

Introduce trees and the tree ADT (Time permitting) ~~X~~

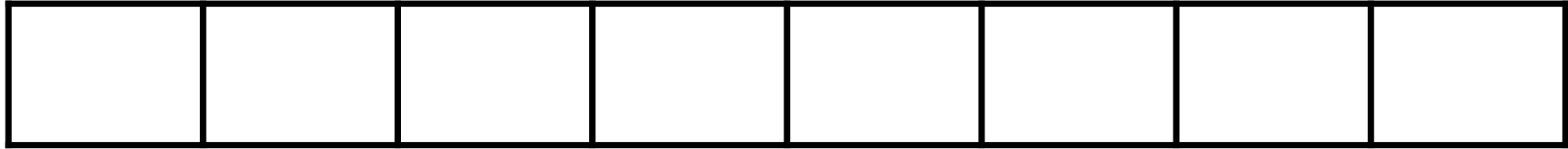
Stack and Queue

Interface
Stack

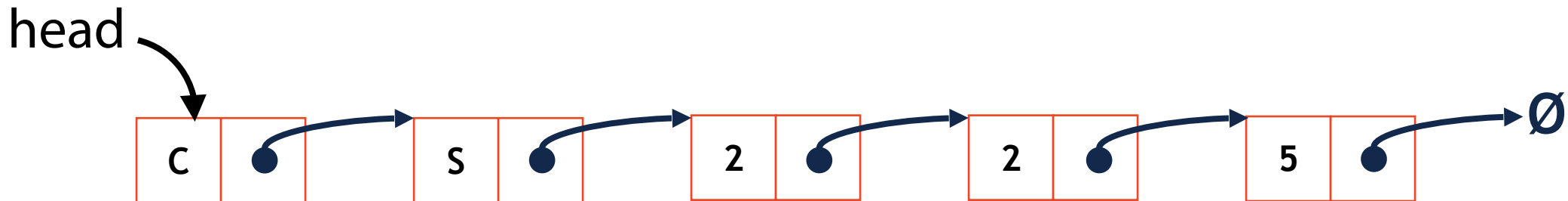
Implement
as array

Taking advantage of special cases in lists / arrays

insertBack $O(1)$



insertFront $O(1)$ Stack as linked list



Stack ADT

- [Order]: LIFO (Last in first out)



- [Implementation]: Array (such as `std::vector`)

Back insert / Remove

*O(1)**

O(1)

Linked List also works using insert / remove Front

- [Runtime]: $O(1)$ Push, Pop, Top

If using array, $O(1)^*$ if we need to resize.

Queue Data Structure

A **queue** stores an ordered collection of objects (like a list)

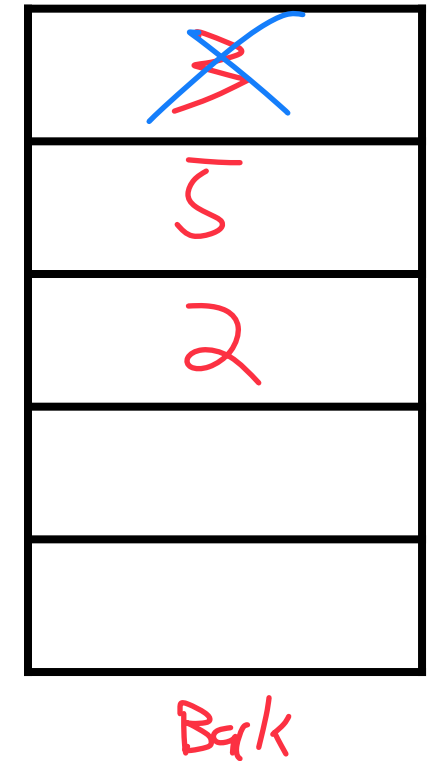
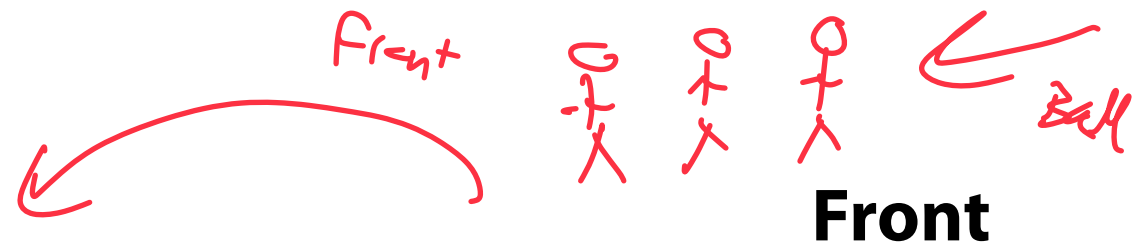
However you can only do three* operations:

Enqueue: Put an item at the back of the queue

Dequeue: Remove the front item of the queue

Front: Return the front item of the queue

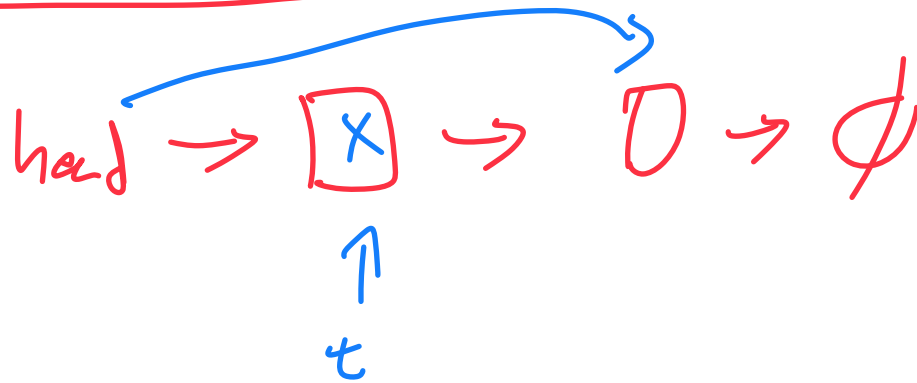
enqueue (3) ; enqueue (5) ; dequeue () ; enqueue (2)



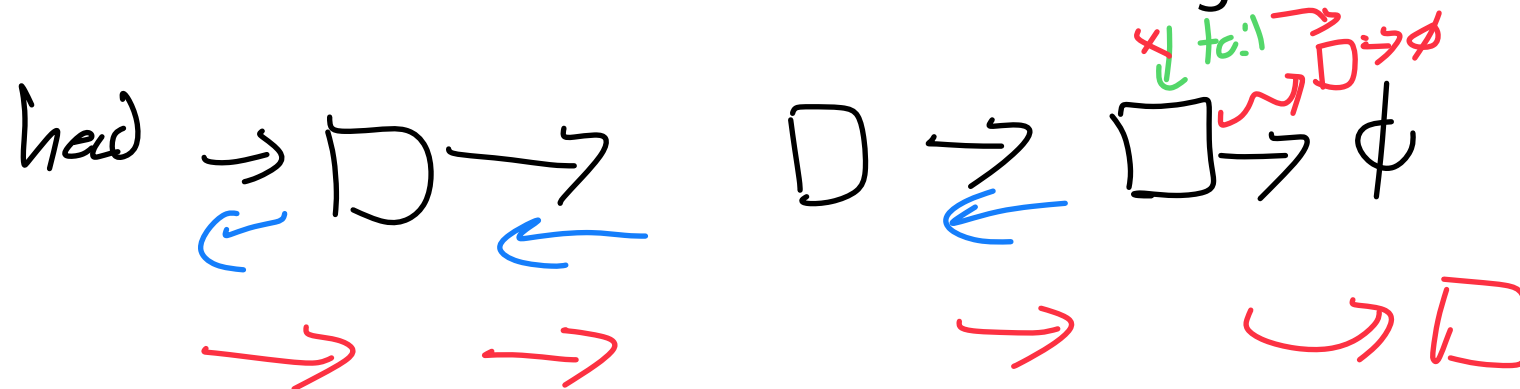
Queue Data Structure

The queue is a **first in — first out** data structure (FIFO)

What data structure excels at removing from the front? *Linked List*



Can we make that same data structure good at inserting at the end?



Doubly linked list does not give $O(1)$ insert Back

tail pointer makes $O(1)$ insert Back

Queue Data Structure

The C++ implementation of a queue is also a vector or deque — why?

1) Engineering

↳ The time to follow pointers is not free

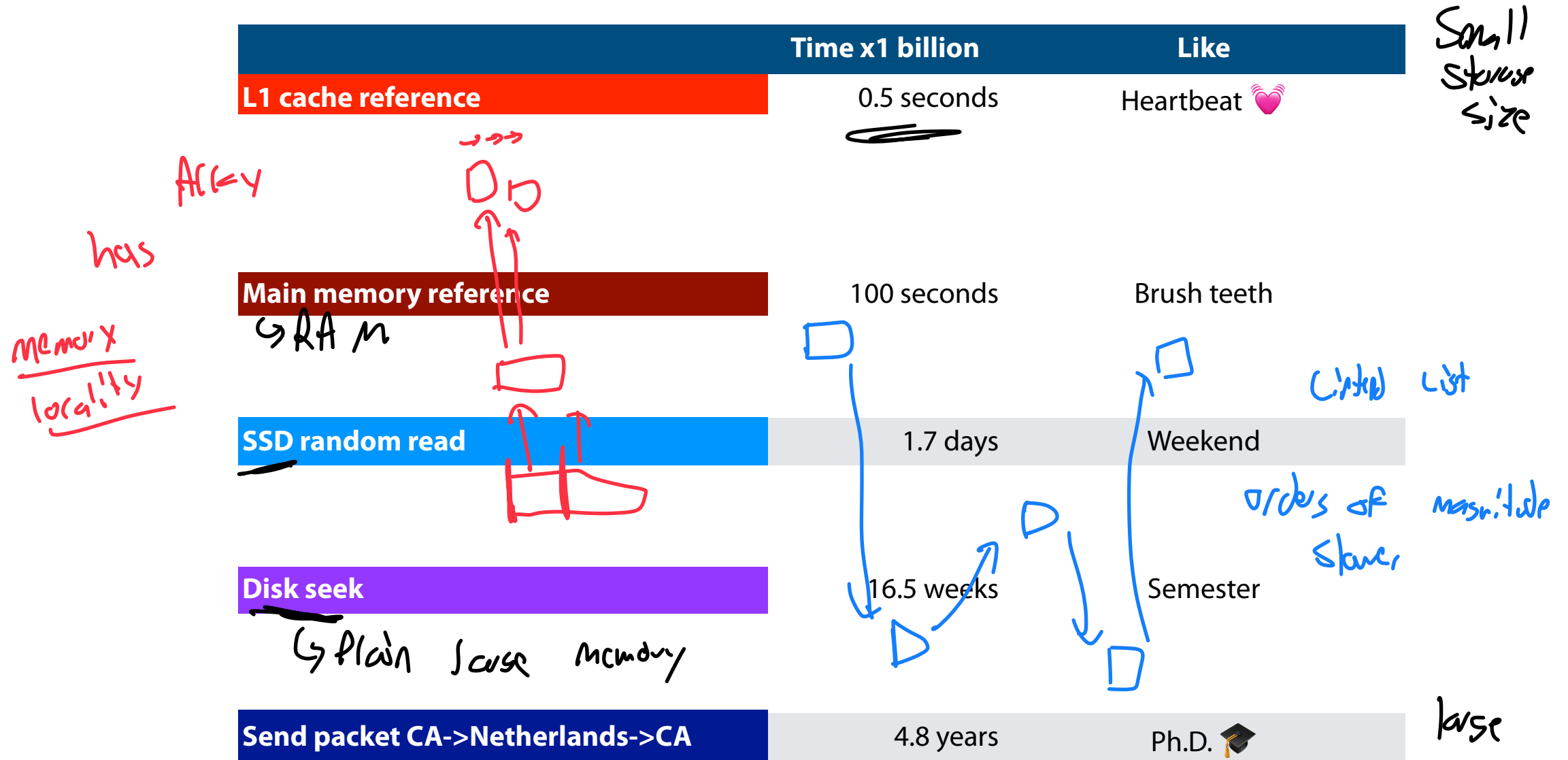
2) We can improve the array (we have the tech...)

Engineering vs Theory Efficiency

	Time x1 billion	Like
L1 cache reference	0.5 seconds	Heartbeat 💖
Branch mispredict	5 seconds	Yawn 🥱
L2 cache reference	7 seconds	Long yawn 🥱 🥱 🥱
Mutex lock/unlock	25 seconds	Make coffee ☕
Main memory reference	100 seconds	Brush teeth
Compress 1K bytes	50 minutes	TV show 📺
Send 2K bytes over 1 Gbps network	5.5 hours	(Brief) Night's sleep 🛌
SSD random read	1.7 days	Weekend
Read 1 MB sequentially from memory	2.9 days	Long weekend
Read 1 MB sequentially from SSD	11.6 days	2 weeks for delivery 📦
Disk seek	16.5 weeks	Semester
Read 1 MB sequentially from disk	7.8 months	Human gestation 🐣
Above two together	1 year	🌍 ☀️
Send packet CA->Netherlands->CA	4.8 years	Ph.D. 🎓

(Care of <https://gist.github.com/hellerbarde/2843375>)

Engineering vs Theory Efficiency



(Care of <https://gist.github.com/hellerbarde/2843375>)

Queue Data Structure

```
q.enqueue(8);  
q.enqueue(4);  
q.dequeue();
```

What do we need to track to maintain a queue with an array list?

lets add a new variable front

Front

↳ data is the first allocated mem giving
↳ No longer guaranteed to be front item

data



size

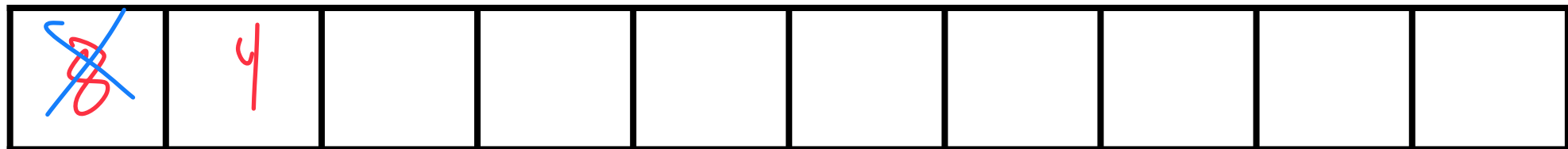


Capacity



Front

Back

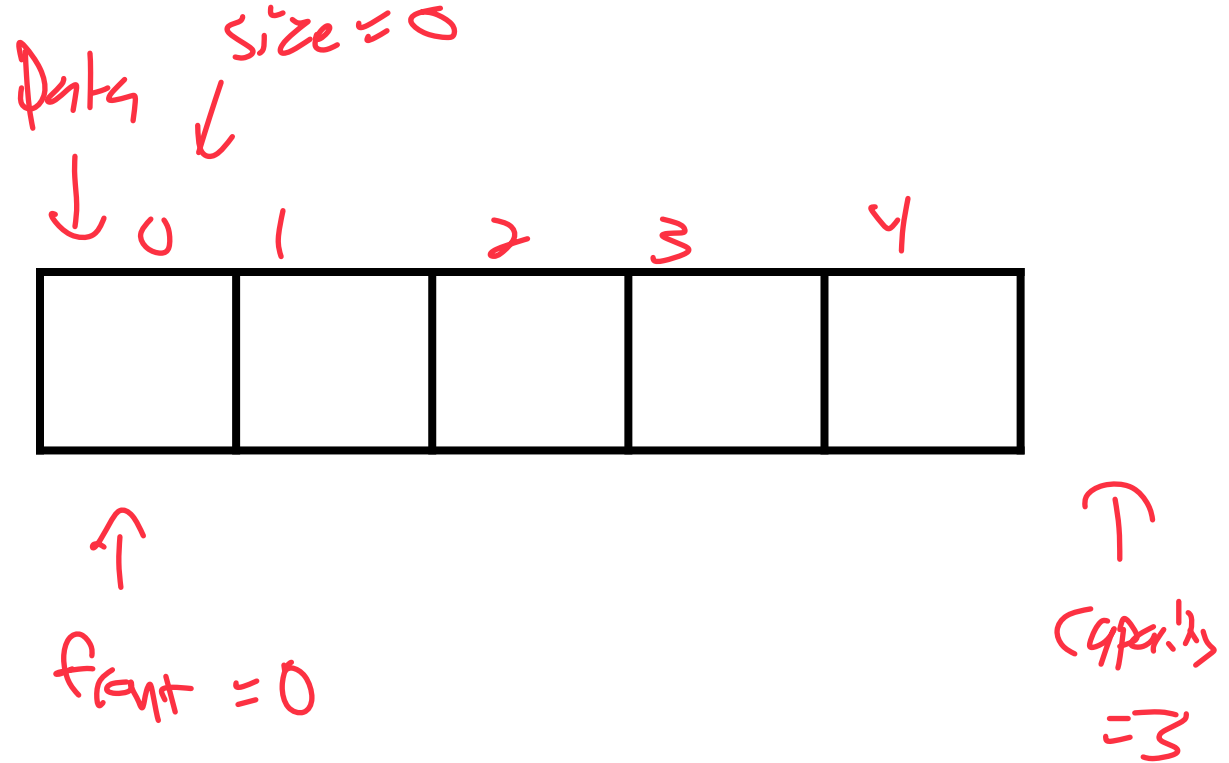


Queue Data Structure

Unlike the array list, it is easier to implement a Queue using unsigned ints

Queue.h

```
1 #pragma once
2
3 template <typename T>
4 class Queue {
5     public:
6         void enqueue(T e);
7         T dequeue();
8         bool isEmpty();
9
10    private:
11        T *data_;
12        unsigned size_;
13        unsigned capacity_;
14        unsigned front_;
15};
```



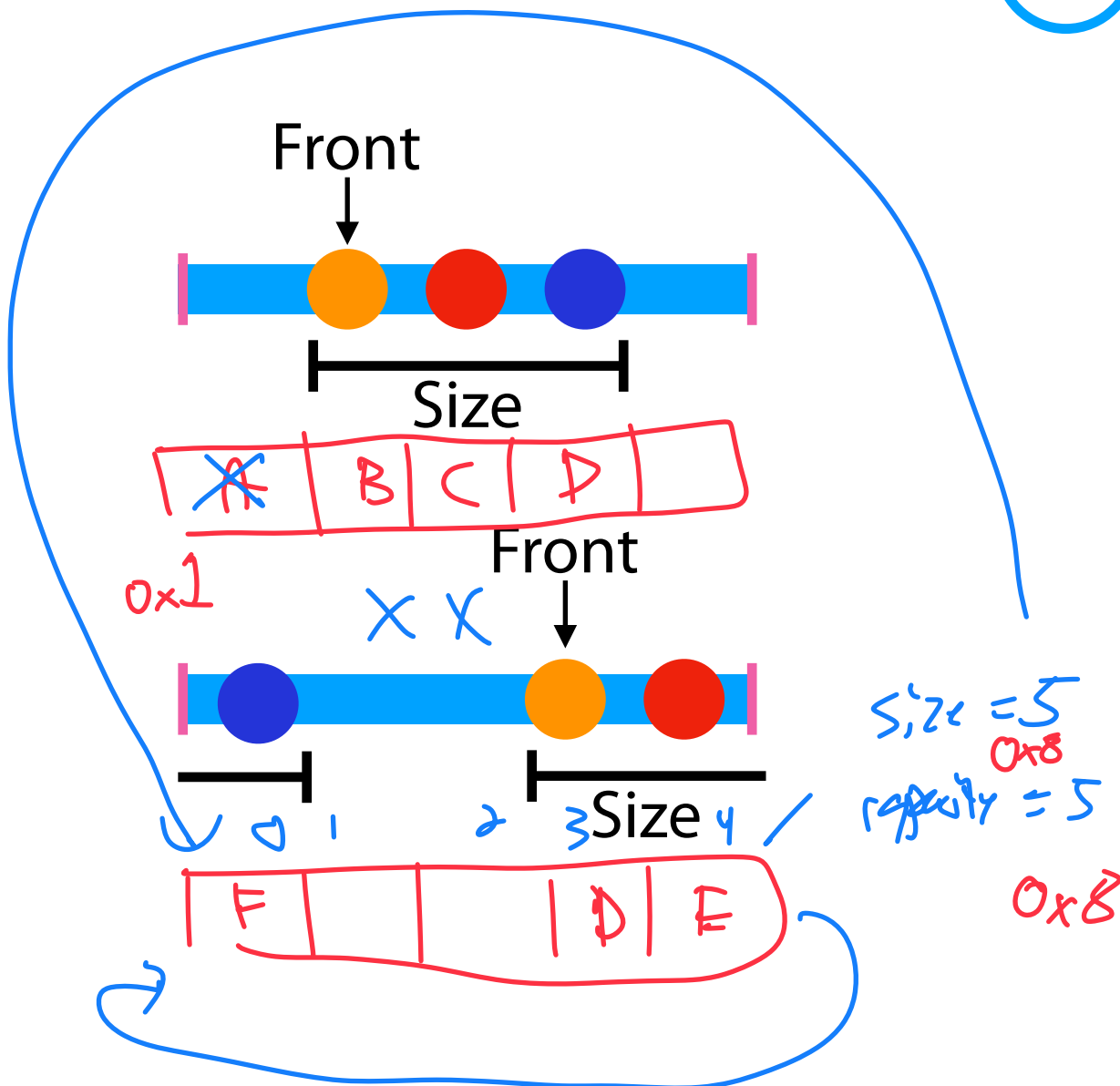
(Circular) Queue Data Structure

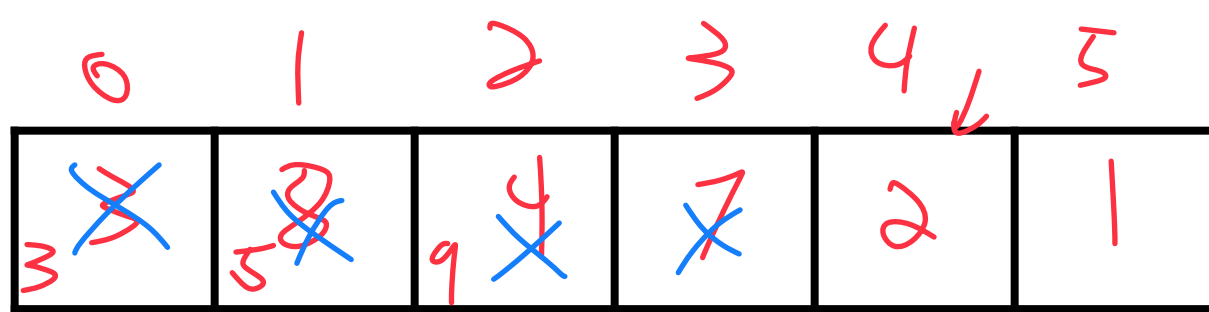


Queue.h

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1 #pragma once
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3 template <typename T>
4 class Queue {
5     public:
6         void enqueue(T e);
7         T dequeue();
8         bool isEmpty();
9
10    private:
11        T *data_;
12        unsigned capacity_;
13        unsigned size_;
14        unsigned front_;
15 };
```

into ints





Enqueue(D): Add to "Back" of array : $(\text{Front} + \text{size}) \% \text{capacity}$
 Size ++

Dequeue(): Remove the front item (@ index front)
 front ++ % capacity
 Size --

Size: ~~0~~ 1 ~~2~~ 3 ~~4~~ 5 ~~6~~ 7

Front: ~~0~~ 1 2 3

6%6 stick figure ← stick figure

```

Front, Size
0, 0 - q.enqueue(3);
0, 1 - q.enqueue(8);
0, 2 - q.enqueue(4);
      q.dequeue();
      q.enqueue(7);
      q.dequeue();
      q.dequeue();
      q.enqueue(2);
      q.enqueue(1);
      q.enqueue(3);
      q.enqueue(5);
      q.dequeue();
      q.enqueue(9);
  
```

Capacity: 6

Both actions have required checks

Enqueue(D): (as long as size $\neq$ capacity)

if Σ

Add data to 'back' of queue

Insert D at index **(size+front) % capacity**

size++

Σ
else {

resize $\rightarrow$

new Array = $\times$ size

for i in range((capacity) Σ

newA[i] = oldA(i+front % capacity);

Add one new data

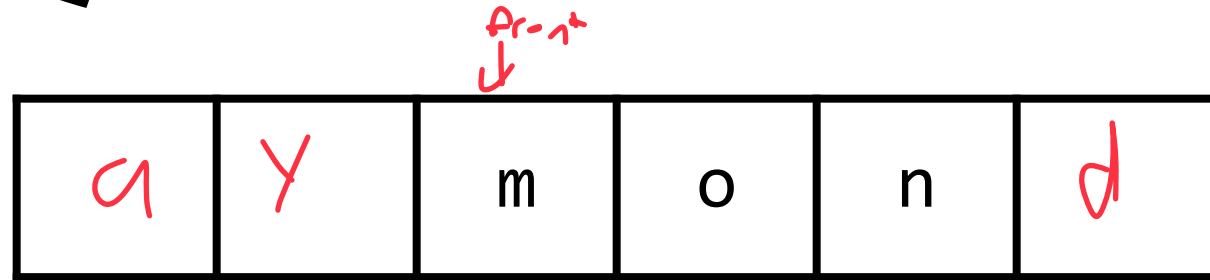
Dequeue(): (as long as size $\neq$ 0)

Remove data at index front

front = **(front+1) % capacity**

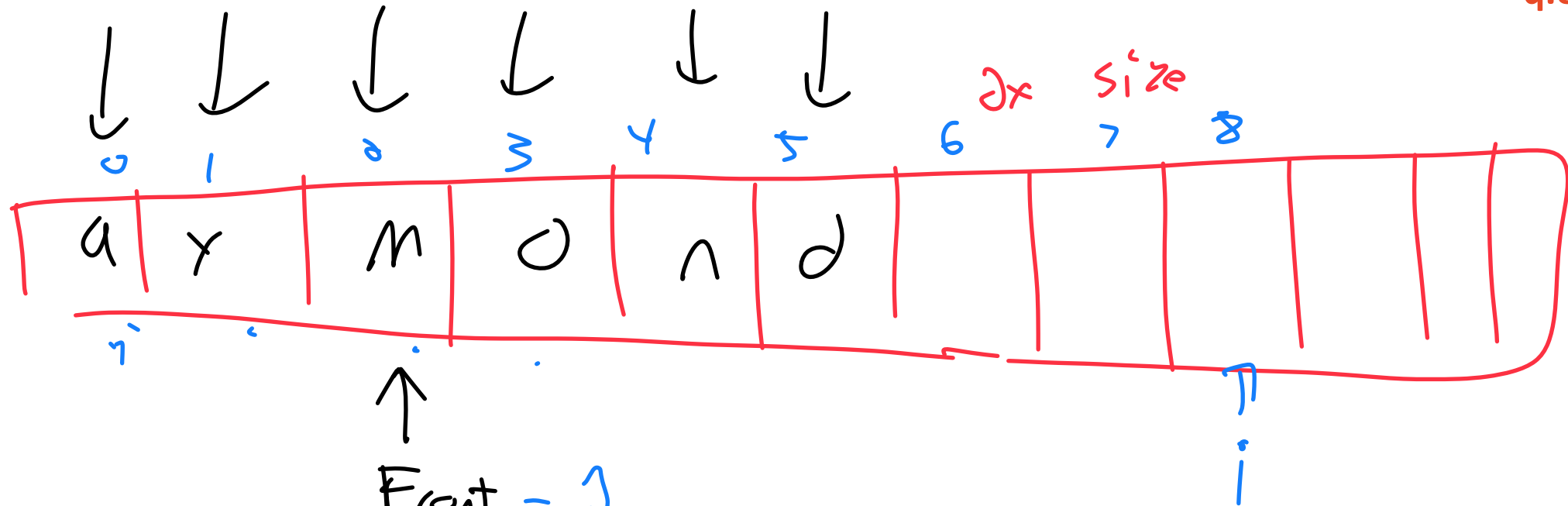
size--

Queue Data Structure: Resizing



capacity = 6

```
Queue<char> q;  
...  
q.enqueue(d);  
q.enqueue(a);  
q.enqueue(y);  
→ q.enqueue(i);  
q.enqueue(s);
```



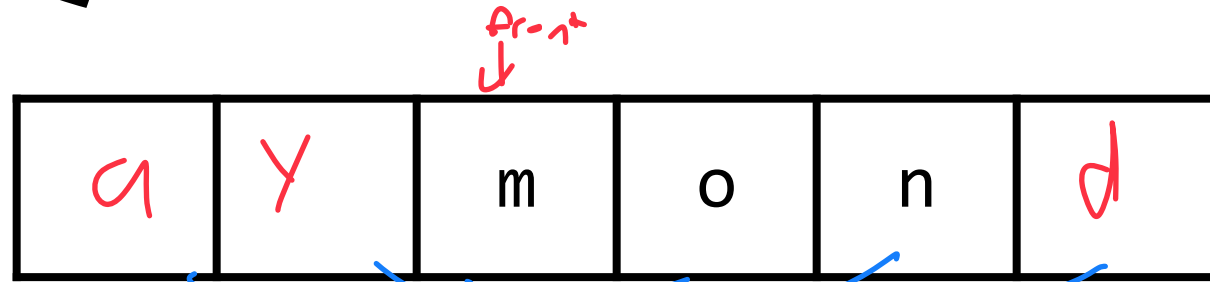
Front = 2

size = 6

capacity = 12

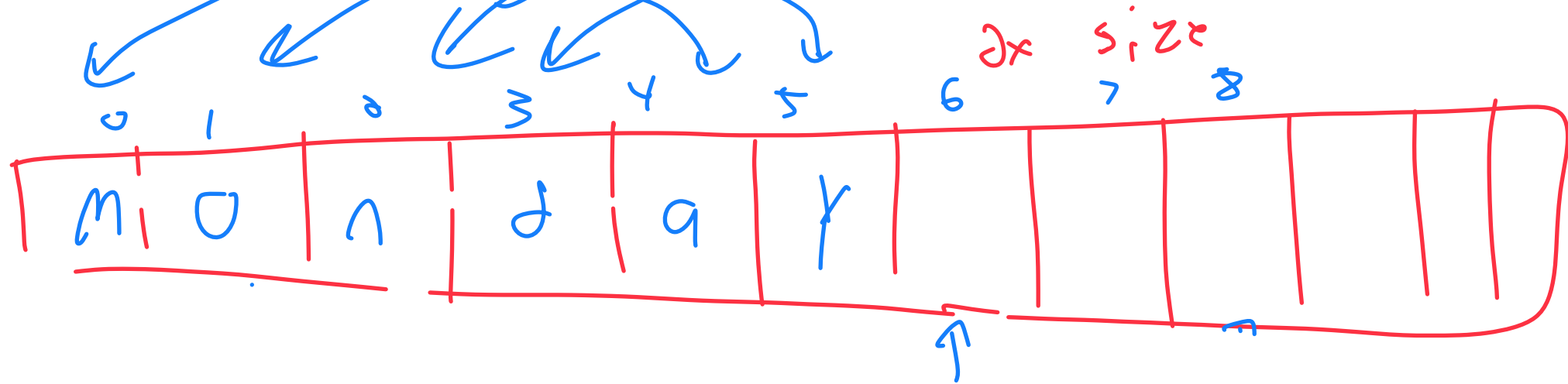
Oh no! Many problems!

Queue Data Structure: Resizing



capacity = 6

```
Queue<char> q;  
...  
q.enqueue(d);  
q.enqueue(a);  
q.enqueue(y);  
→ q.enqueue(i);  
q.enqueue(s);
```

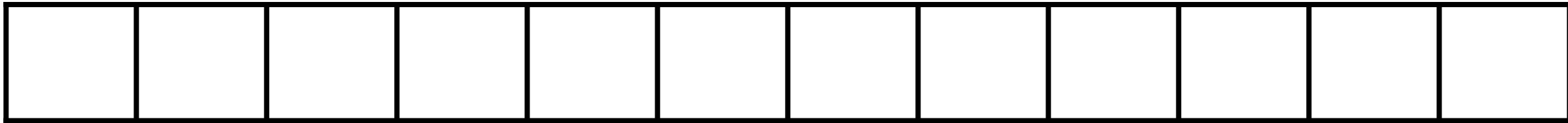
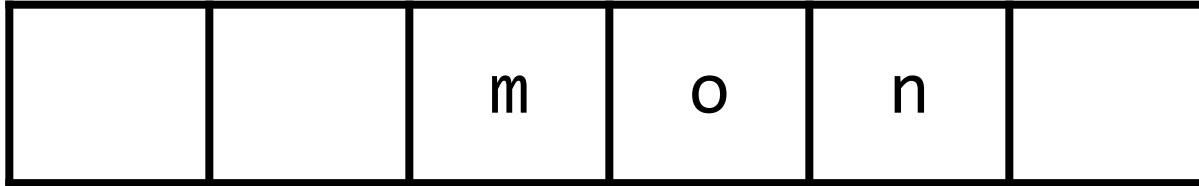


capacity = 12

Front = 0
size = 6

$O(n)$ to resize
(same as array)
 $O(1)^*$ amortized

Queue Data Structure: Resizing



```
Queue<char> q;
```

```
...
```

```
q.enqueue(d);
```

```
q.enqueue(a);
```

```
q.enqueue(y);
```

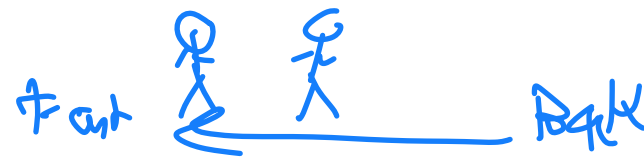
```
q.enqueue(i);
```

```
q.enqueue(s);
```

Queue ADT



- [Order]: First in - First out (FIFO)



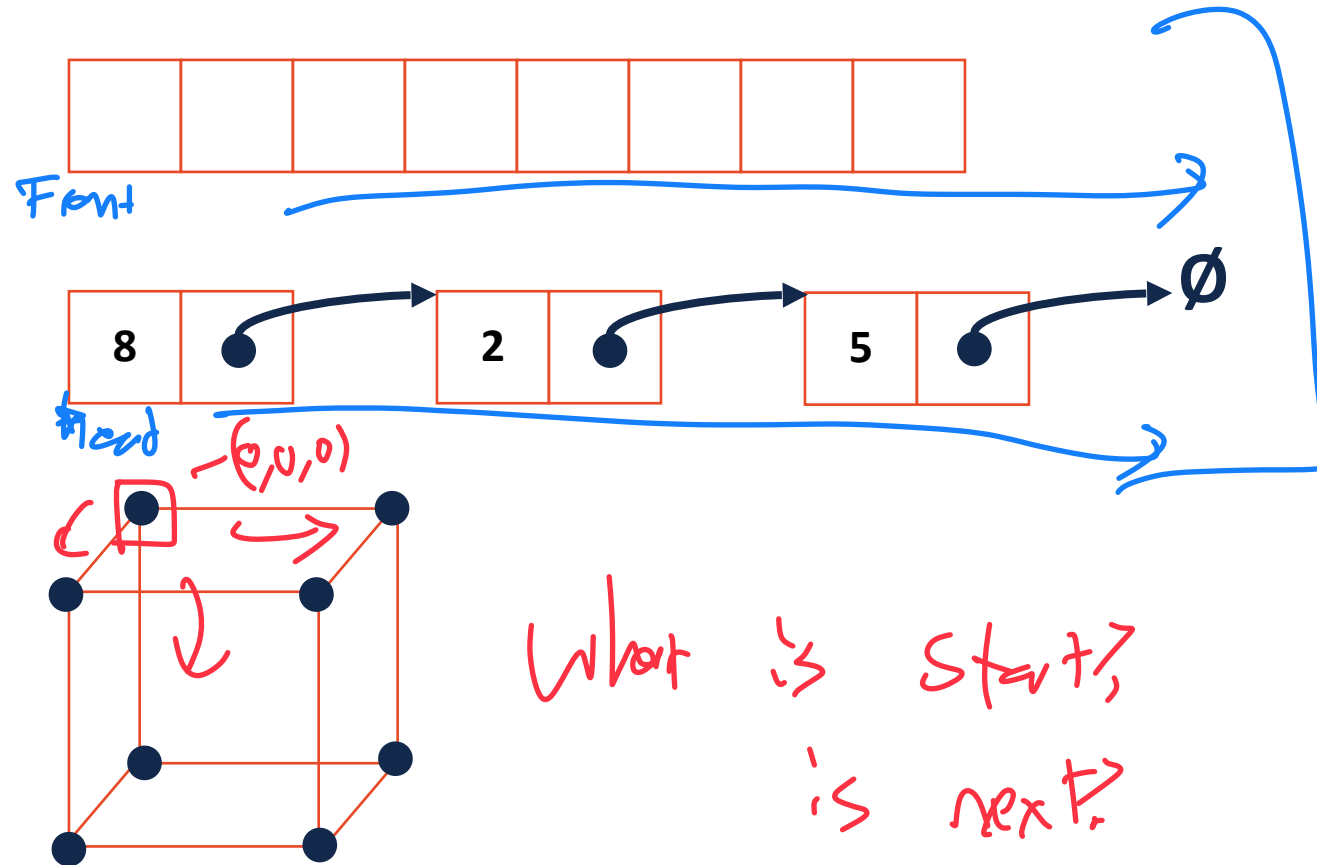
- [Implementation]: trivial as linked list w/ tail pointer

Array as circular queue

- [Runtime]: $O(1)$ / $O(1)^*$ if array resize

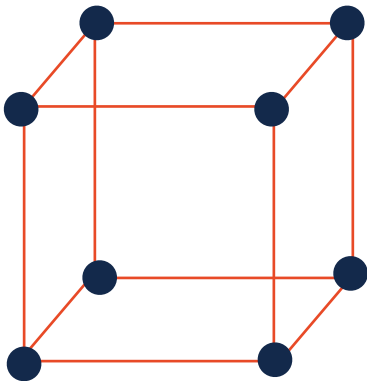
Iterators

We want to be able to loop through all elements for any underlying implementation in a systematic way



Iterators

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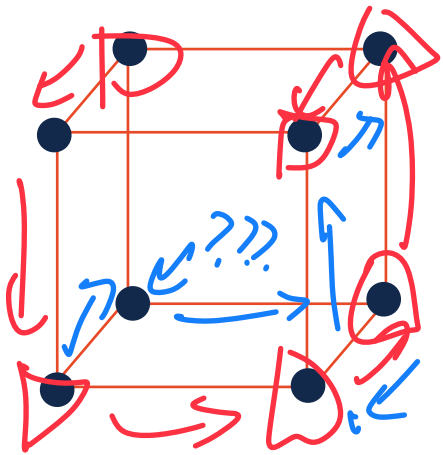


Cur. Location	Cur. Data	Next
ListNode * curr	<i>(curr -> data)</i>	<i>(curr -> next)</i>
unsigned index	<i>A[i]</i>	<i>A[i+1]</i>
Some form of (x, y, z)	<i>???</i>	<i>???</i>

Iterators

Our target interface

Iterators provide a way to access items in a container without exposing the underlying structure of the container



```
1 Cube::Iterator it = myCube.begin();  
2  
3 while (it != myCube.end()) {  
4     std::cout << *it << " ";  
5     it++;  
6 }  
7
```

first item

last (one past end)

get next

get current "data"

Iterators

(like class (as example))

For a class to implement an iterator, it needs two functions:


Iterator begin()

Must return type Iterator

Iterator end()

Iterators

For a class to implement an iterator, it needs two functions:

Iterator begin()

Returns an Iterator object pointing at the 'first item'

Iterator end()

Returns an Iterator object pointing one entry past end of dataset

Iterators

The actual iterator is defined as a class **inside** the outer class:

1. It must be of base class **std::iterator**

2. It must implement at least the following operations:

Iterator& operator ++() "next"

const T & operator *() "dereference"

bool operator !=(const Iterator &) "not equals"

*



Iterators

Here is a (truncated) example of an iterator:

```
1 template <class T>
2 class List {
3
4     class ListIterator : public
5     std::iterator<std::bidirectional_iterator_tag, T> {
6     public:
7
8         ListIterator& operator++();
9
10        ListIterator& operator--();
11
12        bool operator!=(const ListIterator& rhs);
13
14        const T& operator*(); deref
15    };
16
17    ListIterator begin() const;
18
19    ListIterator end() const;
20};
```

next
→
prev
←
B'directional

≠

Specific
iterator
is written
by me

MP_List Iterator

```
1 class ListIterator {  
2     private:  
3         // @TODO: graded in mp_lists part 1  
4         ListNode* position_;  
5     public:  
6         ...  
7  
8         ListIterator() : position_(NULL) { }  
9         ListIterator(ListNode* x) : position_(x) { }
```



You write implementation for List in iterator

```

1  #include <list>
2  #include <string>
3  #include <iostream>
4
5  struct Animal {
6      std::string name, food;
7      bool big;
8      Animal(std::string name = "blob", std::string food = "you", bool big = true) :
9          name(name), food(food), big(big) { /* nothing */ }
10 };
11
12 int main() {
13     Animal g("giraffe", "leaves", true), p("penguin", "fish", false), b("bear");
14     std::vector<Animal> zoo;
15
16     zoo.push_back(g);
17     zoo.push_back(p);    // std::vector's insertAtEnd
18     zoo.push_back(b);
19
20     for ( std::vector<Animal>::iterator it = zoo.begin(); it != zoo.end(); ++it ) {
21         std::cout << (*it).name << " " << (*it).food << std::endl;
22     }
23
24     return 0;
25 }

```

! make you define



call deref operator

You write what this is



```
1
2 std::vector<Animal> zoo;
3
4
5 /* Full text snippet */
6
7 for ( std::vector<Animal>::iterator it = zoo.begin(); it != zoo.end(); ++it ) {
8     std::cout << (*it).name << " " << (*it).food << std::endl;
9 }
10
11
12 /* Auto Snippet */
13
14 for ( auto it = zoo.begin(); it != zoo.end(); ++it ) {
15     std::cout << (*it).name << " " << (*it).food << std::endl;
16 }
17
18 /* For Each Snippet */ For each animal in zoo
19
20 for ( const Animal & animal : zoo ) {
21     std::cout << animal.name << " " << animal.food << std::endl;
22 }
23
24
25
```

Exam 1 Review

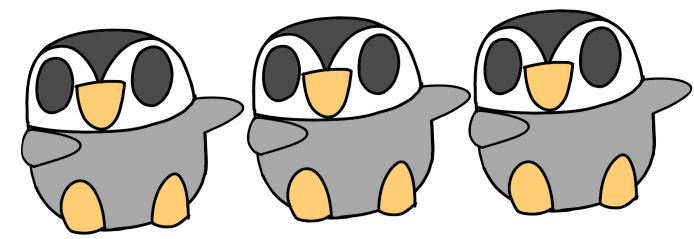
Exam content focuses on fundamentals and lists

To prepare for MCQ:

- Have a clear understanding of the list ADT
- Understand the Big O details behind linked list and array list
- Be able to apply this knowledge in new contexts

e.g. 'What if I had a linked list and made this modification...'

Lists



The not-so-secret underlying implementation for many things

	Singly Linked List	Array
Look up arbitrary location	$O(n)$	$O(1)$
Insert after given element	$O(1)$	$O(n)$
Remove after given element	$O(1)$	$O(n)$
Insert at arbitrary location	$O(n)$	$O(n)$
Remove at arbitrary location	$O(n)$	$O(n)$
Search for an input value	$O(n)$	$O(n)$

Special Cases:

insertFront

insertBack (not full)

Exam 1 Review

To prepare for the free response question:

- Have a clear understanding of the list ADT
- Understand the Big O details behind linked list and array list
- **Be able to apply this knowledge in new contexts**

Brainstorm: What are some variations on linked lists / arrays?

FRQ Example Rubric

Points	Rubric Item	Description
2	Overall Big O is correct	The answer has the best correct Big O for the overall problem
3	Minimum necessary steps described	The explanation describes all steps needed to solve the question (even if the Big O of one or more steps is wrong).
1	Individual step Big O is correct	The explanation correctly describes the Big O of all steps. (To receive points, you must receive full credit for the minimum necessary steps)
2	Explanation of how the 'twist' affects Big O	Each problem has at least one adjustment beyond the core data structure. Your answer should explicitly state how these changed parameters / structural requirements affected (or did not affect) the runtime.
1	Answer is given in complete sentences	Even when practicing, write in full complete sentences.
1	Answer has appropriate structure	Start with a clear declaration of Big O. On a new line state the steps and their Big Os explicitly. Then start a new paragraph where you explain why. Make it easy for your graders

Exam 1 Review

Preparing for the coding question:

- Make sure you understand the fundamentals of debugging
- Review your coding assignments so far (debug / stickers)
- The coding question will be similar to that of practice exam

Brainstorm: What are some ways we could test your knowledge of PNGs / HSLAPixels in the context of arrays?



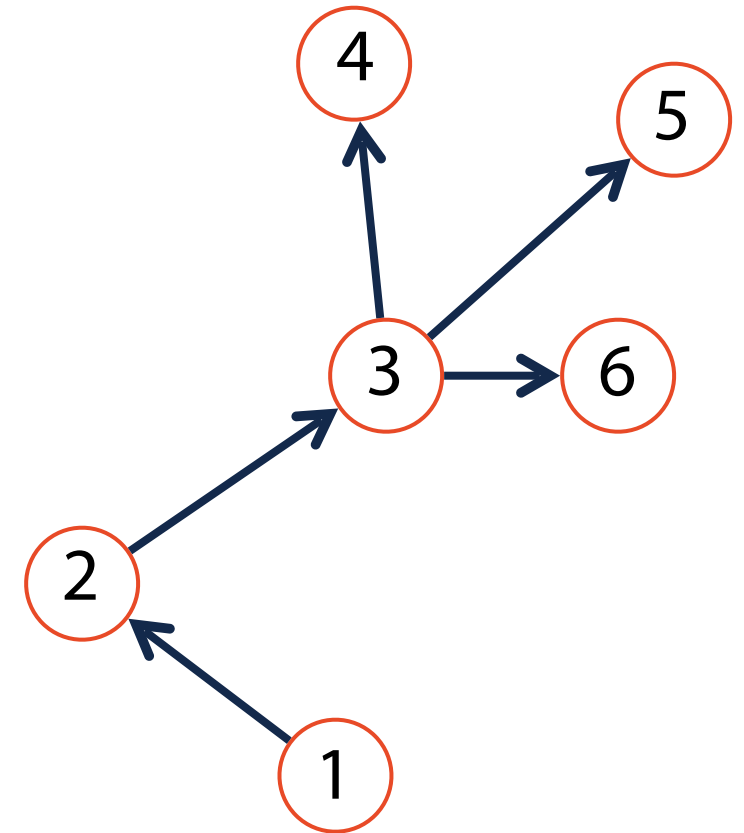
Questions?

Trees

A non-linear data structure defined recursively as a collection of nodes where each node contains a value and zero or more connected nodes.

[In CS 225] a tree is also:

- 1) Acyclic — No path from node to itself
- 2) Rooted — A specific node is labeled root

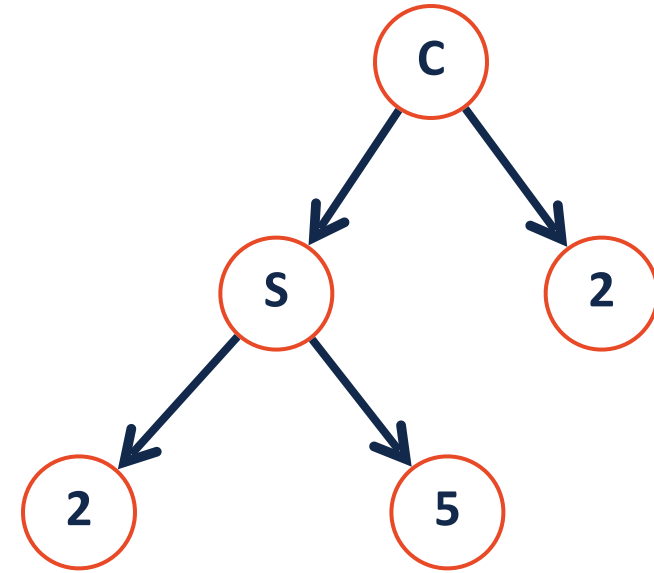


Binary Tree

A **binary tree** is a tree T such that:

1. $T = \emptyset$

2. $T = (data, T_L, T_R)$

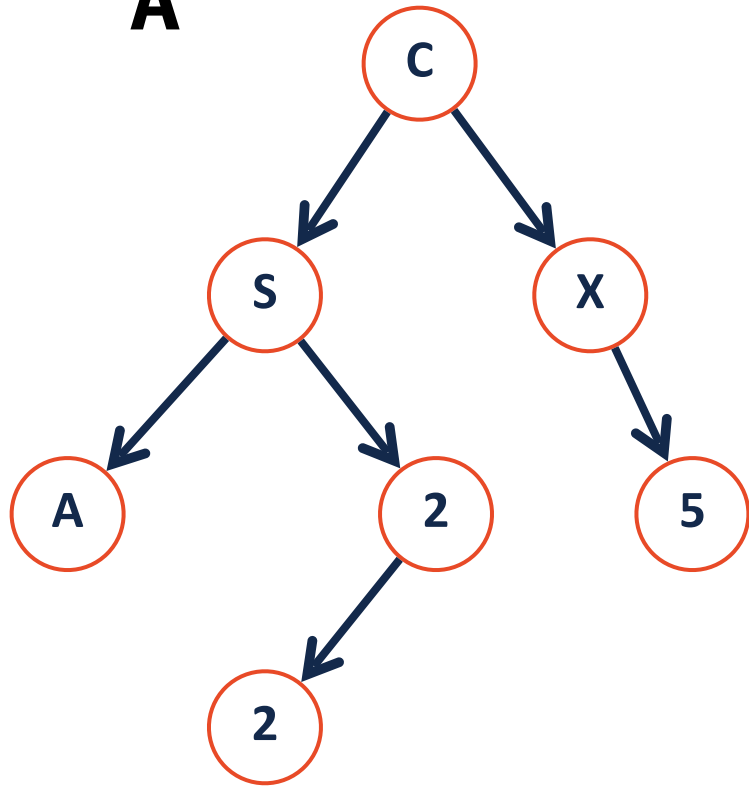


Which of the following are binary trees?

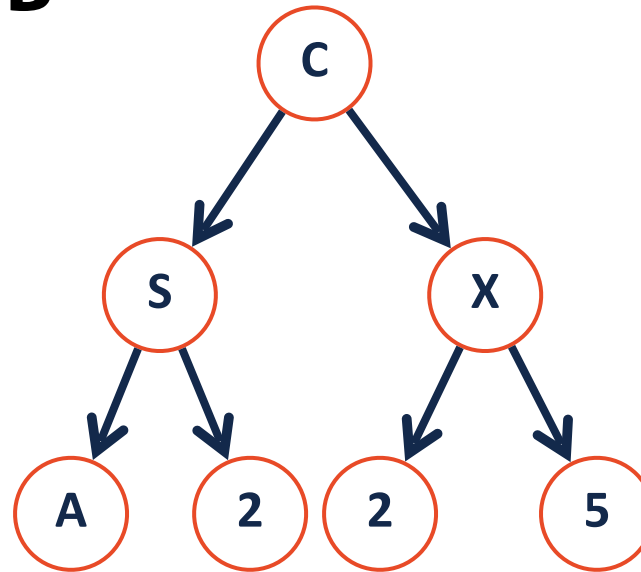


Join Code: 225

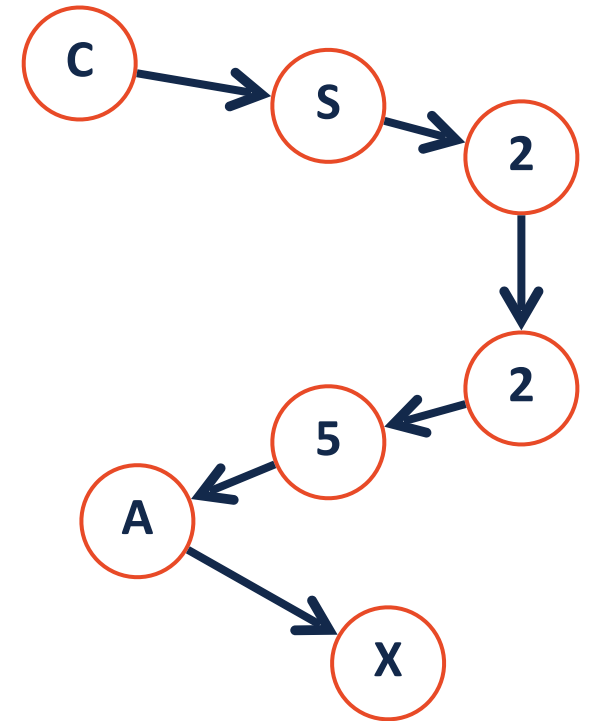
A



B



C





Tree ADT