

Data Structures

Tree Traversal

CS 225

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ILLINOIS
URBANA - CHAMPAIGN

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Extra Credit Reminder

MP submission on PL has two separate submissions

The extra credit portion will only test part 1

Completion of the extra credit portion by the following Monday is worth 4 points

MP_stickers feedback form out now!

Anonymous feedback form worth 2 points **collectively**

Learning Objectives

Review and expand on foundational tree terminology

Discuss the tree ADT

Explore tree implementation details

Binary Tree

Lets define additional terminology for different **types** of binary trees!

1. Full Tree

2. Perfect Tree

3. Complete Tree

Binary Tree: full

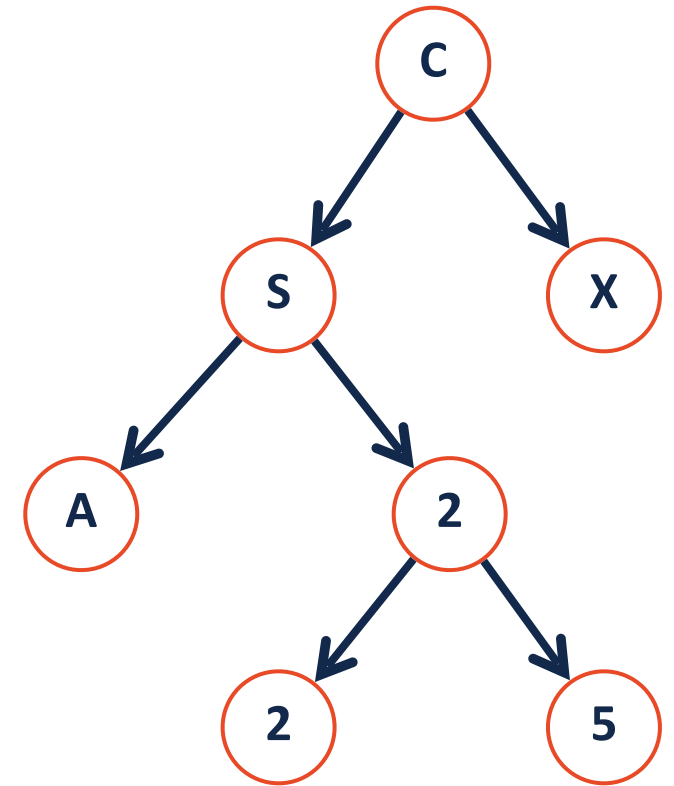
A **full tree** is a binary tree where every node has either 0 or 2 children

A tree **F** is **full** if and only if:

1.

2.

3.



Binary Tree: full

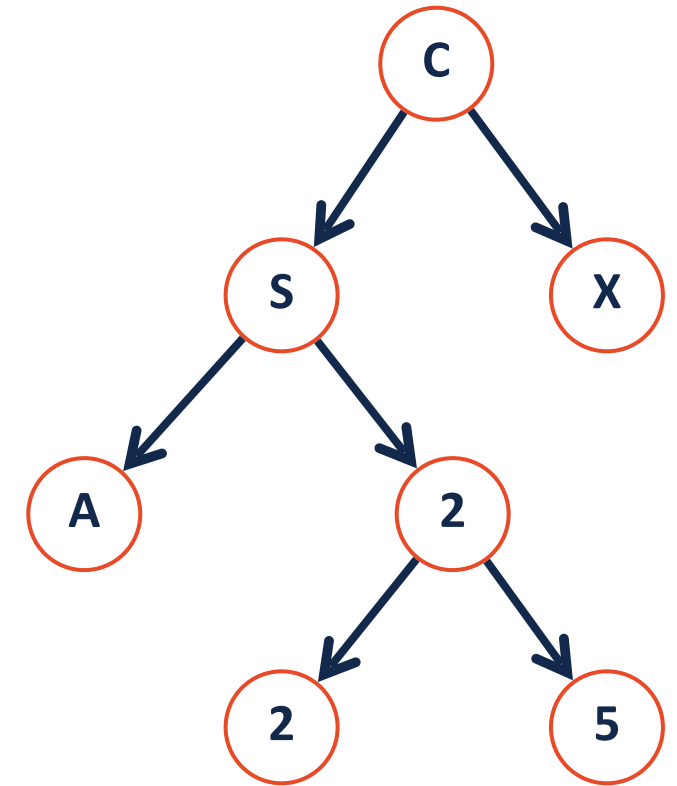
A **full tree** is a binary tree where every node has either 0 or 2 children

A tree **F** is **full** if and only if:

1. $F = \emptyset$

2. $F = (data, \emptyset, \emptyset)$

3. $F = (data, F_l \neq \emptyset, F_r \neq \emptyset)$



Full binary tree : Size

- Question - Which of the following are possible sizes (# of nodes) for a full binary tree?

a) 2 b) 5 c) 7 d) 8

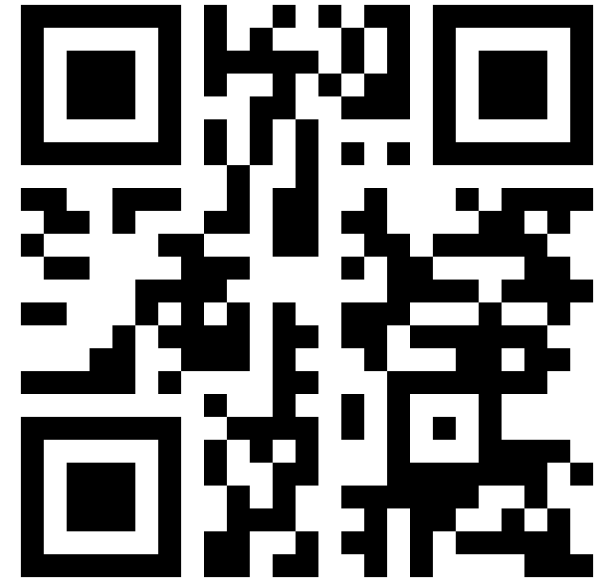
A) 2 and 8

B) 5 and 7

C) 7 only

D) All of these

E) None of these



Join Code : 225

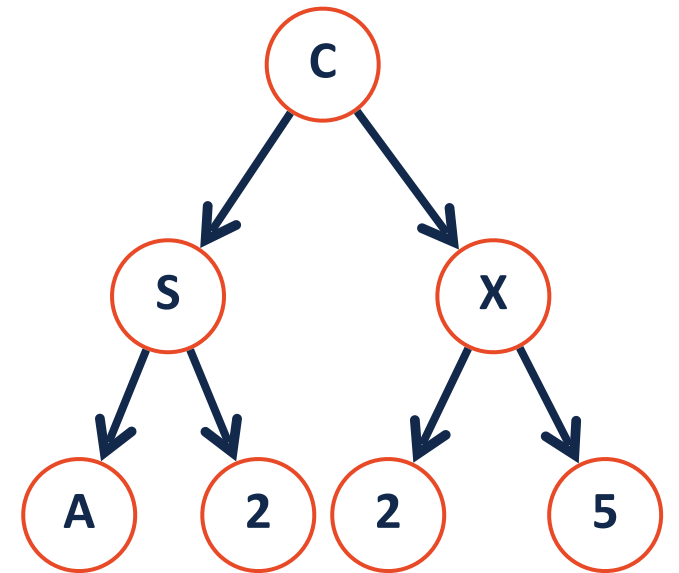
Binary Tree: perfect

A **perfect tree** is a binary tree where...
Every internal node has 2 children and all leaves are at the same level.

A tree **P** is **perfect** if and only if:

1.

2.



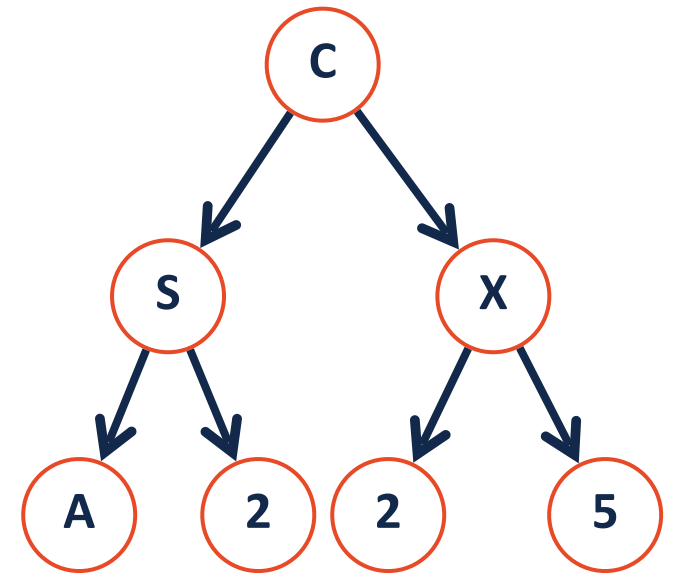
Binary Tree: perfect

A **perfect tree** is a binary tree where...
Every internal node has 2 children and all leaves are at the same level.

A tree **P** is **perfect** if and only if:

$$1. P_h = (data, P_{h-1}, P_{h-1})$$

$$2. P_0 = (data, \emptyset, \emptyset) \equiv P_{-1} = \emptyset$$



Perfect binary tree : Size

- Question - Which of the following are possible sizes (# of nodes) for a perfect binary tree?

a) 2 b) 5 c) 7 d) 8

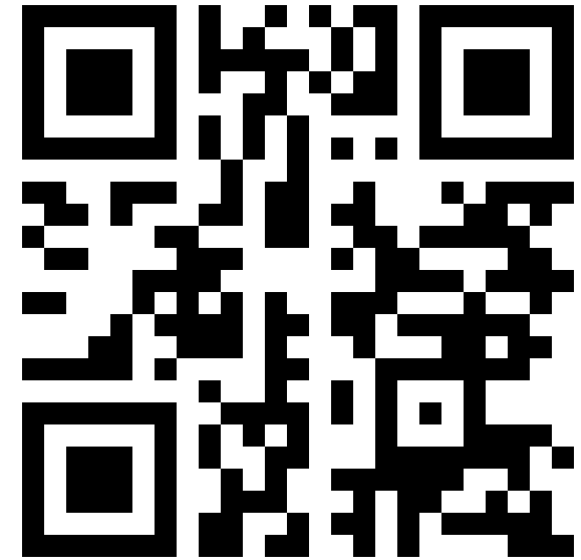
A) 2 and 8

B) 5 and 7

C) 7 only

D) All of these

E) None of these



Join code : 225

Binary Tree: complete

A **complete tree** is a B.T. where...

All levels except the last are completely filled.

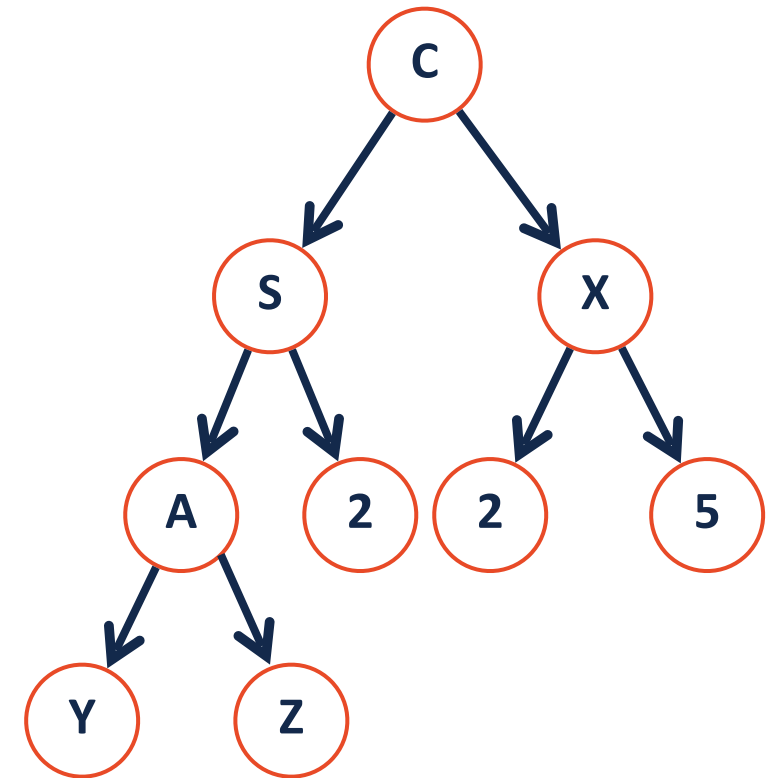
The last level contains at least one node (and is pushed to left)

A tree **C** is **complete** if and only if:

1.

2.

3.



Binary Tree: complete

A **complete tree** is a B.T. where...

All levels except the last are completely filled.

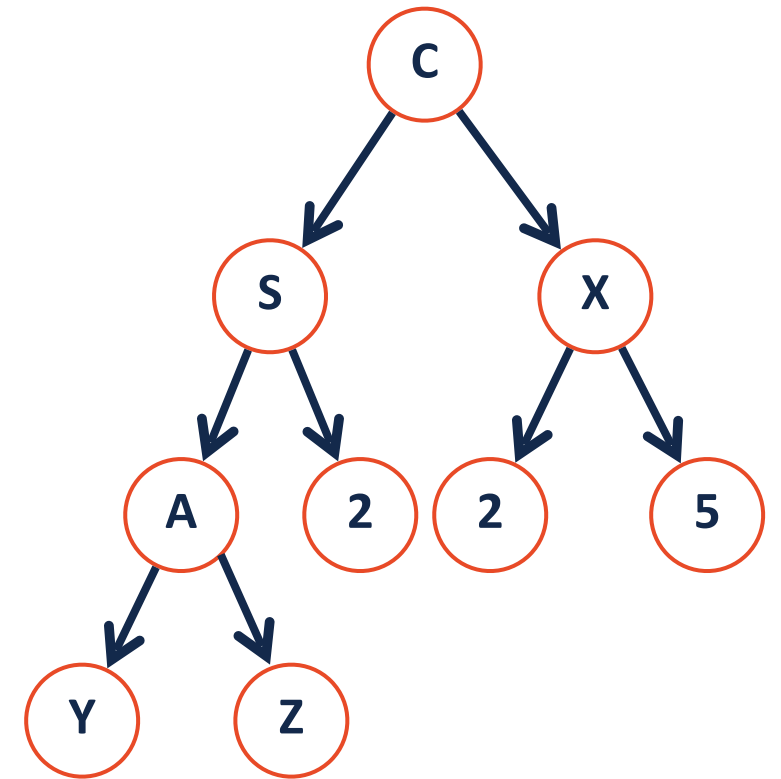
The last level contains at least one node (and is pushed to left)

A tree **C** is **complete** if and only if:

1. $C_h = (data, C_{h-1}, P_{h-2})$

2. $C_h = (data, P_{h-1}, C_{h-1})$

3. $C_{-1} = \emptyset$



Binary Tree: complete

A **complete tree** is a B.T. where...

All levels except the last are completely filled.

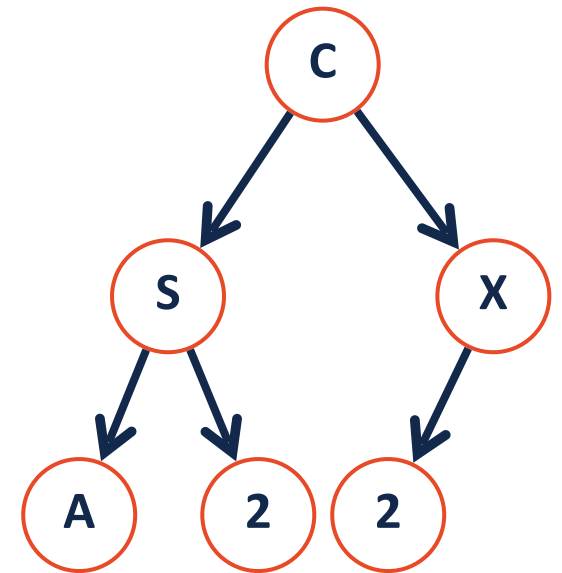
The last level contains at least one node (and is pushed to left)

A tree **C** is **complete** if and only if:

$$1. C_h = (data, C_{h-1}, P_{h-2})$$

$$2. C_h = (data, P_{h-1}, C_{h-1})$$

$$3. C_{-1} = \emptyset$$



Complete binary tree : Size

- Question - Which of the following are possible sizes (# of nodes) for a complete binary tree?

a) 2 b) 5 c) 7 d) 8

A) 2 and 8

B) 5 and 7

C) 7 only

D) All of these

E) None of these



Join code : 225

Binary Tree



Why do we care?

1. Terminology instantly defines a particular tree structure
2. Understanding how to think 'recursively' is very important.

Binary Tree: Thinking with Types

Is every **full** tree **complete**?

Is every **complete** tree **full**?



Tree ADT

Insert

Remove

Traverse

Find

Constructor

List.h

```
1 #pragma once
2
3 template <typename T>
4 class List {
5     public:
6         /* ... */
7     private:
8         class ListNode {
9             T & data;
10
11             ListNode * next;
12
13
14             ListNode(T & data) :
15                 data(data), next(NULL) { }
16         };
17
18         ListNode *head_;
19         /* ... */
20
21 };
```

Tree.h

```
1 #pragma once
2
3 template <typename T>
4 class BinaryTree {
5     public:
6         /* ... */
7     private:
8
9
10
11
12
13
14
15
16
17
18
19 };
20
21     TreeNode *root_;
22     /* ... */
23 };
```

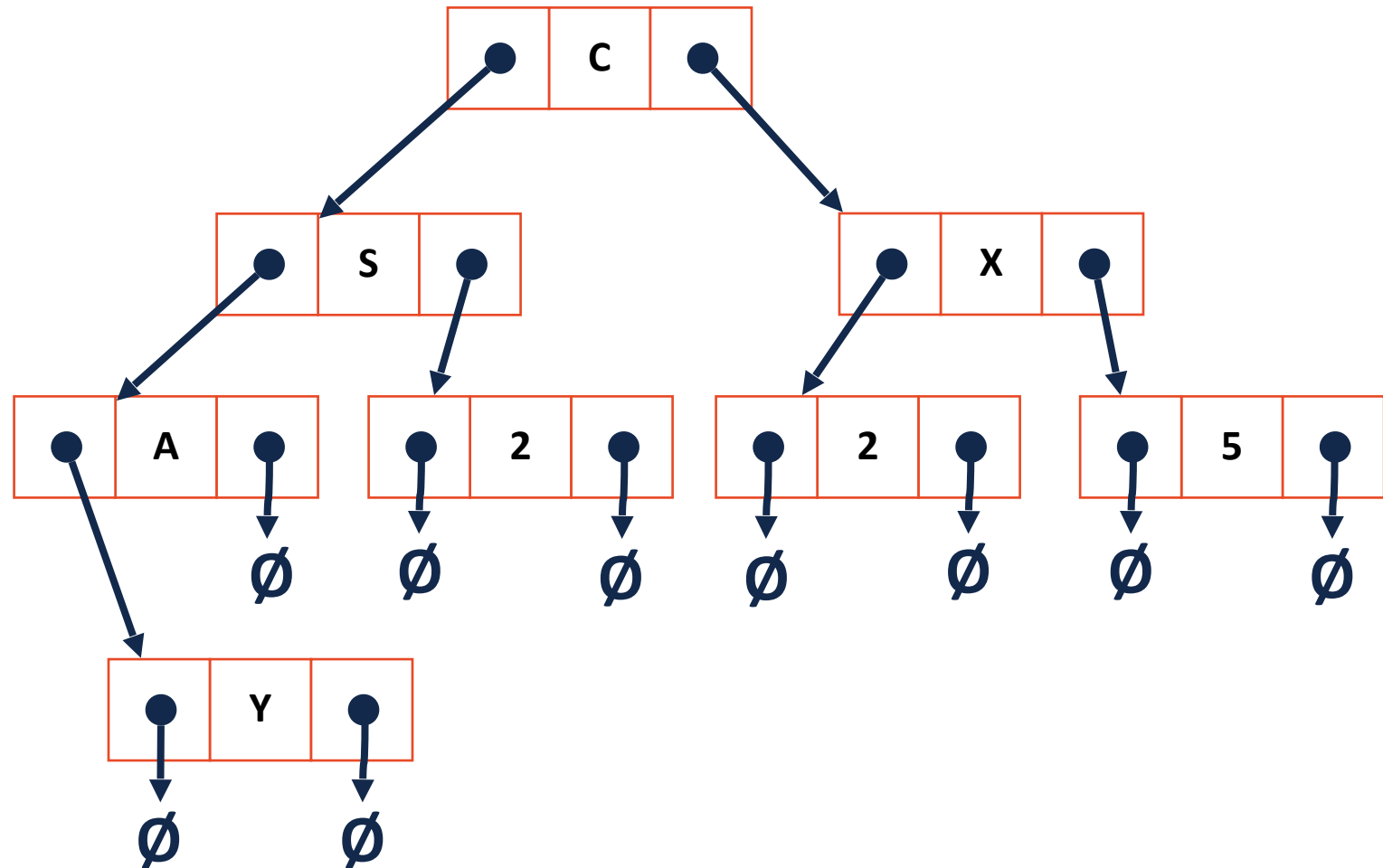
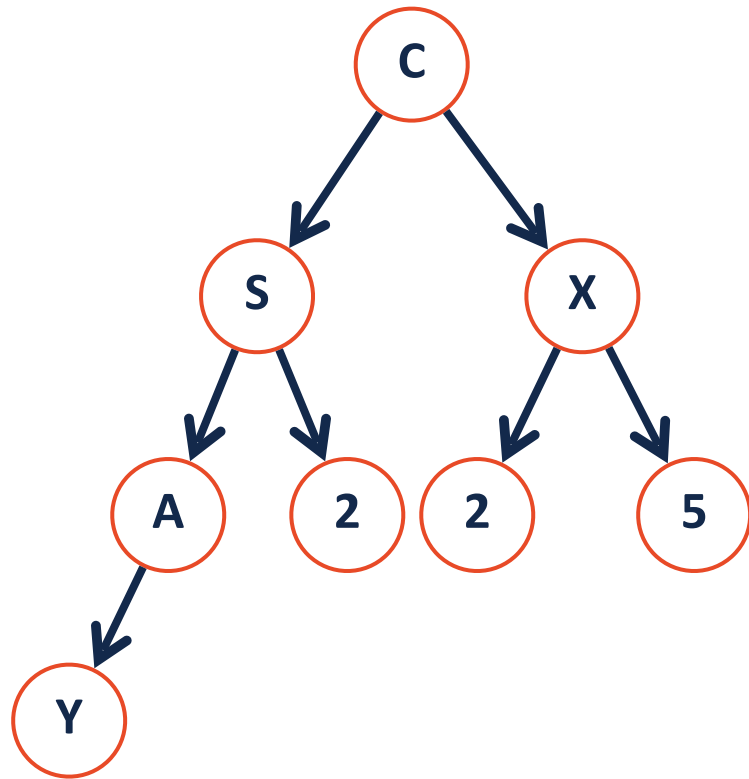
List.h

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11             ListNode * next;
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14             ListNode(T & data) :
15                 data(data), next(NULL) { }
16         };
17
18
19         ListNode *head_;
20         /* ... */
21 };
22
23
```

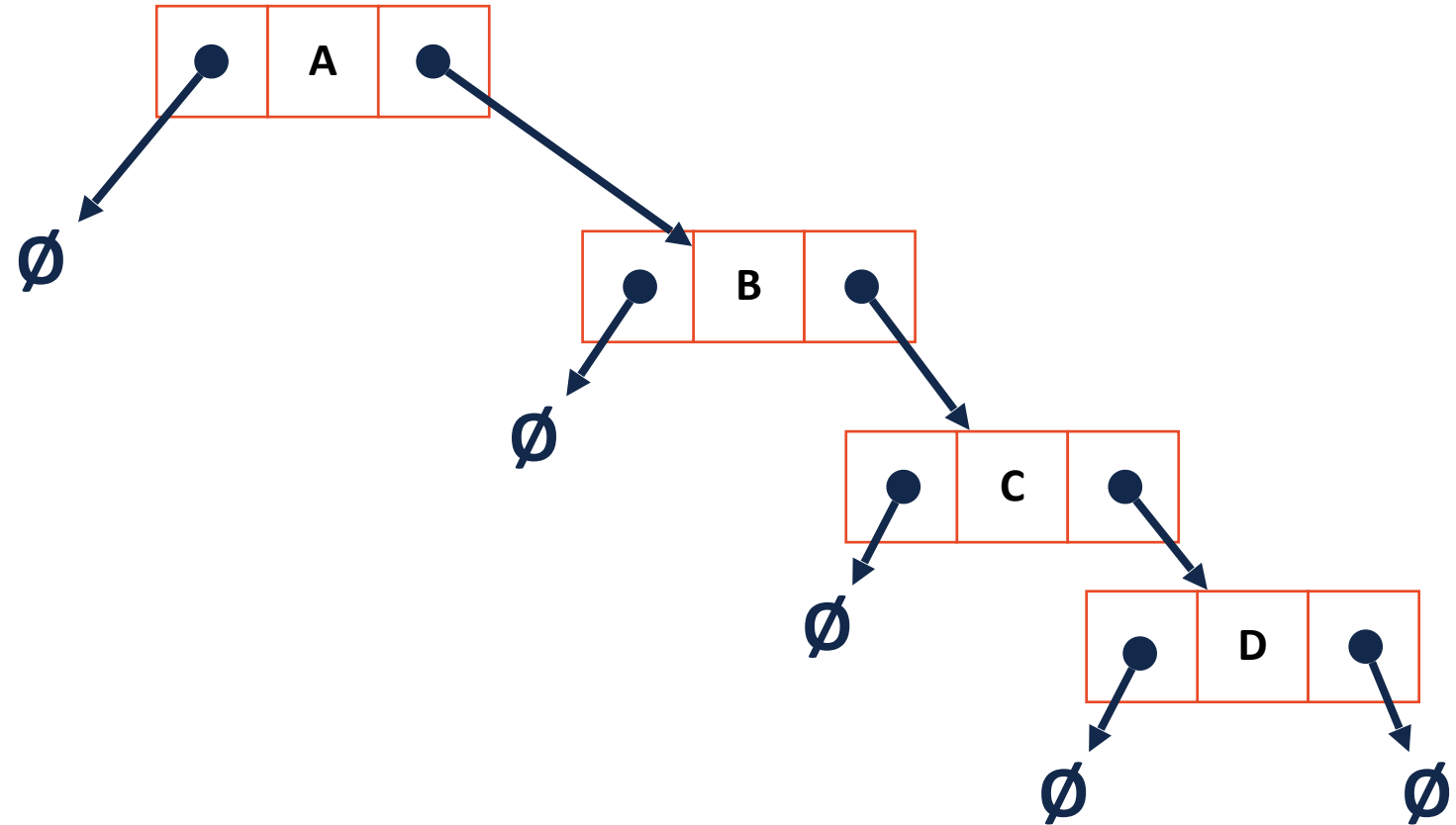
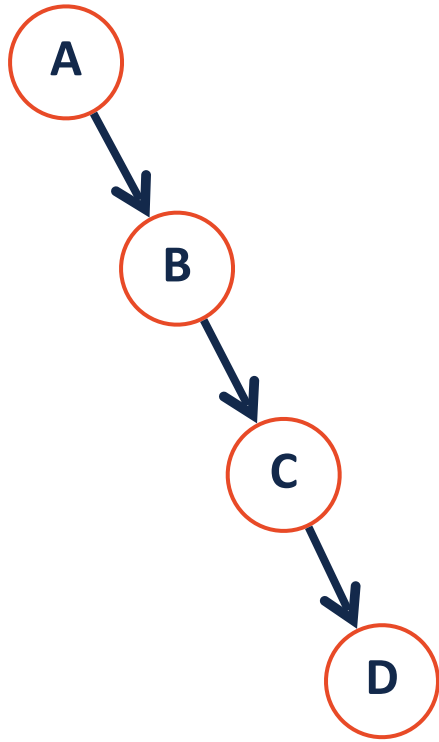
Tree.h

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4 class BinaryTree {
5     public:
6         /* ... */
7     private:
8         class TreeNode {
9             T & data;
10
11             TreeNode * left;
12
13             TreeNode * right;
14
15             TreeNode(T & data) :
16                 data(data), left(NULL),
17                 right(NULL) { }
18
19             };
20
21         TreeNode *root_;
22         /* ... */
23 };
24
```

Visualizing trees

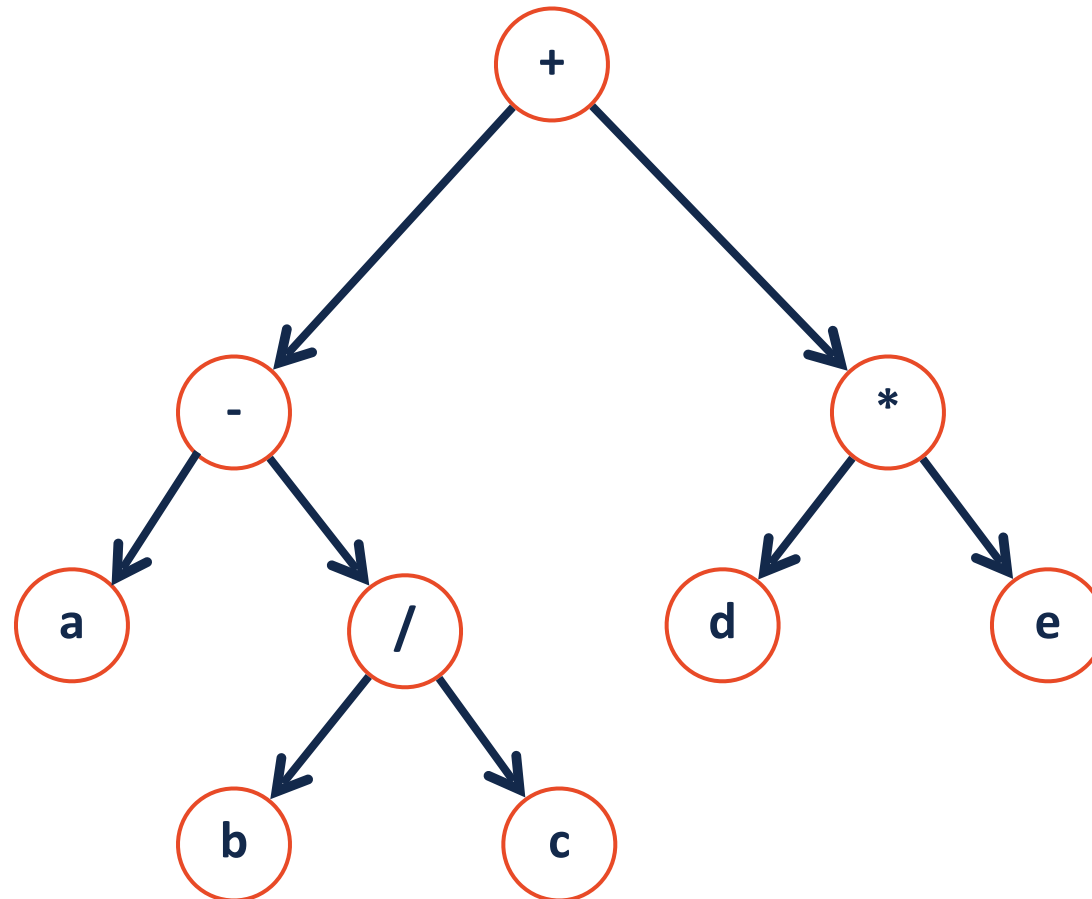


Tree Insert / Remove acts like Linked List

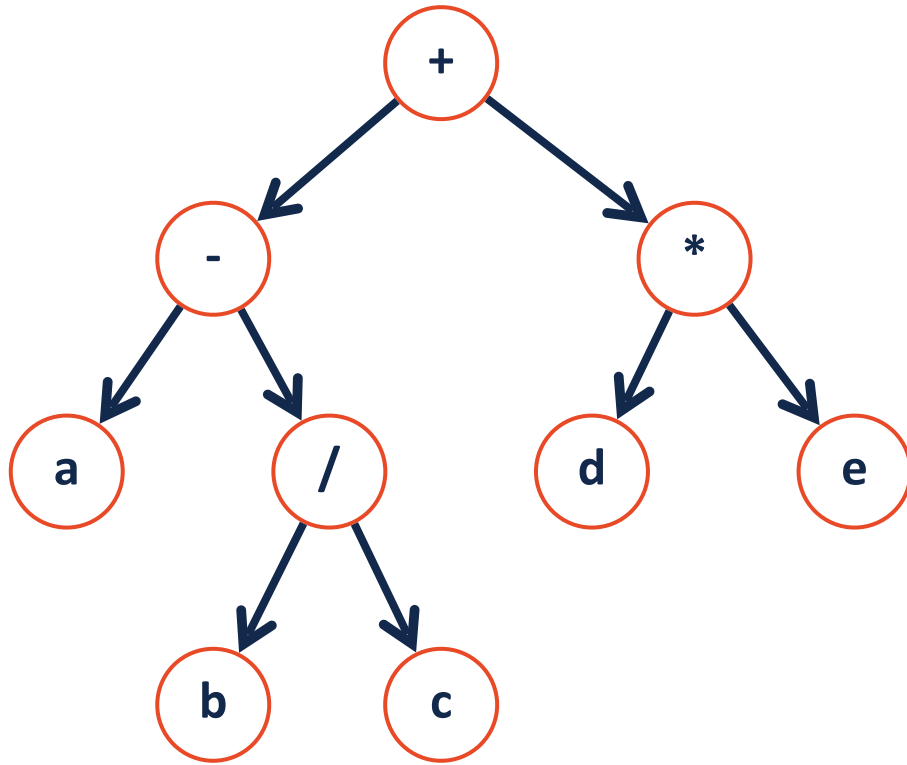


Tree Traversal

A **traversal** of a tree T is an ordered way of visiting every node once.

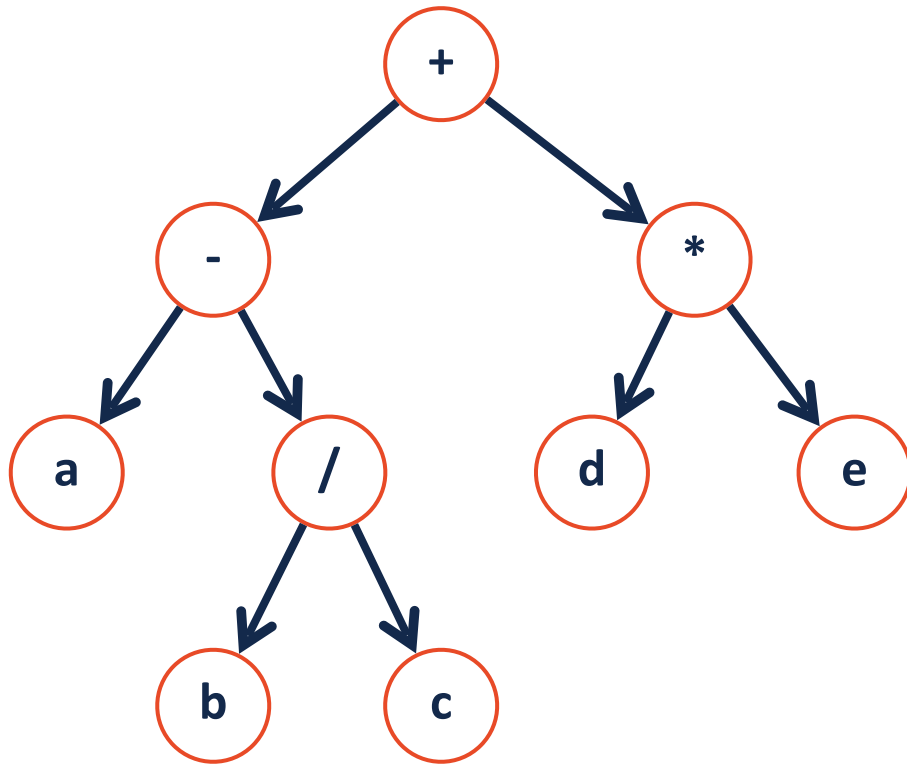


Traversals



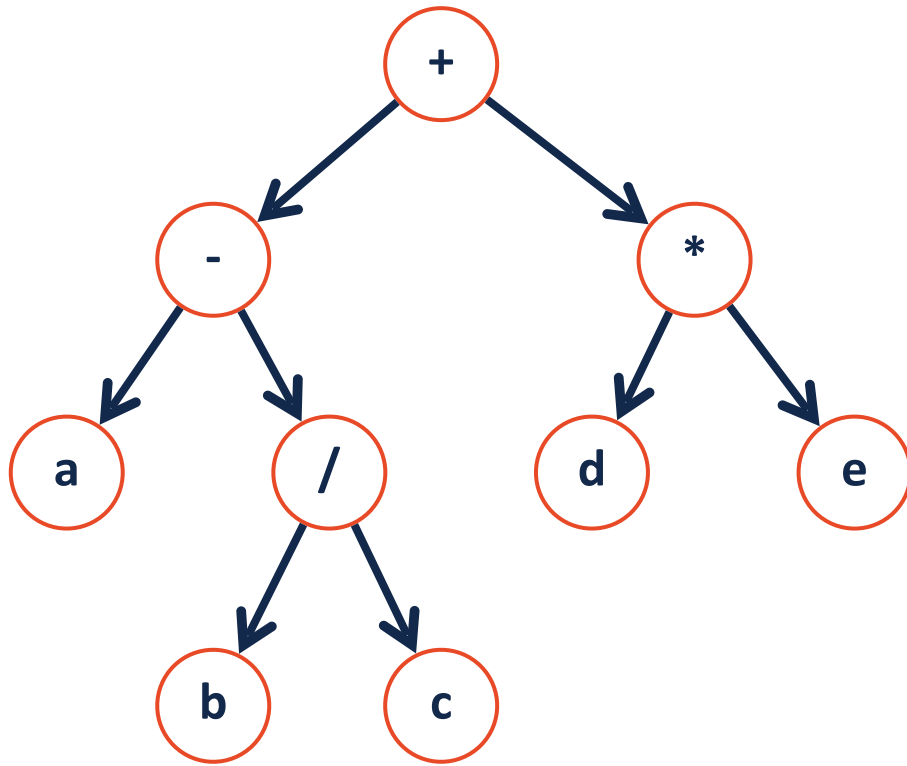
```
1  template<class T>
2  void BinaryTree<T>::____Order(TreeNode * root)
3  {
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21 }
```

Traversals



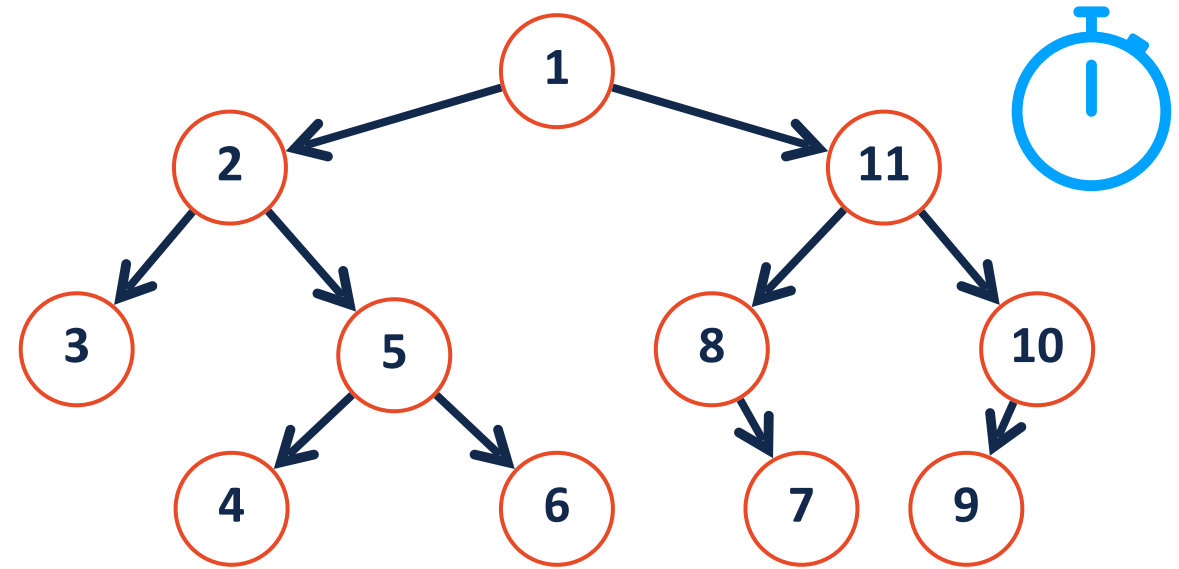
```
1 template<class T>
2 void BinaryTree<T>::____Order(TreeNode * root)
3 {
4
5     if (root) {
6
7         _____;
8
9         ____Order(root->left) ;
10
11         _____;
12
13         ____Order(root->right) ;
14
15         _____;
16
17     }
18
19
20
21 }
```

Traversals



```
1 template<class T>
2 void BinaryTree<T>::____Order(TreeNode * root)
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Tree Traversals

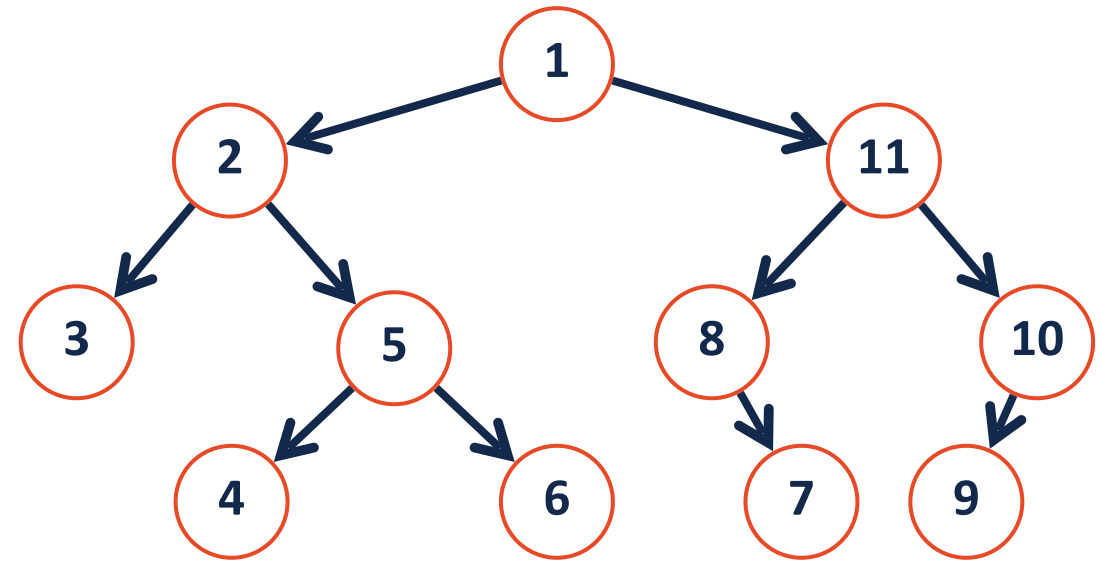


Pre-order:

In-order:

Post-order:

Tree Traversals (Solution)



Pre-order: **1** 2 3 5 4 6 11 8 7 10 9

Tip: Preorder always starts with root!

In-order: **3** 2 4 5 6 **1** 8 7 11 9 10

Tip: Inorder always starts with leftmost node. Root is after all left nodes!

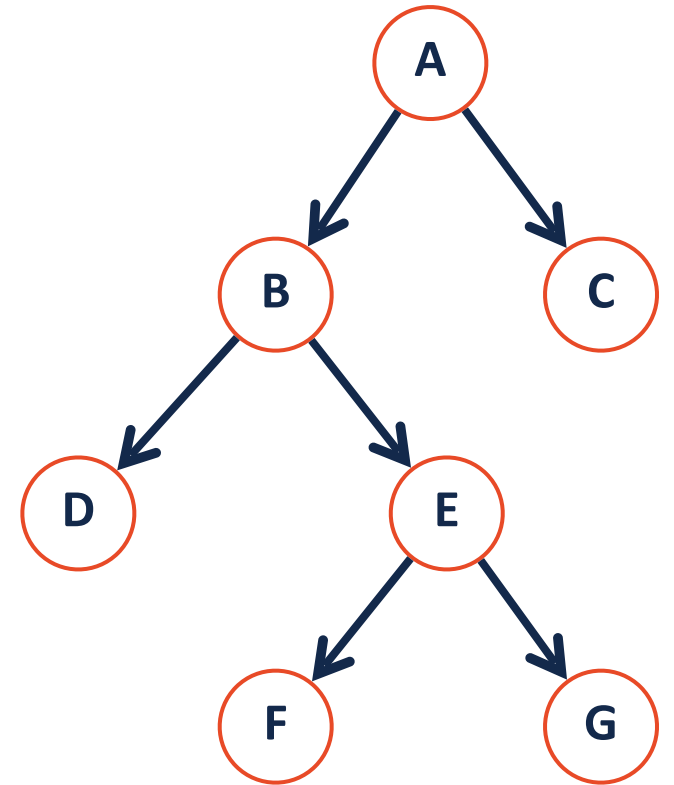
Post-order: **3** 4 6 5 2 7 8 9 10 11 **1**

Tip: Post always starts with leftmost node. Root is always LAST node!

Traversal vs Search

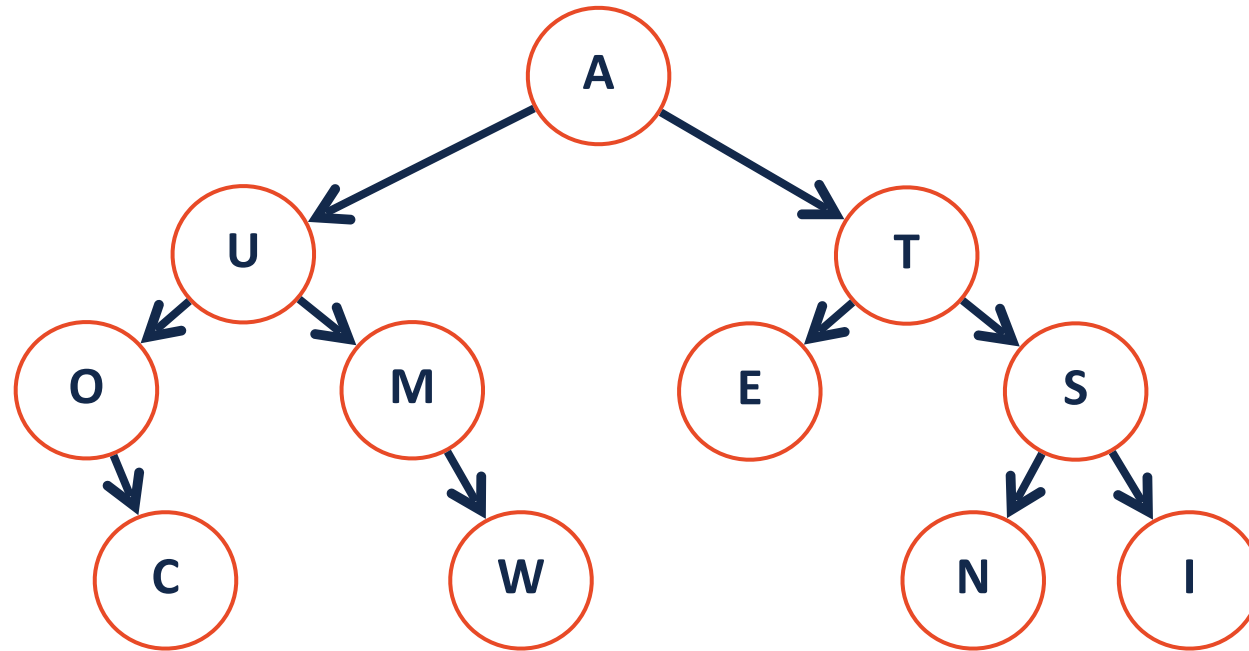
Traversal

Search



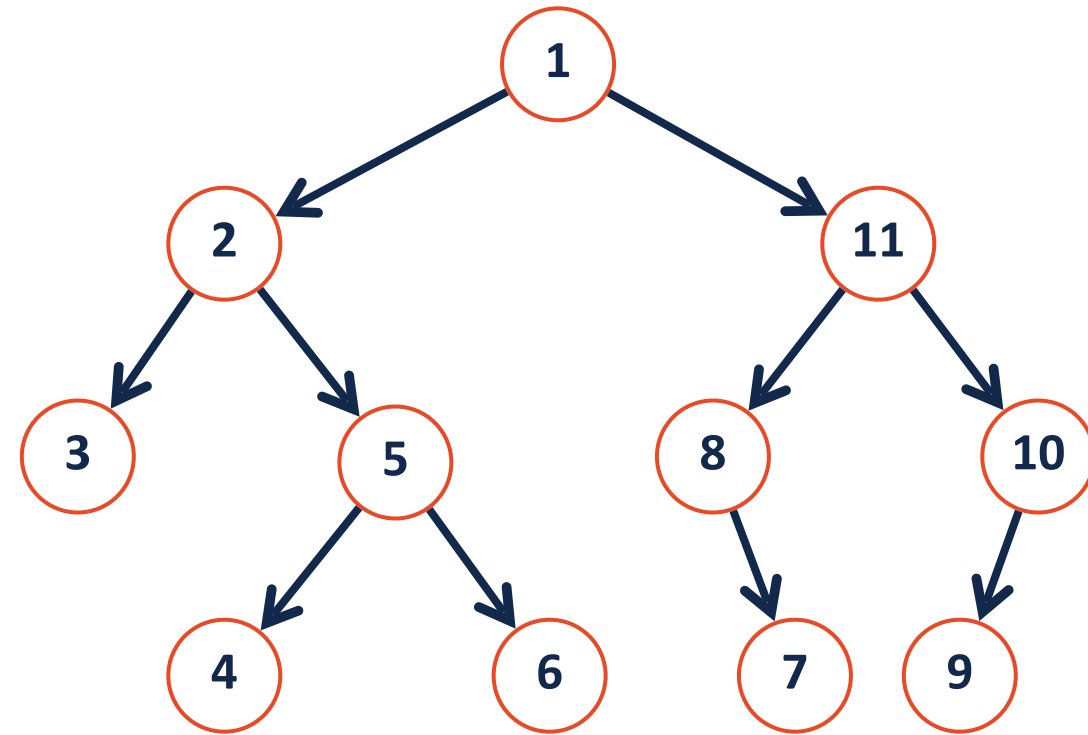
Tree Search

There are two main approaches to searching a binary tree:



Depth First Search

Explore as far along one path as possible before backtracking



Depth First Search

Explore as far along one path as possible before backtracking

Make a stack (initialized with root)

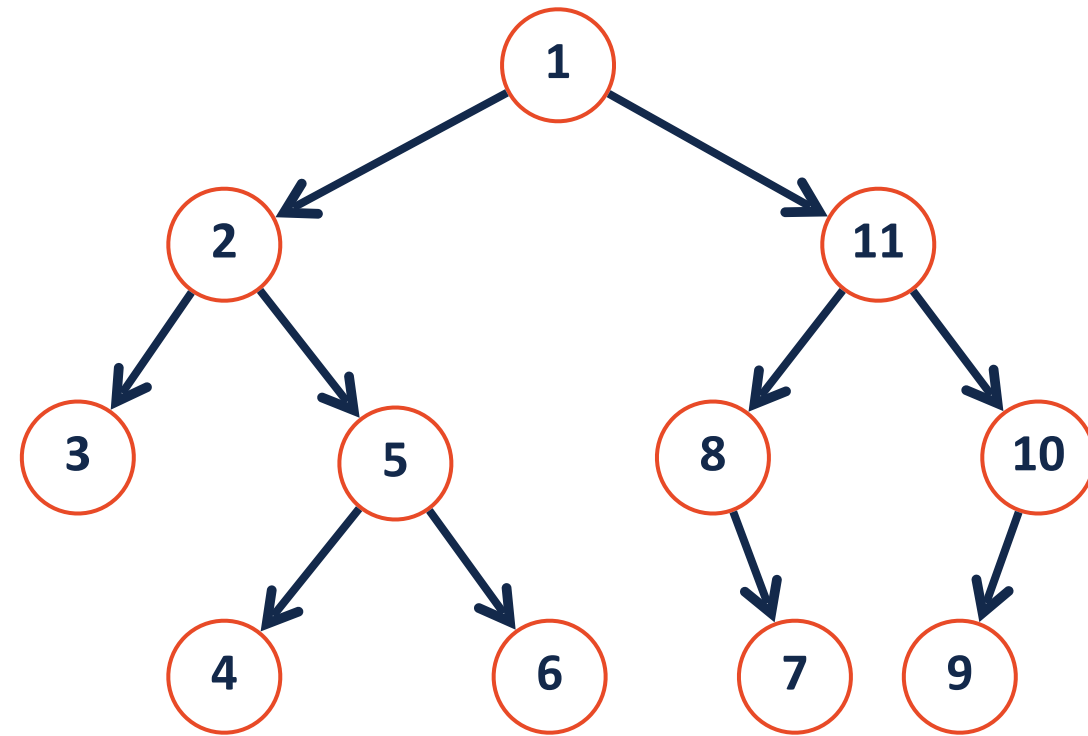
While stack not empty:

```
tmp = stack.pop()
```

```
print(tmp)
```

```
stack.push(tmp->right)
```

```
stack.push(tmp->left)
```

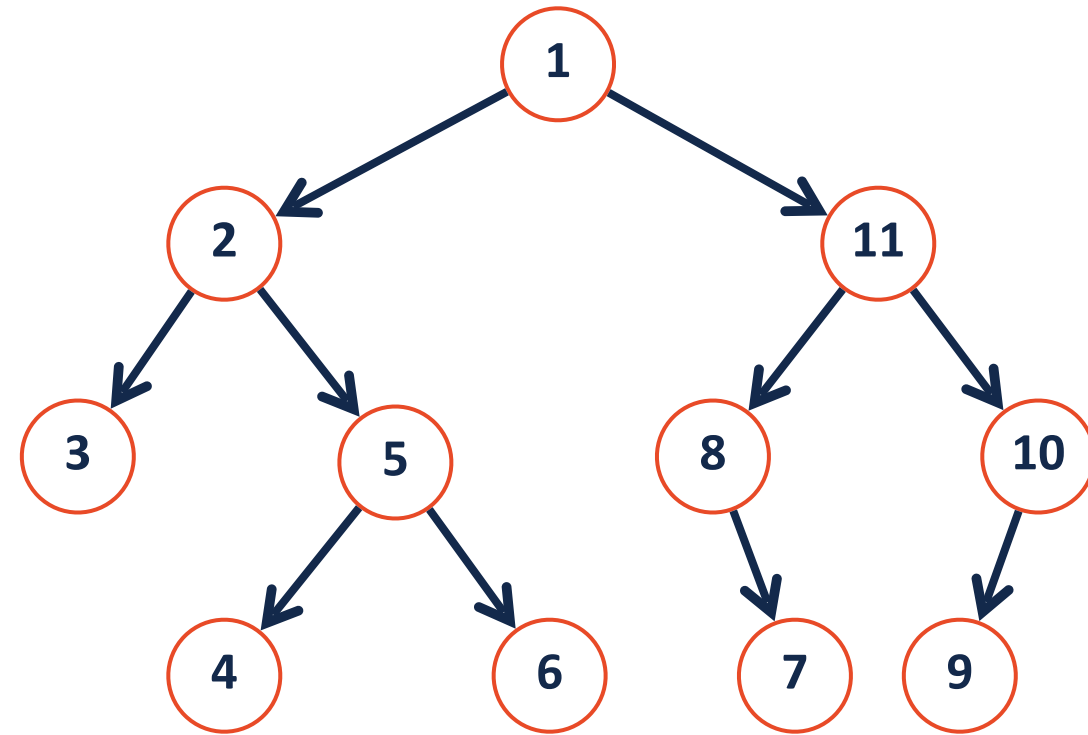


Stack:

Print:

Breadth First Search

Fully explore depth i before exploring depth $i+1$



Breadth First Search

Fully explore depth i before exploring depth $i+1$

Make a queue (initialized with root)

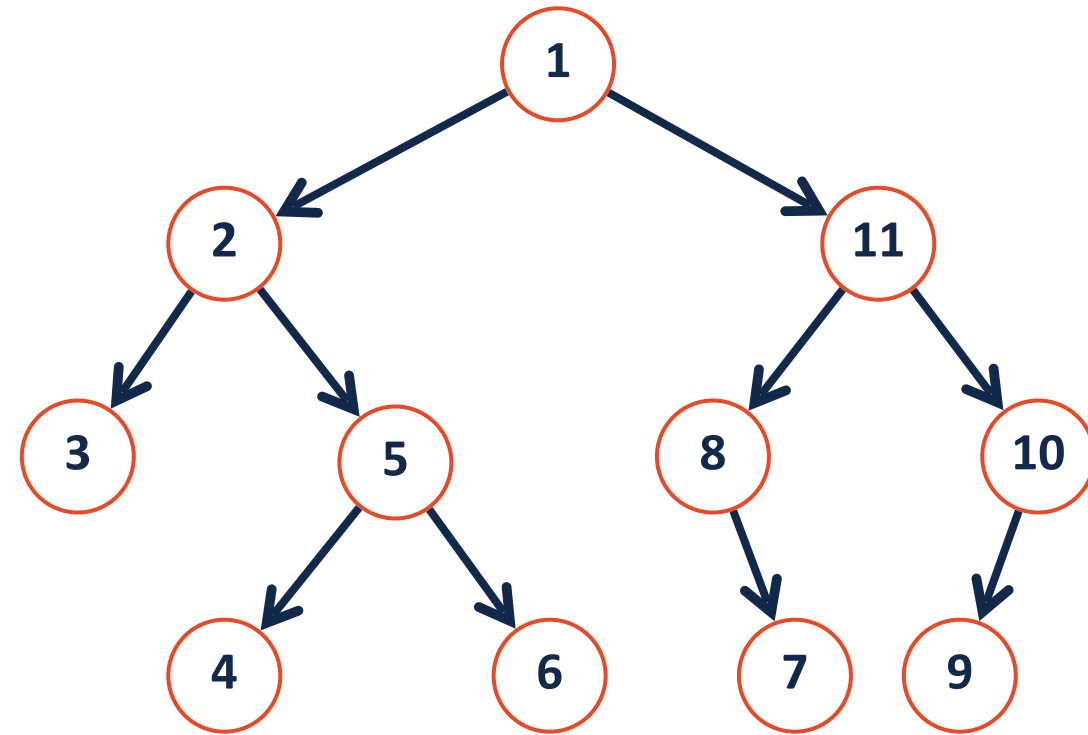
While queue not empty:

```
tmp = queue.dequeue()
```

```
print(tmp)
```

```
queue.enqueue(tmp->left)
```

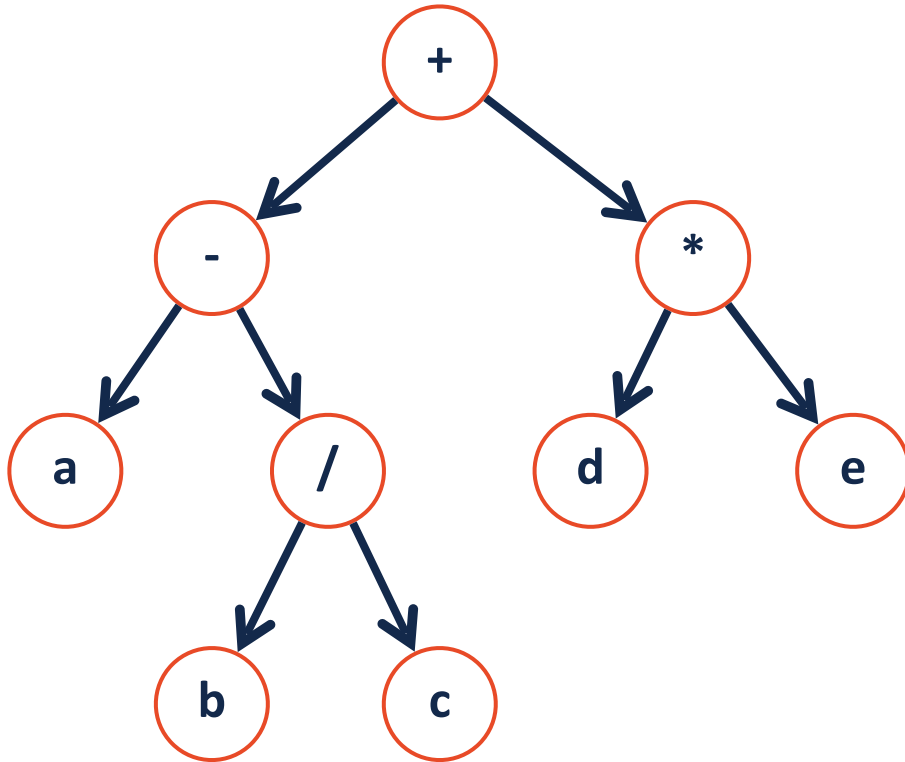
```
queue.enqueue(tmp->right)
```



Queue:

Print:

Level-Order Traversal



```
1 template<class T>
2 void BinaryTree<T>::lOrder(TreeNode * root)
3 {
4
5     Queue<TreeNode*> q;
6     q.enqueue(root);
7
8     while( q.empty() == False){
9
10        TreeNode* temp = q.head();
11        process(temp);
12
13        q.dequeue();
14
15        q.enqueue(temp->left);
16        q.enqueue(temp->right);
17
18    }
19 }
```

Tree Search

How can we improve our ability to search a binary tree?

What do we trade in order to do so?