

Data Structures

Graph Fundamentals

CS 225

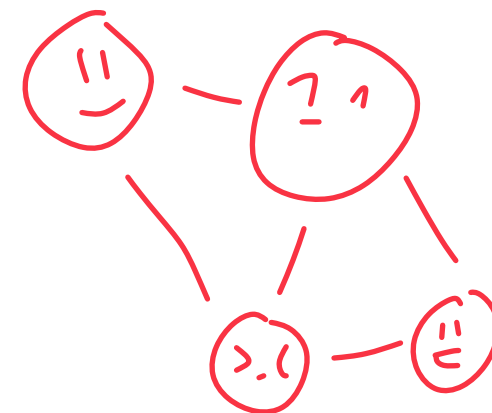
October 21, 2024

Brad Solomon



UNIVERSITY OF
ILLINOIS
URBANA - CHAMPAIGN

Department of Computer Science



Learning Objectives

Define graph vocabulary

Discuss graph implementation and storage strategies

Whats next?

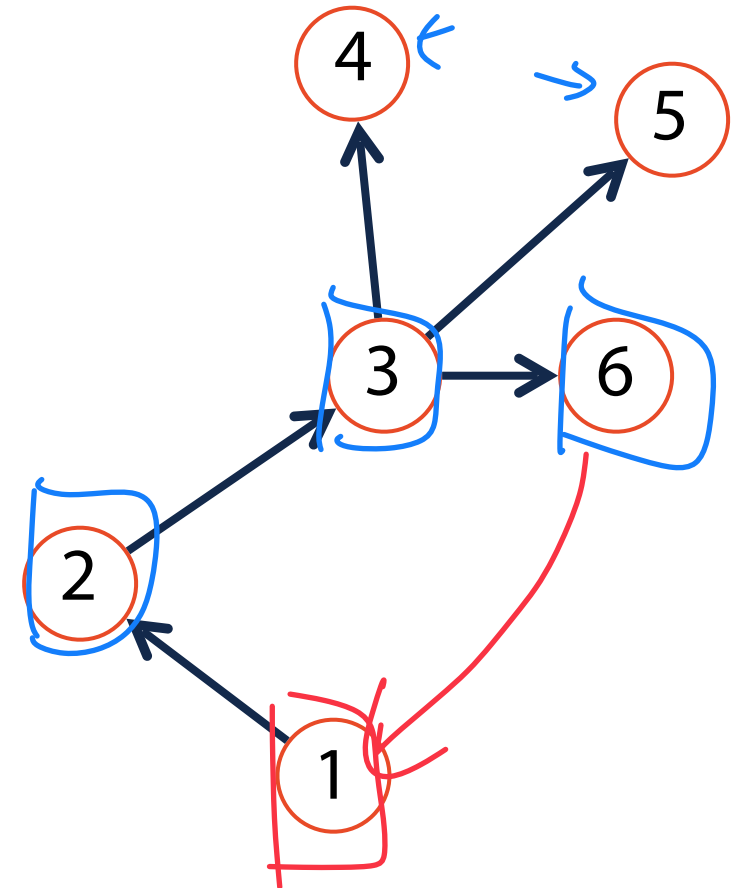
A non-linear data structure defined recursively as a collection of nodes where each node contains a value and zero or more connected nodes.

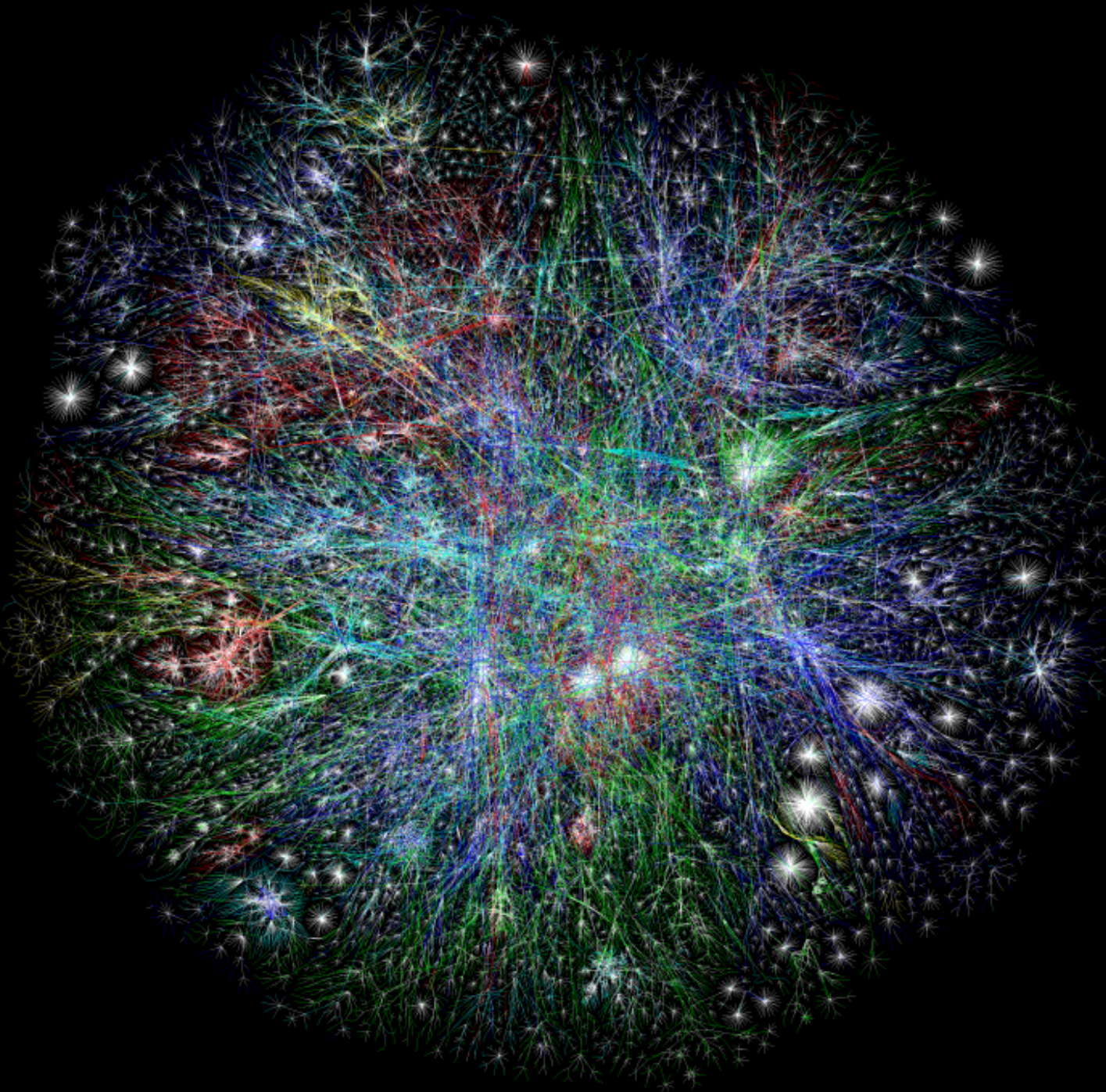
graph

(In CS 225) a ~~tree~~ is also:

1) Acyclic — contains no cycles ~~X~~

2) Rooted — root node connected to all nodes ~~X~~





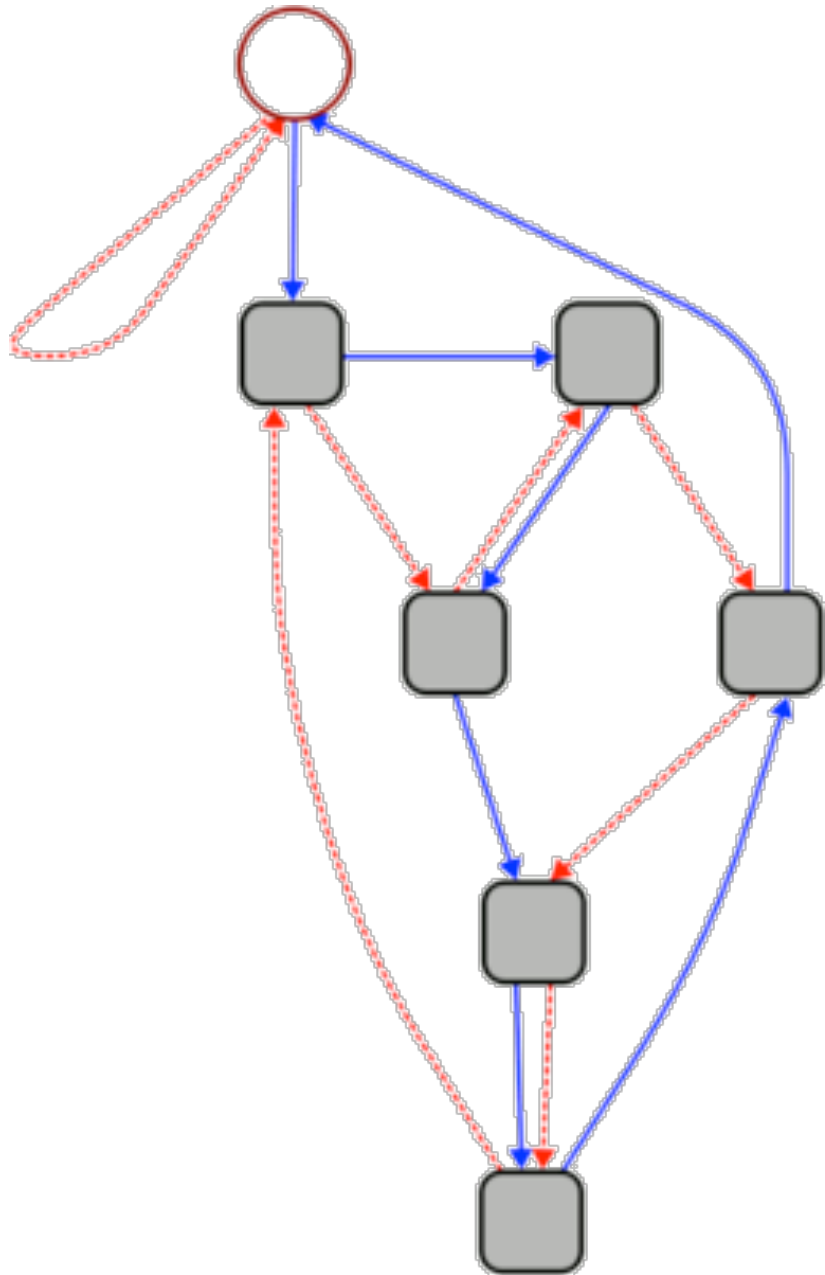
Nodes: Routers and servers

Edges: Connections

Big Data!

The Internet 2003

[The OPTÉ Project \(2003\)](#)



This graph can be used to quickly calculate whether a given number is divisible by 7.

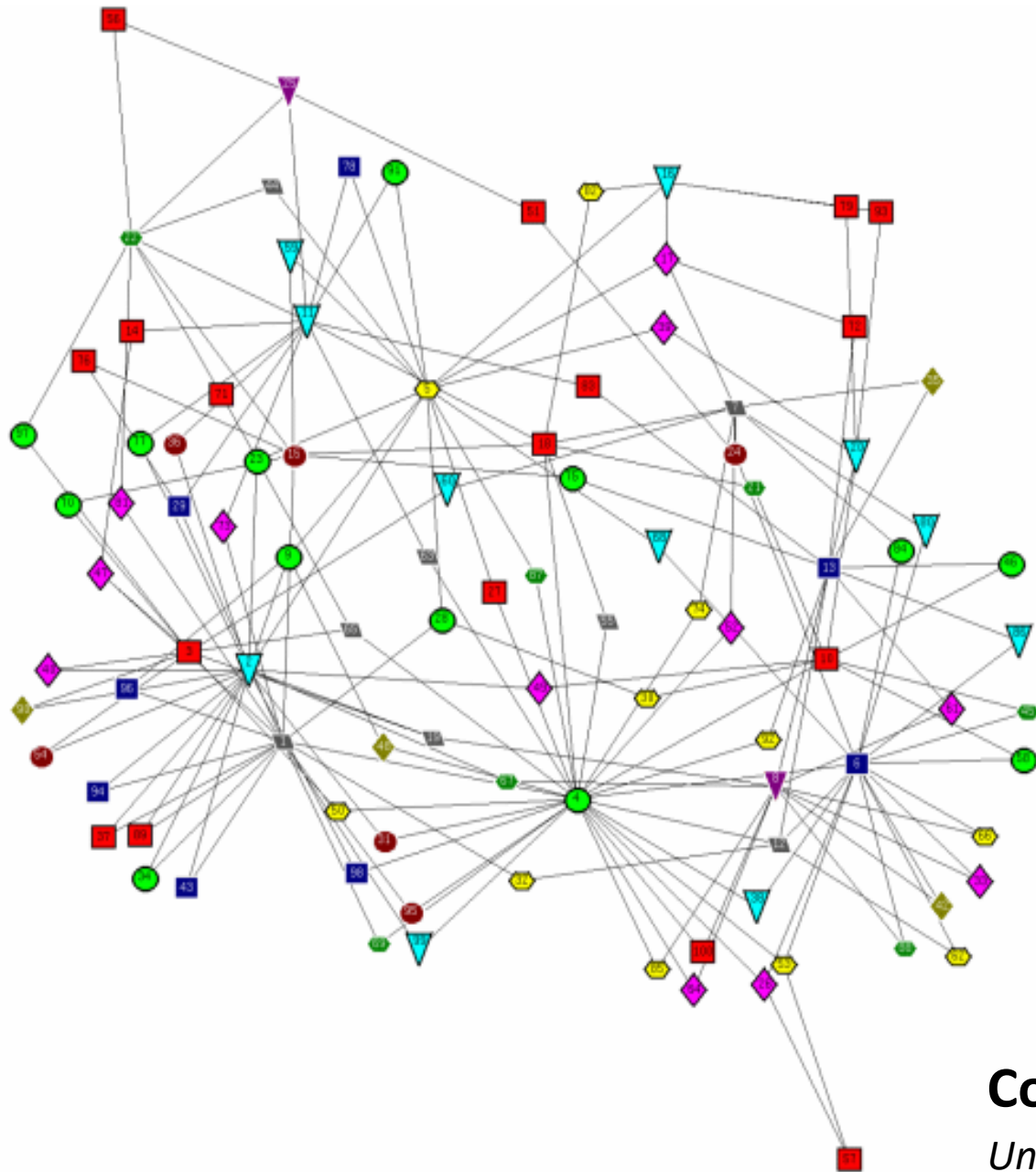
1. Start at the circle node at the top.
2. For each digit **d** in the given number, follow **d blue (solid) edges** in succession. As you move from one digit to the next, follow **1 red (dashed) edge**.
3. If you end up back at the circle node, your number is divisible by 7.

3703 — divisible by 7!

“Rule of 7”

Unknown Source

Presented by Cinda Heeren, 2016



Course as node
edge is shared student

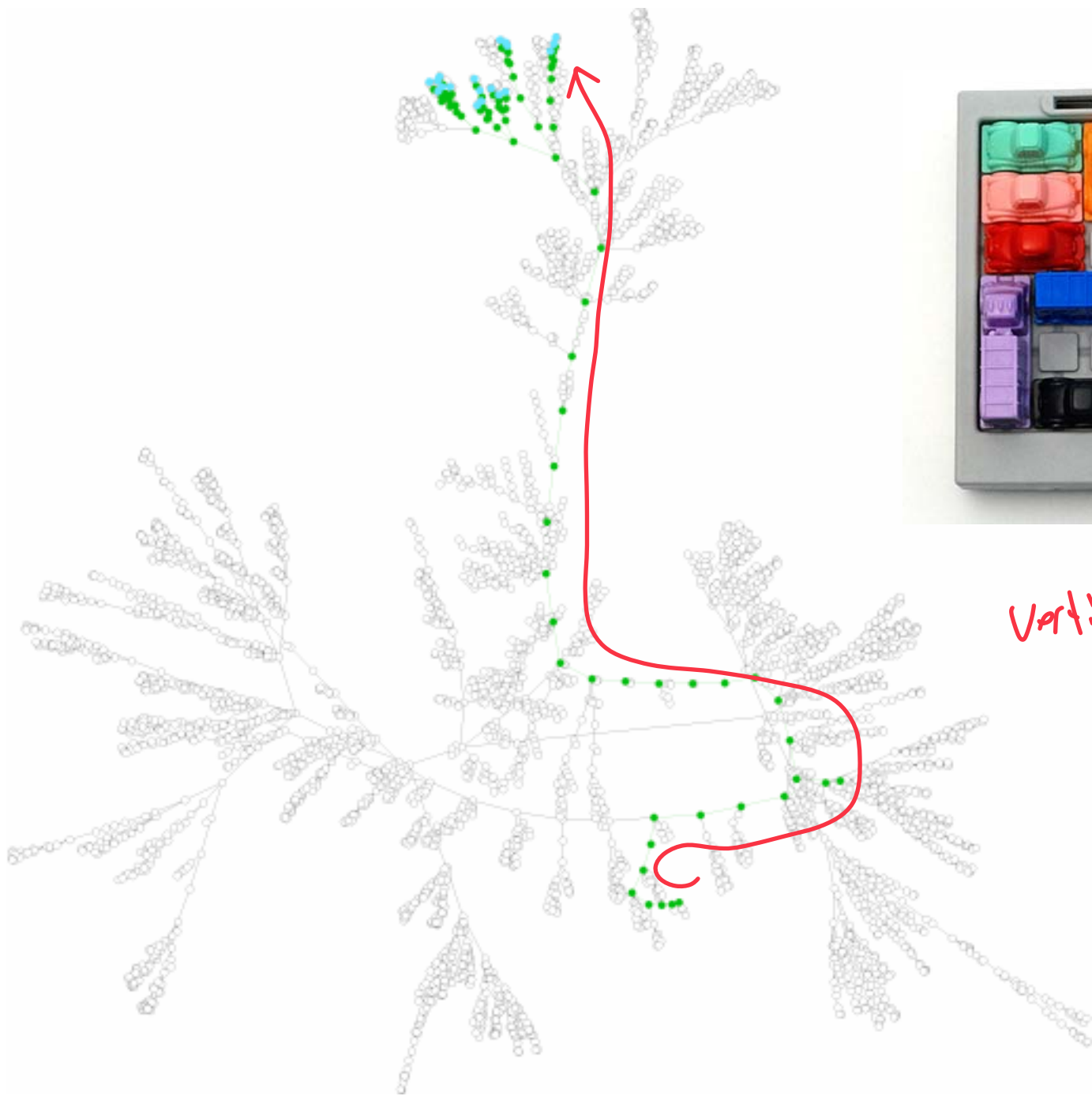
Color = exam time

↳ Oops one too many
exam times!

Conflict-Free Final Exam Scheduling Graph

Unknown Source

Presented by Cinda Heeren, 2016

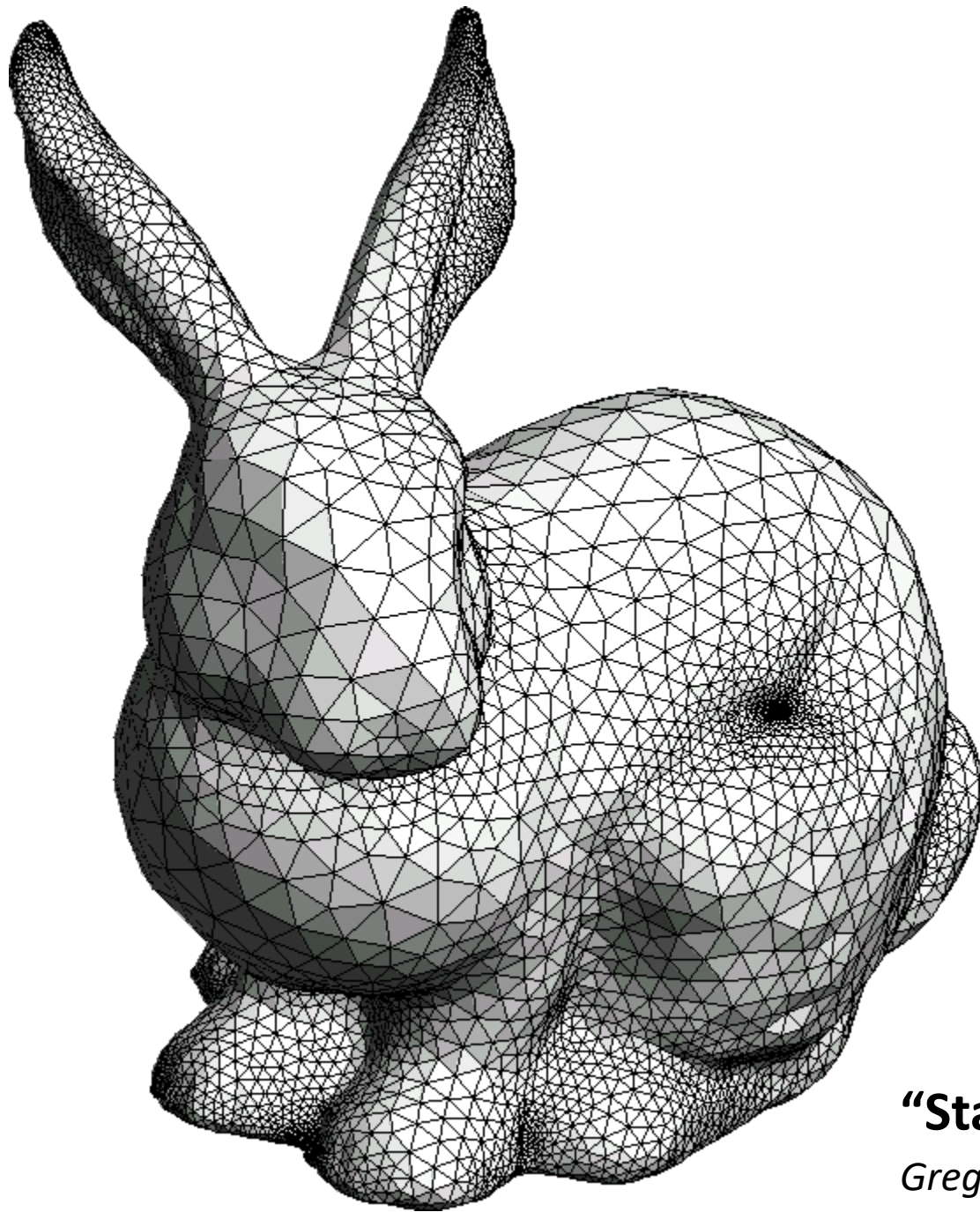


Vertices: States of a System

“Rush Hour” Solution

Unknown Source

Presented by Cinda Heeren, 2016

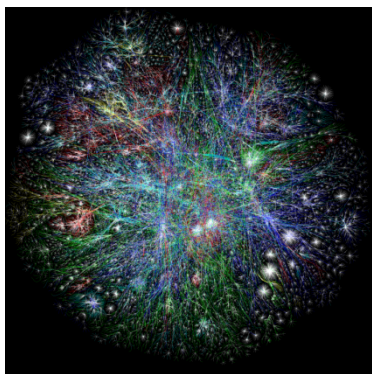


graphics!

“Stanford Bunny”

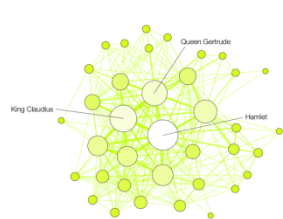
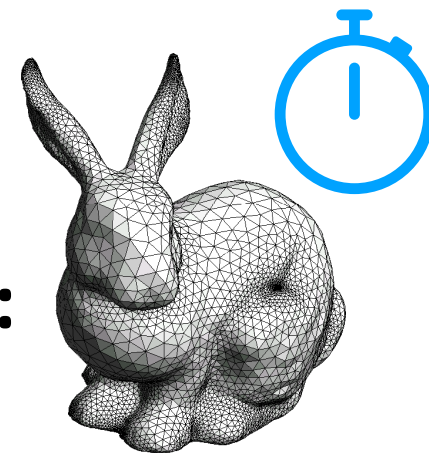
Greg Turk and Mark Levoy (1994)

Graphs

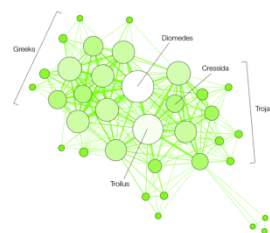


To study all of these structures:

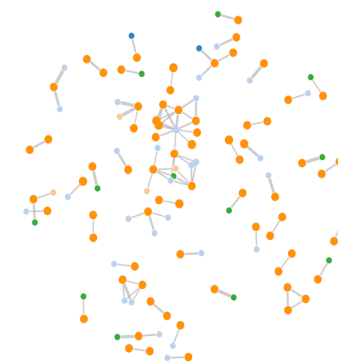
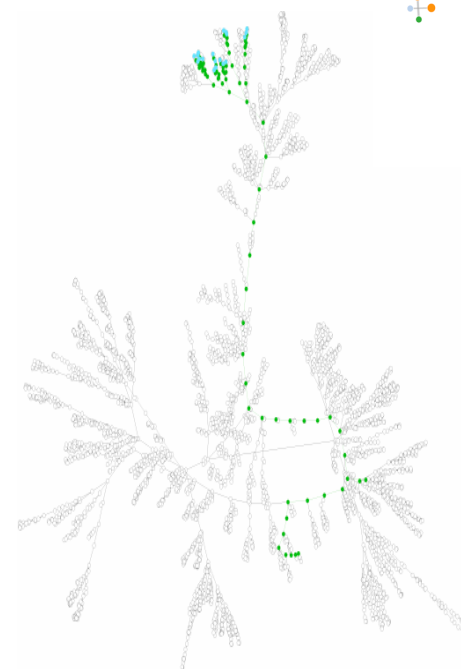
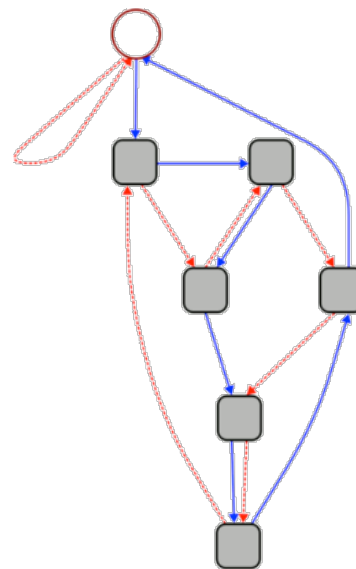
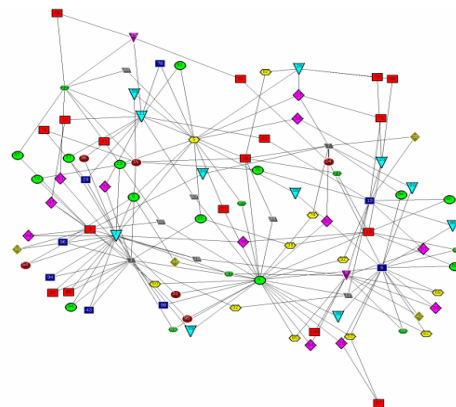
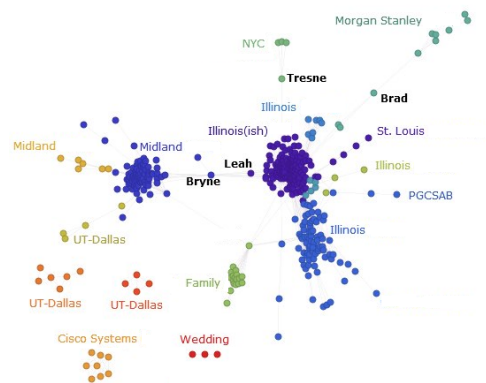
1. A common vocabulary ← Today
2. Graph implementations ← This week
3. Graph traversals ← next week
4. Graph algorithms ←



HAMLET



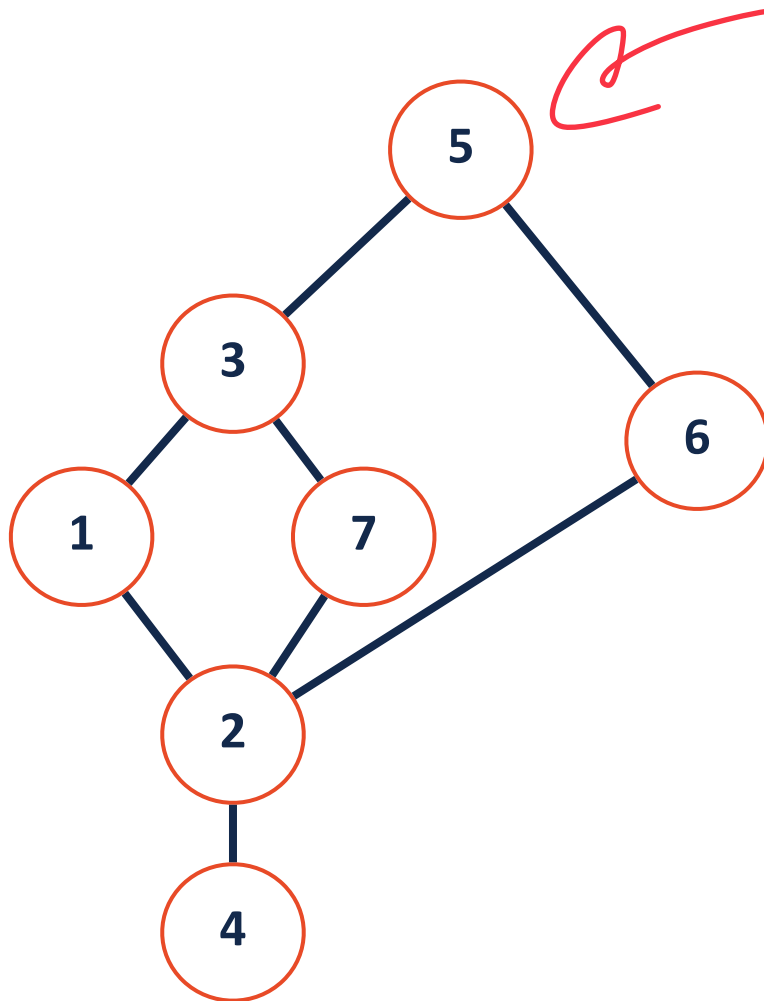
TROILUS AND CRESSIDA



Graph Vocabulary

$$G = (V, E)$$

A **graph** is a data structure containing a set of vertices and a set of edges



Vertex: Nodes of the graph

- ↳ Data
- ↳ Nothing!
- ↳ Key, Value
- ↳ State of System

Edges: The connections between nodes

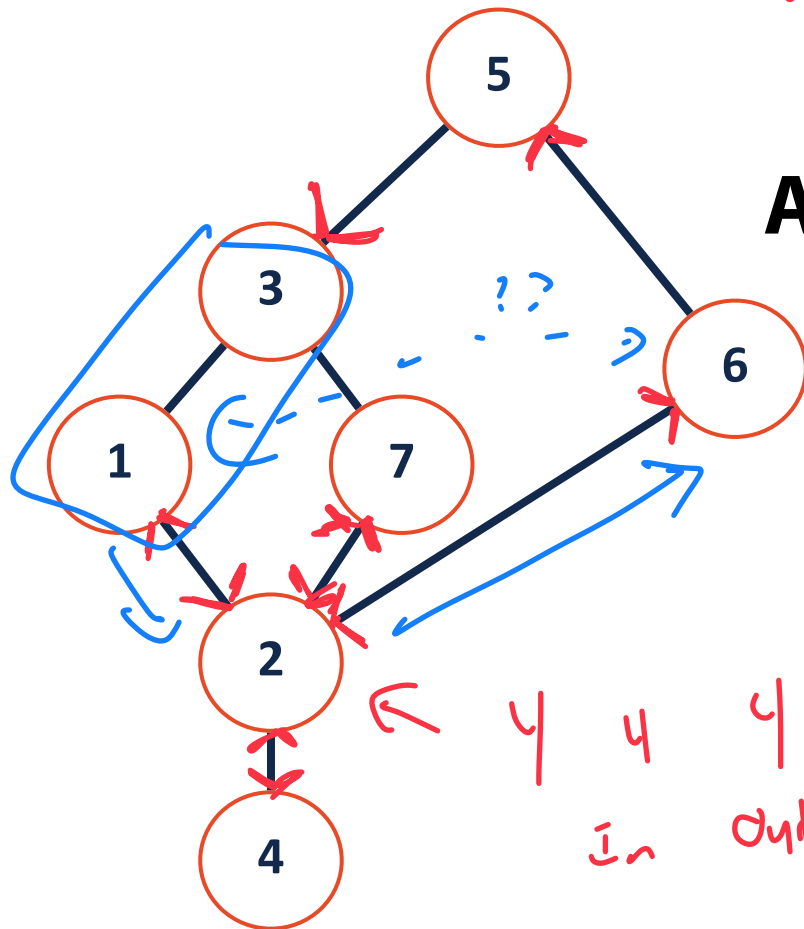
Defined by two endpoints

- ↳ Direction
- ↳ Weight

Graph Vocabulary

Degree: # of edges touching a vertex

Directed
↳ In-degree: # edges in
↳ Out-degree: # edges out



Adjacency: Two vertices are adjacent if they are connected by an edge

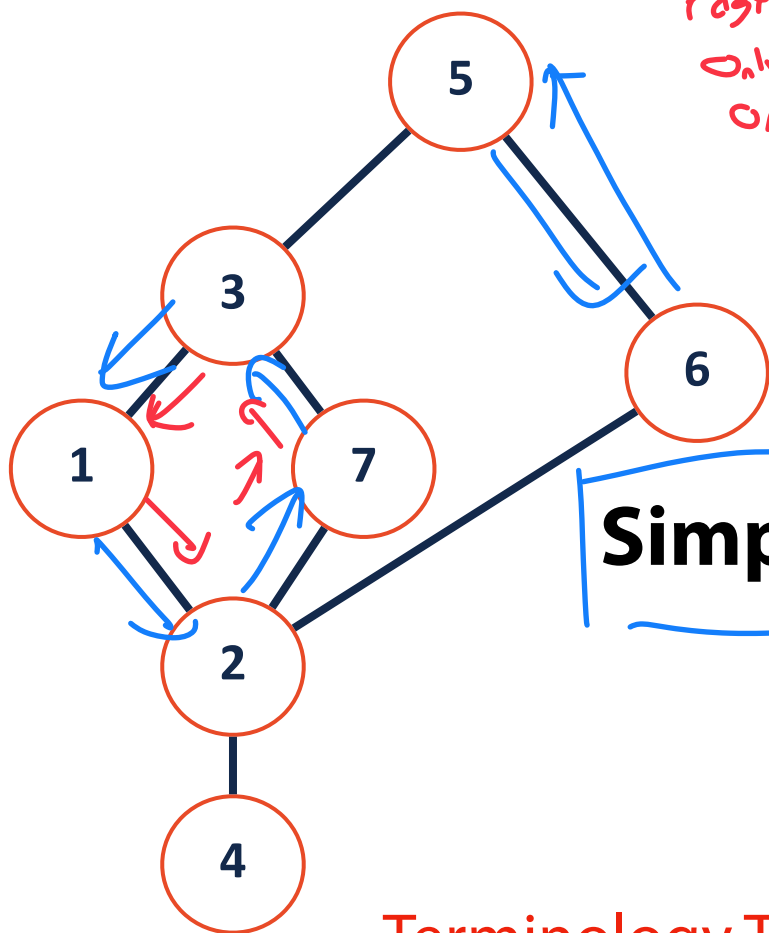
1-3
1-2

Path: A sequence of vertices (or edges) between two nodes

1 → 2 → 6

Graph Vocabulary

A graph has **no root** and **may contain cycles**



using
edges
only
once?

Cycle: A path from a node to itself

$1 \rightarrow 2 \rightarrow 7 \rightarrow 3$

$(1,2) \rightarrow (2,7) \rightarrow (7,3) \rightarrow (3,1)$

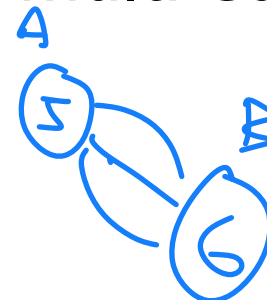
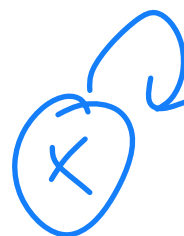
Arg: using at least 2 edges

vertices

edges

Simple Graph:

No self-loops or multi-edges

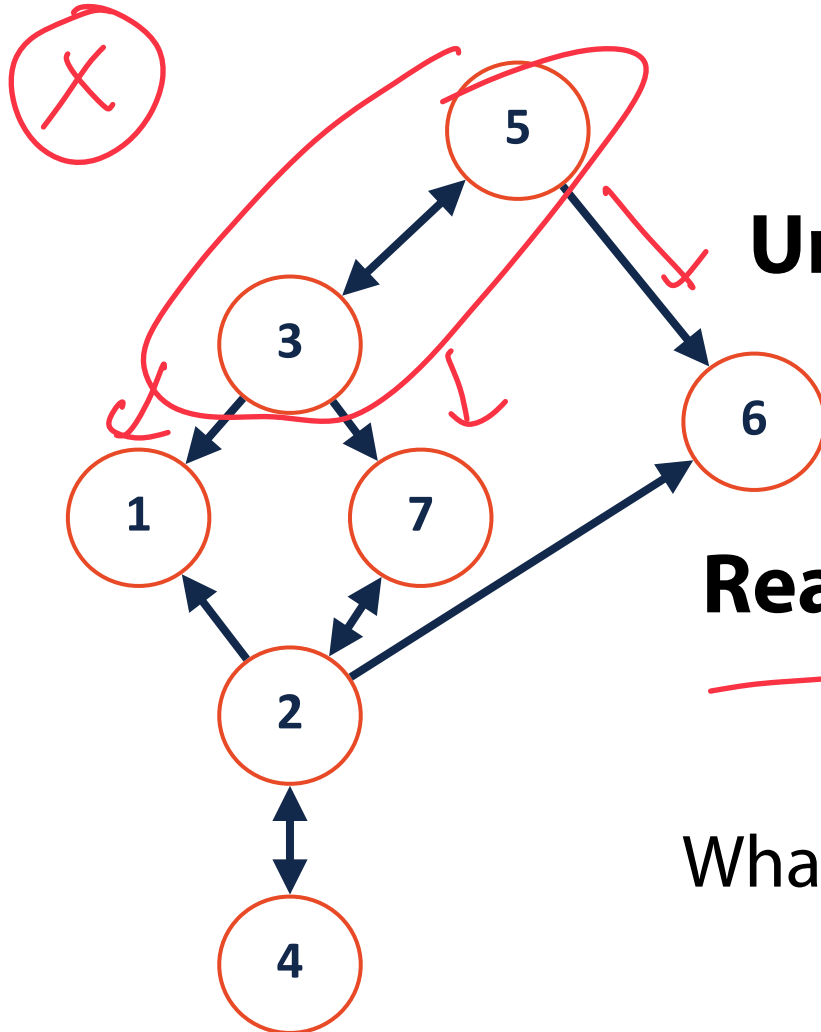


Terminology Trivia: Every tree is a graph but not every graph is a tree

Graph Vocabulary

A graph may be **directed** or **undirected**

$1 \rightarrow 2$
 $2 \rightarrow 1$



Directed: Edges are one way connections

Undirected: Traversable in either direction
 $1 \leftrightarrow 2$

Reachability: v_2 is reachable from v_1 if there is a path from v_1 to v_2

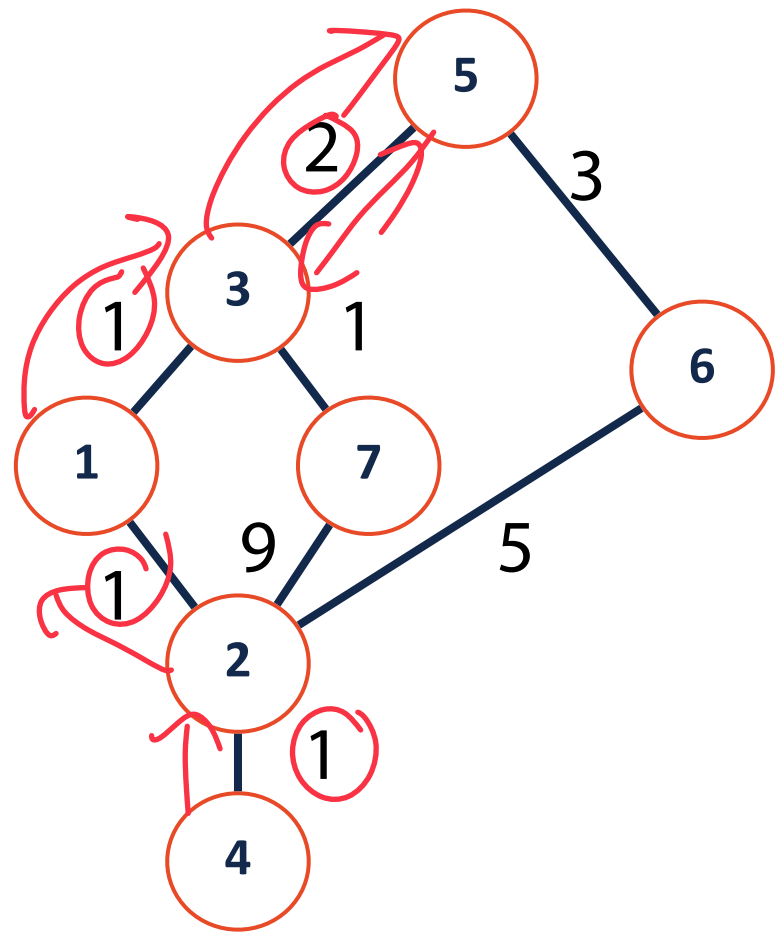
What nodes are not reachable from 4?

3, 5

Graph Vocabulary

Start	End	Weight
3	5	2

A graph may be **weighted** or **unweighted**



Weights: A value associated with an edge

What is the shortest path from 4 to 5? 5

Sum of weights

Unweighted shortest path: # edges

↳ Every edge weight 1

Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

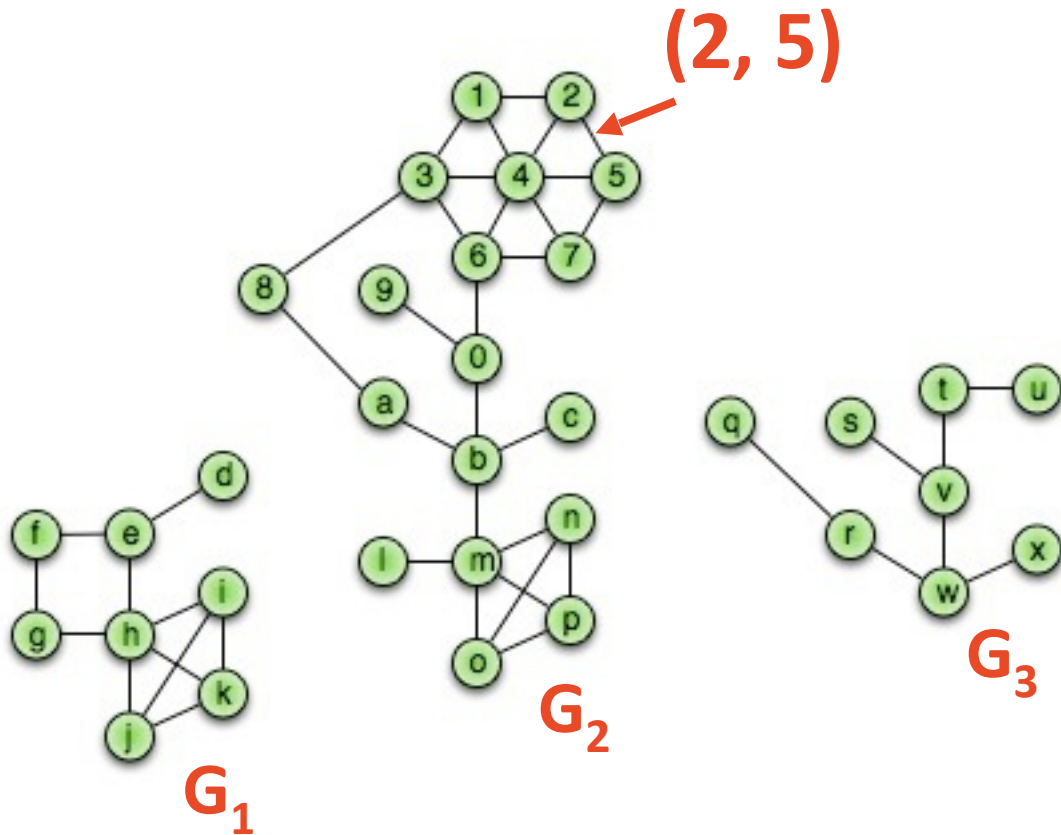
$$|E| = m$$

Subgraph(G):

$$G' = (V', E'):$$

$$V' \subseteq V, E' \subseteq E, \text{ and}$$

$$(u, v) \in E' \rightarrow u \in V', v \in V'$$



Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



Subgraph(G):

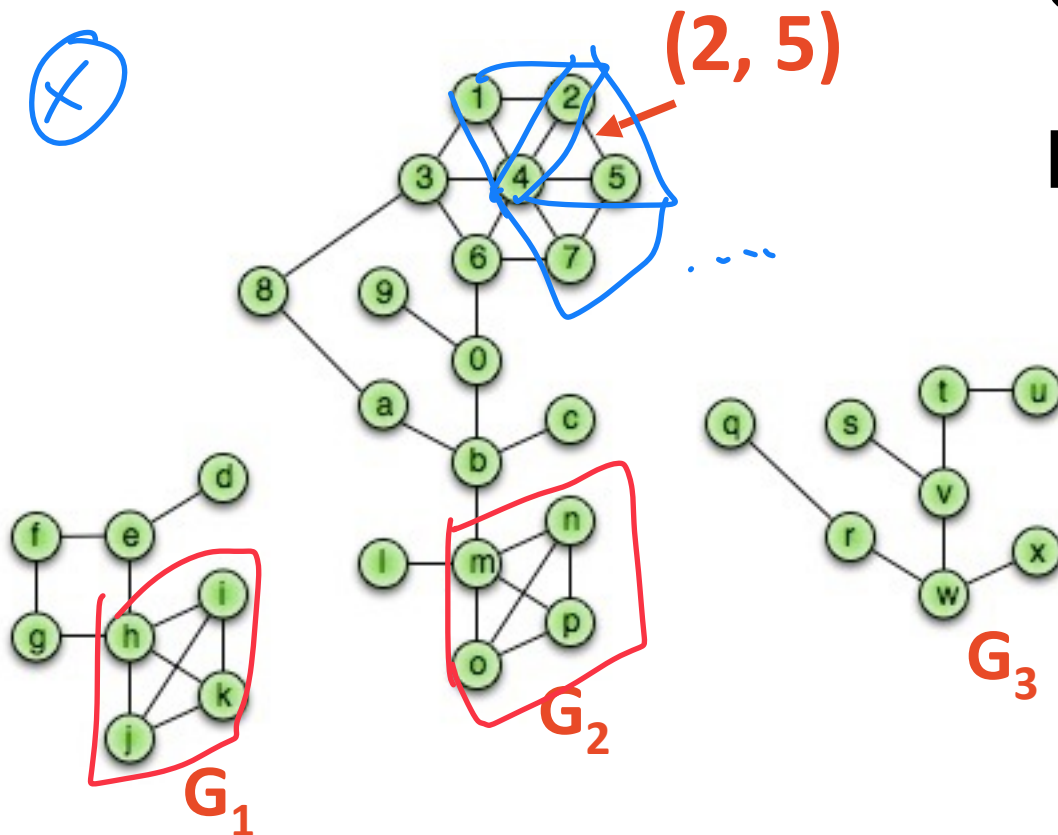
$$G' = (V', E')$$

$V' \subseteq V, E' \subseteq E$, and

$$(u, v) \in E' \rightarrow u \in V', v \in V'$$

Complete Subgraph:

Every pair of vertices are adjacent



Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$

Subgraph(G):

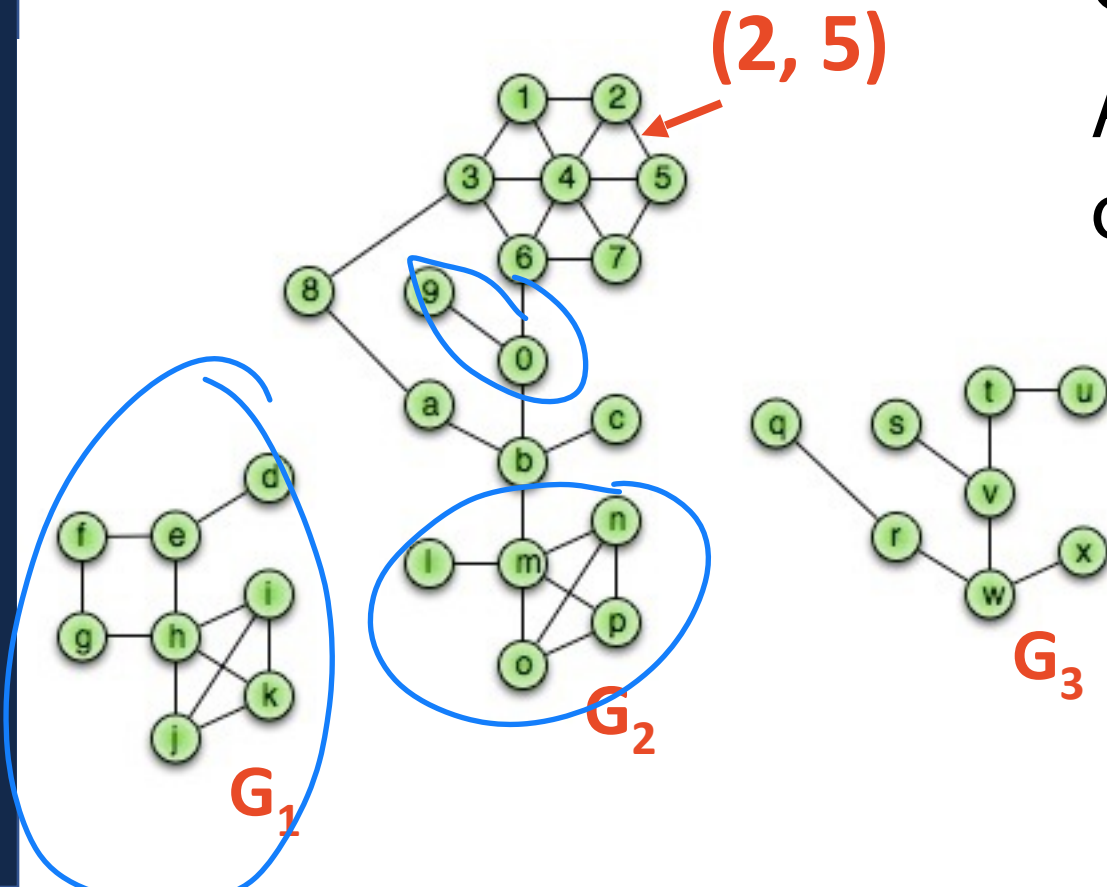
$$G' = (V', E')$$

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Connected Subgraph:

A path exists between every pair of vertices



Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$

Subgraph(G):

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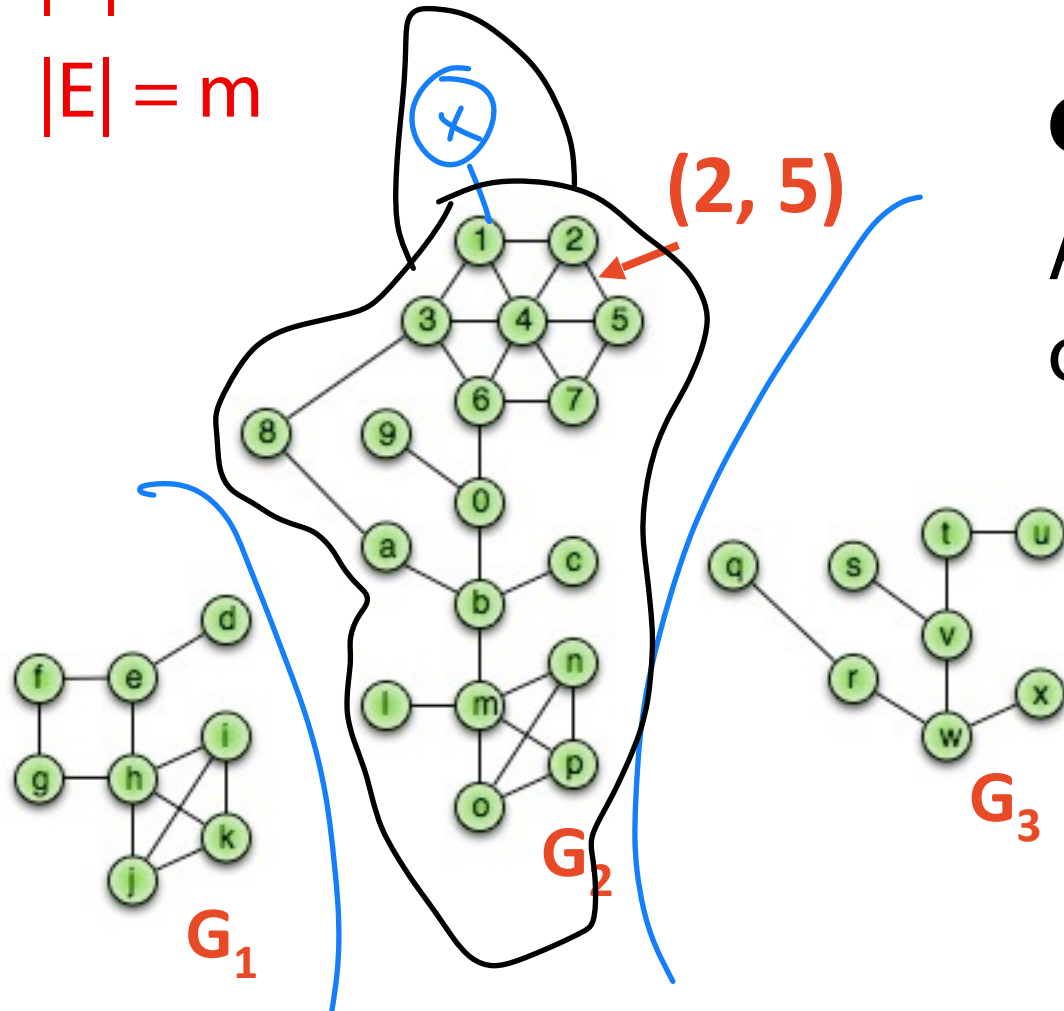
Connected Subgraph:

A path exists between every pair of vertices

Connected Components:

A connected subgraph that is not part of a larger subgraph

↳ The largest connected subgraph

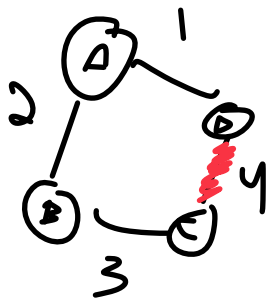


Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



Subgraph(G):

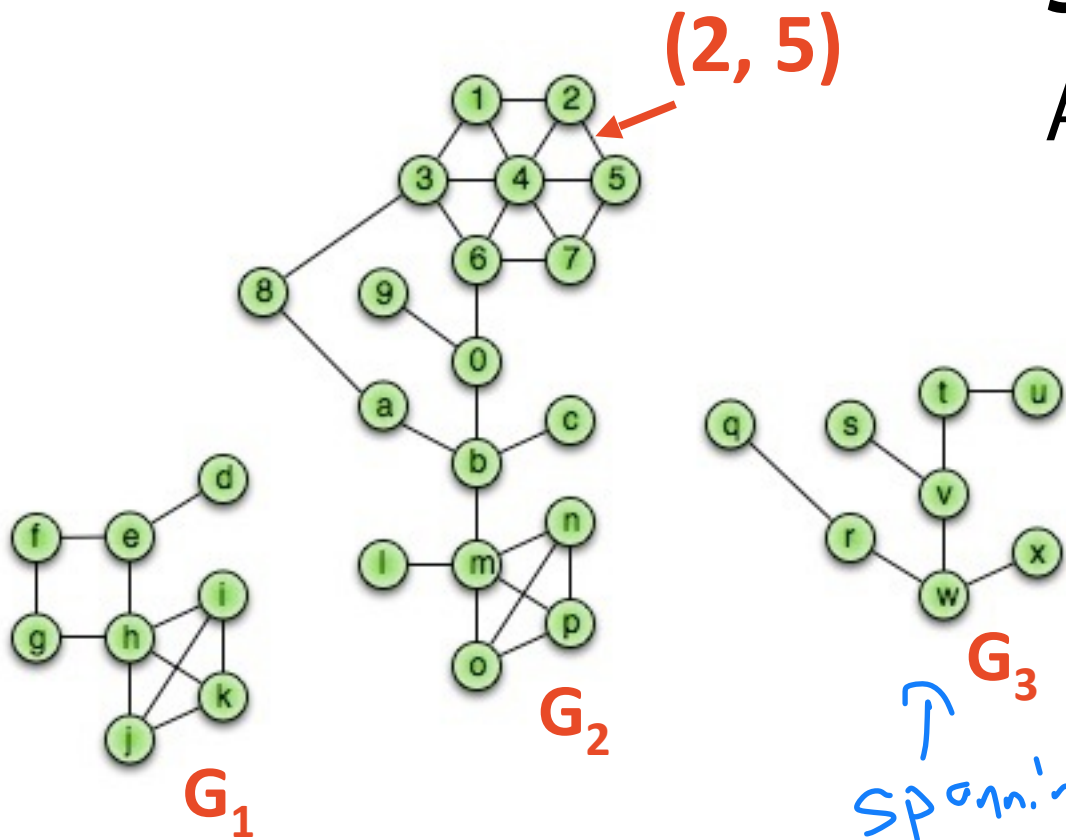
$$G' = (V', E')$$

$V' \subseteq V, E' \subseteq E$, and

$$(u, v) \in E' \rightarrow u \in V', v \in V'$$

Spanning Tree:

A connected graph with no cycles



Minimum Spanning tree

↳ picks lowest weight

$n-1$ edges in spanning tree

Spanning tree



Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$

collection (or set)
of vertices
&
edges

Graph Terminology is very important!

Degree

Weight

Direction

Adjacency

Complete

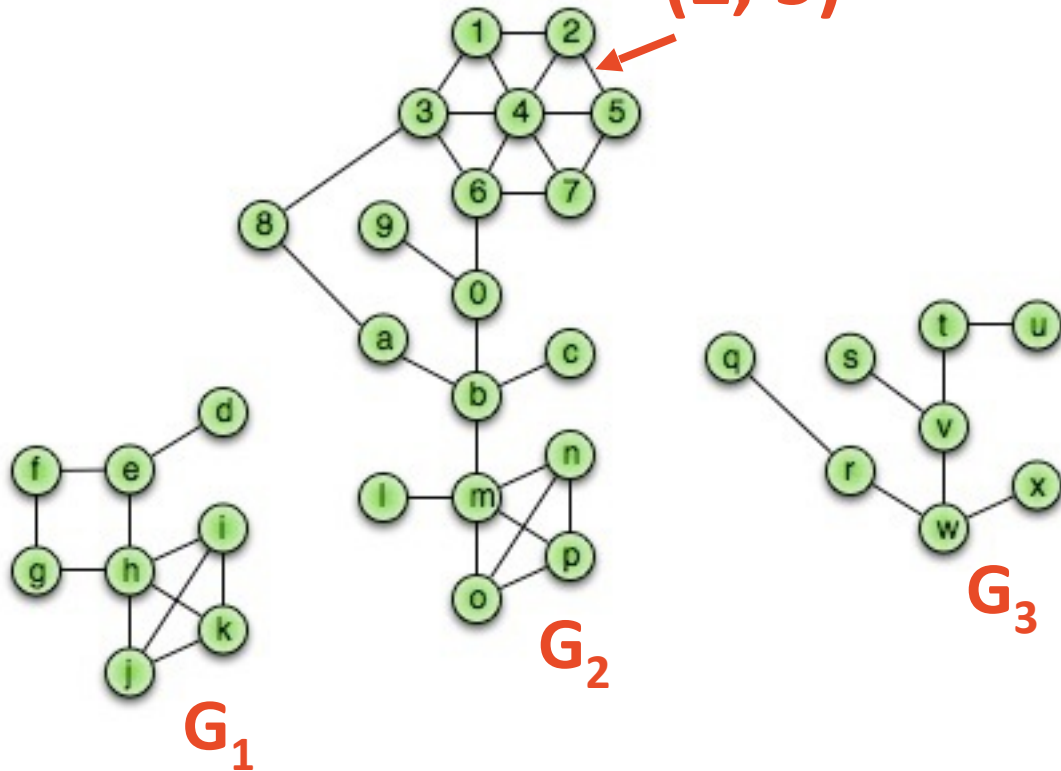
Connected

Acyclic

Spanning

And more...

Edge
(2, 5)



Running times are often reported by n , the number of vertices, but often depend on m , the number of edges.

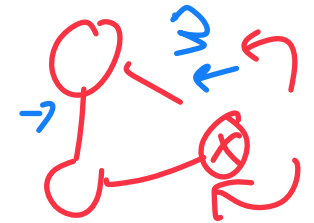
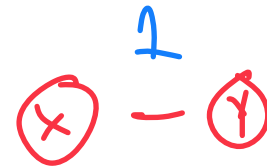
Whats the relationship between n and m ?

Minimum Edges:

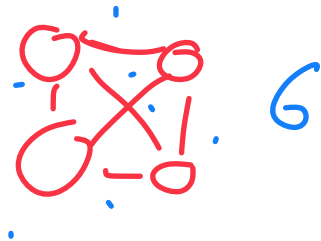
Unconnected Graph: 



Connected (Simple) Graph: $n - 1$



Maximum Edges:

Connected (Simple) Graph: $(n-1) + (n-2) + (n-3) + \dots$ 

$$\sum_{v \in V} \text{deg}(v) = \frac{n(n-1)}{2} \approx O(n^2)$$

Graphs

$A \rightarrow B$

$A \leftarrow B$

Given a collection of individual DMs between individuals, you want to build a graph of connections in a social network.

What is a vertex?

Individual

What is an edge?

A message

Are the edges directed or undirected?

↳ Directed

↳ Depends!

Are the edges weighted or unweighted?

↳ Depends!

Graphs

Given a collection of roads between cities in Illinois, you want to build a graph of the transportation infrastructure in the state.

What is a vertex?

What is an edge?

Are the edges directed or undirected?

Are the edges weighted or unweighted?

Graphs

It is important to be able to describe the structure of a graph given input.

Some other common questions:

Does your graph have cycles?

What is the largest / smallest / average degree in your graph?

What is the total number of edges?

...

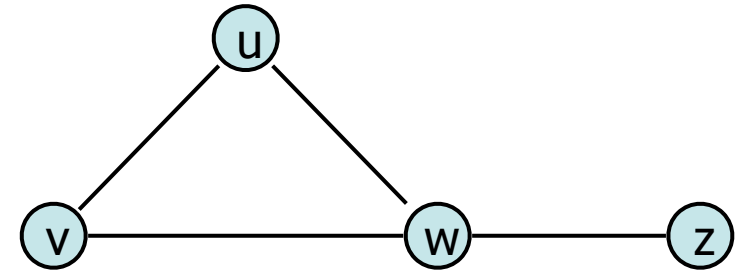
Of course, we also have to understand the graph as a **data structure**

Graph Implementation

What information do we need to store to fully define a graph?

Vertex:

Edge:



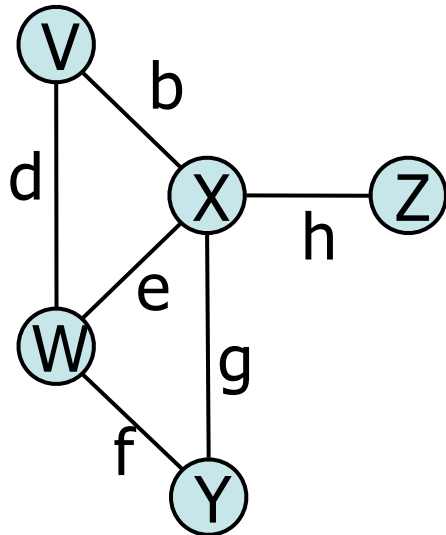
What information do we want to be able to find out quickly?

What operations do we want to prioritize?

Graph ADT

Data:

- Vertices
- Edges
- Some data structure maintaining the structure between vertices and edges.

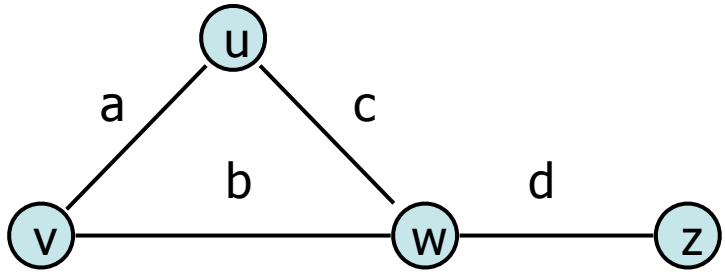


Functions:

- insertVertex(K key);
- insertEdge(Vertex v1, Vertex v2, K key);
- removeVertex(Vertex v);
- removeEdge(Vertex v1, Vertex v2);
- getEdges(Vertex v);
- areAdjacent(Vertex v1, Vertex v2);
- origin(Edge e);
- destination(Edge e);



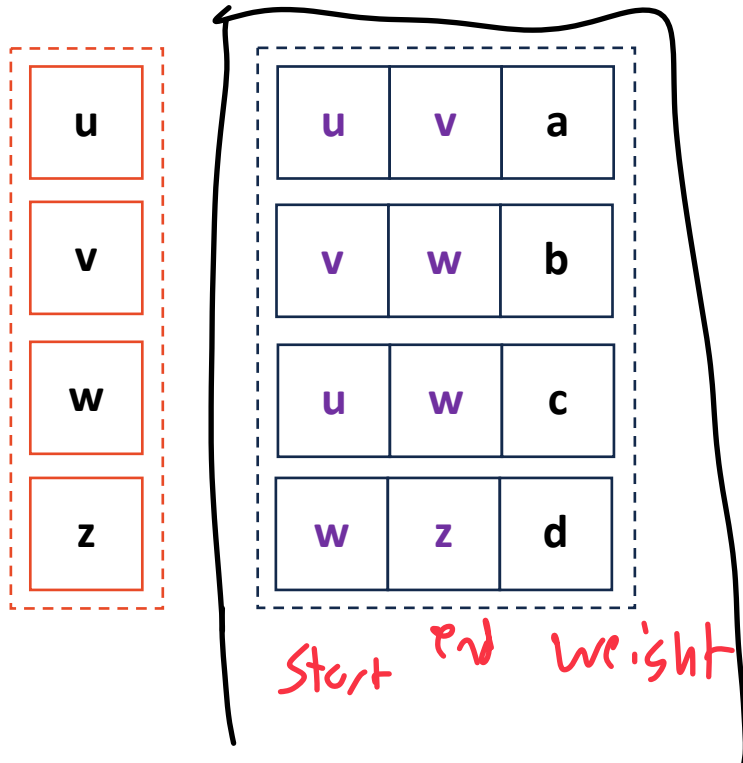
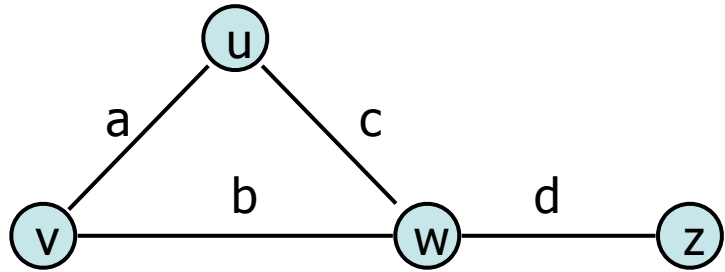
Graph Implementation Idea



↳ Adj List!

Graph Implementation: Edge List $|V| = n, |E| = m$

The equivalent of an 'unordered' data structure



Vertex Storage:

↳ optionally as a list of vertex labels

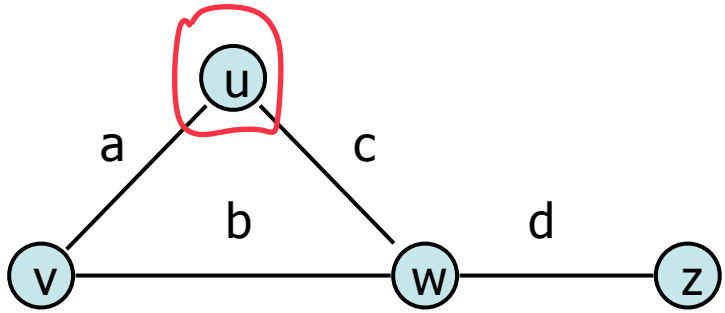
↳ Build this from edge storage

Edge Storage:

↳ A list of three values

↳ Very compact representation!

Graph Implementation: Edge List $|V| = n, |E| = m$

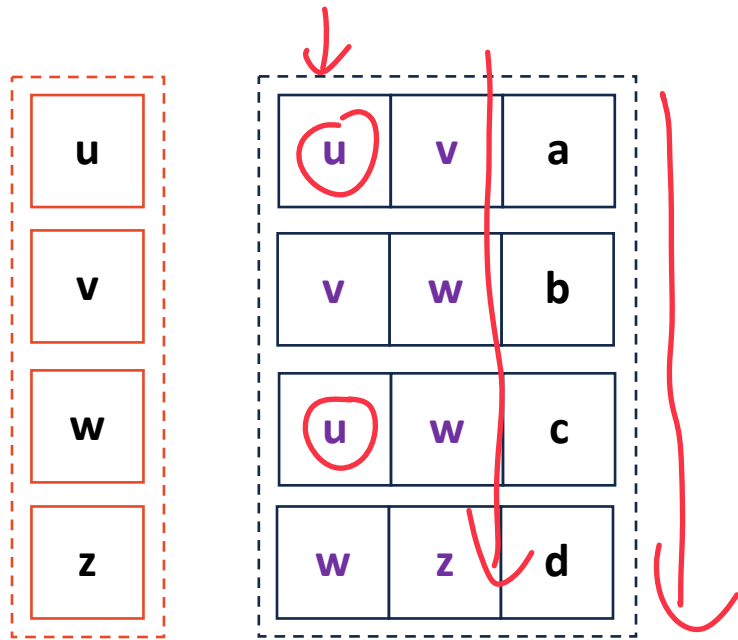


getEdges(Vertex v)

$O(m)$

↳ Traverse edge storage

Look for u in start or end

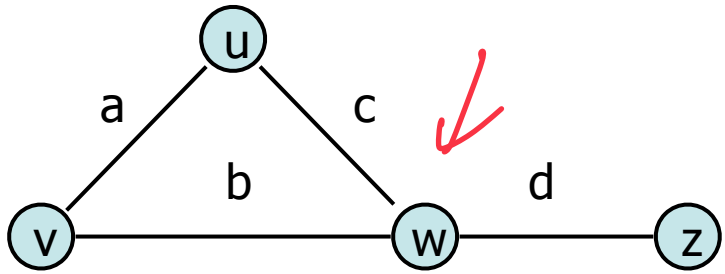


areAdjacent(Vertex v1, Vertex v2) $O(m)$

↳ Do same thing!

w u c.

Graph Implementation: Edge List $|V| = n, |E| = m$

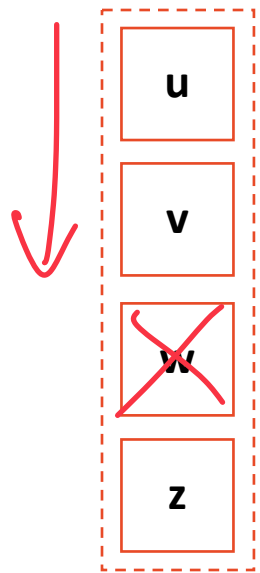


insertVertex(K key)

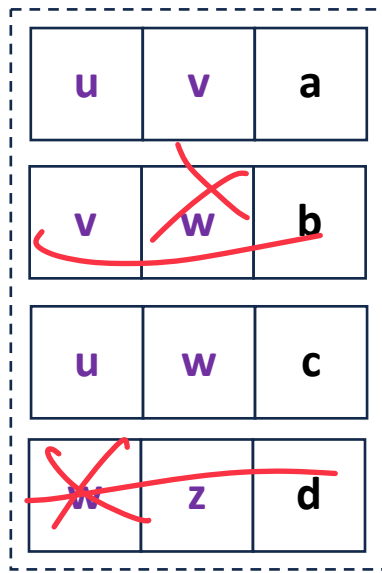
$$O(1)^*$$

↳ Array insert

Vertex



$$O(n)$$



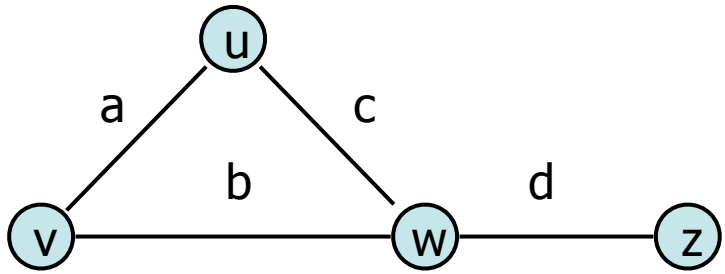
$$O(m)$$

removeVertex(Vertex v)

$$O(n + m)$$

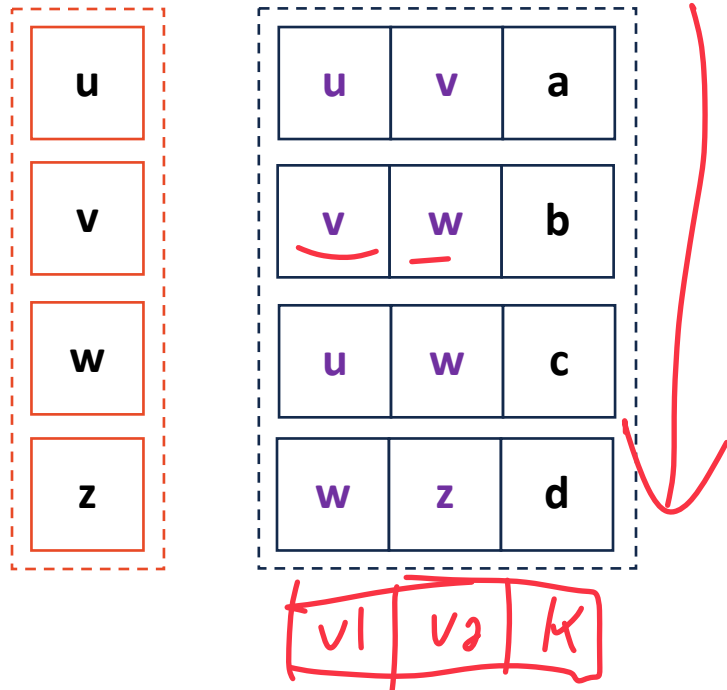
↳ Array removal

Graph Implementation: Edge List $|V| = n, |E| = m$



insertEdge(Vertex v1, Vertex v2, K key)

↳ Array insert $O(1)^*$



removeEdge(Vertex v1, Vertex v2)

↳ $O(m)$

Graph Implementation: Edge List



Pros:

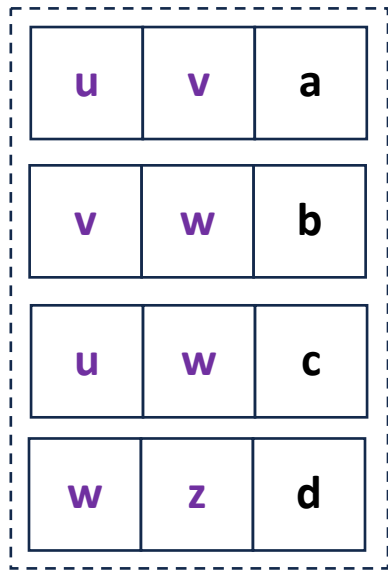
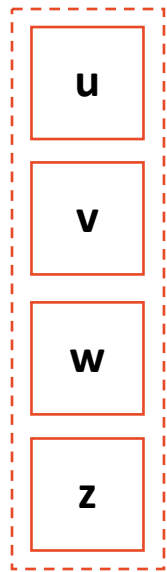
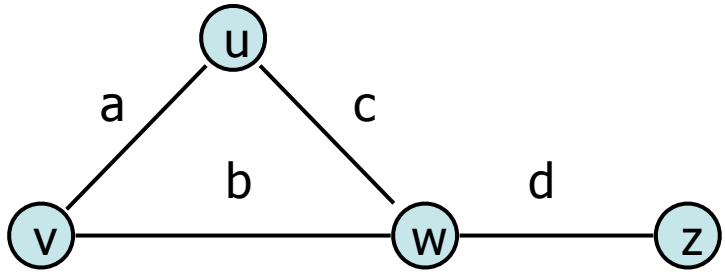
Cons:

Graph Implementation: Brainstorming better

What operations might I want to do very quickly?

What modifications might allow me to do these things faster?

Graph Implementation: Adjacency Matrix



	u	v	w	z
u				
v				
w				
z				