

# Data Structures

## Tree Definitions

CS 225

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**ILLINOIS**  
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# Exam 1 (9/18 — 9/20)

Autograded MC and one coding question

Manually graded short answer prompt

Practice exam will be released on PL

Topics covered can be found on website

**Registration started August 22**

<https://courses.engr.illinois.edu/cs225/fa2024/exams/>

# Learning Objectives

Review trees and binary trees

Practice tree theory with recursive definitions and proofs

Discuss the tree ADT

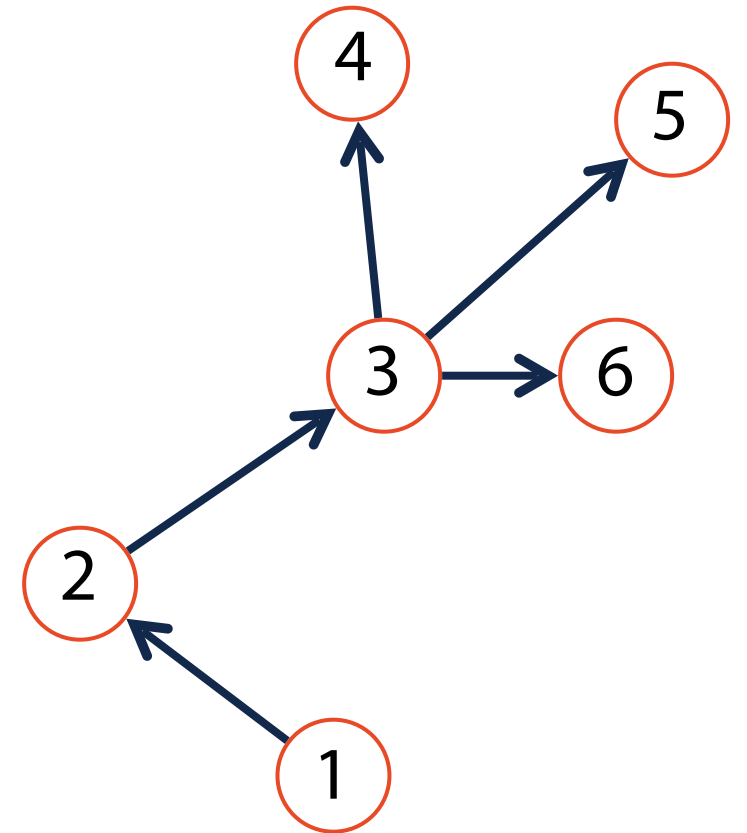
Explore tree implementation details

# Trees

A non-linear data structure defined recursively as a collection of nodes where each node contains a value and zero or more connected nodes.

[In CS 225] a tree is also:

- 1) Acyclic — No path from node to itself
- 2) Rooted — A specific node is labeled root

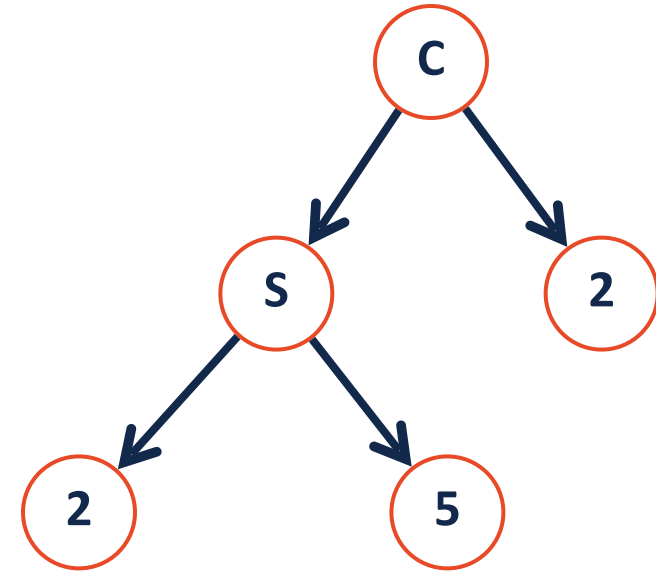


# Binary Tree

A **binary tree** is a tree  $T$  such that:

1.  $T = \emptyset$

2.  $T = (data, T_L, T_R)$



# Binary Tree

Lets define additional terminology for different **types** of binary trees!

1.

2.

3.

# Binary Tree: full

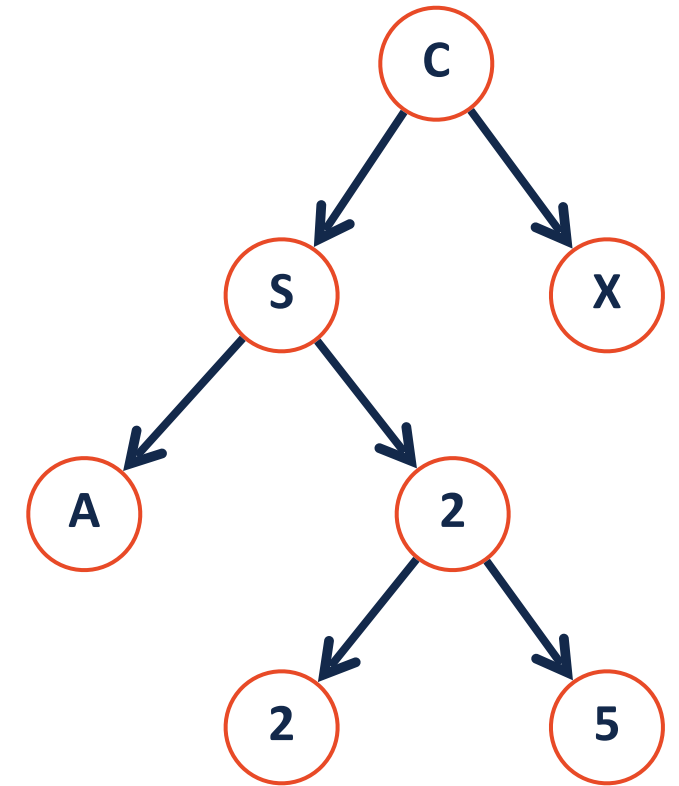
A **full tree** is a binary tree where every node has either 0 or 2 children

A tree **F** is **full** if and only if:

1.

2.

3.



# Binary Tree: full

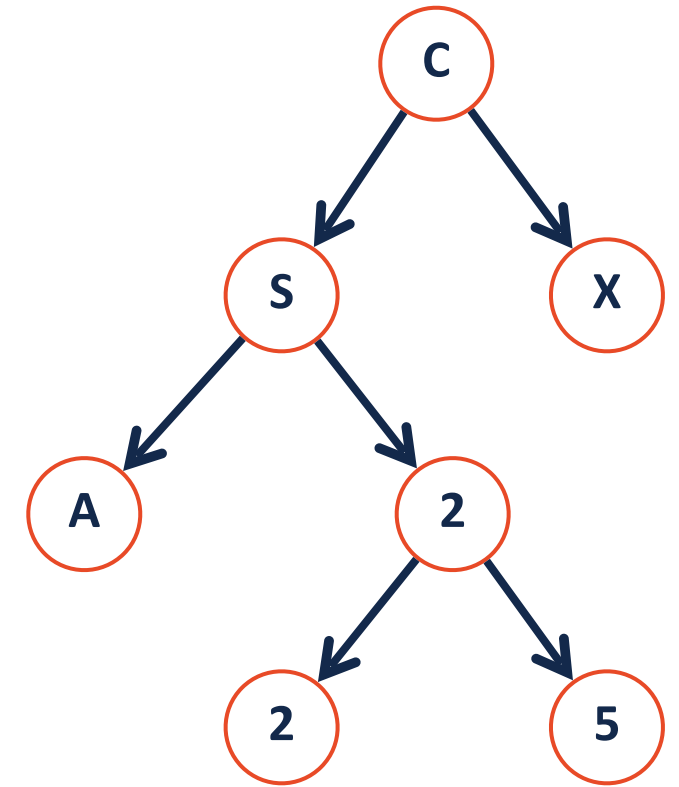
A **full tree** is a binary tree where every node has either 0 or 2 children

A tree **F** is **full** if and only if:

1.  $F = \emptyset$

2.  $F = (data, \emptyset, \emptyset)$

3.  $F = (data, F_l \neq \emptyset, F_r \neq \emptyset)$





# Binary Tree: perfect

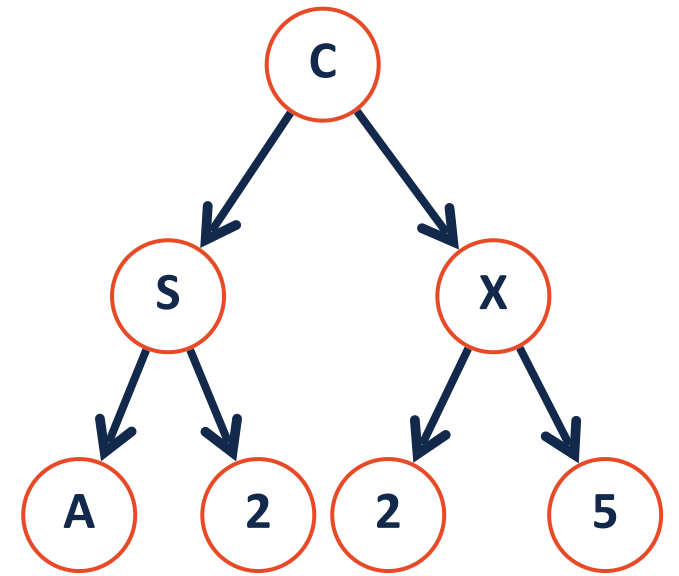
A **perfect tree** is a binary tree where...

Every internal node has 2 children and all leaves are at the same level.

A tree **P** is **perfect** if and only if:

1.

2.



# Binary Tree: perfect

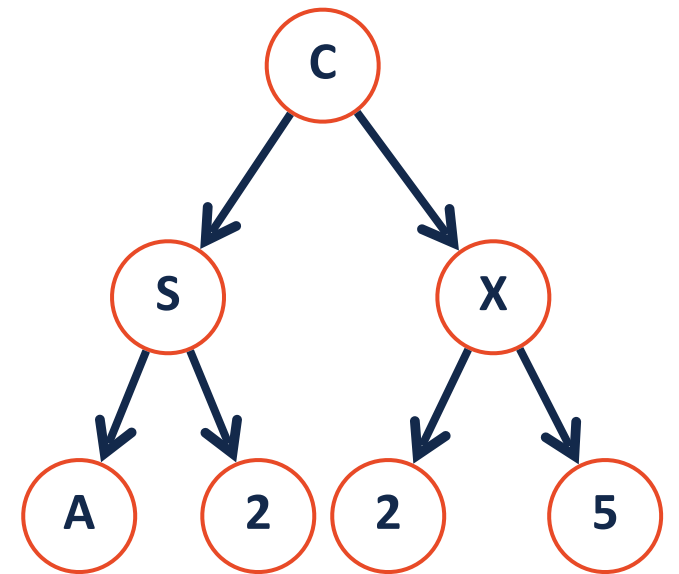
A **perfect tree** is a binary tree where...

Every internal node has 2 children and all leaves are at the same level.

A tree **P** is **perfect** if and only if:

$$1. P_h = (data, P_{h-1}, P_{h-1})$$

$$2. P_0 = (data, \emptyset, \emptyset) \equiv P_{-1} = \emptyset$$



# Binary Tree: complete

A **complete tree** is a B.T. where...

All levels except the last are completely filled.

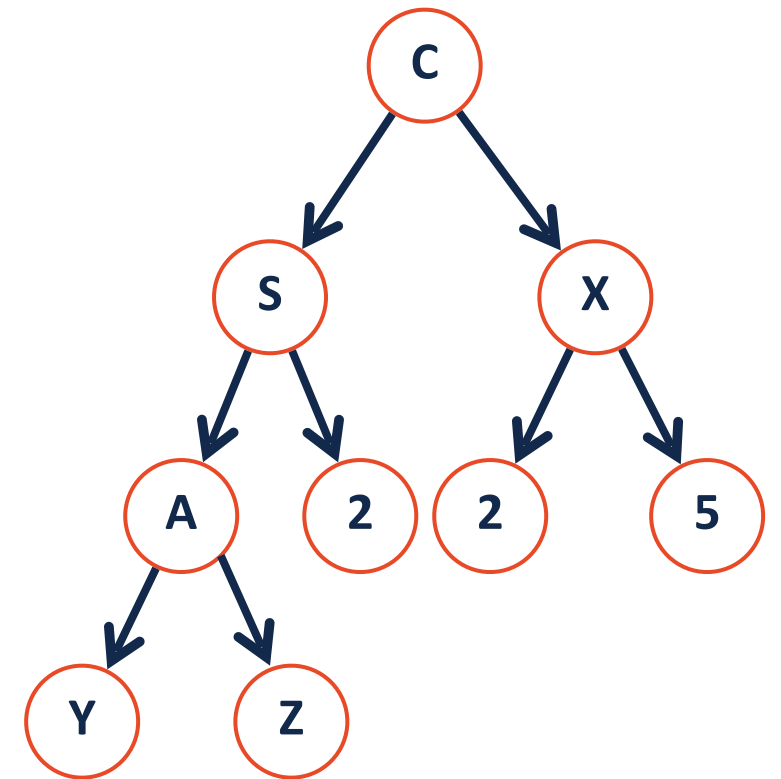
The last level contains at least one node (and is pushed to left)

A tree **C** is **complete** if and only if:

1.

2.

3.



# Binary Tree: complete

A **complete tree** is a B.T. where...

All levels except the last are completely filled.

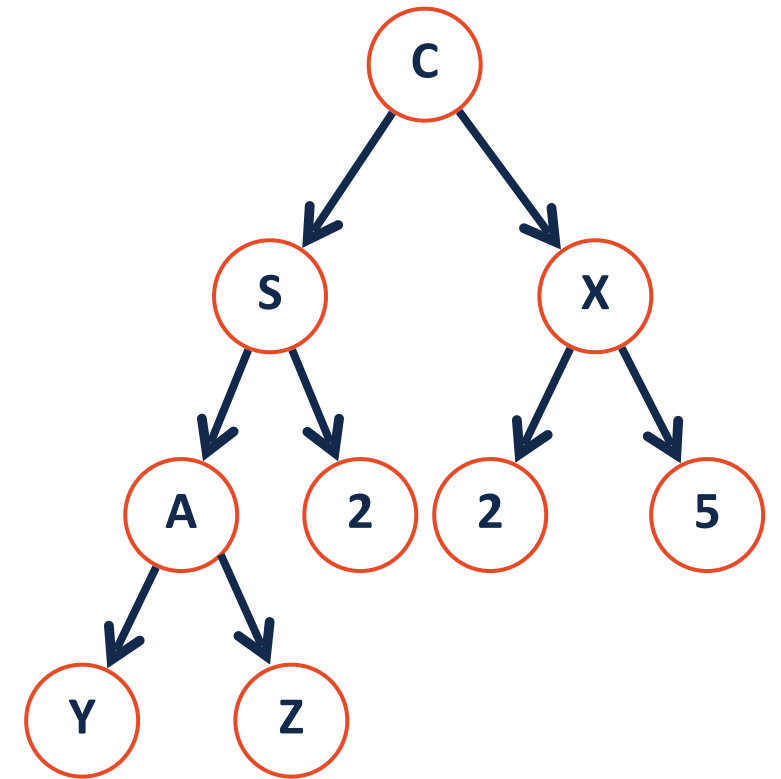
The last level contains at least one node (and is pushed to left)

A tree **C** is **complete** if and only if:

1.  $C_h = (data, C_{h-1}, P_{h-2})$

2.  $C_h = (data, P_{h-1}, C_{h-1})$

3.  $C_{-1} = \emptyset$



# Binary Tree



Why do we care?

1. Terminology instantly defines a particular tree structure
2. Understanding how to think 'recursively' is very important.

# Binary Tree: Thinking with Types

Is every **full** tree **complete**?

Is every **complete** tree **full**?

# Binary Tree: Practicing Proofs

**Theorem:** If there are  $n$  objects in our representation of a binary tree, then there are \_\_\_\_\_ NULL pointers.

# Binary Tree: Practicing Proofs

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Base Case:



# Binary Tree: Practicing Proofs

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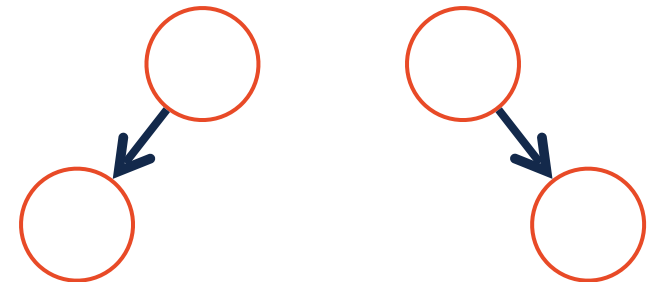
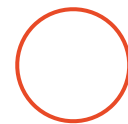
Base Case:

Let  $F(n)$  be the max number of NULL pointers in a tree of  $n$  nodes

$N=0$  has one NULL

$N=1$  has two NULL

$N=2$  has three NULL



**Theorem:** If there are  $n$  objects in our representation of a binary tree, then there are  $n+1$  NULL pointers.

Induction Step:

**Theorem:** If there are  $n$  objects in our representation of a binary tree, then there are  $n+1$  NULL pointers.



**IS:** Assume claim is true for  $|T| \leq k - 1$ , prove true for  $|T| = k$

By def,  $T = r, T_L, T_R$ . Let  $q$  be the # of nodes in  $T_L$

Since  $r$  exists,  $0 \leq q \leq k - 1$ . By IH,  $T_L$  has  $q + 1$  NULL

All nodes not in  $r$  or  $T_L$  exist in  $T_R$ . So  $T_R$  has  $k - q - 1$  nodes

$k - q - 1$  is also smaller than  $k$  so by IH,  $T_R$  has  $k - q$  NULL

Total number of NULL is the sum of  $T_L$  and  $T_R$ :  $q + 1 + k - q = k + 1$



# Tree ADT

Insert

Remove

Traverse

Find

Constructor

# BinaryTree.h

```
1 #pragma once
2
3 template <class T>
4 class BinaryTree {
5     public:
6         /* ... */
7
8     private:
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25 };
```

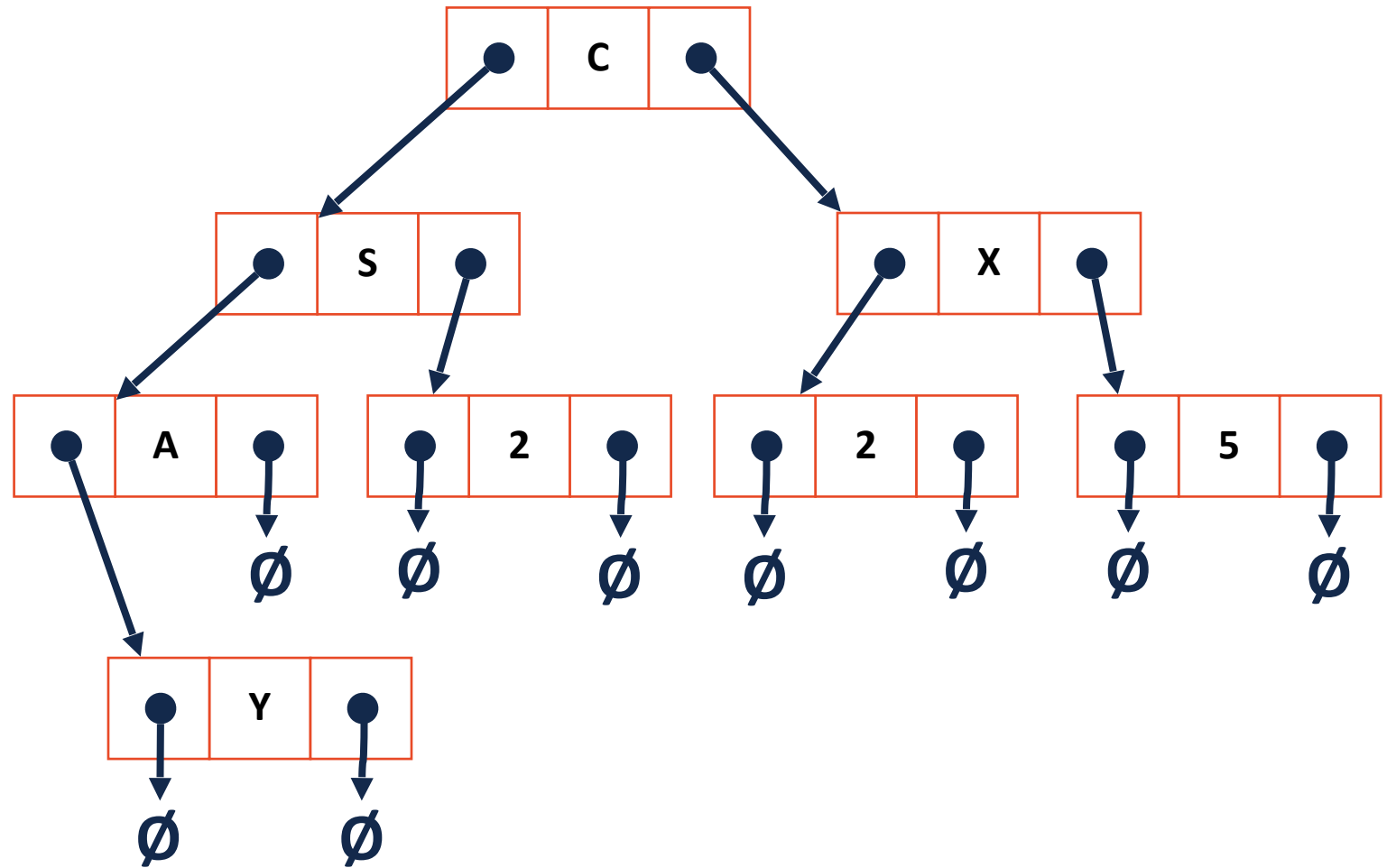
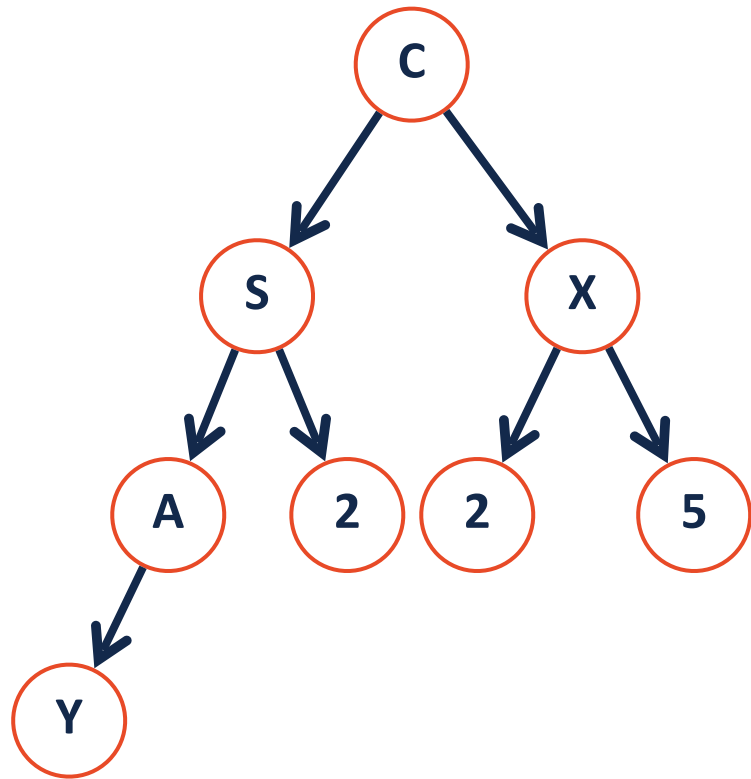
## List.h

```
1 #pragma once
2
3 template <typename T>
4 class List {
5     public:
6         /* ... */
7     private:
8         class ListNode {
9             T & data;
10
11             ListNode * next;
12
13
14             ListNode(T & data) :
15                 data(data), next(NULL) { }
16         };
17
18
19         ListNode *head_;
20         /* ... */
21 };
```

## Tree.h

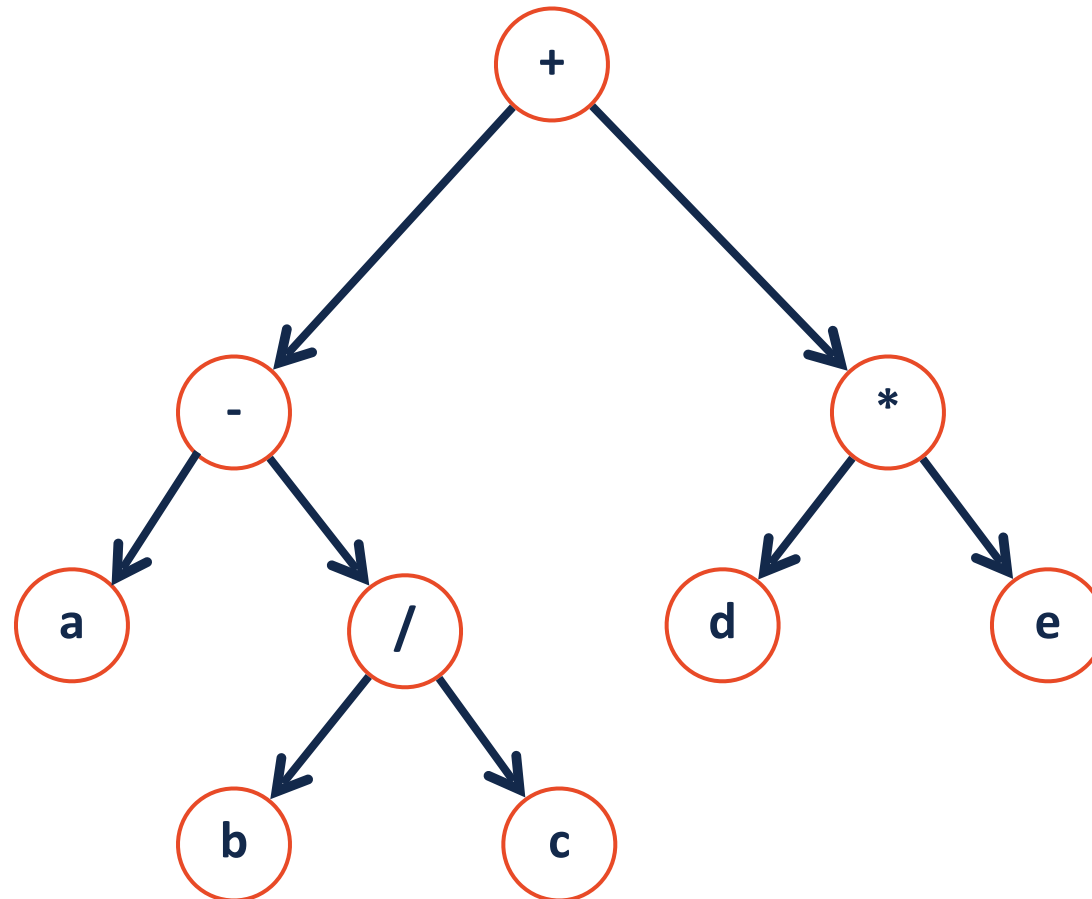
```
1 #pragma once
2
3 template <typename T>
4 class BinaryTree {
5     public:
6         /* ... */
7     private:
8         class TreeNode {
9             T & data;
10
11             TreeNode * left;
12
13             TreeNode * right;
14
15             TreeNode(T & data) :
16                 data(data), left(NULL),
17                 right(NULL) { }
18         };
19
20         TreeNode *root_;
21         /* ... */
22 };
```

# Visualizing trees



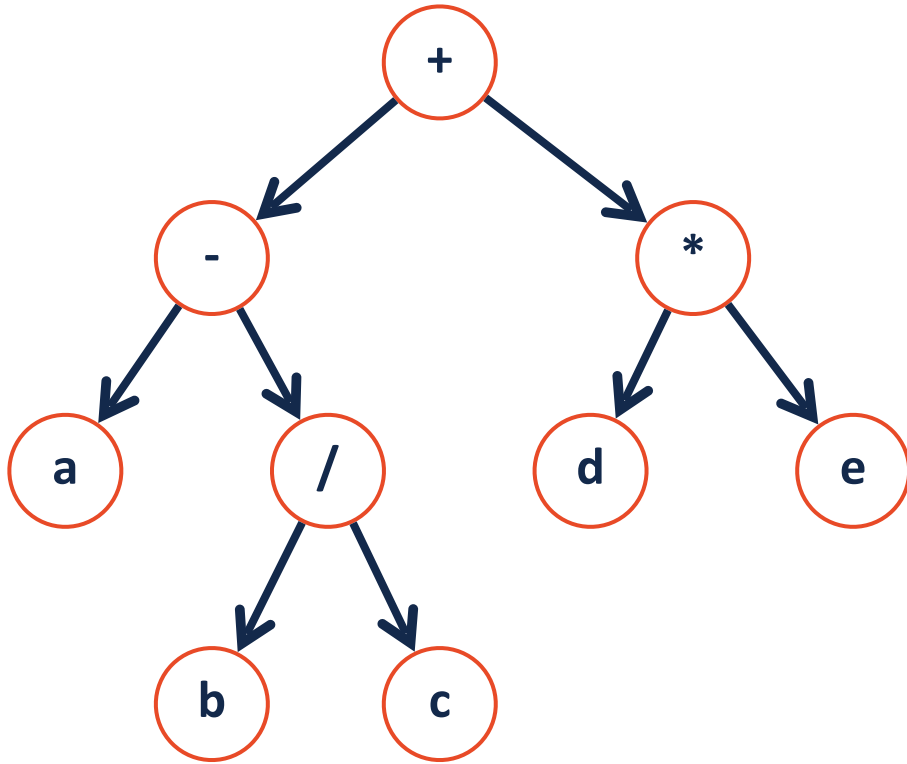
# Tree Traversal

A **traversal** of a tree  $T$  is an ordered way of visiting every node once.





# Traversals



```
1 template<class T>
2 void BinaryTree<T>::_____Order(TreeNode * root)
3 {
4
5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21 }
```