Data Structures

Linked List 3 & Array Lists

CS 225

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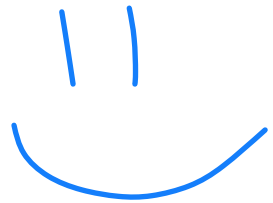
Department of Computer Science



Learning Objectives

Review the importance of index in a linked list

Finish implementing the List ADT (as a linked list)

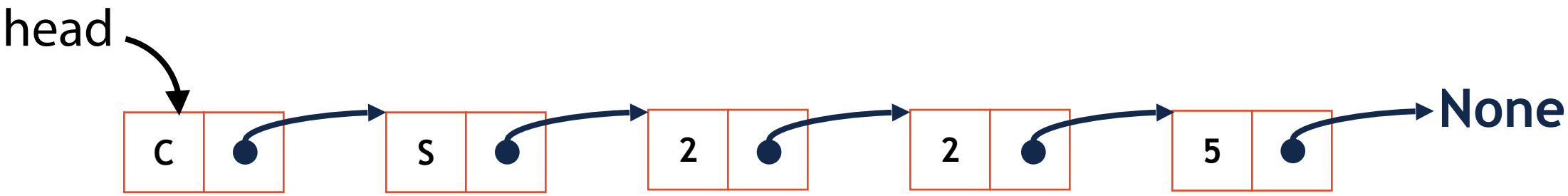
Discuss data variables for implementing array lists

Explore the List ADT (as an array list)

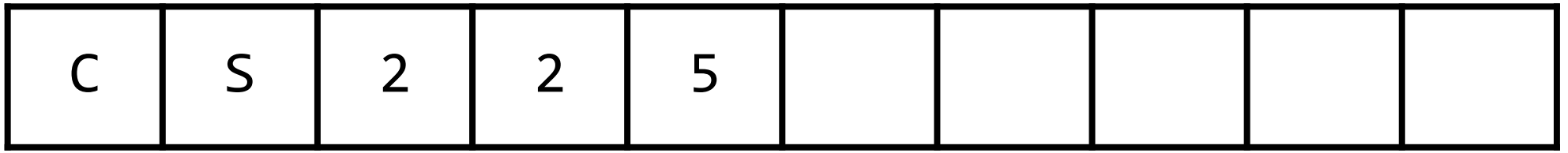
↳ Big O

List Implementations

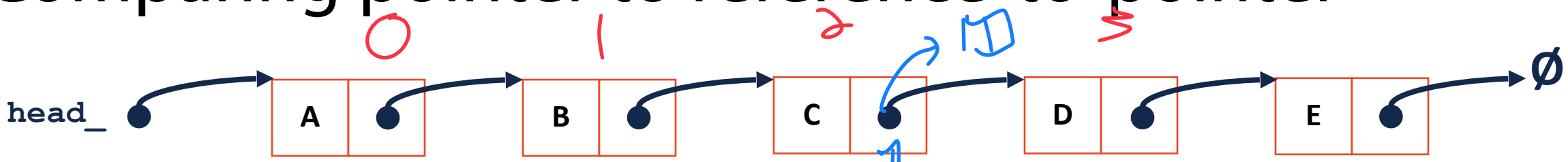
1. Linked List



2. Array List



Comparing pointer to reference-to-pointer



2 diff imp, see Wednesday!

```
ListNode * curr = _index(3);
```

↳ Not our implementation

↳ gives us address to Node @ 3

But doesn't let us change what node it is 3

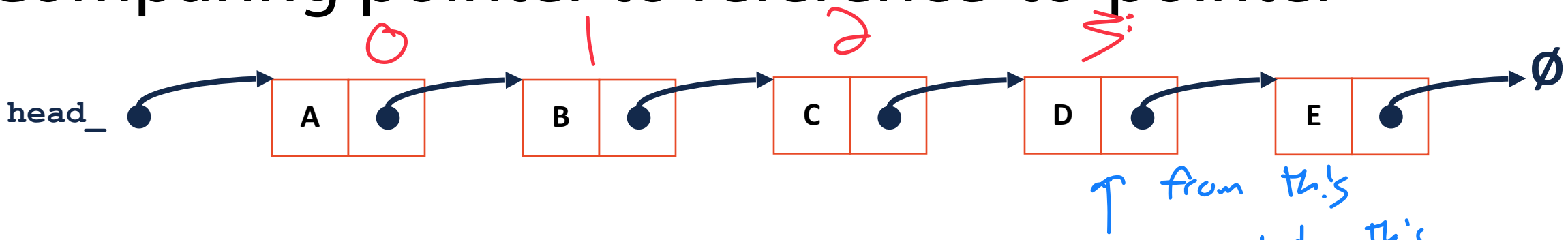
We chose not to do this!

```
ListNode *& curr = _index(3);
```

alias of

C → next

Comparing pointer to reference-to-pointer



curr

from this

to this

X

```
ListNode * curr = _index(3);
```

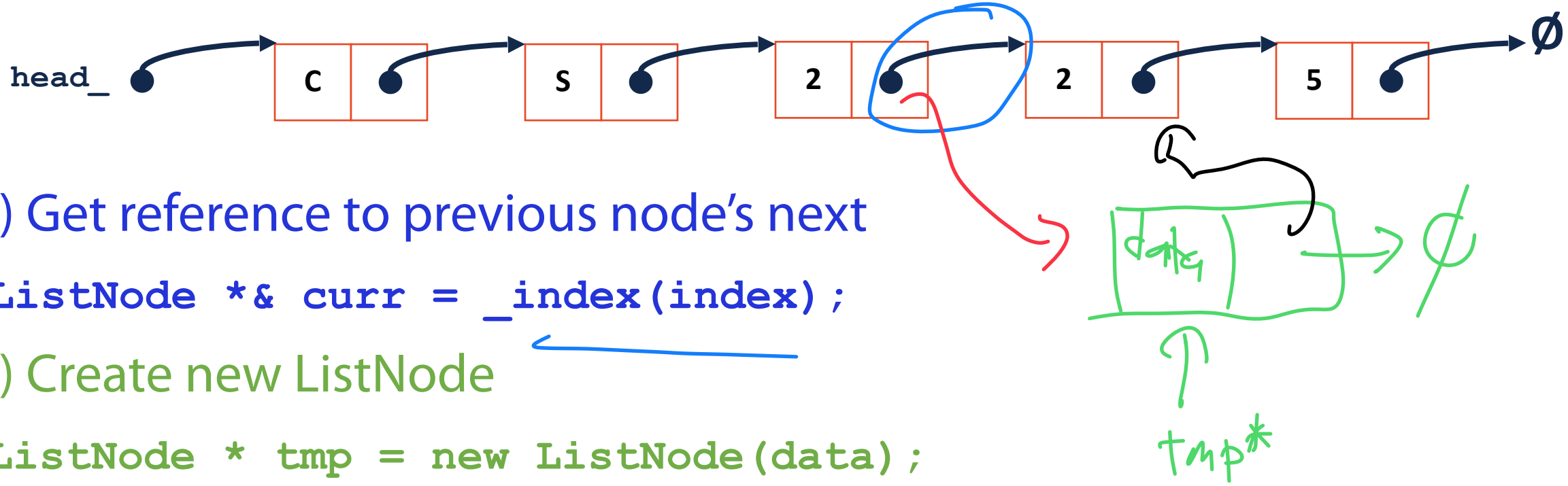
curr = new ListNode(X)

```
ListNode *& curr = _index(3);
```

alias of

C → next

Linked List: `insert(data, index)`



1) Get reference to previous node's next

```
ListNode *& curr = _index(index);
```

2) Create new ListNode

```
ListNode * tmp = new ListNode(data);
```

3) Update new ListNode's next

```
tmp->next = curr;
```

4) Modify the previous node to point to new ListNode

```
curr = tmp;
```

Lets compare...

```
1
2  template <typename T>
3  void List<T>::insertAtFront(const T& t)
4  {
5    ListNode *tmp = new ListNode(t);
6
7    tmp->next = head_;
8
9    head_ = tmp;
10
11 }
12
13
14
15
16
17
18
19
20
21
22
```

List.hpp

```
1
2  template <typename T>
3  void List<T>::insert(const T & data,
4  unsigned index) {
5
6
7
8    ListNode *& curr = _index(index);
9
10
11
12   ListNode * tmp = new ListNode(data);
13
14
15
16   tmp->next = curr;
17
18
19
20   curr = tmp;
21 }
22
```

List Random Access [ ]

L[7] → item @ 1

Given a list L, what operations can we do on L[ ]?

↳ Get Value

↳ Set (or change value)

What return type should this function have?

↳ pointer would work / reference will work

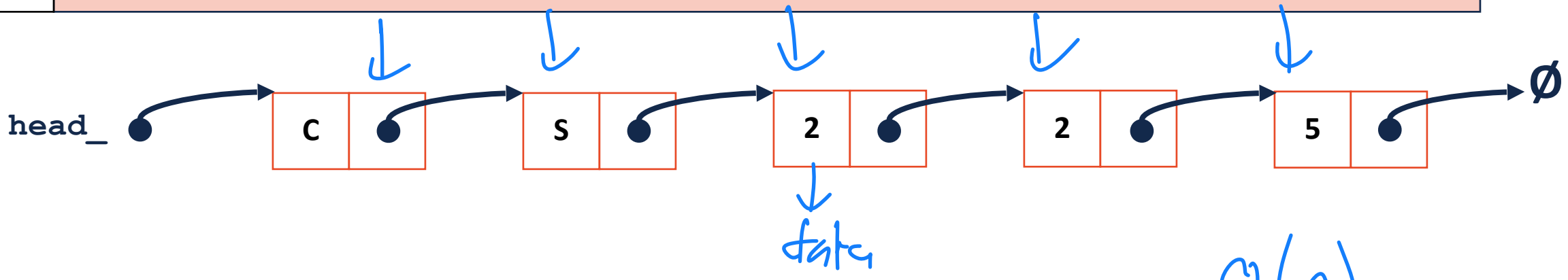
List.hpp

```
48 template <typename T>
49 T & List<T>::operator[](unsigned index) {
50
51
52
53
54
55
56
57
58 }
```

ListNode *val = _index(index);

O(n)

return val->data;

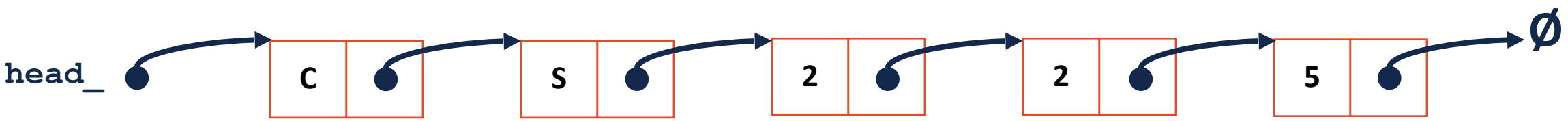


O(n)

List.hpp

```
48 template <typename T>
49 T & List<T>::operator[](unsigned index) {
50
51
52 ListNode *&new_node = _index(index);
53
54
55 return new_node->data;
56
57
58 }
```

Join Code: 225



What is the Big O of random access?

Linked List: remove(<parameters>)

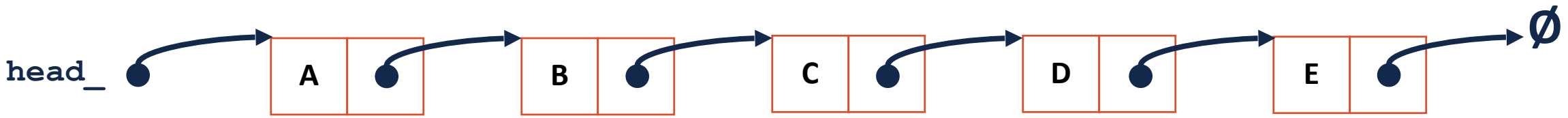
What input parameters make sense for remove?

↳ unsigned index

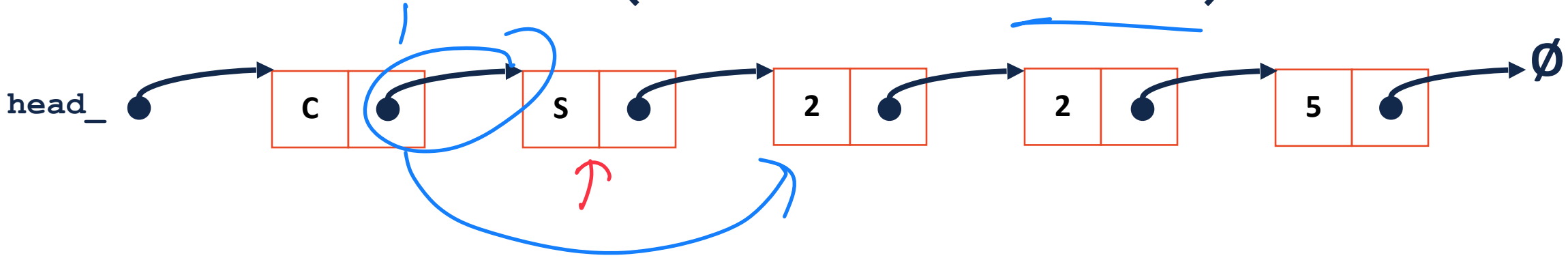
↳ T & value

↳ ListNode *& ref pointer

↳ remove bulk



Linked List: `remove(ListNode *& n)`



ListNode * tmp = n;

n = n -> next;

delete tmp;

List.hpp

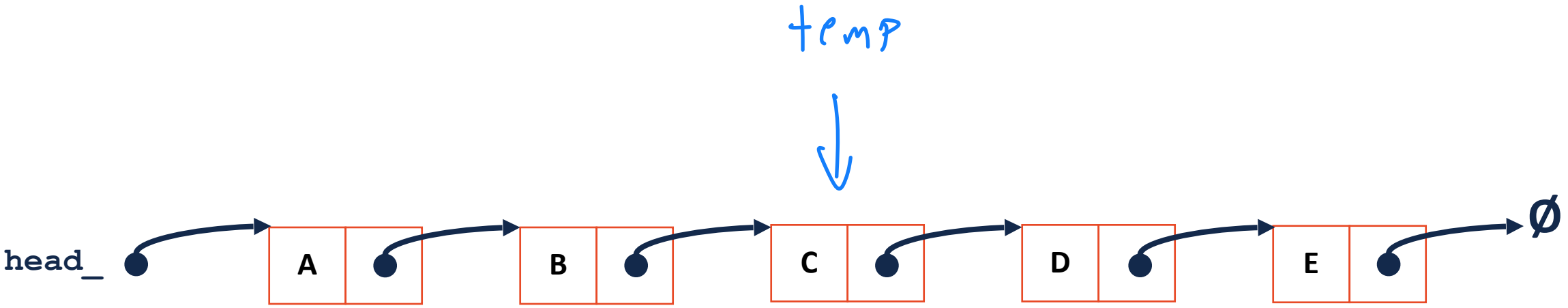
```
103 template <typename T>
104 T List<T>::remove(ListNode *& node) {
105
106 ListNode * temp = node;
107 node = node->next;
108 T data = temp->data;
109 delete temp;
110 return data;
111
112 }
```

// O(1)
// O(1)
// O(1)

ListNode *& t;

remove(t);

O(1)

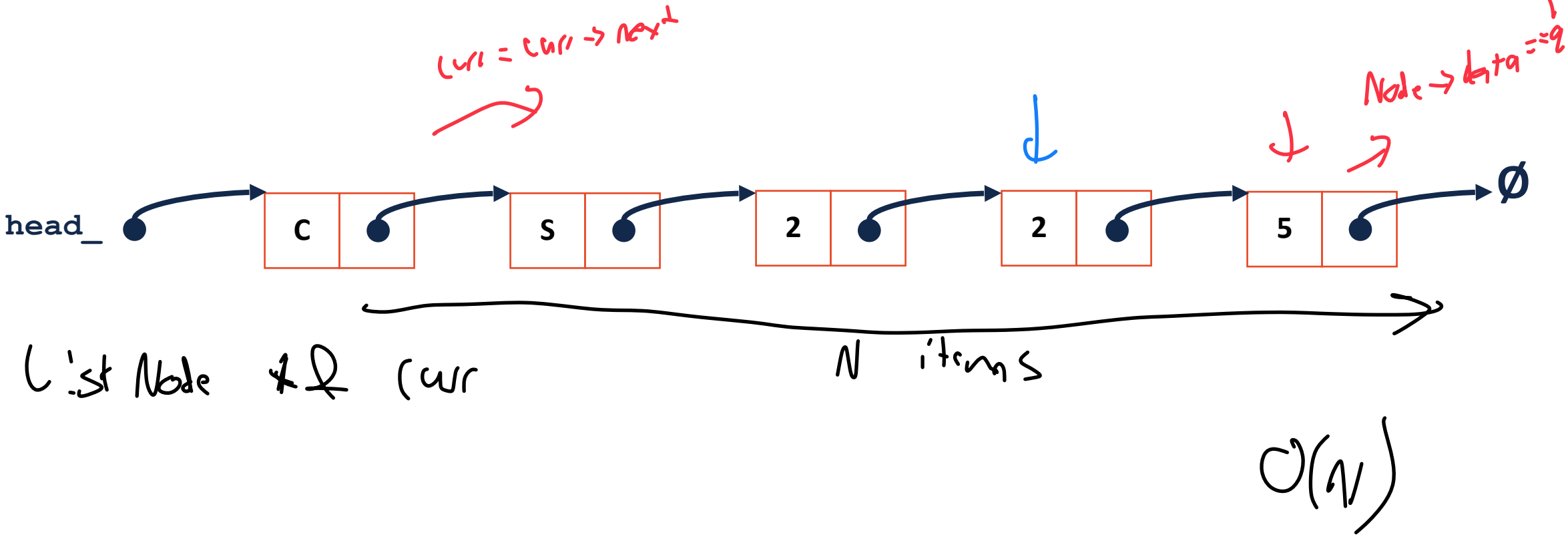


Linked List: remove(T & data)

List Node * prev

List Node * curr

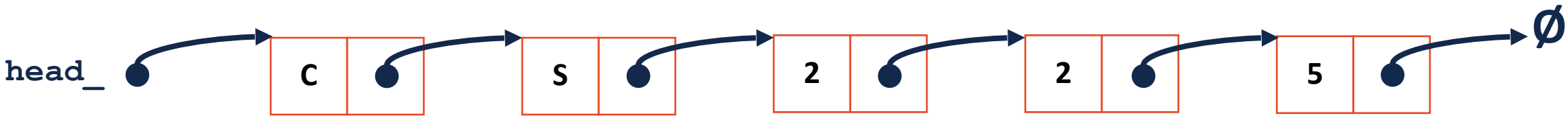
remove(5)



List Node *& curr

O(N)

Linked List: `remove`



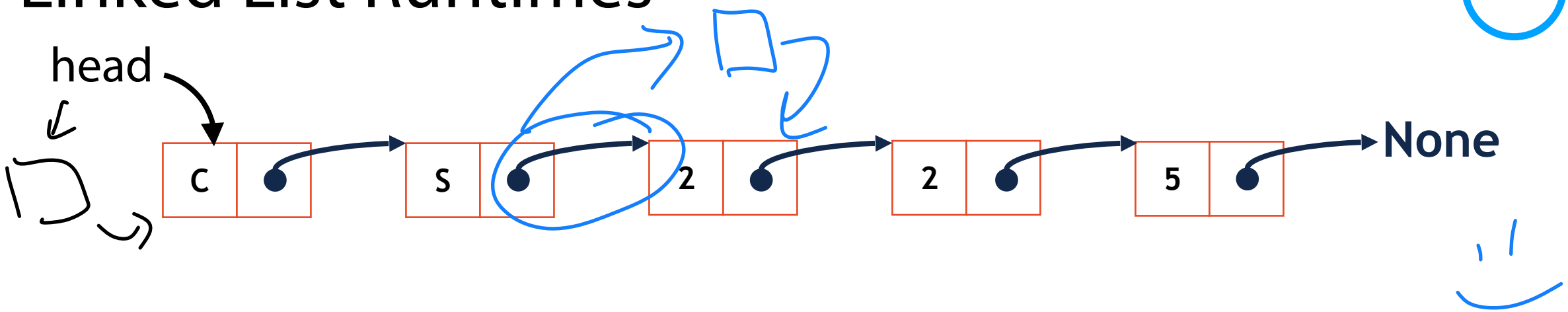
Running time for **`remove(ListNode *&)`**

↳ $O(1)$

Running time for **`remove(T & data)`**

↳ $O(n)$

Linked List Runtimes



	@Front	@RefPointer	@Index
Insert	O(1)	O(1)	O(n)
Delete	O(1)	O(1)	O(n)

Thinking critically about linked lists…

When would we use insert/delete on a reference to a pointer?

↳ If we have already run find()

↳ Loop over list

↳ O(n)

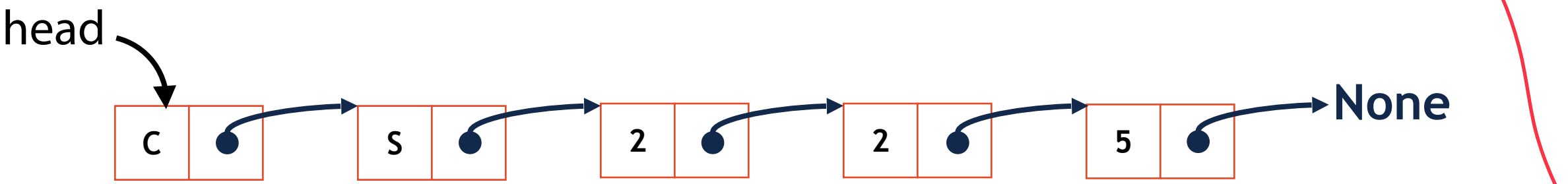
O(1) for each mod

What is the runtime to find an item of interest?

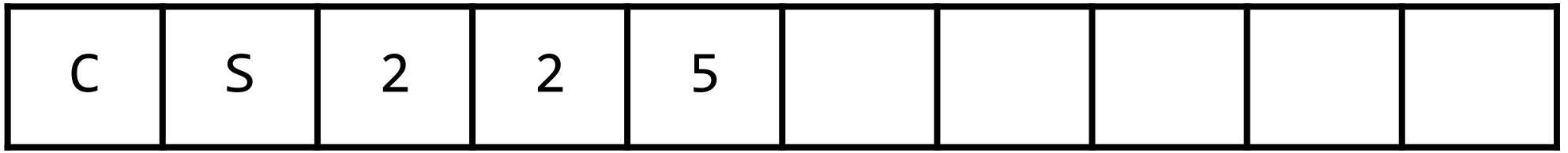
↳ O(n)

List Implementations

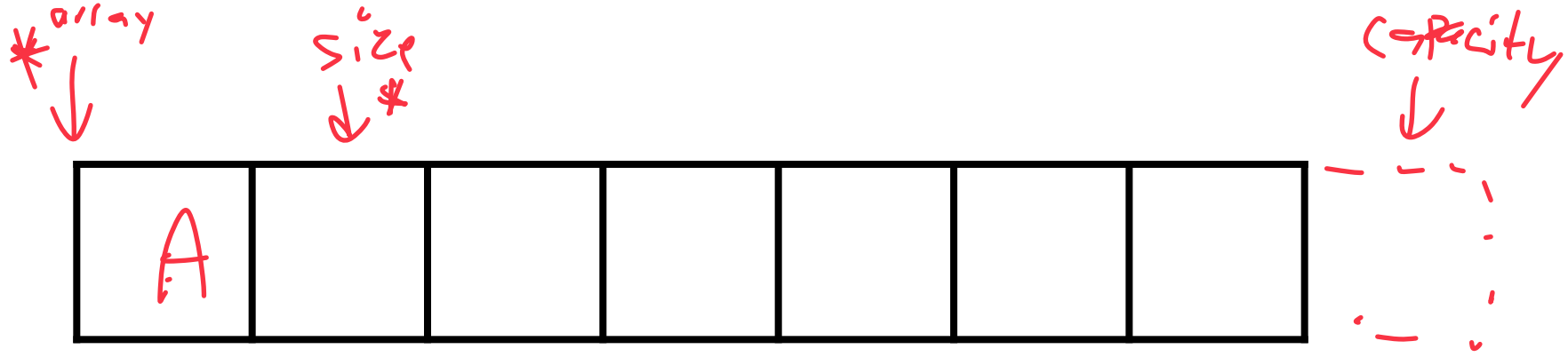
1. Linked List



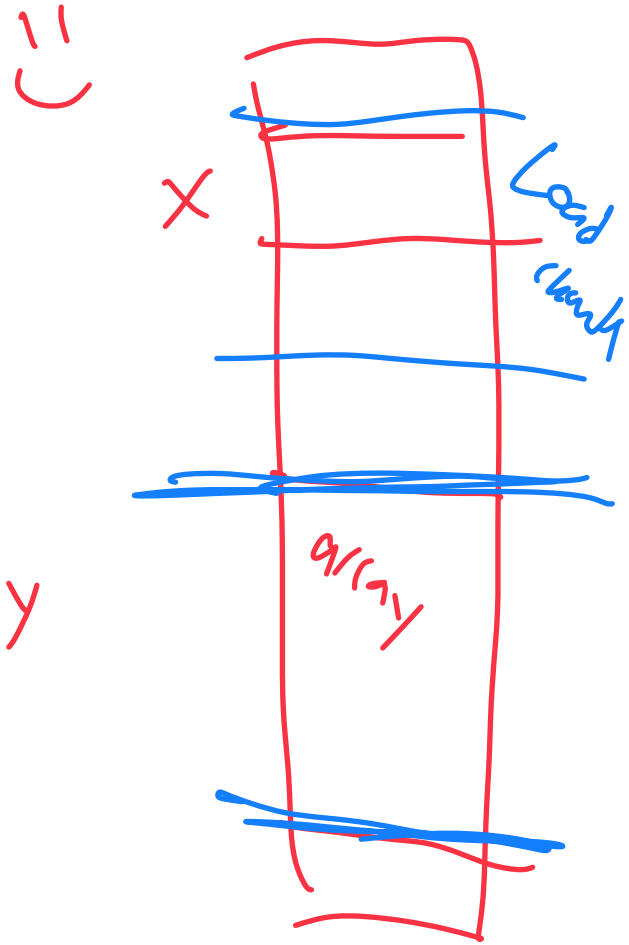
2. Array List



Array List



An array is allocated as continuous memory.



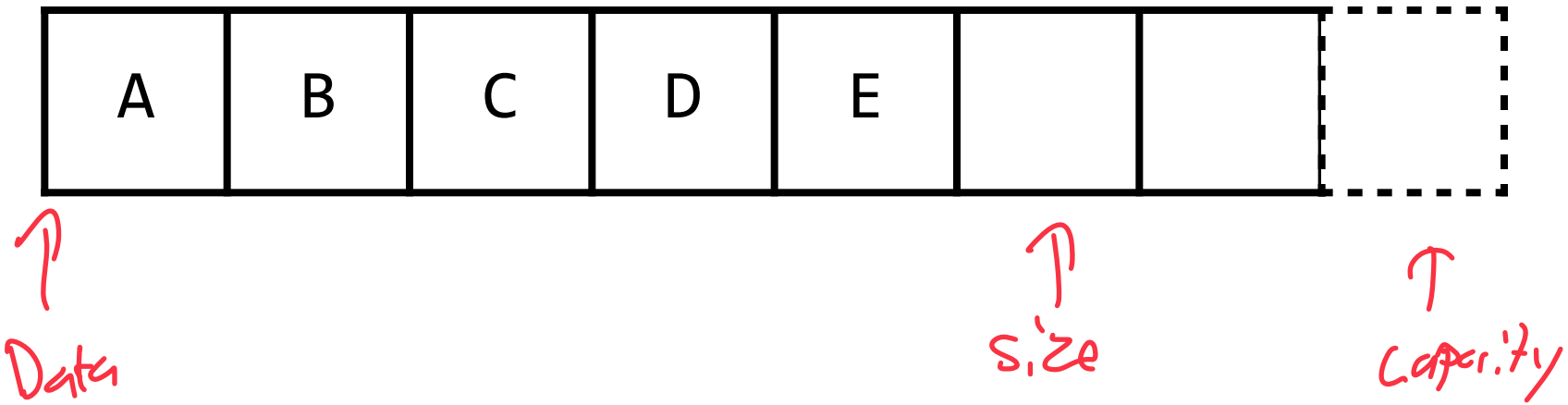
Three values are necessary for efficient array usage:

1) Data : the start location of my array

2) Size : How many things are in our array

3) Capacity : The max space allocated

Array List



In C++, vector is implemented as:

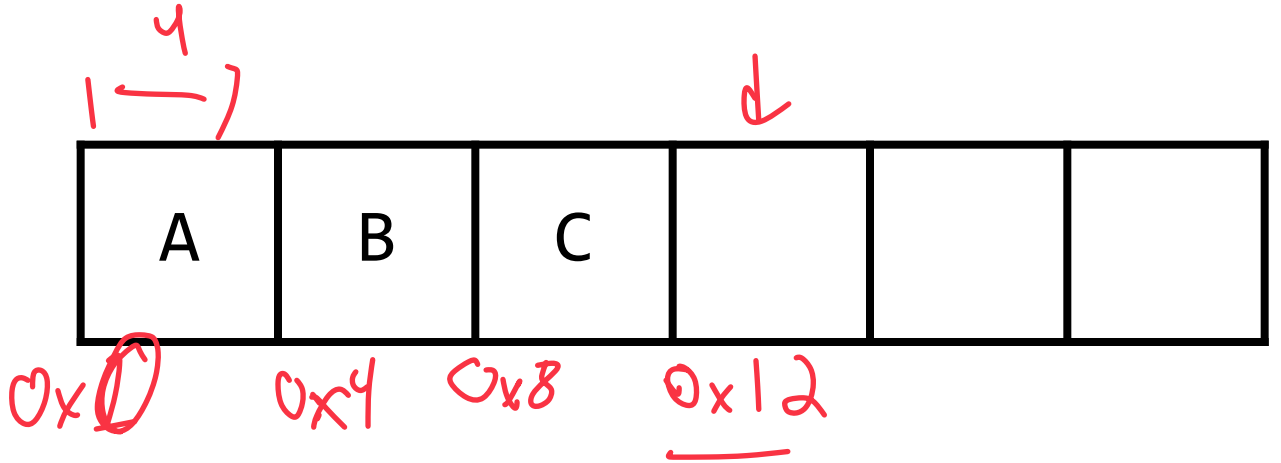
1) **Data:** Stored as a pointer to array start

2) **Size:** Stored as a pointer to the next available space

3) **Capacity:** Stored as a pointer past the end of the array

List.h

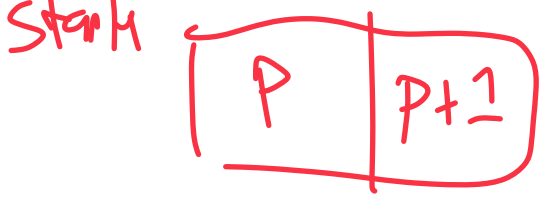
```
 1  #pragma once
 2
 3  template <typename T>
 4  class List {
 5  public:
 …      /* --- */
25  private:
26    T *data_;
27
28    T *size_;
29
30    T *capacity_;
 …
        /* --- */
    };
```



size of (4) = 4

int * p

p++;



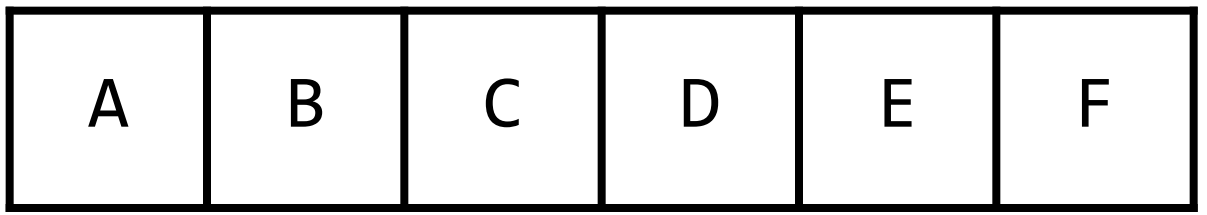
If I want to know the number of items in the array:

size - data = 3 ← O(1)

0x12 0x0

List.h

```
1   #pragma once
2
3   template <typename T>
4   class List {
5   public:
        /* --- */
…
25  private:
26    T *data_;
27
28    T *size_;
29
30    T *capacity_;
…
        /* --- */
    };
```



capacity

size

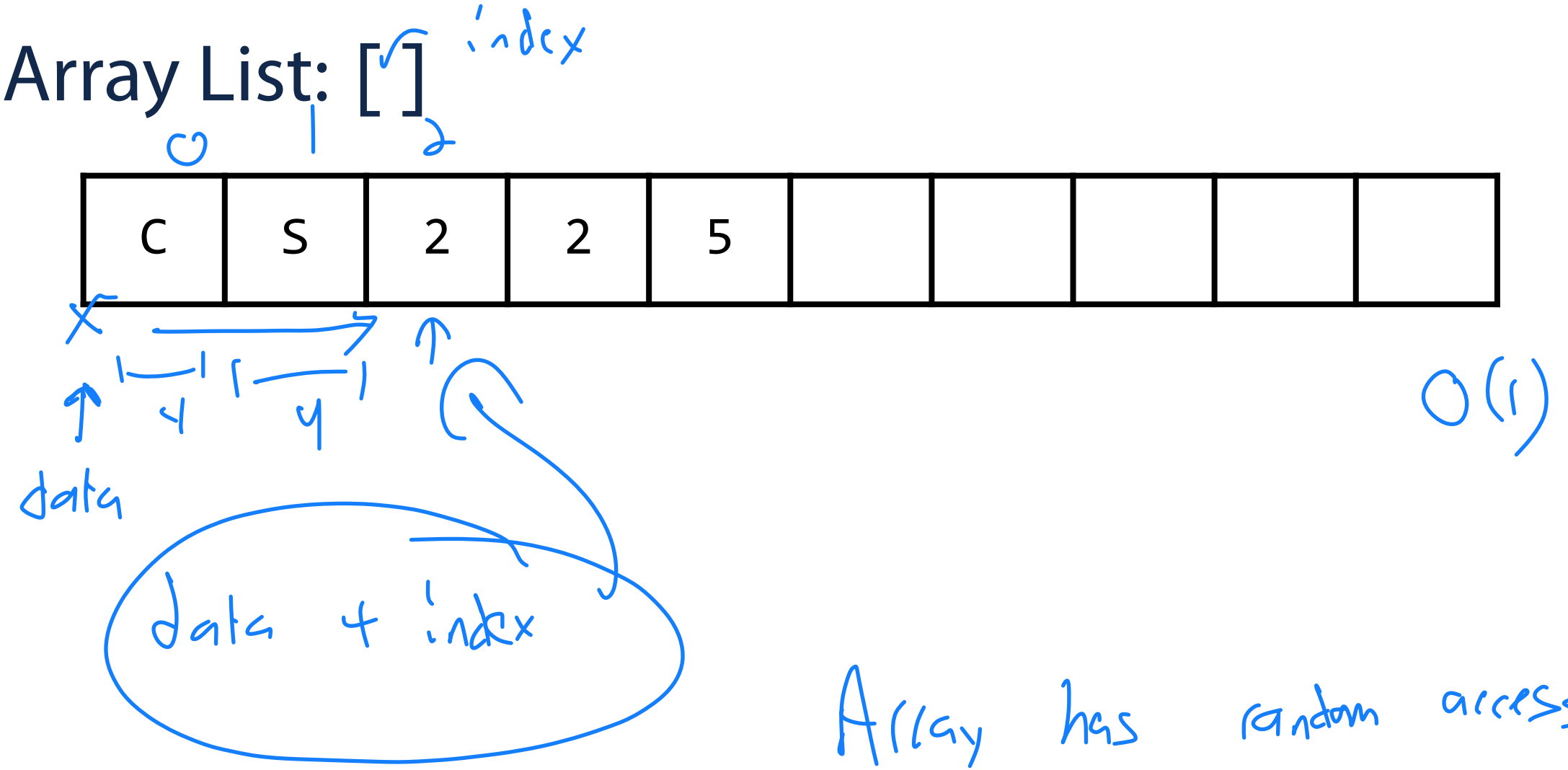
How do I know if I'm at capacity?

size == capacity, my array is full

↳ Dont insert

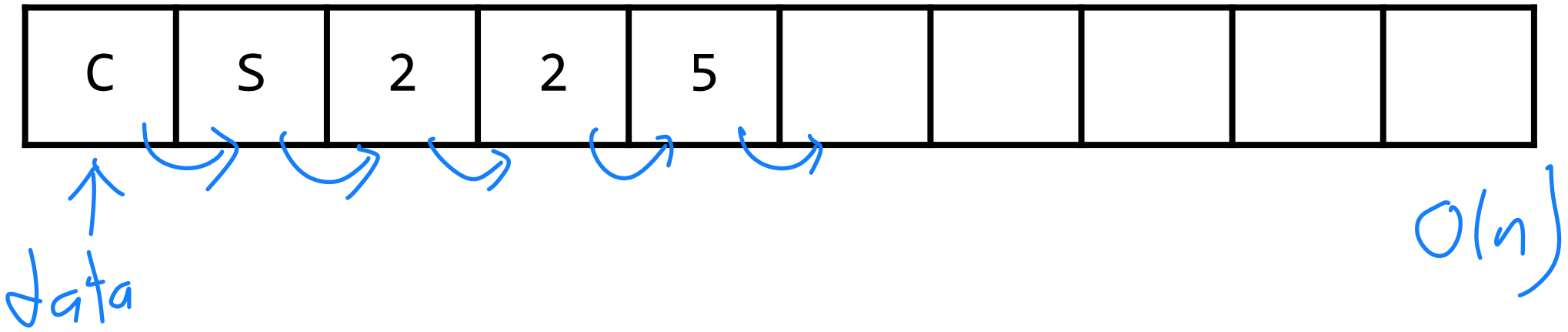
O(1)

Array List: [ ]



Array has random access

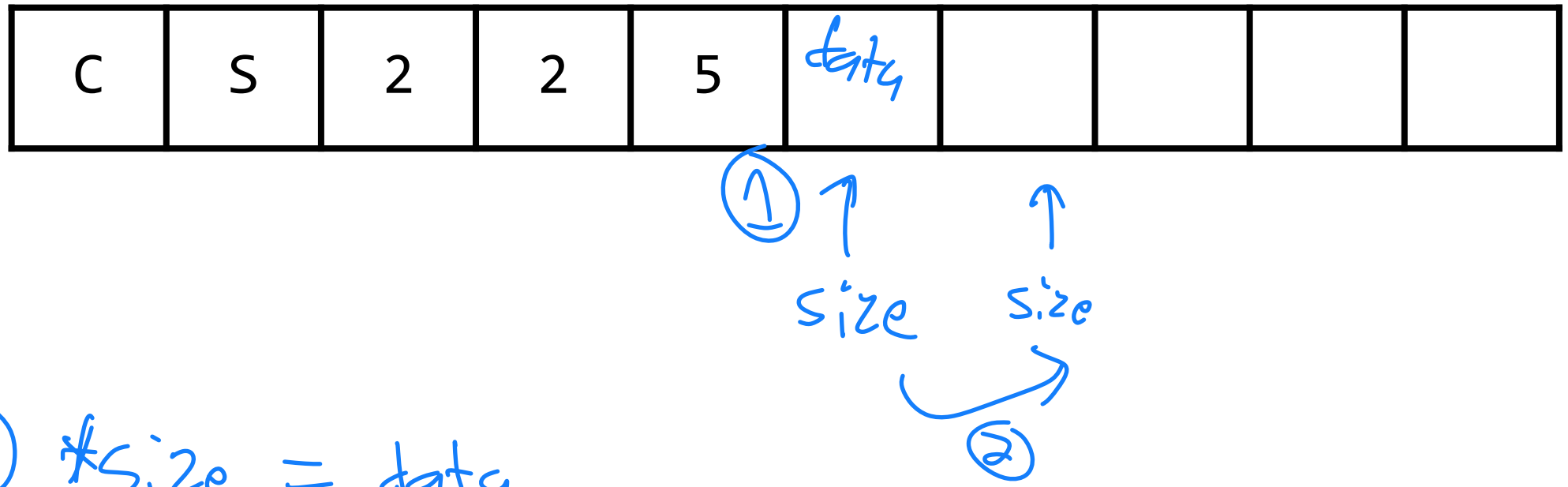
Array List: `insertFront(data)`



insert at front moves all items

Array List: `insertBack(data)`

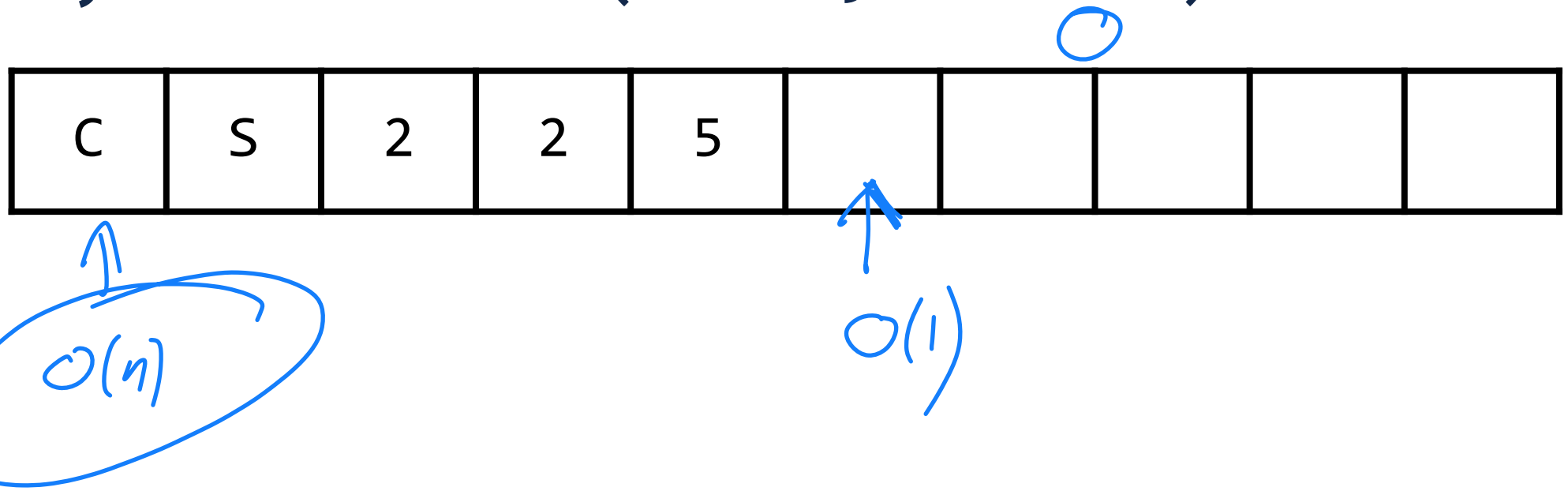
\# Not at capacity



1. `*size = data`
2. `size++;`

O(1)

Array List: `insert(data, index)`



Array List: addspace(data)