

# Data Structures and Algorithms

## Hashing 2

CS 225

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# Learning Objectives

Review fundamentals of hash tables

Introduce closed hashing approaches to hash collisions

Determine when and how to resize a hash table

Justify when to use different index approaches

# A Hash Table based Dictionary

**User Code (is a map):**

```
1 Dictionary<KeyType, ValueType> d;  
2 d[k] = v;
```

A **Hash Table** consists of three things:

1. A hash function
2. A data storage structure
3. A method of addressing *hash collisions*

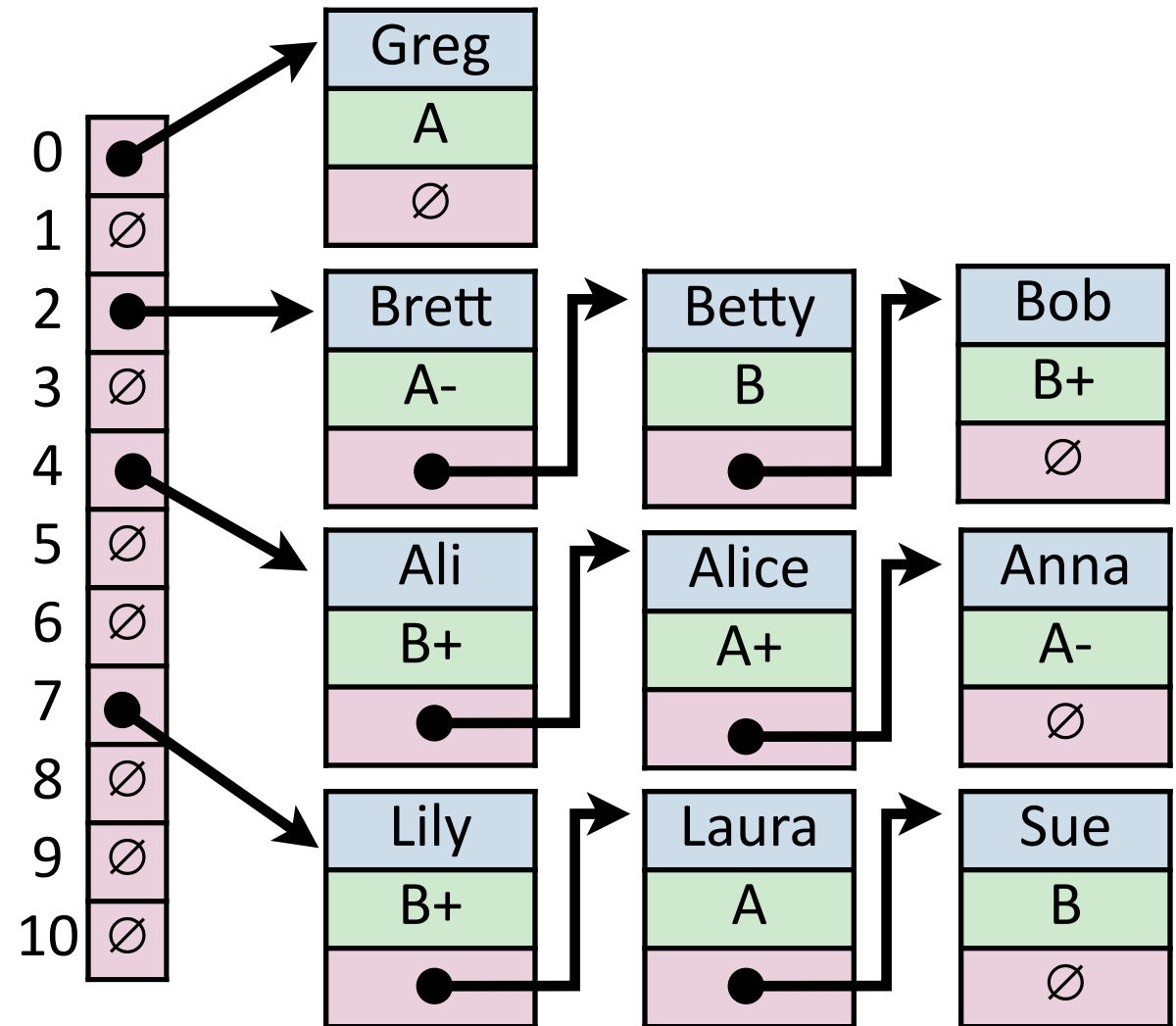
# Open vs Closed Hashing

Addressing hash collisions depends on your storage structure.

- **Open Hashing:** store  $k, v$  pairs externally
  
  
  
  
  
  
  
  
  
  
- **Closed Hashing:** store  $k, v$  pairs in the hash table

# Hash Table (Separate Chaining)

Key	Value	Hash
Bob	B+	2
Anna	A-	4
Alice	A+	4
Betty	B	2
Brett	A-	2
Greg	A	0
Sue	B	7
Ali	B+	4
Laura	A	7
Lily	B+	7



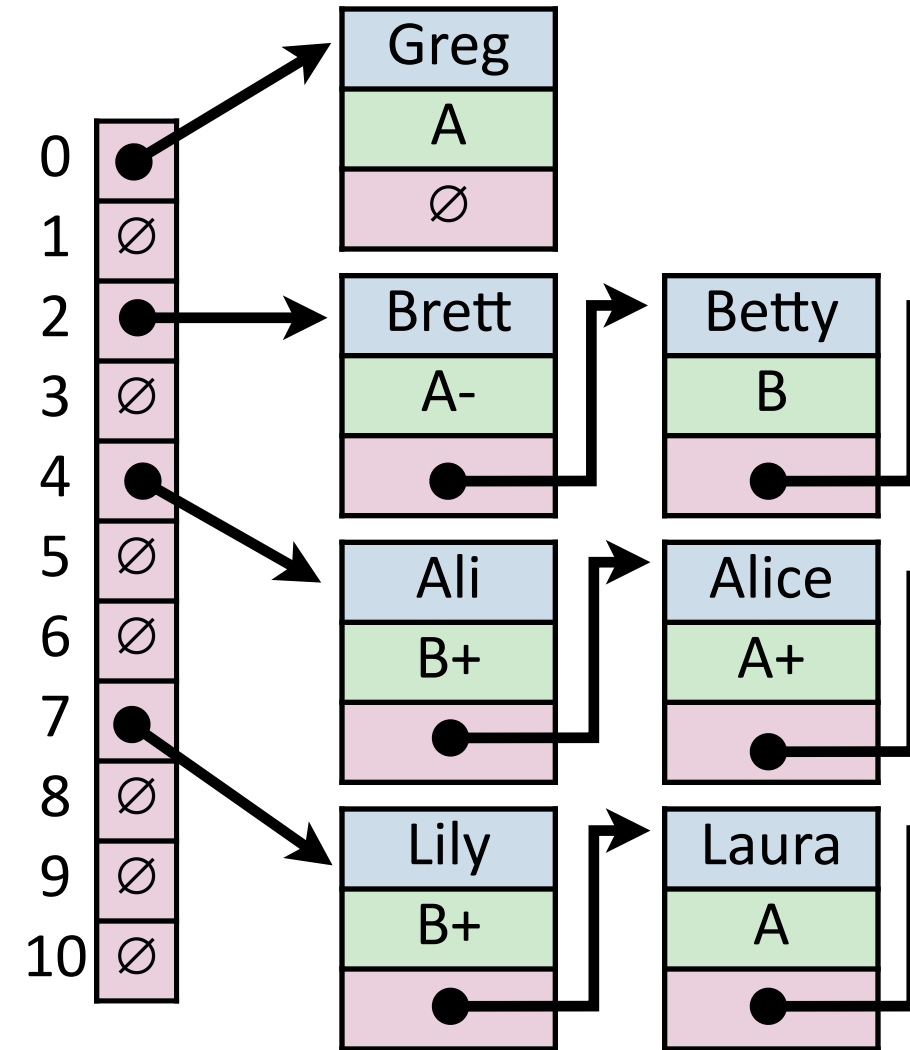
# Hash Table (Separate Chaining)

For hash table of size  $m$  and  $n$  elements:

Find runs in: \_\_\_\_\_

Insert runs in: \_\_\_\_\_

Remove runs in: \_\_\_\_\_



# Hash Table

Worst-Case behavior is bad — but what about randomness?

1) **Fix  $h$** , our hash, and assume it is good for *all keys*:

2) Create a *universal hash function family*:

# Simple Uniform Hashing Assumption

Given table of size  $m$ , a simple uniform hash,  $h$ , implies

$$\forall k_1, k_2 \in U \text{ where } k_1 \neq k_2, \Pr(h[k_1] = h[k_2]) = \frac{1}{m}$$

**Uniform:**

**Independent:**



# Separate Chaining Under SUHA

Given table of size  $m$  and  $n$  inserted objects

**Claim:** Under SUHA, expected length of chain is  $\frac{n}{m}$

# Separate Chaining Under SUHA

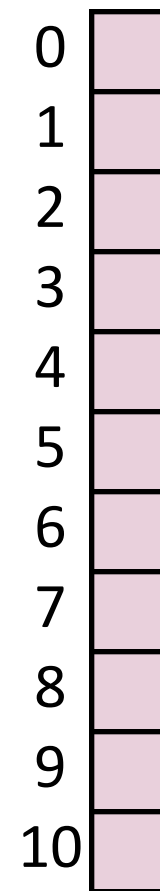


**Under SUHA, a hash table of size  $m$  and  $n$  elements:**

Find runs in: \_\_\_\_\_.

Insert runs in: \_\_\_\_\_.

Remove runs in: \_\_\_\_\_.



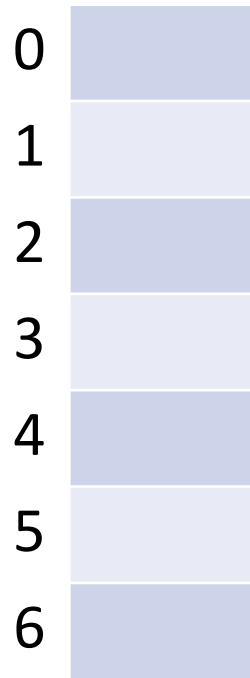
# Collision Handling: Probe-based Hashing

$$S = \{ 1, 8, 15 \}$$

$$h(k) = k \% 7$$

$$|S| = n$$

$$|\text{Array}| = m$$

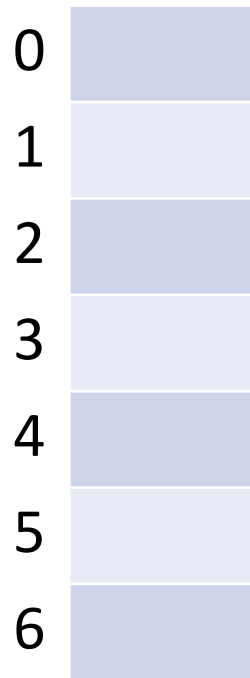


# Collision Handling: Linear Probing

$S = \{ 16, 8, 4, 13, 29, 11, 22 \}$        $|S| = n$

$h(k) = k \% 7$

$|\text{Array}| = m$



$h(k, i) = (k + i) \% 7$

Try  $h(k) = (k + 0) \% 7$ , if full...

Try  $h(k) = (k + 1) \% 7$ , if full...

Try  $h(k) = (k + 2) \% 7$ , if full...

Try ...

# Collision Handling: Linear Probing

$S = \{ 16, 8, 4, 13, 29, 11, 22 \}$        $|S| = n$

$h(k, i) = (k + i) \% 7$        $|Array| = m$

0	22
1	8
2	16
3	29
4	4
5	11
6	13

`_find(29)`

# Collision Handling: Linear Probing

$S = \{ 16, 8, 4, 13, 29, 11, 22 \}$        $|S| = n$

$h(k, i) = (k + i) \% 7$        $|\text{Array}| = m$

0	22
1	8
2	16
3	29
4	4
5	11
6	13

\_remove(16)

# A Problem w/ Linear Probing

**Primary clustering:**



**Description:**

**Remedy:**

# Collision Handling: Quadratic Probing

$S = \{ 16, 8, 4, 13, 29, 12, 22 \}$        $|S| = n$

$h(k) = k \% 7$

$|\text{Array}| = m$

0	
1	8
2	16
3	
4	4
5	
6	13

$h(k, i) = (k + i*i) \% 7$

Try  $h(k) = (k + 0) \% 7$ , if full...

Try  $h(k) = (k + 1*1) \% 7$ , if full...

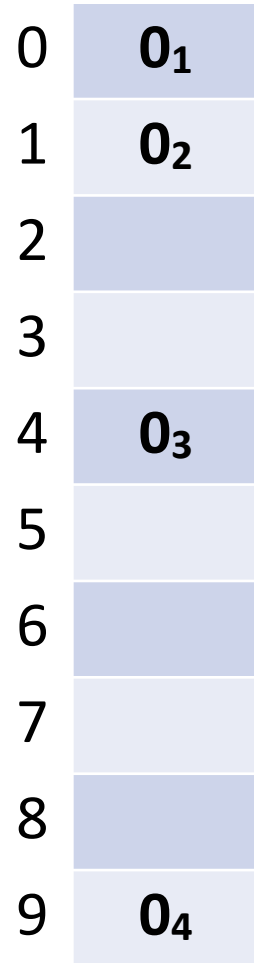
Try  $h(k) = (k + 2*2) \% 7$ , if full...

Try ...



# A Problem w/ Quadratic Probing

**Secondary clustering:**



**Description:**

**Remedy:**

# Collision Handling: Double Hashing

$$S = \{ 16, 8, 4, 13, 29, 11, 22 \} \quad |S| = n$$

$$h_1(k) = k \% 7$$

$$|\text{Array}| = m$$

$$h_2(k) = 5 - (k \% 5)$$

$$h(k, i) = (h_1(k) + i * h_2(k)) \% 7$$

Try  $h(k) = (k + 0 * h_2(k)) \% 7$ , if full...

Try  $h(k) = (k + 1 * h_2(k)) \% 7$ , if full...

Try  $h(k) = (k + 2 * h_2(k)) \% 7$ , if full...

Try ...

0	
1	8
2	16
3	
4	4
5	
6	13

# Running Times *(Don't memorize these equations, no need.)*

*(Expectation under SUHA)*

## Open Hashing:

insert: \_\_\_\_\_.

find/ remove: \_\_\_\_\_.

## Closed Hashing:

insert: \_\_\_\_\_.

find/ remove: \_\_\_\_\_.

# Running Times *(Don't memorize these equations, no need.)*

*The expected number of probes for find(key) under SUHA*

## Linear Probing:

- Successful:  $\frac{1}{2}(1 + 1/(1-\alpha))$
- Unsuccessful:  $\frac{1}{2}(1 + 1/(1-\alpha))^2$

## Double Hashing:

- Successful:  $1/\alpha * \ln(1/(1-\alpha))$
- Unsuccessful:  $1/(1-\alpha)$

## Separate Chaining:

- Successful:  $1 + \alpha/2$
- Unsuccessful:  $1 + \alpha$

**Instead, observe:**

- **As  $\alpha$  increases:**

- **If  $\alpha$  is constant:**

# Running Times

*The expected number of probes for find(key) under SUHA*

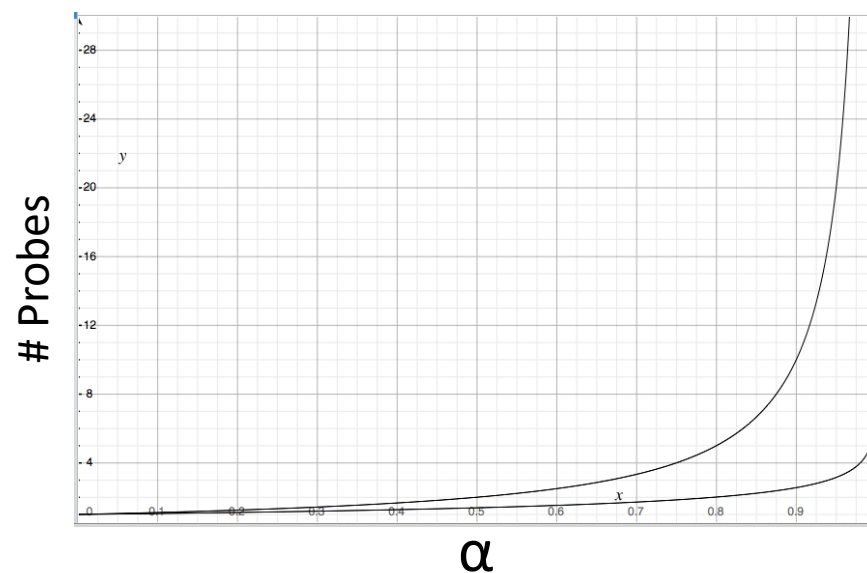
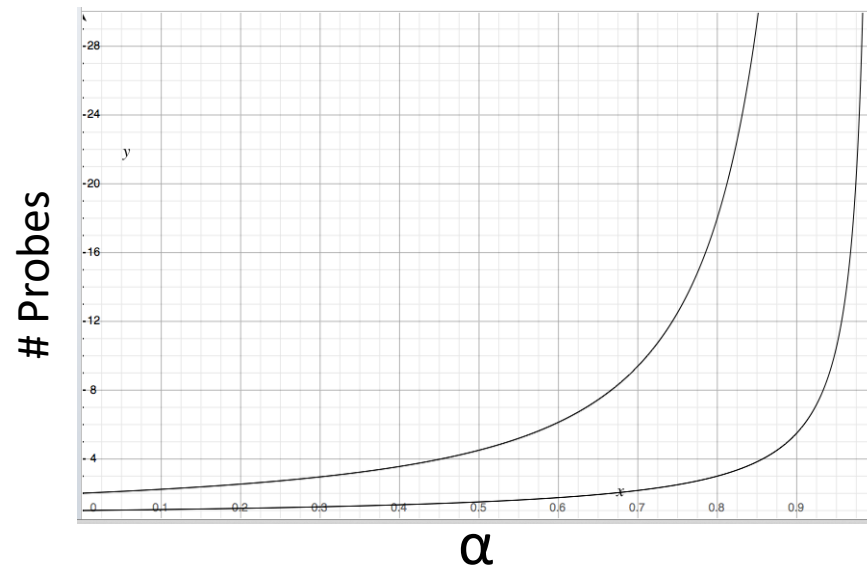
## Linear Probing:

- Successful:  $\frac{1}{2}(1 + 1/(1-\alpha))$
- Unsuccessful:  $\frac{1}{2}(1 + 1/(1-\alpha))^2$

## Double Hashing:

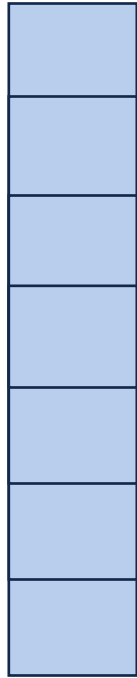
- Successful:  $1/\alpha * \ln(1/(1-\alpha))$
- Unsuccessful:  $1/(1-\alpha)$

**When do we resize?**



# Resizing a hash table

How do you resize?



## **Which collision resolution strategy is better?**

- Big Records:
- Structure Speed:

## **What structure do hash tables implement?**

## **What constraint exists on hashing that doesn't exist with BSTs?**

## **Why talk about BSTs at all?**

# Running Times

	Hash Table	AVL	Linked List
<b>Find</b>	Expectation*:  Worst Case:		
<b>Insert</b>	Expectation*:  Worst Case:		
<b>Storage Space</b>			