Data Structures and Algorithms Cardinality Sketches CS 225 Dec 2, 2022

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Reminder: Final Exam Scheduling

You can sign up now for the final exam

There are no extensions or make-ups for the final exam

Learning Objectives



Introduce the concept of cardinality and cardinality estimation

Demonstrate the relationship between cardinality and similarity

Introduce the MinHash Sketch for set similarity detection

Bloom Filters

A probabilistic data structure storing a set of values

Has three key properties:

k, number of hash functions n, expected number of insertions m, filter size in bits

Expected false positive rate:

$$\left(1 - \left(1 - \frac{1}{m}\right)^{nk}\right)^k \approx \left(1 - e^{\frac{-nk}{m}}\right)^k$$

Optimal accuracy when:

$$k^* = \ln 2 \cdot \frac{m}{n}$$

 $h_{\{1,2,3,...,k\}}$

Count-Min Sketch

A probabilistic data structure storing a set of values

Has **four** key properties:

k, number of hash functions
n, expected number of insertions
m, filter size in *registers*b, number of bits per register

Minimal increase reduces overcounting by identifying collisions.

(Count returned by sketch) \geq (True count of the query)

$h_{\{1,2\}}$	2,3,,k
1	0

1

0	3	1	0
0	0	2	3
2	0	1	0
3	1	0	1
0	0	2	2

The hidden problem with sketches...



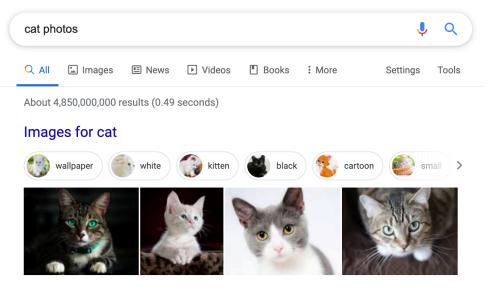
Cardinality

Cardinality is a measure of how many unique items are in a set

2
4
9
3
7
9
7
8
5
6

Cardinality

Sometimes its not possible or realistic to count all objects!



Estimate: 60 billion — 130 trillion

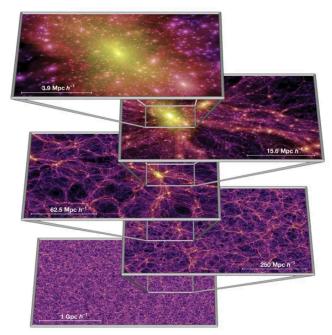
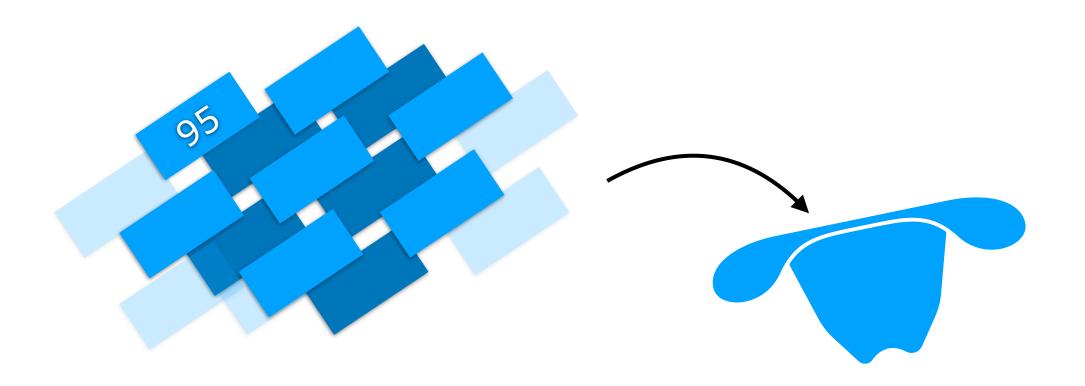


Image: https://doi.org/10.1038/nature03597

5581
8945
6145
8126
3887
8925
1246
8324
4549
9100
5598
8499
8970
3921
8575
4859
4960
42
6901
4336
9228
3317
399
6925
2660
2314

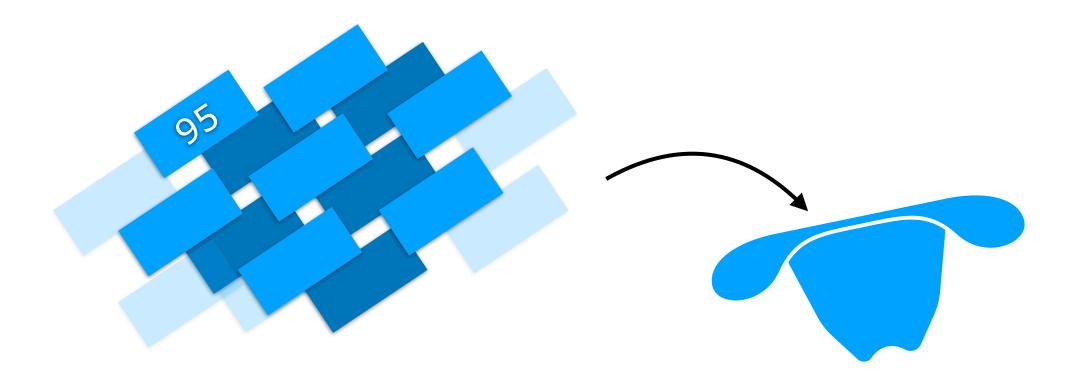
Imagine I fill a hat with numbered cards and draw one card out at random.

If I told you the value of the card was 95, what have we learned?



Imagine I fill a hat with a random subset of numbered cards from 0 to 999

If I told you that the **minimum** value was 95, what have we learned?



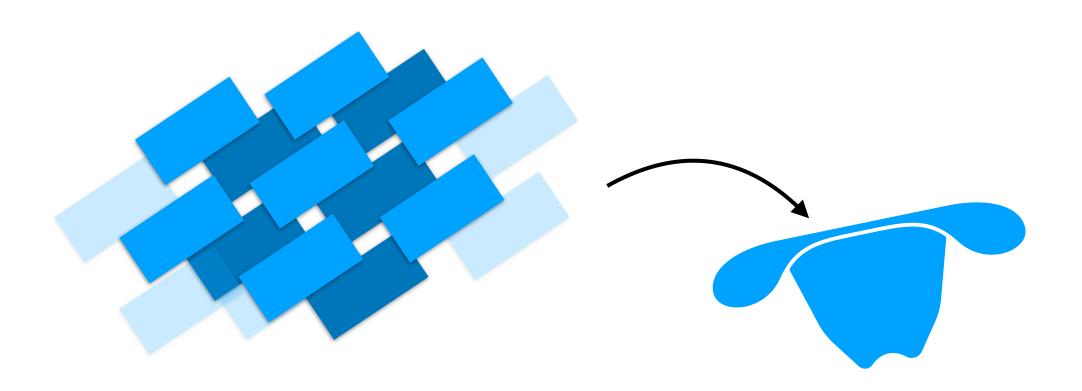
Imagine we have multiple sets (multiple minimums).



Let min = 95. Can we estimate *N*, the cardinality of the set?



Why do we care about "the hat problem"?



Why do we care about "the hat problem"?

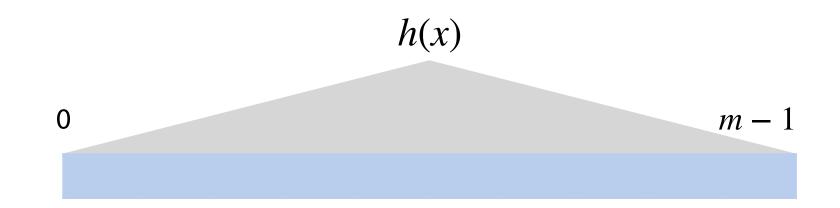
m possible minima Key Value Universe of card sets

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Now imagine we have a SUHA hash *h* over a range *m*.

Here a hash insert is equivalent to adding a card to our hat!

Now storing only the minimum hash value is a **sketch!**



Let $M = min(X_1, X_2, ..., X_N)$ where each $X_i \in [0, 1]$ is an independent random variable

Claim:
$$\mathbf{E}[M] = \frac{1}{N+1}$$

0

Claim:
$$E[M] = \frac{1}{N+1}$$
 $N \approx \frac{1}{M} - 1$

Attempt 20.2530.8390.3270.6550.491

Attempt 3

|--|

Consider an N + 1 draw:

 $\mathbf{0}$

$$X_1 X_2 X_3 \dots X_N X_{N+1} \qquad M = \min_{1 \le i \le N} X_N$$

Claim:
$$\mathbf{E}[M] = Pr(X_{N+1} < \min_{1 \le i \le N} X_i)$$

Cardinality Sketch $I_1 \quad I_2 \quad I_3 \quad I_N \quad I_{N+1}$ Image: Image:

Claim: $\mathbf{E}[M] = Pr(X_{N+1} < \min_{1 \le i \le N} X_i)$ $I_i = \begin{cases} 1 & \text{if } X_i < \min_{j \ne i} X_j \\ 0 & \text{otherwise} \end{cases}$

$$\Pr\left(X_{N+1} < \min_{1 \le i \le N} X_i\right) = \mathbf{E}[I_{N+1}] = \frac{1}{N+1} = \mathbf{E}[M]$$

$$0 \qquad M \qquad 1$$

0

The minimum hash is a valid sketch of a dataset but can we do better?

Claim: Taking the k^{th} -smallest hash value is a better sketch!

Claim:
$$\mathbf{E}[\mathbf{M}_k] = \frac{k}{N+1}$$

$$0 \quad \underbrace{M_1}_{1} \quad \underbrace{M_2}_{1} \quad \underbrace{M_3}_{1} \quad \ldots \quad \underbrace{M_k}_{1}$$

 M_1

Claim: Taking the k^{th} -smallest hash value is a better sketch!

 M_2

Claim:
$$\mathbf{E}[M_k] = \frac{k}{N+1}$$

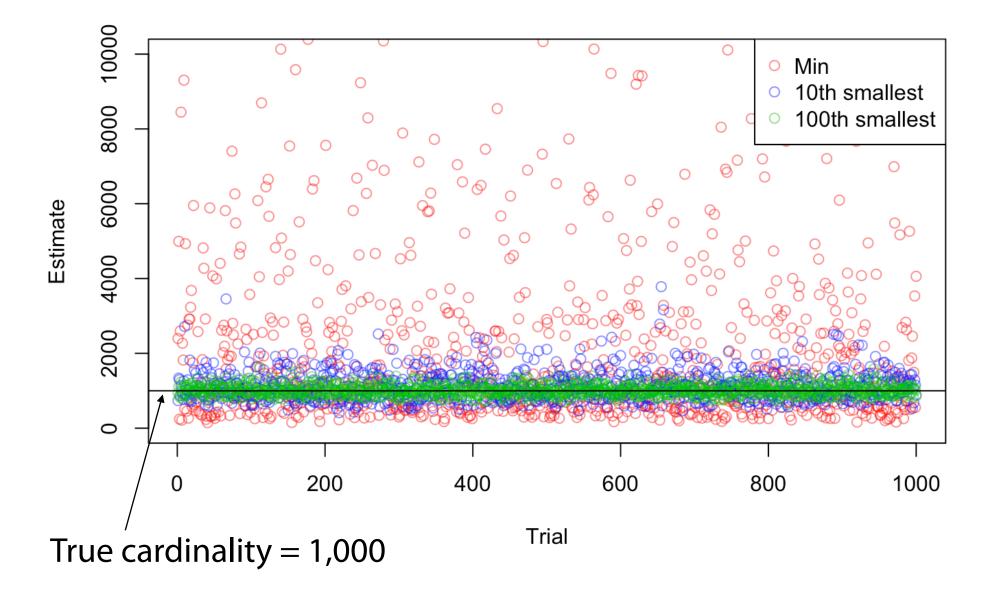
= $\left[\mathbf{E}[M_1] + (\mathbf{E}[M_2] - \mathbf{E}[M_1]) + \dots + (\mathbf{E}[M_k] - \mathbf{E}[M_{k-1}])\right] \cdot \frac{1}{k}$

 M_3 ...

 M_{k-1}

 M_k

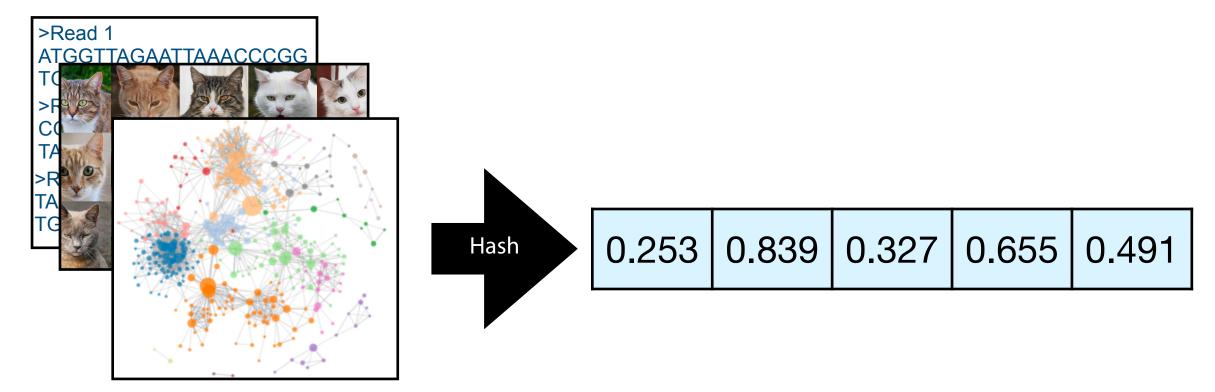
Cardinality



Cardinality



Given any dataset and a SUHA hash function, we can estimate the number of unique items by tracking the minimum hash values.



Applied Cardinalities

Cardinalities

|A|

 $|A \cup B|$

 $|A \cap B|$

Set similarities

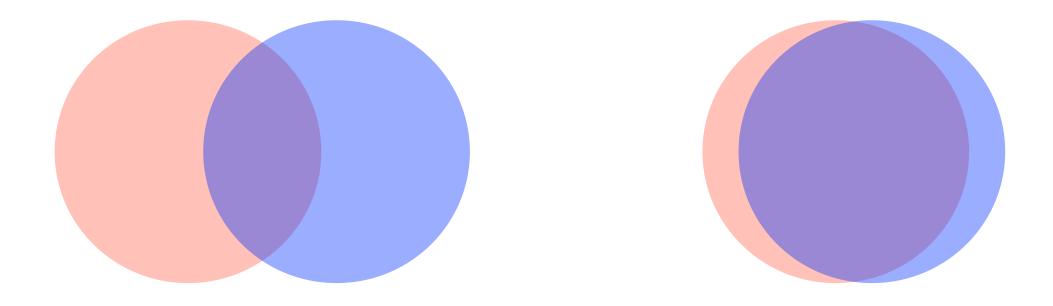
$$O = \frac{|A \cap B|}{\min(|A|, |B|)}$$

 $J = \frac{|A \cap B|}{|A \cup B|}$

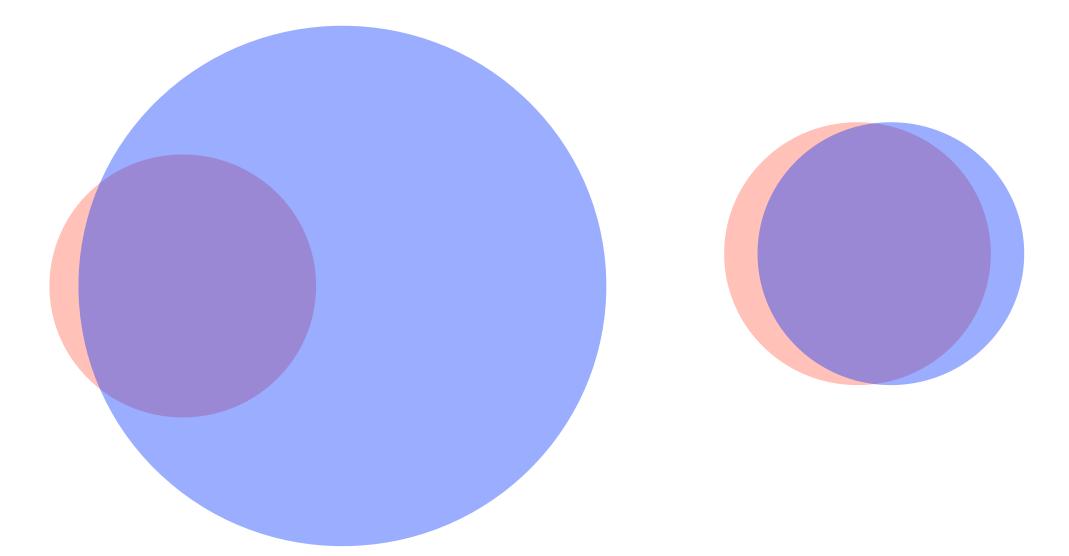
Real-world Meaning Aggccacagtgtattatgactg

GAGG--TCAGATTCACAGCCAC

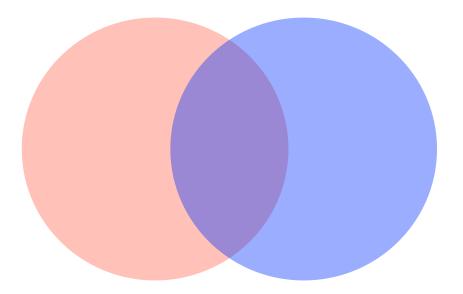
How can we describe how *similar* two sets are?



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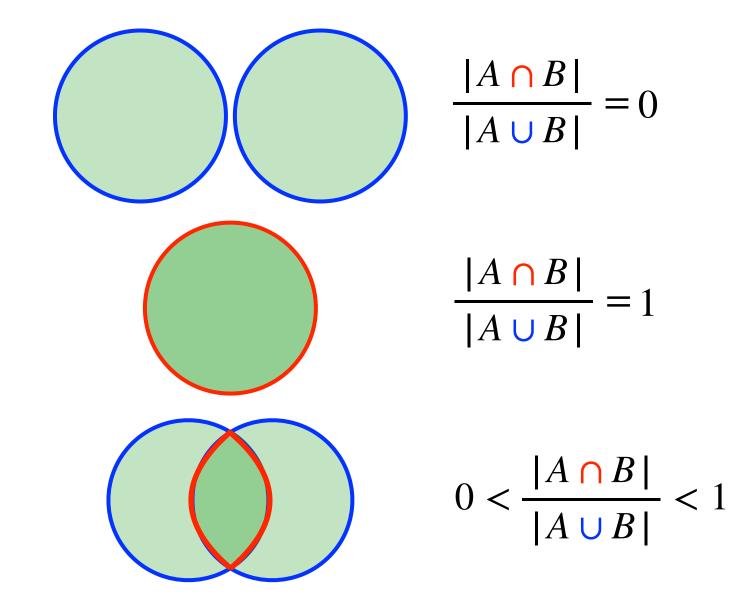


To measure **similarity** of *A* & *B*, we need both a measure of how similar the sets are but also the total size of both sets.



$$J = \frac{|A \cap B|}{|A \cup B|}$$

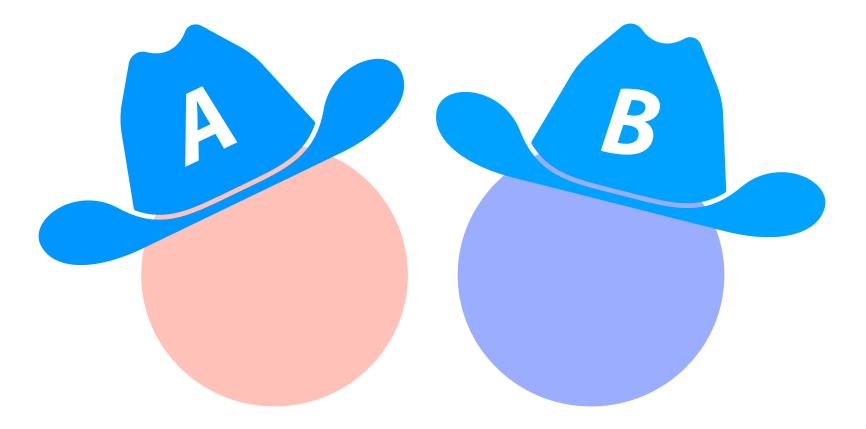
J is the Jaccard coefficient





Similarity Sketches

But what do we do when we only have a sketch?



Similarity Sketches

Imagine we have two datasets represented by their kth minimum values

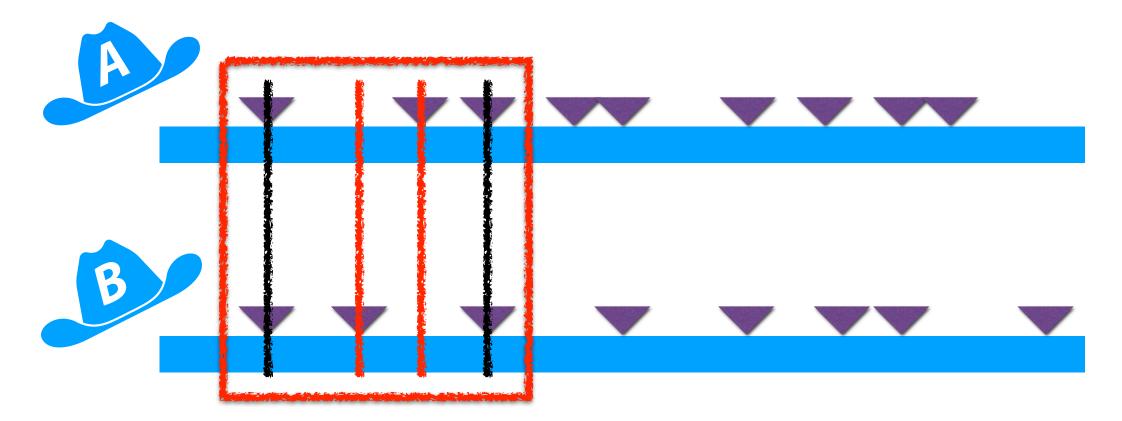


Image inspired by: Ondov B, Starrett G, Sappington A, Kostic A, Koren S, Buck CB, Phillippy AM. **Mash Screen:** high-throughput sequence containment estimation for genome discovery. *Genome Biol* 20, 232 (2019)

Similarity Sketches

Claim: Under SUHA, set similarity can be estimated by sketch similarity!

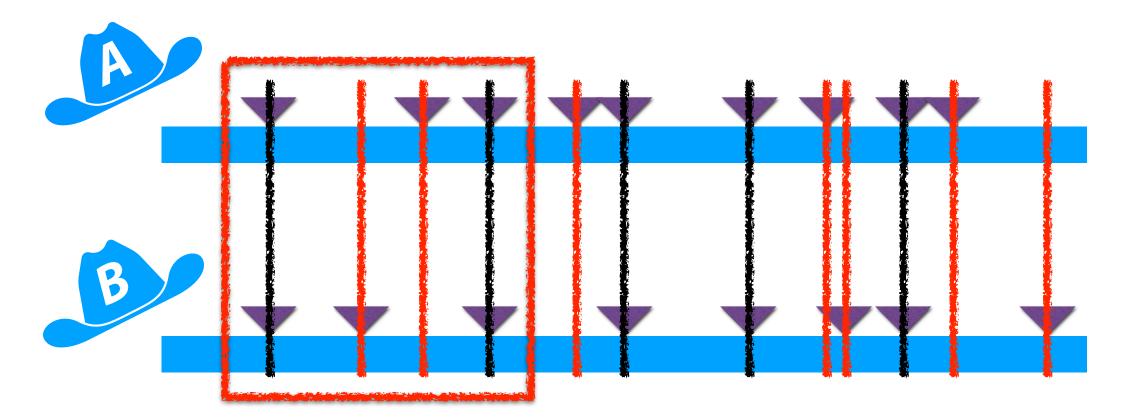
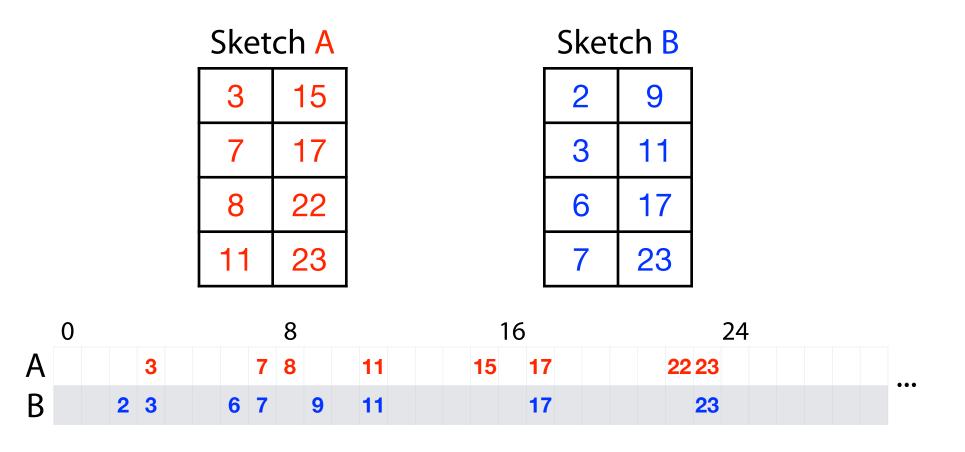


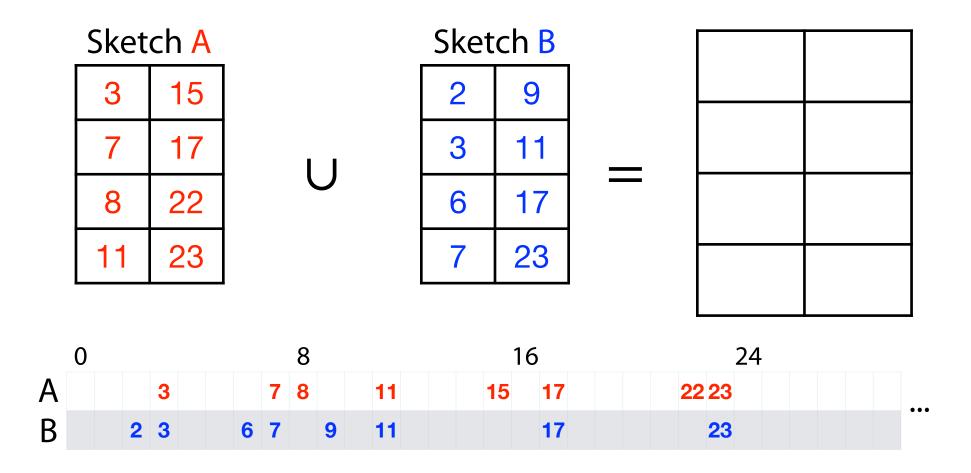
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Let sets A and B be two arbitrary sets of at least 8 elements

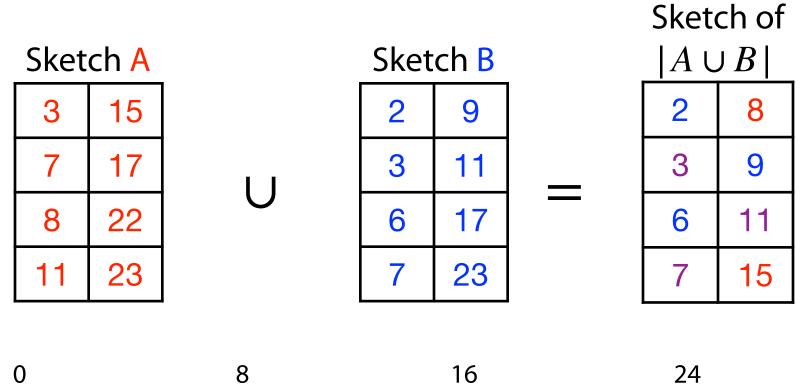
The eight minimum hash values for sets A and B is a MinHash Sketch



To get similarity, we want to estimate $|A \cup B|$ and $|A \cap B| \dots$

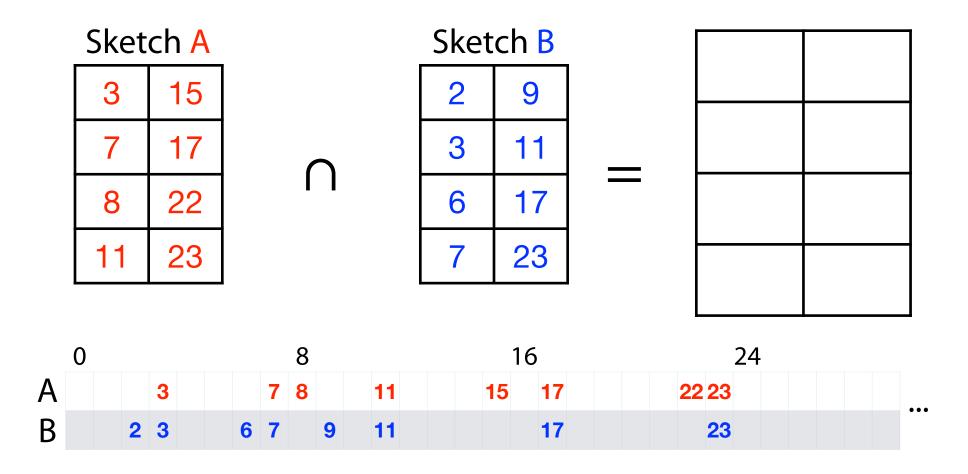


To get similarity, we want to estimate $|A \cup B|$ and $|A \cap B| \dots$



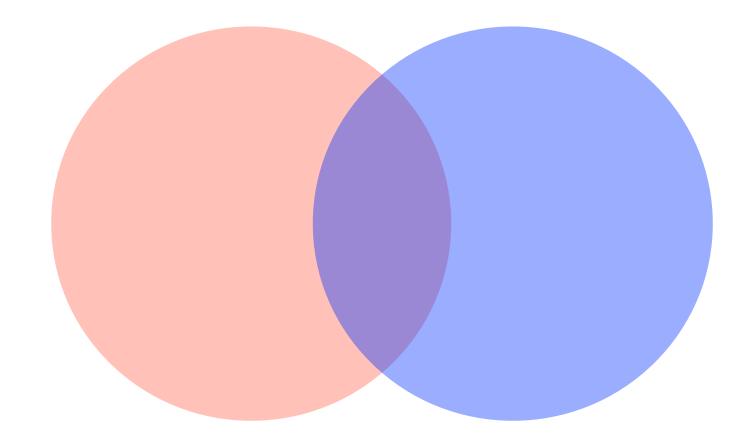


To get similarity, we want to estimate $|A \cup B|$ and $|A \cap B| \dots$

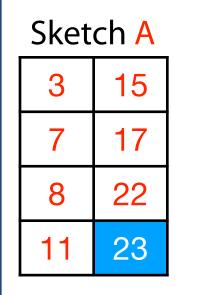


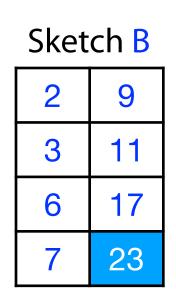
Inclusion-Exclusion Principle

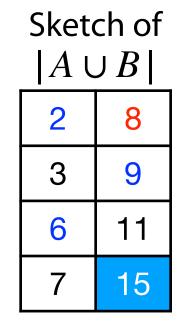
$|A \cap B| =$



Using **inclusion-exclusion** principle and KMV, we can estimate similarity!





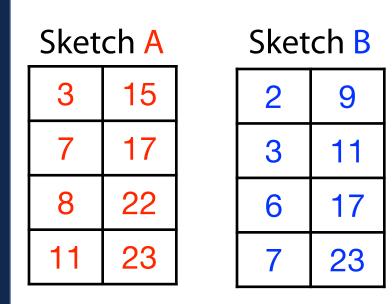


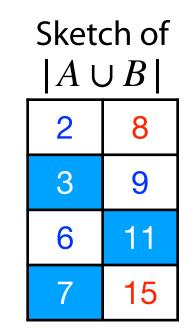
*k*th minimum value (KMV) with k = 8, assuming hash range is integers in [0, 100):

$$= \frac{800/23 - 1 + 800/23 - 1 - 800/15 - 1}{800/15 - 1}$$
$$= \frac{34.782 + 34.782 - 53.333 - 1}{53.333 - 1}$$
$$\approx 0.29$$

$$|A| + |B| - |A \cup B|$$
$$|A \cup B|$$

Claim: Cardinality of the intersection can also be estimated directly!

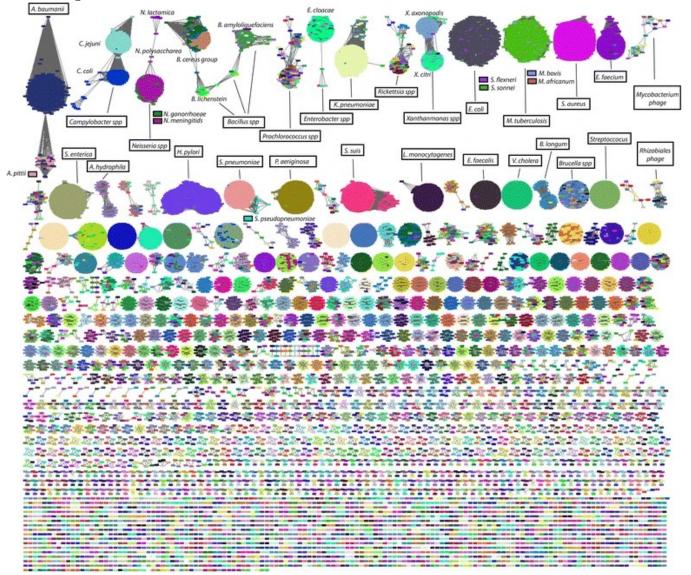




1) Sequence decomposed into **kmers**

Assembling large genomes with single-molecule sequencing and locality-sensitive hashing Berlin et al (2015) *Nature Biotechnology*

MinHash in practice



Mash: fast genome and metagenome distance estimation using MinHash Ondov et al (2016) *Genome Biology* Reviewing probabilistic data sketches



Does a *specific* object exist in my data?

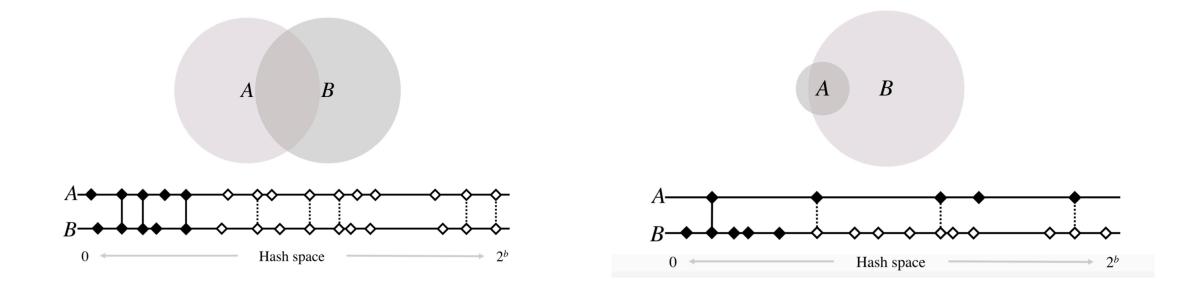
How often is a *specific* object repeated in my data?

How many unique objects do I have in my set?

How similar are two datasets?

Bonus Slides (Taking it one step further...)

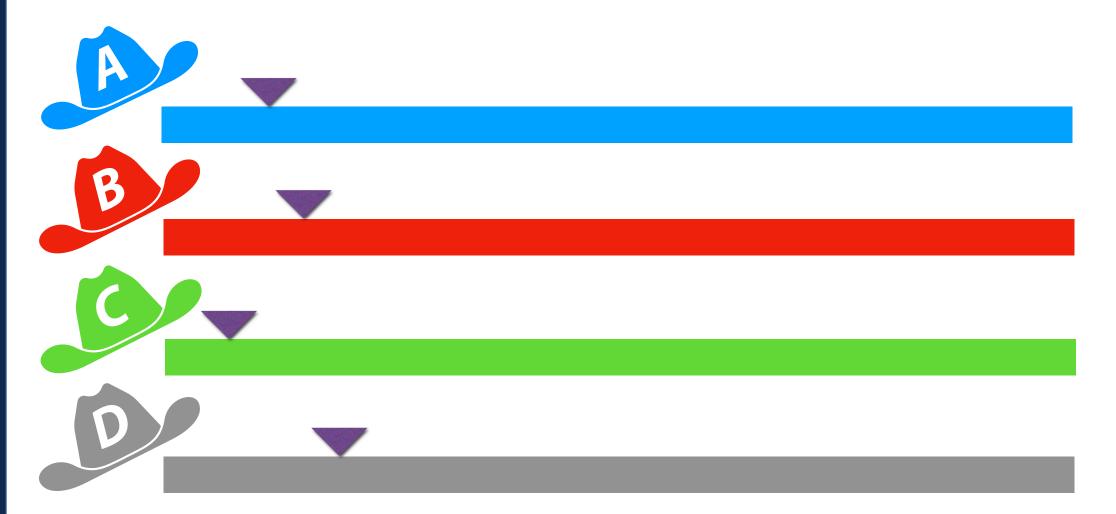
Bottom-k minhash has low accuracy if the cardinality of sets are skewed



Ondov, Brian D., Gabriel J. Starrett, Anna Sappington, Aleksandra Kostic, Sergey Koren, Christopher B. Buck, and Adam M. Phillippy. **Mash Screen: High-throughput sequence containment estimation for genome discovery**. *Genome biology* 20.1 (2019): 1-13.

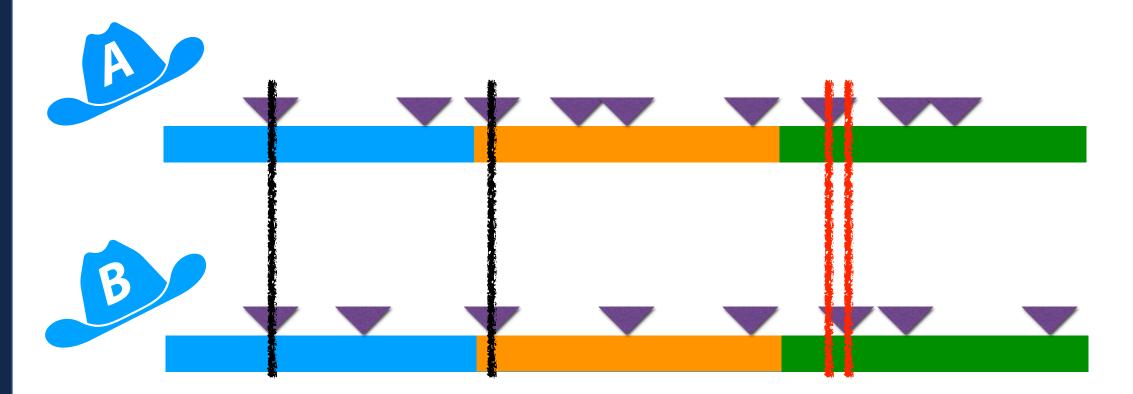
K-Hash Minhash

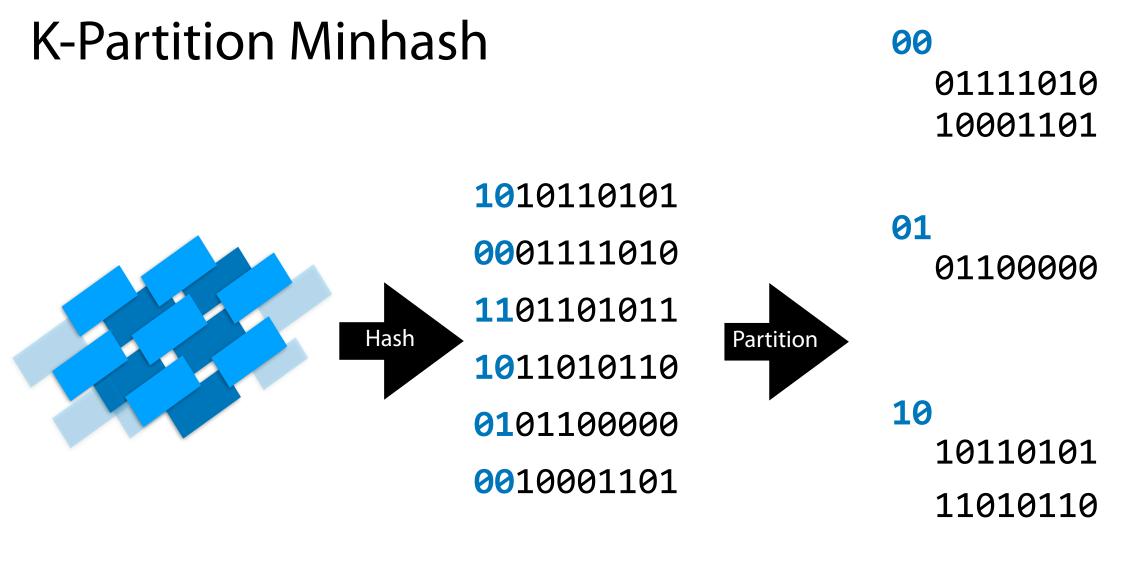
What if instead we used k different hashes and took the min each time?



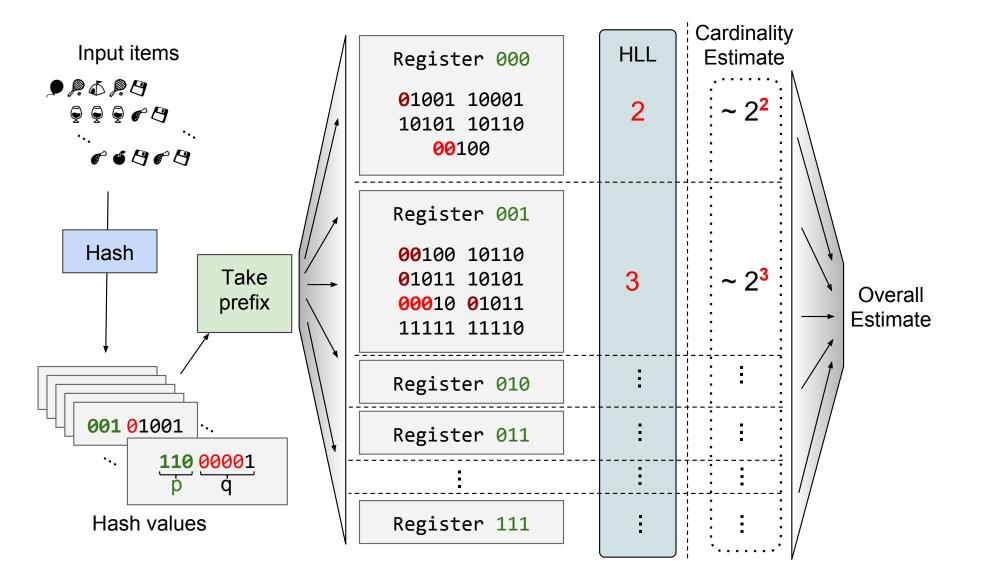
K-Partition Minhash

What if we instead took the minimum of k-partitions?





HyperLogLog



Baker, Daniel et al. "Dashing: fast and accurate genomic distances with HyperLogLog." Genome biology 20.1 (2019): 1-12.