# Data Structures and Algorithms Hashing 2

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## Learning Objectives

Review fundamentals of hash tables

Introduce closed hashing approaches to hash collisions

Determine when and how to resize a hash table

Justify when to use different index approaches

# A Hash Table based Dictionary

#### **Client Code:**

```
1 Dictionary<KeyType, ValueType> d;
2 d[k] = v;
```

### A Hash Table consists of three things:

- 1. A hash function
- 2. A data storage structure
- 3. A method of addressing hash collisions

## **Open vs Closed Hashing**

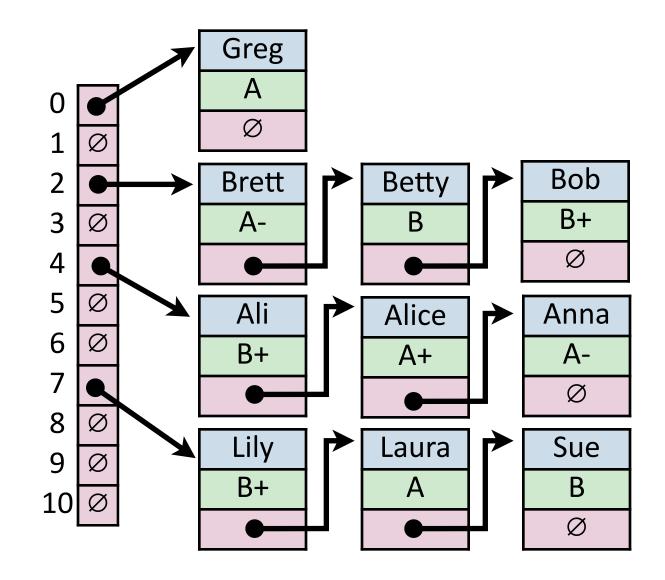
Addressing hash collisions depends on your storage structure.

• **Open Hashing:** store *k*,*v* pairs externally

• **Closed Hashing:** store *k*,*v* pairs in the hash table

## Hash Table (Separate Chaining)

Кеу	Value	lue Hash	
Bob	B+	2	
Anna	A-	4	
Alice	A+	4	
Betty	В	2	
Brett	A-	2	
Greg	А	0	
Sue	В	7	
Ali	B+	4	
Laura	А	7	
Lily	B+	7	



## Simple Uniform Hashing Assumption

Given table of size *m*, a simple uniform hash, *h*, implies

$$\forall k_1, k_2 \in U$$
 where  $k_1 \neq k_2$ ,  $Pr(h[k_1] = h[k_2]) = \frac{1}{m}$ 

Uniform: keys are equally likely to hash to any position

**Independent:** key hash values are independent of other keys

Separate Chaining Under SUHA



Under SUHA, a hash table of size *m* and *n* elements:

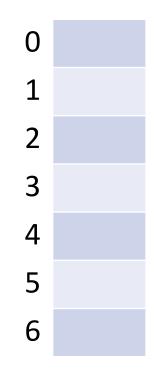
Expected length of chain is \_\_\_\_\_\_.

find runs in: \_\_\_\_\_.

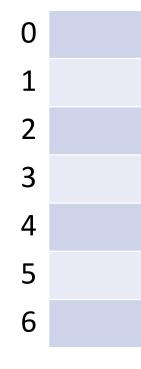
insert runs in: \_\_\_\_\_\_.

remove runs in: \_\_\_\_\_\_.

#### Collision Handling: Probe-based Hashing S = { 1, 8, 15} h(k) = k % 7 Array| = m



Collision Handling: Linear Probing S = { 16, 8, 4, 13, 29, 11, 22 } |S| = n h(k) = k % 7 |Array| = m



h(k, i) = (k + i) % 7 Try h(k) = (k + 0) % 7, if full... Try h(k) = (k + 1) % 7, if full... Try h(k) = (k + 2) % 7, if full... Try ...

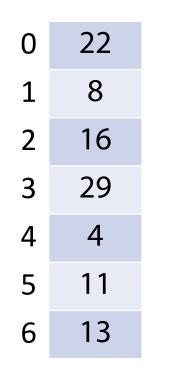
# Collision Handling: Linear Probing S = { 16, 8, 4, 13, 29, 11, 22 } |S| = n h(k, i) = (k + i) % 7 |Array| = m

find(29)

0	22
1	8
2	16
3	29
4	4
5	11
6	13

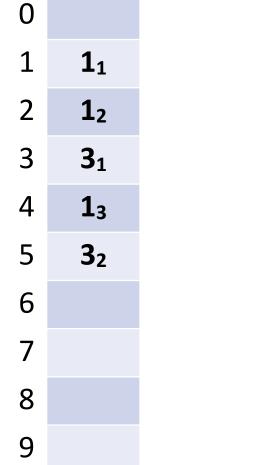
Collision Handling: Linear Probing S = { 16, 8, 4, 13, 29, 11, 22 } |S| = n h(k, i) = (k + i) % 7 |Array| = m

\_remove(16)



A Problem w/ Linear Probing

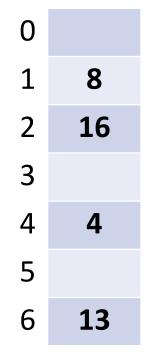
#### **Primary clustering:**



#### **Description:**

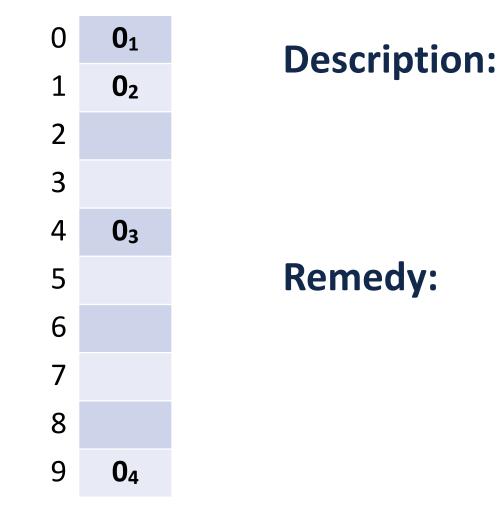
**Remedy:** 

Collision Handling: Quadratic Probing **S** = { 16, 8, 4, 13, 29, 12, 22 } |**S**| = n h(k) = k % 7 |Array| = m



h(k, i) = (k + i\*i) % 7 Try h(k) = (k + 0) % 7, if full... Try h(k) = (k + 1\*1) % 7, if full... Try h(k) = (k + 2\*2) % 7, if full... Try ...

# A Problem w/ Quadratic Probing Secondary clustering:



### **Collision Handling: Double Hashing**

 $S = \{ 16, 8, 4, 13, 29, 11, 22 \}$  |S| = n  $h_1(k) = k \% 7$  |Array| = m  $h_2(k) = 5 - (k \% 5)$ 

> $h(k, i) = (h_1(k) + i^*h_2(k)) \% 7$ Try h(k) = (k + 0\*h\_2(k)) % 7, if full... Try h(k) = (k + 1\*h\_2(k)) % 7, if full... Try h(k) = (k + 2\*h\_2(k)) % 7, if full... Try ...



0

1

2

3

4

5

6

### **Running Times** (Don't memorize these equations, no need.) (Expectation under SUHA)

#### **Open Hashing:**

insert: \_\_\_\_\_.

find/ remove: \_\_\_\_\_\_.

**Closed Hashing:** 

insert: \_\_\_\_\_.

find/ remove: \_\_\_\_\_\_.

### **Running Times** (Don't memorize these equations, no need.) The expected number of probes for find(key) under SUHA

### **Linear Probing:**

- Successful: ½(1 + 1/(1-α))
- Unsuccessful: ½(1 + 1/(1-α))<sup>2</sup>

Instead, observe:

- As α increases:

- Double Hashing:
  Successful: 1/α \* ln(1/(1-α))
- Unsuccessful:  $1/(1-\alpha)$
- **Separate Chaining:**
- Successful:  $1 + \alpha/2$
- Unsuccessful:  $1 + \alpha$

- If  $\alpha$  is constant:

# **Running Times**

The expected number of probes for find(key) under SUHA

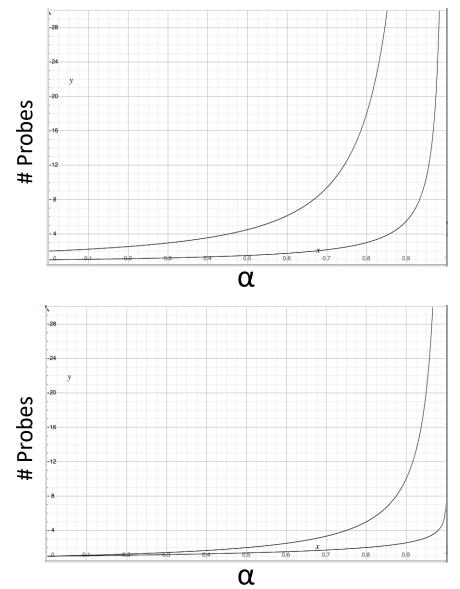
### **Linear Probing:**

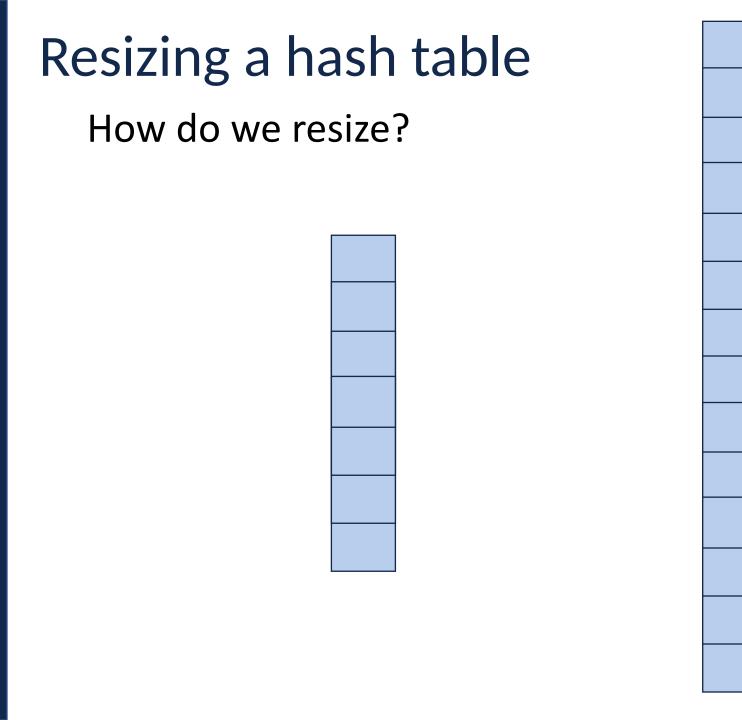
- Successful: ½(1 + 1/(1-α))
- Unsuccessful: ½(1 + 1/(1-α))<sup>2</sup>

### **Double Hashing:**

- Successful: 1/α \* ln(1/(1-α))
- Unsuccessful: 1/(1-α)

#### When do we resize?





#### Which collision resolution strategy is better?

- Big Records:
- Structure Speed:

What structure do hash tables implement?

What constraint exists on hashing that doesn't exist with BSTs?

Why talk about BSTs at all?

# Running Times

	Hash Table	AVL	Linked List
Find	Amortized: Worst Case:		
Insert	Amortized: Worst Case:		
Storage Space			