## Data Structures and Algorithms

## Skip List

CS 225
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## Final Project Reminders / Updates

Regrade late penalties accumulate weekly (Can see penalties on rubric)
"Weekly" is one week after you received feedback on your submission

Mid-project checkins will be available either Nov 16-18 or Nov 28-30

Your mentor will communicate their schedule to you directly

Final project final deadline extended to December 12th

## Learning Objectives

Motivate and introduce the skip list ADT

Conceptualize and code core functions

Analyze efficiency of skip list

## Linked List with ‘Checkpoints’

With some small overhead costs, we can store checkpoints.


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## Linked List with Perfect Checkpoints

For optimal checkpoints, we want half the number of items at each level.


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Maintaining this while inserting and deleting is too costly!


## Linked List with Random Checkpoints

Instead of having exactly half each level, let's have approximately half!


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To analyze runtimes we use: $\qquad$


## The Skip List

An ordered linked list where each node has variable size
Each node has at most one key but an arbitrary number of pointers
The decision for height is randomized
Claim: The expected time to insert, search, or delete is $O(\log n)$


## Skip List

```
template <class T>
class SkipList{
    public:
        class SkipNode{
            public:
                SkipNode() {
                    next.push_back(nullptr);
                }
                    SkipNode(int h, T & d) {
                    data = d;
                    for(int i = 0; i <= h; i++) {
                    next.push_back(nullptr);
                    }
            }
                T data;
                std::vector<SkipNode*> next;
        };
        int max; // max height
        float c; //update constant
        SkipNode* head;
```

Skip List Find Find (9)


Skip List Find Find (7)


Skip List Find Find (1)


## Skip List Find

```
template <class T>
bool SkipList<T>::find(T data){
    SkipNode* curr = head;
    for(int i = max; i >= 0; i--){
        while(curr->next[i]!= nullptr && curr->next[i]->data < data ){
            curr = curr->next[i];
        }
    }
    curr = curr->next[0];
    if (curr != nullptr && curr->data == data){
        return true;
    }
    return false;
}
```



Skip List Insert


Skip List Insert


## Skip List Insert



```
void SkipList<T>::insert(T data) {
    int h = randHeight();
    SkipNode* n = new SkipNode(h, data);
    SkipNode* curr = head;
    for(int i = max; i >= 0; i--) {
        while(curr->next[i]!= nullptr && curr->next[i]->data < data){
            curr = curr->next[i];
        }
        if (h >= i){
            curr->next[i]=n;
            n->next[i]=nextNode;
        }
    }
    if (h > max){
        int diff = h-max;
        for(int i = 0; i < diff; i++) {
            (head->next) .push_back (n) ;
        }
        max =h;
    }
}
```



Skip List Remove


Skip List Remove


Skip List Remove


## Skip List Expectation

Lets assume our skip list uses a coin flip for randomness ( $c=0.5$ )
Claim: Expected size of a node is 2.

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Claim: Expected size of skip list is 2 n .

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Claim: Expected height of skip list is $O(\log n)$

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Claim: Expected height of skip list is $O(\log n)$

$$
E[h]=\sum_{l=0}^{\lceil\log n\rceil} E\left[I_{l}\right]+\sum_{l=\lceil\log n\rceil+1}^{\infty} E\left[I_{l}\right] \quad I_{l}=\left\{\begin{array}{l}
1 l \text { th level contains a node } \\
0 l \text { th level empty }
\end{array}\right.
$$

## Skip List Expectation

Claim: Expected length of search of skip list is $O(\log n)$


