CS 225

Data Structures

October 19 – Graphs
G Carl Evans

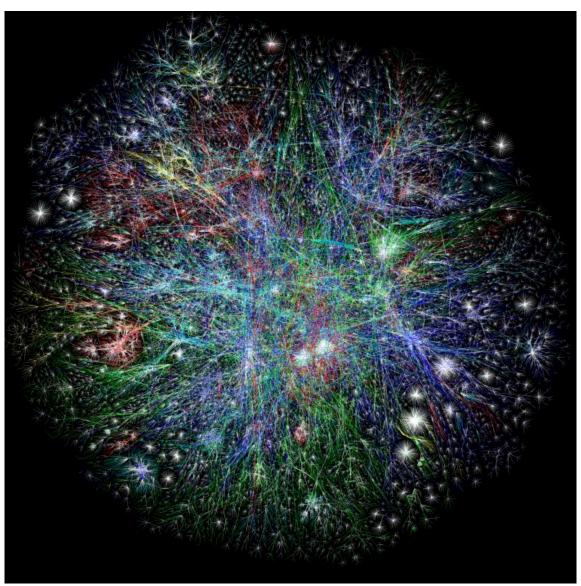
In Review: Data Structures

Array

- Sorted Array
- Unsorted Array
 - Stacks
 - Queues
 - Priority Queues
 - Heaps
 - Disjoint Sets
 - UpTrees

Linked

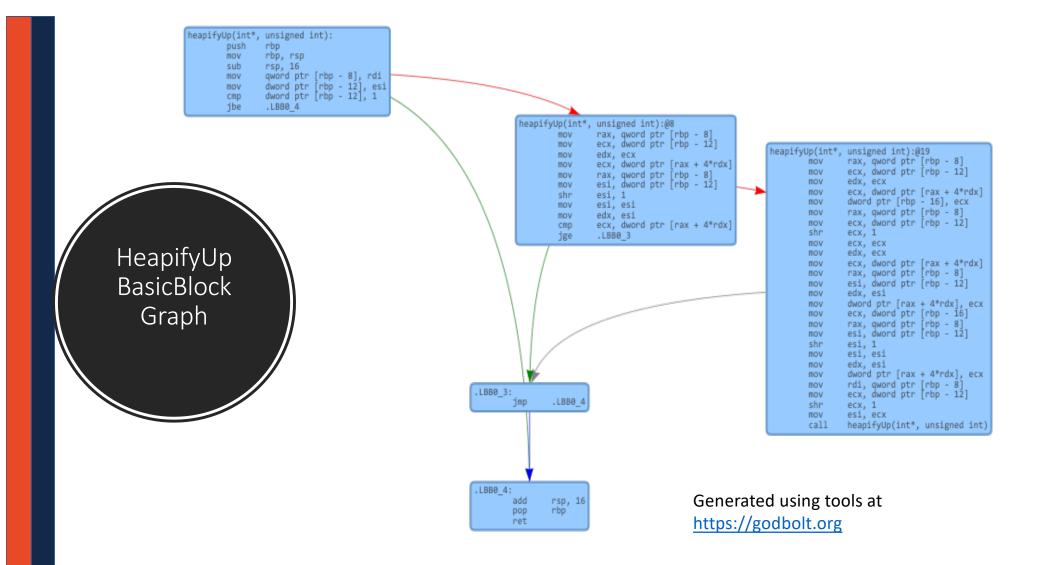
- Doubly Linked List
- Trees
 - BTree
 - Binary Tree
 - Huffman Encoding
 - kd-Tree
 - AVL Tree

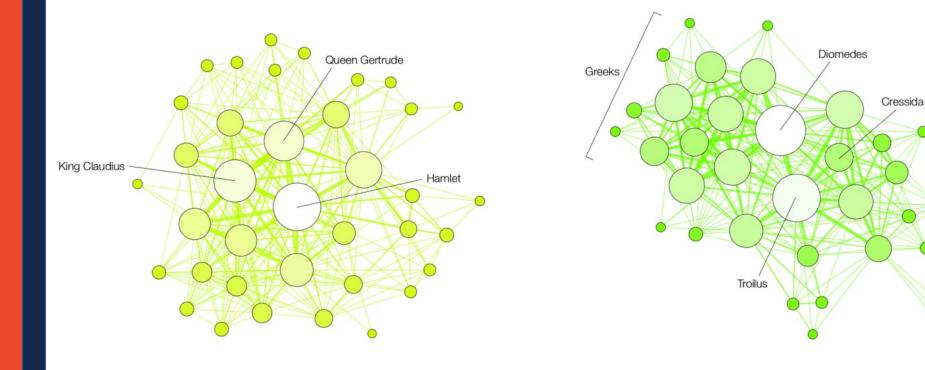


The Internet 2003

The OPTE Project (2003)

Map of the entire internet; nodes are routers; edges are connections.





HAMLET

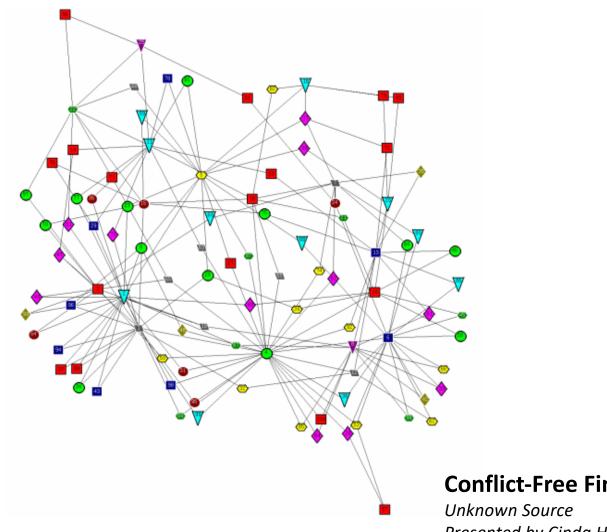
TROILUS AND CRESSIDA

Trojans

Who's the real main character in Shakespearean tragedies?

Martin Grandjean (2016)

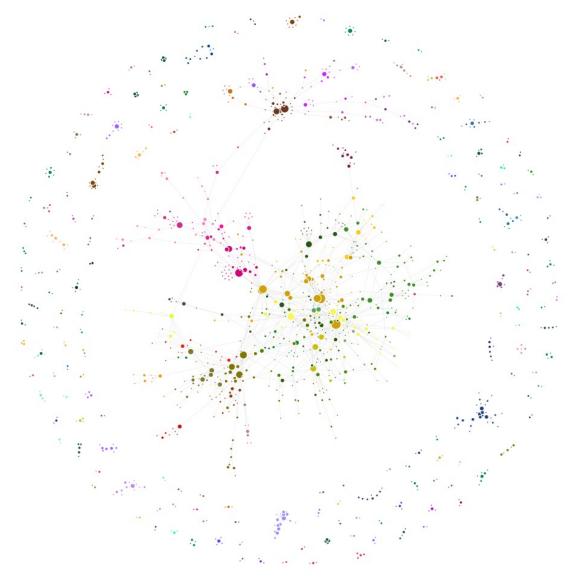
https://www.pbs.org/newshour/arts/whos-the-real-main-character-in-shakespearen-tragedies-heres-what-the-data-say

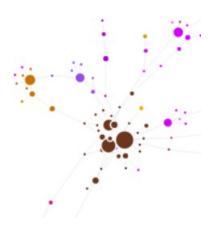


Conflict-Free Final Exam Scheduling Graph

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Presented by Cinda Heeren, 2016





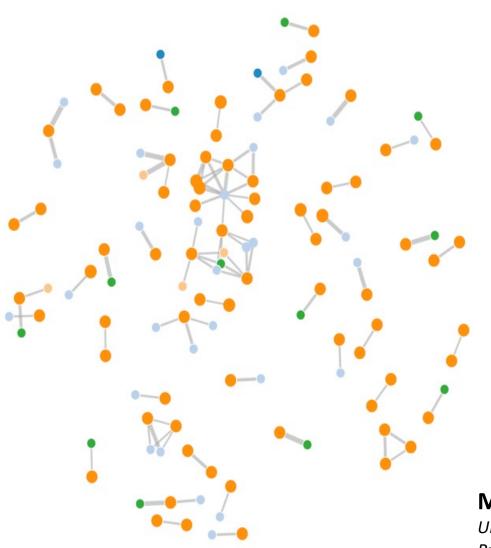


Class Hierarchy At University of Illinois Urbana-Champaign

A. Mori, W. Fagen-Ulmschneider, C. Heeren

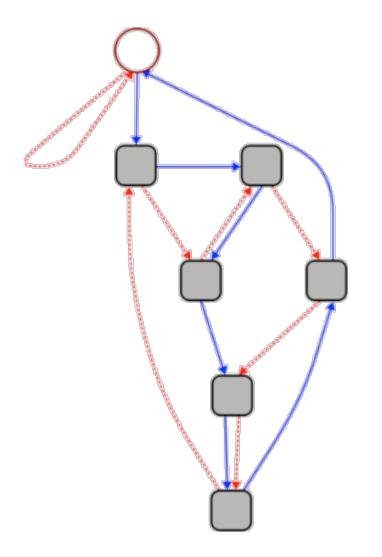
Graph of every course at UIUC; nodes are courses, edges are prerequisites

http://waf.cs.illinois.edu/discovery/class_hi
erarchy_at_illinois/



MP Collaborations in CS 225

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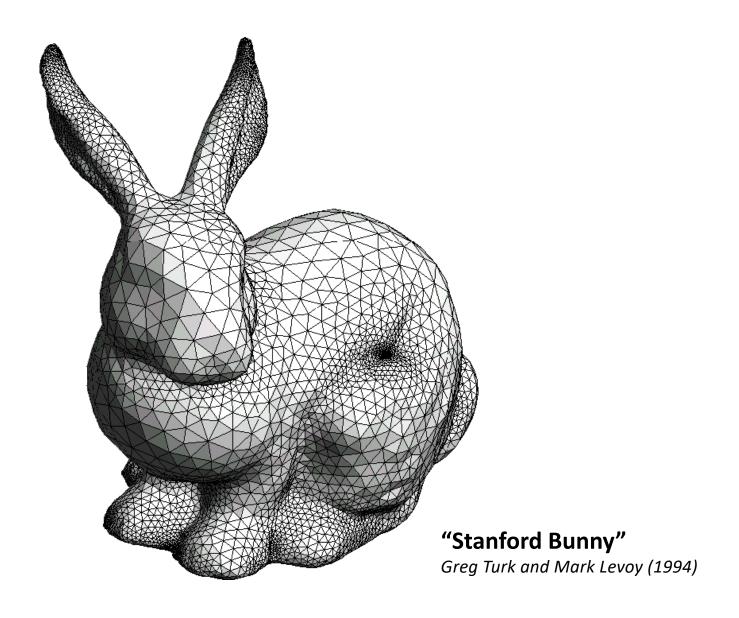
This graph can be used to quickly calculate whether a given number is divisible by 7.

- 1. Start at the circle node at the top.
- 2. For each digit **d** in the given number, follow **d** blue (solid) edges in succession. As you move from one digit to the next, follow **1** red (dashed) edge.
- 3. If you end up back at the circle node, your number is divisible by 7.

3703

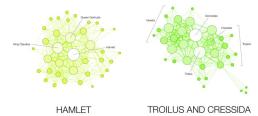
"Rule of 7"

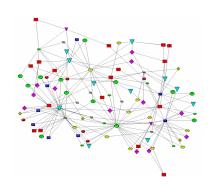
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Graphs

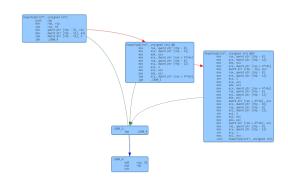


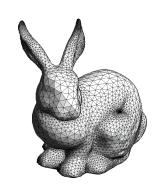


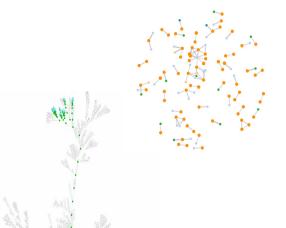


To study all of these structures:

- 1. A common vocabulary
- 2. Graph implementations
- 3. Graph traversals
- 4. Graph algorithms







Graph Vocabulary

```
G = (V, E)
|V| = n
|E| = m
                     (2, 5)
```

Degree(v): ||

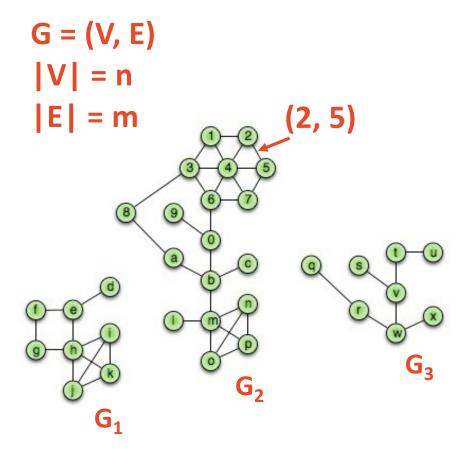
Adjacent Vertices: A(v) = { x : {x, v} in E }

Path(G₂): Sequence of vertices connected by edges

Cycle(G₁): Path with a common begin and end vertex with at least 3 vertices.

Simple Graph(G): A graph with no self loops or multi-edges.

Graph Vocabulary



```
Subgraph(G):

G' = (V', E'):

V' \in V, E' \in E, and

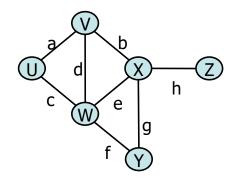
(u, v) \in E' \rightarrow u \in V', v \in V'
```

Complete subgraph(G)
Connected subgraph(G)
Connected component(G)
Acyclic subgraph(G)
Spanning tree(G)

Running times are often reported by **n**, the number of vertices, but often depend on **m**, the number of edges.

How many edges? Minimum edges:

Not Connected:



Connected*:

Maximum edges:

Simple:

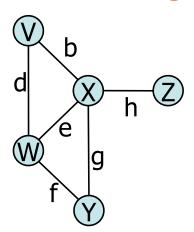
Not simple:

$$\sum_{v \in V} \deg(v) =$$

Graph ADT

Data:

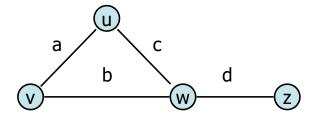
- Vertices
- Edges
- Some data structure maintaining the structure between vertices and edges.



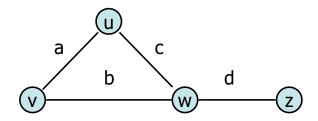
Functions:

- insertVertex(K key);
- insertEdge(Vertex v1, Vertex v2, K key);
- removeVertex(Vertex v);
- removeEdge(Vertex v1, Vertex v2);
- incidentEdges(Vertex v);
- areAdjacent(Vertex v1, Vertex v2);
- origin(Edge e);
- destination(Edge e);

Graph Implementation Idea



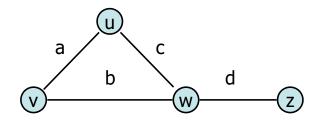
Vertex Collection:



u v a
v w b
u w c
w z d

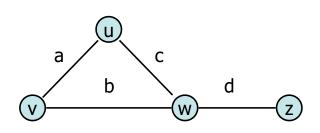
Edge Collection:

insertVertex(K key):

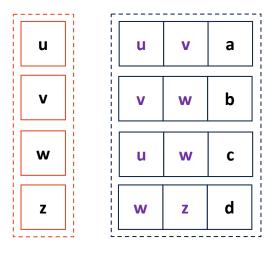


u v a
v w b
u w c
z w z d

removeVertex(Vertex v):

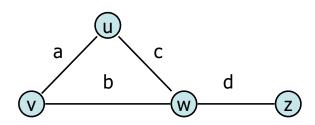


incidentEdges(Vertex v):

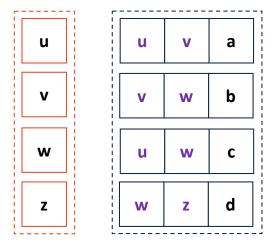


areAdjacent(Vertex v1, Vertex v2):

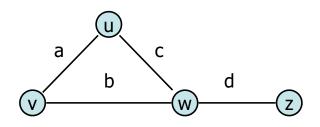
G.incidentEdges(v1).contains(v2)



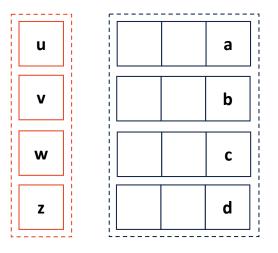
insertEdge(Vertex v1, Vertex v2, K key):



Graph Implementation: Adjacency Matrix



insertVertex(K key);
removeVertex(Vertex v);
areAdjacent(Vertex v1, Vertex v2);
incidentEdges(Vertex v);



	u	V	w	Z
u				
v				
w				
Z				