CS 225

Data Structures

October 12 – Disjoint Sets and kD-tree G Carl Evans

Implementation #2

- We will continue to use an array where the index is the key
- The value of the array is:
 - -1, if we have found the representative element
 - The index of the parent, if we haven't found the rep. element
- We will call theses UpTrees:



Disjoint Sets



0	1	2	3	4	5	6	7	8	9

Disjoint Sets – Smart Union





Both guarantee the height of the tree is: _____

Disjoint Sets Find

```
1 int DisjointSets::find(int i) {
2    if ( s[i] < 0 ) { return i; }
3    else { return find( s[i] ); }
4 }</pre>
```

```
void DisjointSets::unionBySize(int root1, int root2) {
 1
 2
     int newSize = arr [root1] + arr [root2];
 3
 4
     // If arr [root1] is less than (more negative), it is the larger set;
 5
     // we union the smaller set, root2, with root1.
 6
     if ( arr [root1] < arr [root2] ) {</pre>
 7
       arr [root2] = root1;
 8
       arr [root1] = newSize;
 9
     }
10
11
     // Otherwise, do the opposite:
     else {
12
13
       arr [root1] = root2;
       arr [root2] = newSize;
14
15
     }
16
```

Path Compression



MP Mosaics and One More Tree

kD-Trees



kD-Trees





kD-Trees

kD-Tree Constructor

How to construct your own kD-tree?

kD-Tree

- Data structures (trees) which are often used to find the *nearest neighbor* of a *k*-dimensional point
 - Why you should care: actually very applicable to real world scenarios!
- kD-Trees are used to organize Points in k-dimensional space, for any k > 0

kD-Tree Constructor

Nearest Neighbors

Some suggestions to keep in mind

Backtracking: start recursing backwards -- store "best" possibility as you trace back

On ties, use smallerDimVal to determine which point remains curBest

