



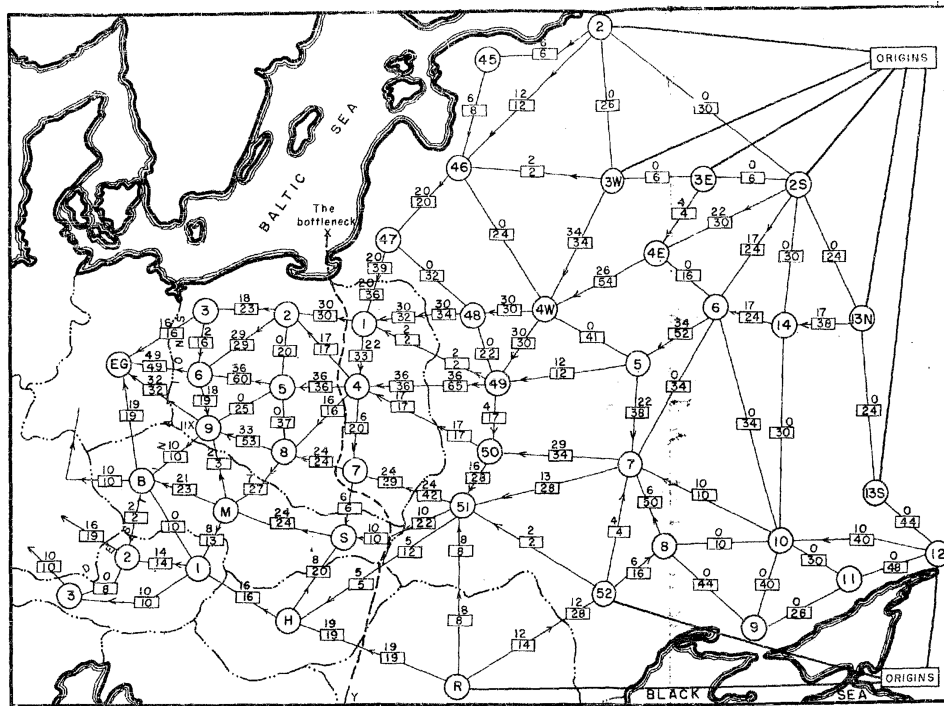
# CS 225

## Data Structures

*December 6 – Maximum Flow*

*G Carl Evans*

# Origin of Maxflow Problem



SECRET <sup>RM-3773</sup>  
10-24-55  
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Fig. 7 — Traffic pattern: entire network available

Legend:

— International boundary

⊙ Railway operating division

←  $\frac{a}{b}$  → Capacity:  $a$  each way per day. Required flow of  $b$  per day toward destinations (in direction of arrow) with equivalent number of returning trains in opposite direction

All capacities in  $\sqrt{1000}$ 's of tons each way per day

Origins: Divisions 2, 3W, 3E, 2S, 13N, 13S, 12, 52 (USSR), and Roumania

Destinations: Divisions 3, 6, 9 (Poland); B (Czechoslovakia); and 2, 3 (Austria)

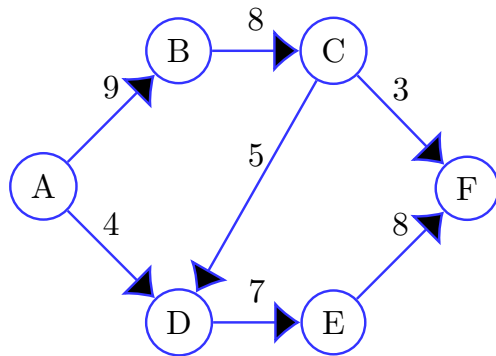
Alternative destinations: Germany or East Germany

Note IIX at Division 9, Poland

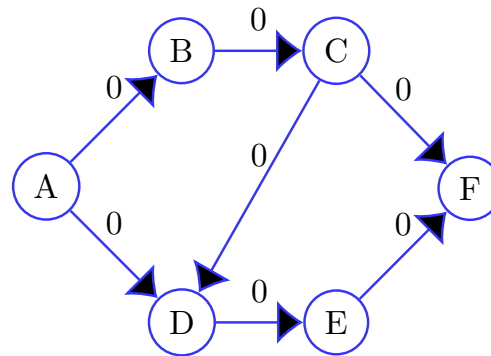
We are given as input graph  $G$ .

We create two new graphs: a *flow graph*  $F$  and a *residual graph*  $R$ .

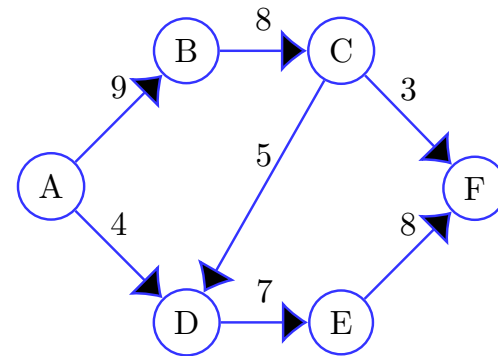
Graph  $G$



Flow Graph  $F$



Residual Graph  $R$

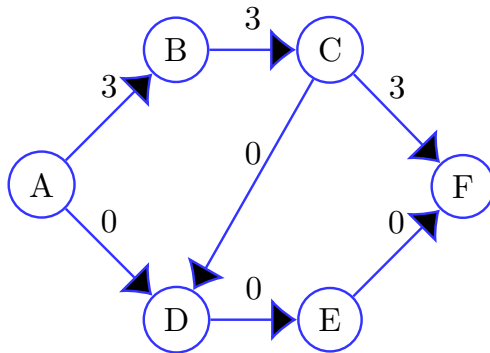


**Problem 1.**

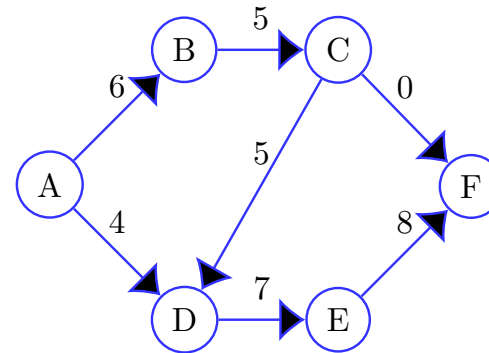
The algorithm works by selecting paths from the residual graph  $R$ . The first path selected is  $A \rightarrow B \rightarrow C \rightarrow F$  in graph  $R$ . This path's flow capacity is 3. What do you think determines the flow capacity?

The algorithm uses the path to modify graphs  $F$  and  $R$ . Here is the result.

Graph  $F$



Graph  $R$



**Problem 2.**

Examine the new versions of  $F$  and  $R$  above. What is being done with the path selected from  $R$  to modify these graphs?

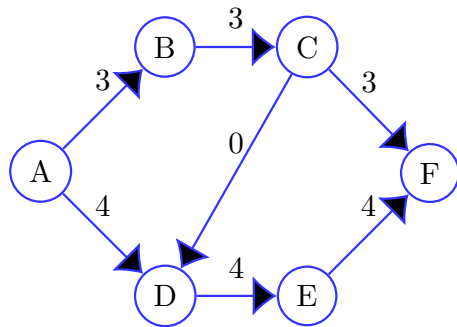
**Problem 3.**

The next path selected was  $A \rightarrow D \rightarrow E \rightarrow F$  in graph  $R$ . What is the flow capacity of that path?

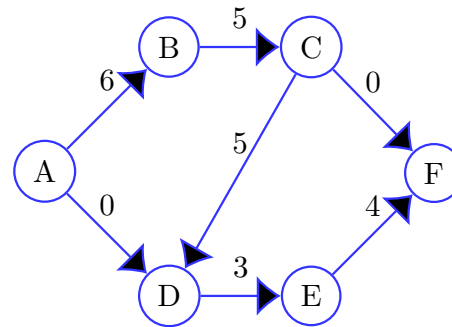
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The resulting working graphs are these:

Graph  $F$



Graph  $R$



**Problem 4.**

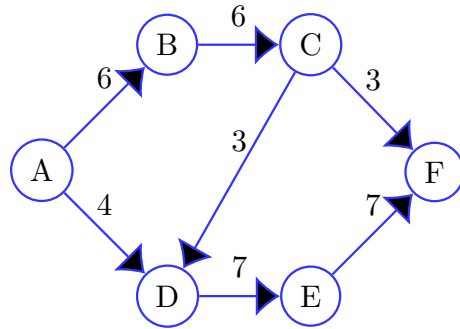
We select path  $A \rightarrow B \rightarrow C \rightarrow D \rightarrow E \rightarrow F$ . What is the flow capacity of that path?

**Problem 5.**

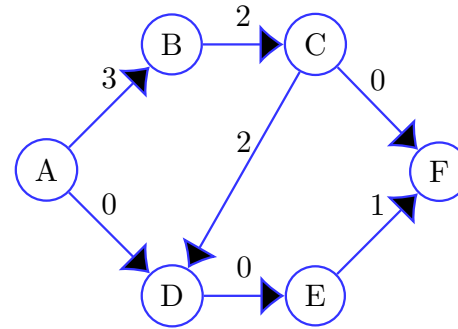
The paths selected always start from node  $A$  and end with node  $F$ . What is different about these nodes compared to the others?

Here are the final working graphs  $F$  and  $R$ .

Graph  $F$



Graph  $R$



**Problem 6.**

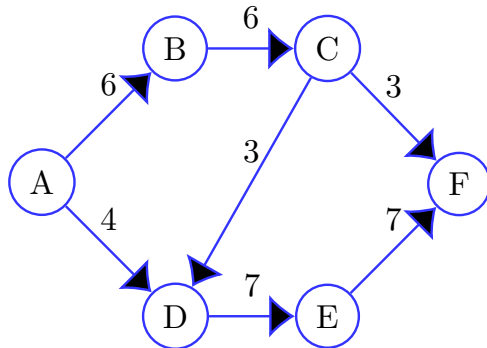
At this point, the algorithm is finished. How can we know the algorithm is done by examining graph  $R$ ?

**Problem 7.**

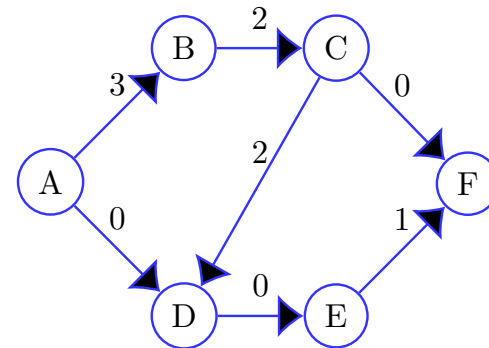
For nodes  $B$ ,  $C$ ,  $D$ , and  $E$ , what is the relationship between the in-flows and the out-flows? Why does that relationship have to exist?

Here are the final working graphs  $F$  and  $R$ .

Graph  $F$



Graph  $R$



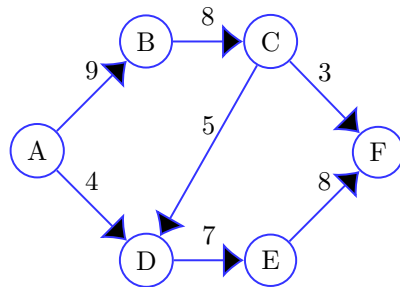
**Problem 8.**

Using the final flow graph  $F$  above, determine the maximum flow of graph  $G1$ .

**Problem 9.**

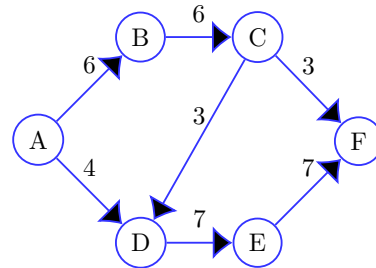
In graph  $F$ , the outflow of  $A$  is equal to the inflow of  $F$ . Should that always be the case?

Graph  $G$

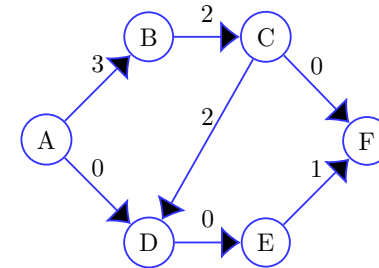


Here are the final working graphs  $F$  and  $R$ .

Graph  $F$



Graph  $R$



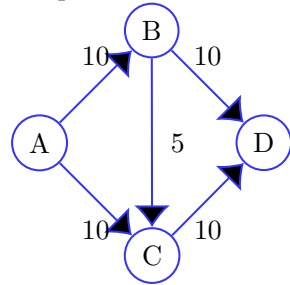
**Problem 10.**

Node  $A$  is called a *source node* and node  $F$  is called a *sink node*. Would this technique work if there were multiple source and sink nodes? Why or why not?

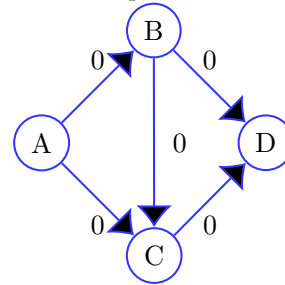


Now we are going to look at a case that messes up the algorithm.

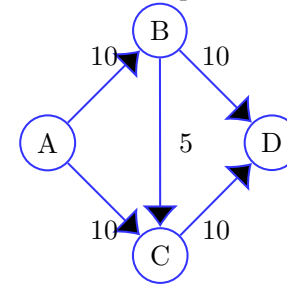
Graph  $G_2$



Flow Graph



Residual Graph



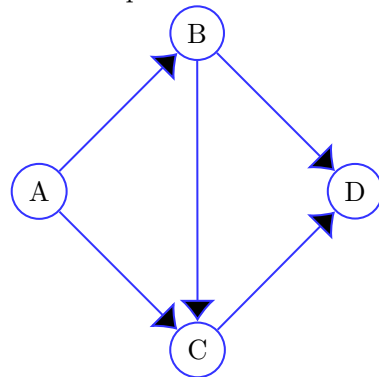
**Problem 11.**

The algorithm picks path  $A \rightarrow B \rightarrow C \rightarrow D$ . What is the capacity of that path?

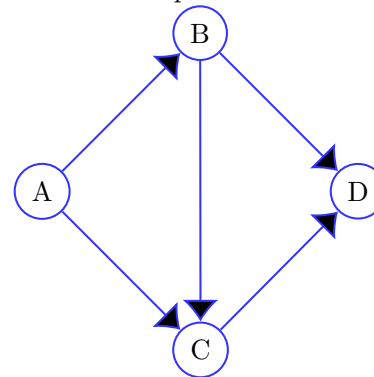
**Problem 12.**

Update the flow and residual graphs as a result of selecting this path.

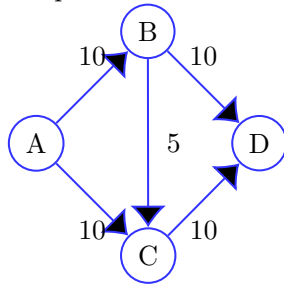
Flow Graph



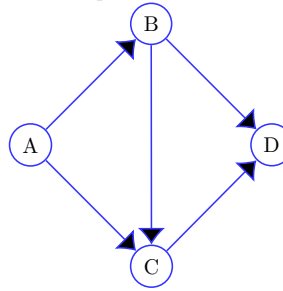
Residual Graph



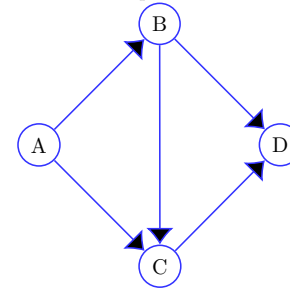
Graph  $G_2$



Flow Graph



Residual Graph



**Problem 17.**

What is the maximum network flow of  $G_2$ , according to the algorithm?

**Problem 18.**

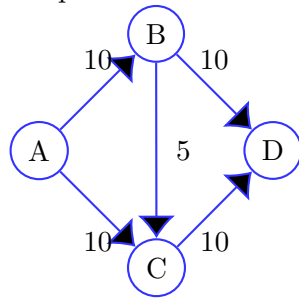
Is this number correct? Why or why not? Examine  $G_2$  to verify your answer.

**Problem 19.**

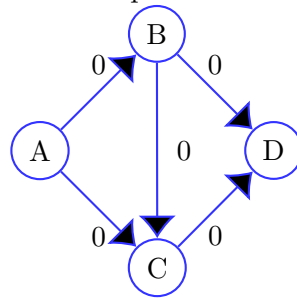
Suppose  $G_2$  modeled a network of water pipes. What would happen on edge  $B \rightarrow C$  in this situation? Would it change the total flow of  $G_2$  if we deleted that edge?

We are going to modify the algorithm. Starting again with the previous graph, we make a new kind of residual graph. The dotted edges are added, and are legal edges to be traversed in the residual graph.

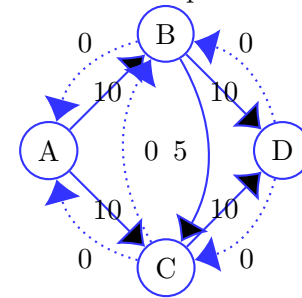
Graph  $G_3$



Flow Graph



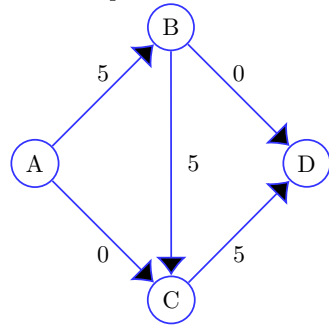
Residual Graph



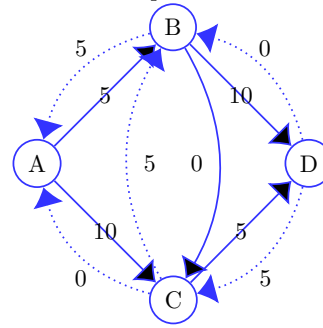
**Problem 20.**

Select path  $A \rightarrow B \rightarrow C \rightarrow D$ . What is the capacity of that path?

Flow Graph

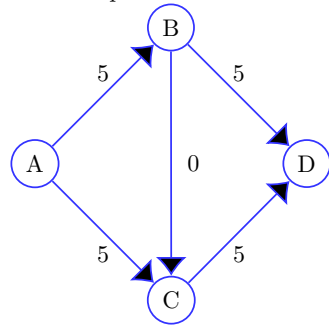


Residual Graph

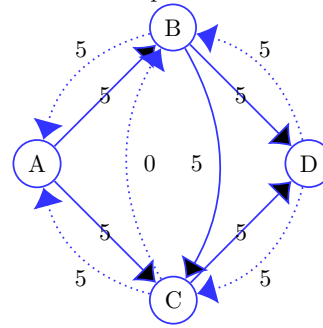


Now we select path  $A \rightarrow C \rightarrow B \rightarrow D$ .  
Here are the updated flow and residual graphs:

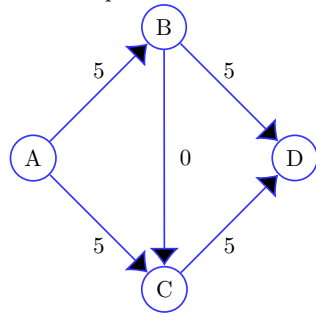
Flow Graph



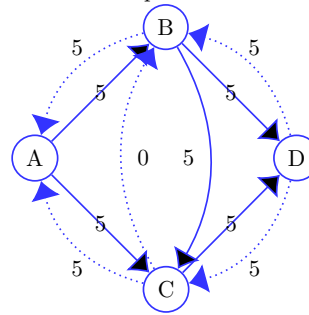
Residual Graph



Flow Graph



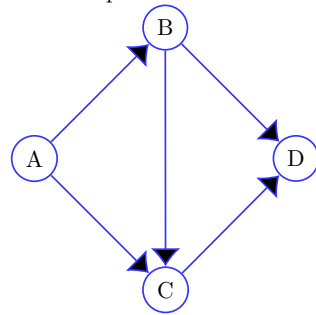
Residual Graph



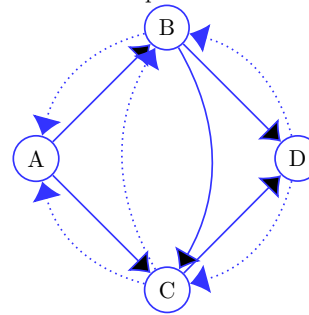
**Problem 22.**

Select path  $A \rightarrow B \rightarrow C \rightarrow D$ . (Yes, we are repeating this path.) What are the resulting flow and residual graphs?

Flow Graph



Residual Graph

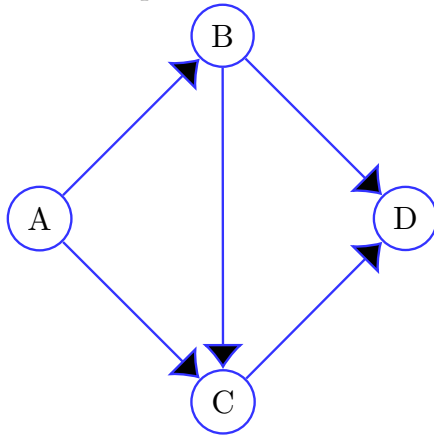


**Problem 23.**

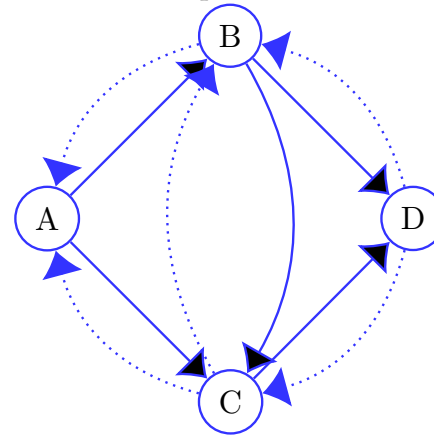
Now we select path  $A \rightarrow C \rightarrow B \rightarrow D$ .

What are the updated flow and residual graphs?

Flow Graph



Residual Graph



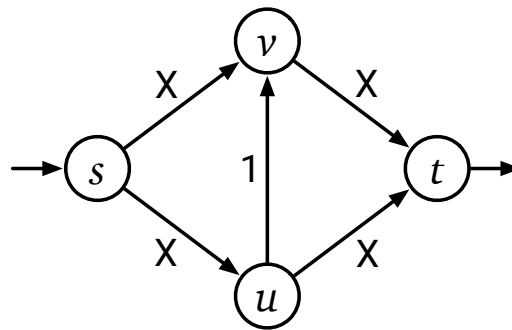
**Problem 24.**

At this point, the algorithm should be done. Is the final network flow accurate now?



# Ford Fulkerson Requirements and Runtime

# Ford Fulkerson Problems



**Figure 10.7.** Edmonds and Karp's bad example for the Ford-Fulkerson algorithm.

Image from <https://jeffe.cs.illinois.edu/teaching/algorithms/book/10-maxflow.pdf>





# Edmonds and Karp's Algorithms