



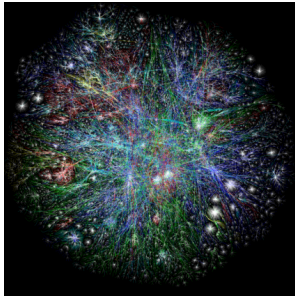
# CS 225

## Data Structures

*November 12 – Graphs*

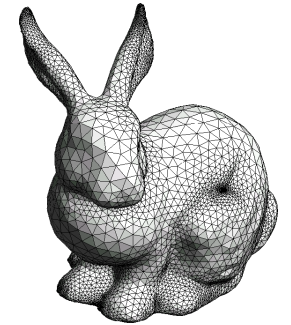
*Wade Fagen-Ulmschneider*

# Graphs

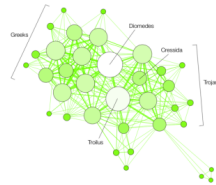


**To study all of these structures:**

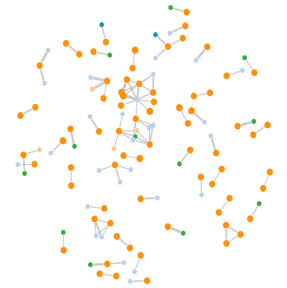
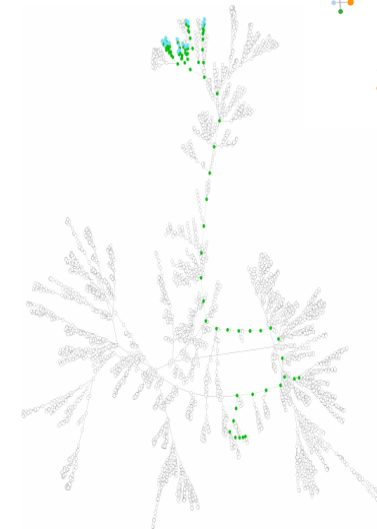
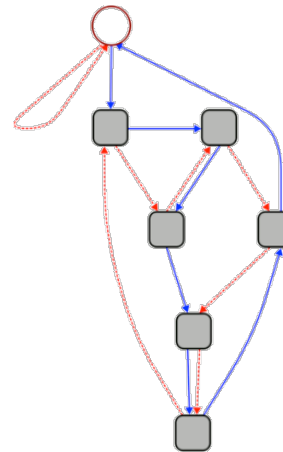
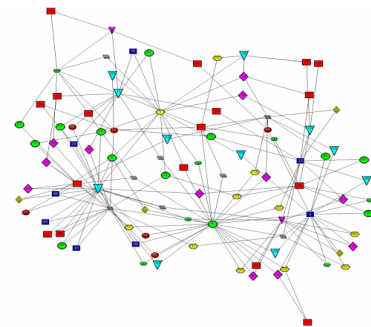
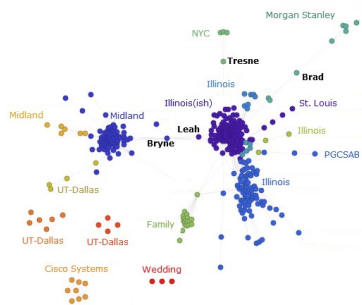
1. A common vocabulary
2. Graph implementations
3. Graph traversals
4. Graph algorithms



HAMLET



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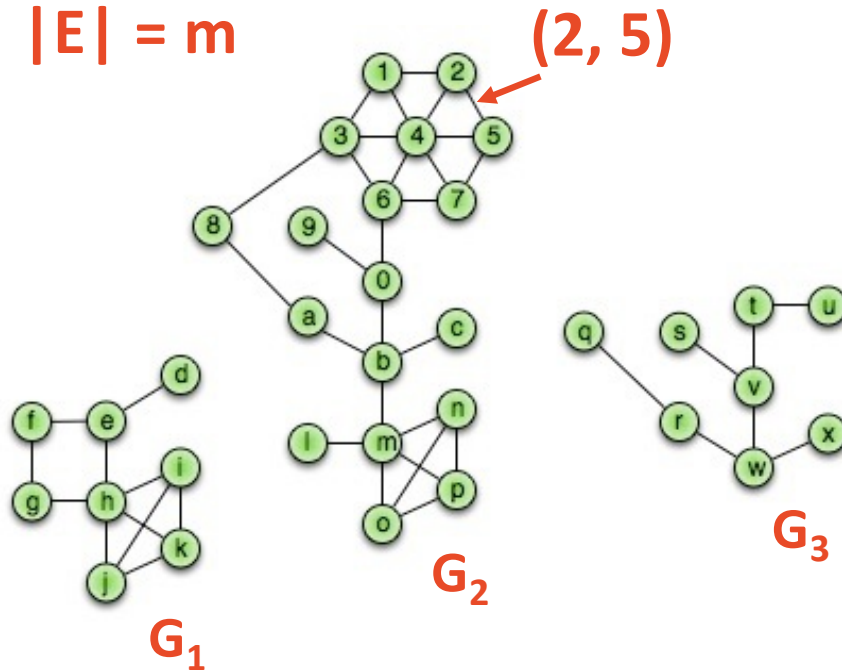


# Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



Incident Edges:

$$I(v) = \{ (x, v) \text{ in } E \}$$

Degree(v):  $|I(v)|$

Adjacent Vertices:

$$A(v) = \{ x : (x, v) \text{ in } E \}$$

Path( $G_2$ ): Sequence of vertices connected by edges

Cycle( $G_1$ ): Path with a common begin and end vertex.

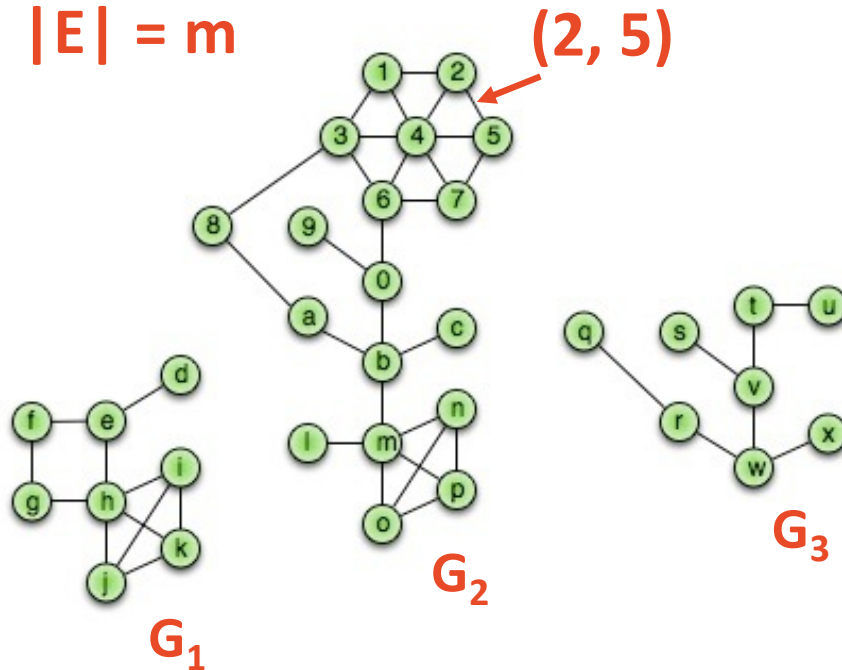
Simple Graph( $G$ ): A graph with no self loops or multi-edges.

# Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



**Subgraph(G):**

$$G' = (V', E')$$

$V' \subseteq V, E' \subseteq E$ , and

$$(u, v) \in E \rightarrow u \in V', v \in V'$$

**Complete subgraph(G)**

**Connected subgraph(G)**

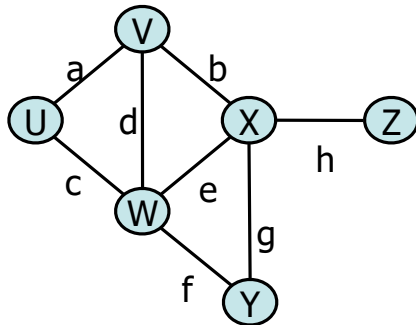
**Connected component(G)**

**Acyclic subgraph(G)**

**Spanning tree(G)**

Running times are often reported by  $n$ , the number of vertices, but often depend on  $m$ , the number of edges.

How many edges? **Minimum edges:**  
Not Connected:



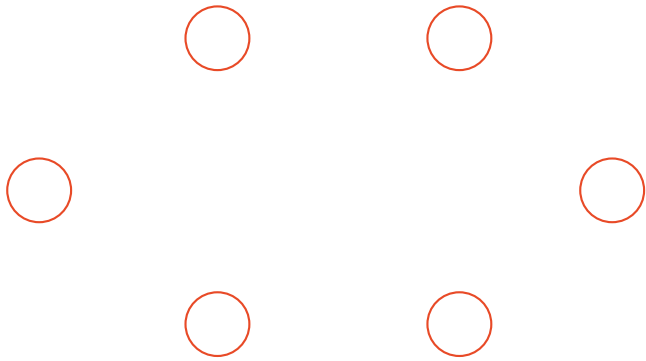
Connected\*:

**Maximum edges:**  
Simple:

Not simple:

$$\sum_{v \in V} \deg(v) =$$

# Connected Graphs






## Proving the size of a minimally connected graph

### **Theorem:**

Every connected graph  $G=(V, E)$  has at least  $|V|-1$  edges.



**Thm:** Every connected graph  $G=(V, E)$  has at least  $|V|-1$  edges.

**Proof:** Consider an arbitrary, connected graph  $G=(V, E)$ .





**Suppose  $|V| = 1$ :**

**Definition:** A connected graph of 1 vertex has 0 edges.

**Theorem:**  $|V| - 1$  edges  $\rightarrow 1 - 1 = 0$ .

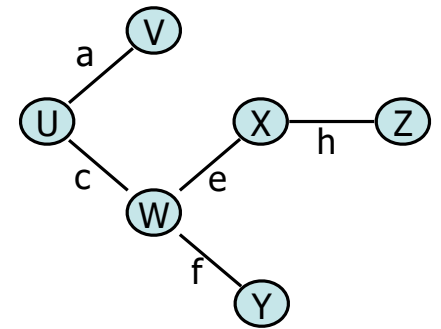


**Inductive Hypothesis:** For any  $j < |V|$ , any connected graph of  $j$  vertices has at least  $j-1$  edges.

**Suppose  $|V| > 1$ :**

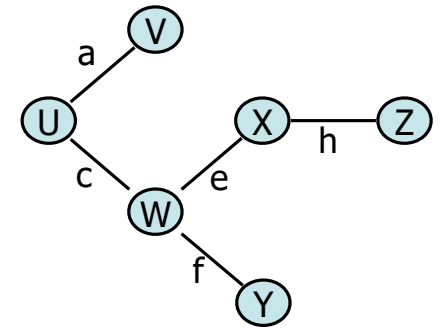
1. Choose any vertex:

2. Partition:



**Suppose  $|V| > 1$ :**

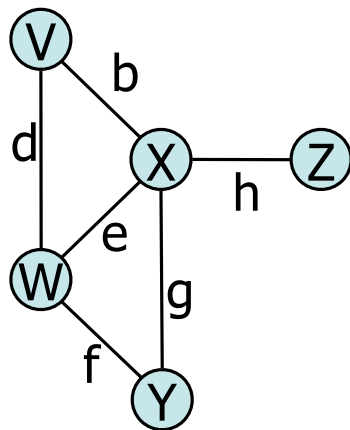
3. Count the edges



# Graph ADT

## Data:

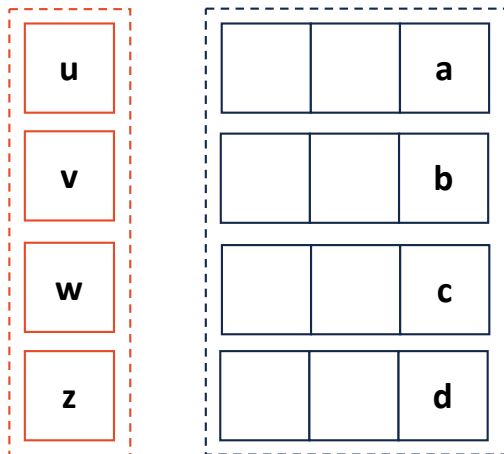
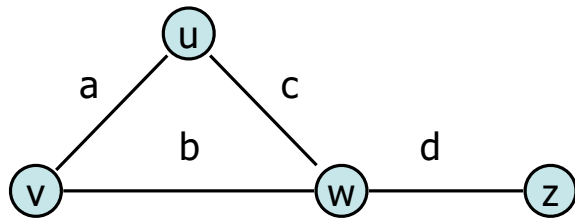
- Vertices
- Edges
- Some data structure maintaining the structure between vertices and edges.



## Functions:

- insertVertex(K key);
- insertEdge(Vertex v1, Vertex v2, K key);
- removeVertex(Vertex v);
- removeEdge(Vertex v1, Vertex v2);
- incidentEdges(Vertex v);
- areAdjacent(Vertex v1, Vertex v2);
- origin(Edge e);
- destination(Edge e);

# Graph Implementation: Edge List



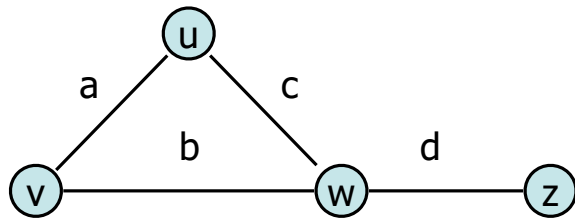
**insertVertex(K key);**

**removeVertex(Vertex v);**

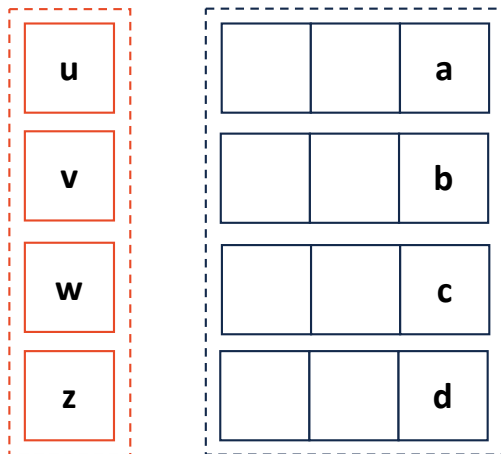
**areAdjacent(Vertex v1, Vertex v2);**

**incidentEdges(Vertex v);**

# Graph Implementation: Adjacency Matrix



**insertVertex(K key);**  
**removeVertex(Vertex v);**  
**areAdjacent(Vertex v1, Vertex v2);**  
**incidentEdges(Vertex v);**



	u	v	w	z
u				
v				
w				
z				