Grammar Trees

lan Ludden

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• Given a grammar *G*, give examples of trees that are/aren't generated by *G* and determine whether a given tree could be generated by *G*.

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- Build a tree matching grammar *G* with a specific terminal string *s*. When multiple trees are possible, build more than one or describe the set of possible trees.

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Parse Trees, a.k.a. Grammar Trees

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Definition

A *parse tree* is a visualization of the generation of a string by a context-free grammar.

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Example 1: Binary Strings

Start symbol is *S*, terminals are 0, 1, and ε , rules are:

 $S \rightarrow 0S$ $S \rightarrow 1S$ $S \rightarrow \varepsilon$

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Example 2: Simple Arithmetic Expressions

Start symbols are *E* and *V* (also non-terminals), terminals are *a*, *b*, +, and \times , rules are:

$$E \rightarrow E + V \mid E \times V \mid V + V \mid V \times V$$

 $V \rightarrow a \mid b$

Example 3: Where are the implicit parentheses?

$$E \to E + V \mid E \times V \mid \mathbf{V} \times \mathbf{E} \mid V + V \mid V \times V$$
$$V \to a \mid b$$

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Designing Parse Trees, and Impossible Strings

$$S \rightarrow AB \mid C$$
$$A \rightarrow Aa \mid a$$
$$B \rightarrow Bb \mid b$$
$$C \rightarrow C$$

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Can *G* generate the following strings? If so, give a parse tree. If not, explain why not.

aabca

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- aabca
- b

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