

Big-O

Part c: (Dis)Proving Big-O Relationships

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Learning Objectives

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- Given functions f and g , (dis)prove that f is $O(g)$ and/or $\Theta(g)$.

Example 0: Warm-up

Prove $17n$ is $O(n^2)$.

$\exists c, k \in \mathbb{R}^+$
Fix some particular values

$\forall n \geq k$
Prove universal claim

$|f(n)| \leq c |g(n)|$
 $0 \leq f(n) \leq cg(n)$

Sketch:

$$17n \leq cn^2 \quad \forall n \geq k$$

$$17 \leq cn$$

$$c = 1$$

$$k = 20$$

$$17n \leq 17n^2$$
$$= c \cdot n^2$$

$$c = 17$$
$$k = 1$$

Proof. Fix $c = 17$ and $k = 1$.
Let $n \geq k$ be arbitrary.

Then,

$$17n \leq 17n^2 \quad (\text{since } n \geq k=1)$$
$$= c \cdot n^2$$

So $17n$ is $O(n^2)$. \square

Example 0: Warm-up

Prove $17n$ is $O(n^2)$.

Example 1: The bigger, the better

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Prove $450n^2 + 25n + 2$ is $O(n^3)$.

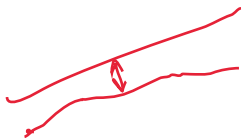
Scratch:

$$450n^2 + 25n + 2 \leq c \cdot n^3 \quad \forall n \geq k$$

$$450n^3 + 25n^3 + 2n^3 \leq c \cdot n^3$$

$$c = 500$$

$$k = 1$$



Proof: Fix $c=500$, $k=1$.

Let $n \geq k$ be arb.

Then,

$$450n^2 + 25n + 2 \leq 450n^3 + 25n^3 + 2n^3$$

(since $n \geq k=1$)

$$= 477n^3$$

$$\leq 500n^3$$

$$= c \cdot n^3$$

So $450n^2 + 25n + 2$ is $O(n^3)$.

□

Example 2: I've heard it both ways

$$\left[\begin{array}{l} f \text{ is } O(g) \\ g \text{ is } O(f) \end{array} \right. \text{ and} \Rightarrow \begin{array}{l} f \text{ is } \Theta(g) \\ \text{and} \\ g \text{ is } \Theta(f). \end{array}$$

Example 2: I've heard it both ways

Prove $\overbrace{2n \log n + 3n}^f$ is $\Theta(\overbrace{n \log n}^g)$.

Scratch. $2n \log n + 3n$ is $O(n \log n)$

$$2n \log n + 3n \log n = 5n \log n$$

$\underbrace{\hspace{10em}}_{k=2} \quad \uparrow_{c=5}$

Scratch. $n \log n$ is $O(2n \log n + 3n)$
 $c=1, k=1$

Pf. Fix $c=5, k=2$.

Let $n \geq k$ be arb.

Then,

$$2n \log n + 3n \leq 2n \log n + 3n \log n$$

(since $n \geq k=2$)

$$= 5n \log n$$

$$= c \cdot n \cdot \log n. \quad \square$$

Pf. Fix $c=1, k=1$.

Let $n \geq k$ be arb.

Then,

$$n \log n \leq 2n \log n$$

$$\leq 2n \log n + 3n$$

$$= c \cdot (2n \log n + 3n). \quad \square$$

Example 3: Trading places

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Prove $|\sin(n\pi/2)|$ is **not** $O(|\cos(n\pi/2)|)$.

Proof. Let $c, k \in \mathbb{R}^+$ be arb.

Construct $n = 2m+1$, where $m \in \mathbb{Z}$
and $m \geq k$.

Then, n is odd \rightarrow n is odd.

$$|\sin(\frac{n\pi}{2})| = 1 > 0 = c \cdot |\cos(\frac{n\pi}{2})|.$$

So $|\sin(\frac{n\pi}{2})|$ is not
 $O(|\cos(\frac{n\pi}{2})|)$. \square

$$\exists c, k \in \mathbb{R}^+ \quad \forall n \geq k \quad |f(n)| \leq c|g(n)|$$
$$\forall c, k \in \mathbb{R}^+ \quad \exists n \geq k \quad |f(n)| > c|g(n)|$$

Scratch.

Need $n \geq k$, and n odd.

$n = 2m+1$, where
 $m \in \mathbb{Z}$ and $m \geq k$.

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