

Introduction to Trees

Ian Ludden

Learning Objectives

By the end of this lesson, you will be able to:

Learning Objectives

By the end of this lesson, you will be able to:

- Define and use tree terminology.

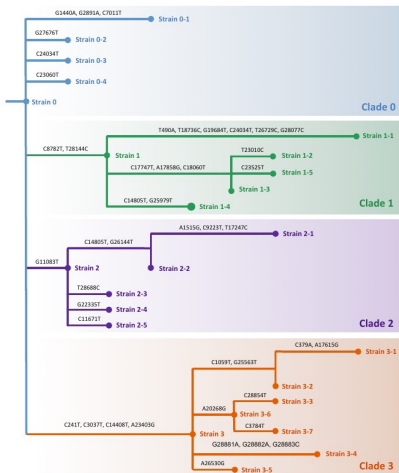
Learning Objectives

By the end of this lesson, you will be able to:

- Define and use tree terminology.
- Define and identify various tree properties.

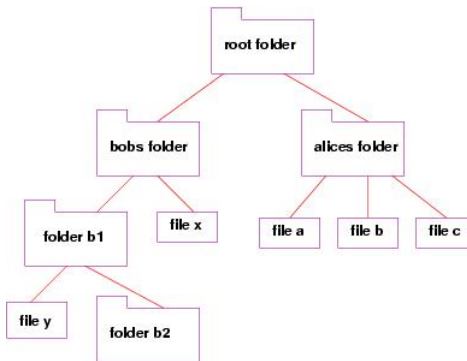
Why do we care about trees?

Why do we care about trees?



A phylogenetic tree (Source)

Why do we care about trees?



A file tree (Source)

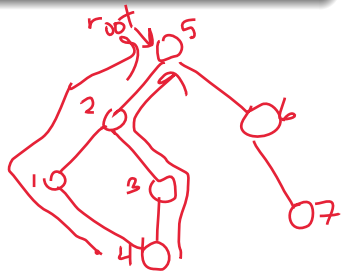
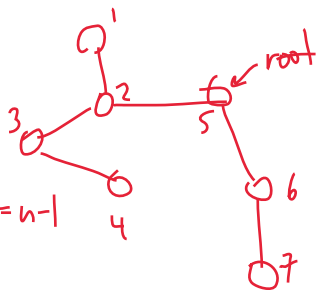
What is a tree?

Definition

A **tree** is a connected acyclic graph. A rooted tree has a special vertex called a **root**.

$T = (V, E)$
↑
nodes

$|V| = n \rightarrow |E| = n - 1$



What is a tree?

Definition

A **tree** is a connected acyclic graph. A **rooted** tree has a special vertex called a **root**.

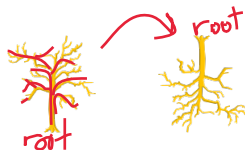
- Root goes at the top by convention

What is a tree?

Definition

A **tree** is a connected acyclic graph. A **rooted** tree has a special vertex called a **root**.

- Root goes at the top by convention
- Terms borrowed from biological trees and family trees

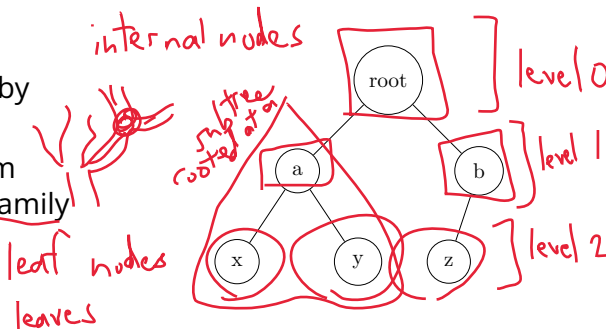


What is a tree?

Definition

A **tree** is a connected acyclic graph. A **rooted** tree has a special vertex called a **root**.

- Root goes at the top by convention
- Terms borrowed from biological trees and family trees



parent/child/sibling
ancestor/descendant

height := max level 0
"proper"

Definition

An **m -ary tree** is a tree in which each node has at most m children.

$m=2$: binary tree
 $m=3$: ternary tree

Definition

An **m -ary tree** is a tree in which each node has at most m children.

Some special cases (shown for $m = 2$):

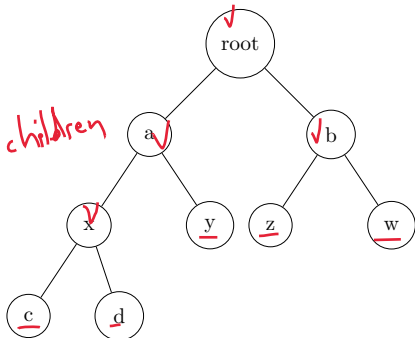
Definition

An **m -ary tree** is a tree in which each node has at most m children.

Some special cases (shown for $m = 2$):

- **full** m -ary tree

every node:
0 or m children



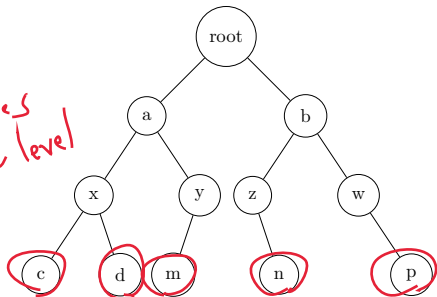
Definition

An **m -ary tree** is a tree in which each node has at most m children.

Some special cases (shown for $m = 2$):

- **full** m -ary tree
- **complete** m -ary tree

*all leaves
at same level*

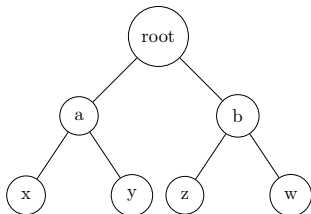


Definition

An **m -ary tree** is a tree in which each node has at most m children.

Some special cases (shown for $m = 2$):

- **full** m -ary tree
- **complete** m -ary tree
- full *and* complete m -ary tree



Counting Nodes

Counting Nodes

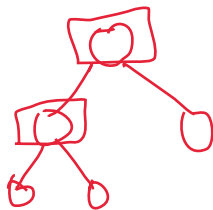
Fact

A full m -ary tree with i internal nodes has $mi + 1$ nodes total.

Proof: Ask everyone how many kids they have (then add the root).

Internal nodes: m each
Leaf nodes: 0 each

$$mi + 1$$



Counting Nodes

Fact

A full m -ary tree with i internal nodes has $mi + 1$ nodes total.

Proof: Ask everyone how many kids they have (then add the root).

Fact

A binary tree of height h has at least $h + 1$ and at most $2^{h+1} - 1$ nodes.

Proof: Consider a path of length h and a full, complete binary tree of height h .

Handwritten mathematical derivations and diagrams illustrating the proof for counting nodes in a binary tree.

Left side (General m -ary tree):

$$\sum_{k=0}^h 2^k = 2^{h+1} - 1$$
$$\sum_{k=0}^h m^k = \frac{m^{h+1} - 1}{m - 1}$$

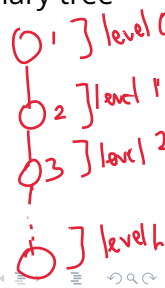
Bottom left (Base 10 example):

10
100
1000
10000
+ 1 = 100000

Center (Binary tree diagram):

```
graph TD
    n0(( )) --- n1L(( ))
    n0 --- n1R(( ))
    n1L --- n2LL(( ))
    n1L --- n2LR(( ))
    n1R --- n2RL(( ))
    n1R --- n2RR(( ))
    n2LL --- n3LLL(( ))
    n2LL --- n3LLR(( ))
    n2LR --- n3LRL(( ))
    n2LR --- n3LRR(( ))
    n2RL --- n3RLL(( ))
    n2RL --- n3RLR(( ))
    n2RR --- n3RRL(( ))
    n2RR --- n3RRR(( ))
```

level	# nodes
0	1
1	2
2	4
3	8
k	$2^k = m^{h+1}$



Counting Nodes

Fact

A full m -ary tree with i internal nodes has $mi + 1$ nodes total.

Proof: Ask everyone how many kids they have (then add the root).

Fact

A binary tree of height h has at least $h + 1$ and at most $2^{h+1} - 1$ nodes.

Proof: Consider a path of length h and a full, complete binary tree of height h .

Fact

The height of a full and complete binary tree with n nodes is proportional to $\log_2 n$.

$$n = 2^{h+1} - 1$$

$$\log_2(n+1) = h+1$$

$$h = \log_2(n+1) - 1 \approx \log_2(n)$$



$[a, b, c, d, e, f]$ $O(n)$

Recap: Learning Objectives

By the end of this lesson, you will be able to:

- Define and use tree terminology.
- Define and identify various tree properties.